

Date

```
public class LL {  
    private Node head;  
    private Node tail;  
    private int size;  
    public LL(){  
        this.size = 0;  
    }  
    // Insert at First  
    public void insertFirst(int val){  
        Node node = new Node(val);  
        node.next = head;  
        head = node;  
        if(tail == null){  
            tail = head;  
        }  
        size++;  
    }  
}
```

```
// Insert at Last  
public void insertLast(int val){  
    Node node = new Node(val);  
    tail.next = node;  
    tail = node;  
    size++;  
    if(tail == null){  
        insertFirst(val);  
    }  
    return;  
}
```

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//Display

```
public void display(){
    Node temp = head;
    while (temp != null) {
        System.out.print (temp.val + " → ");
        temp = temp.next;
    }
    System.out.println("END");
}
```

//Insert at any specific index

```
public void insert (int val, int index) {
```

```
if (index == 0) {
    insertFirst (val);
    return;
}
```

```
if (index == size) {
    insertLast (val);
    return;
}
```

//Otherwise check from head to our index

Node temp = head; //0 itself the loop from its index

//add element at index 3 means iterates from
1 to 2

for (int i=1 ; i< index ; i++) {

```
    temp = temp.next;
}
```

Node node = new Node (int val, int index);
temp.next = node; //created node

& size++;

}

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delete first ele. from linked list.

```
public int deleteFirst() {  
    int val = head.value;  
    head = head.next; // i.e., if list has only 1 element  
    if (head == null) { // then it deleted after head is  
        tail = null; // null because no value is there  
    } // at next.  
    size--;  
    return val;  
}
```

list: {3} → delete first

{3} → next = null

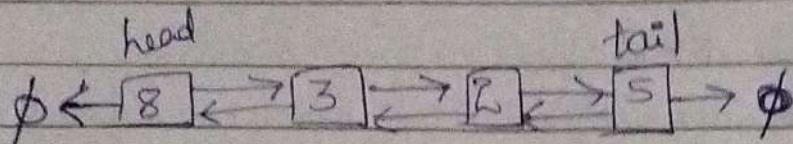
/* by head = tail = null */

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Doubly linked list :- When we traverse back it is useful than S.L.L.

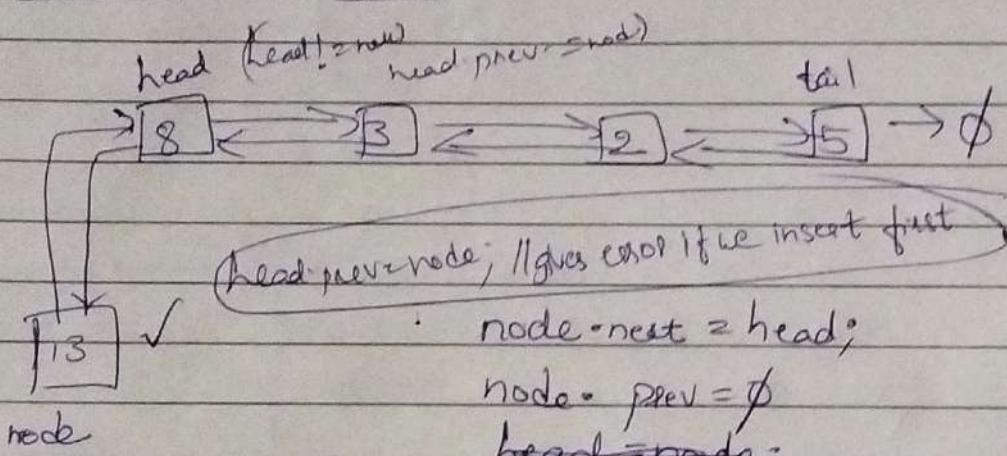


→ Every single box having the structure :-

```

class Node {
    int val;
    Node next;
}
  
```

① insert 13 at first then



Node node = new Node(val);

node.next = head;

node.prev = ϕ; (null)

if (head != null) {

head.prev = node;

}

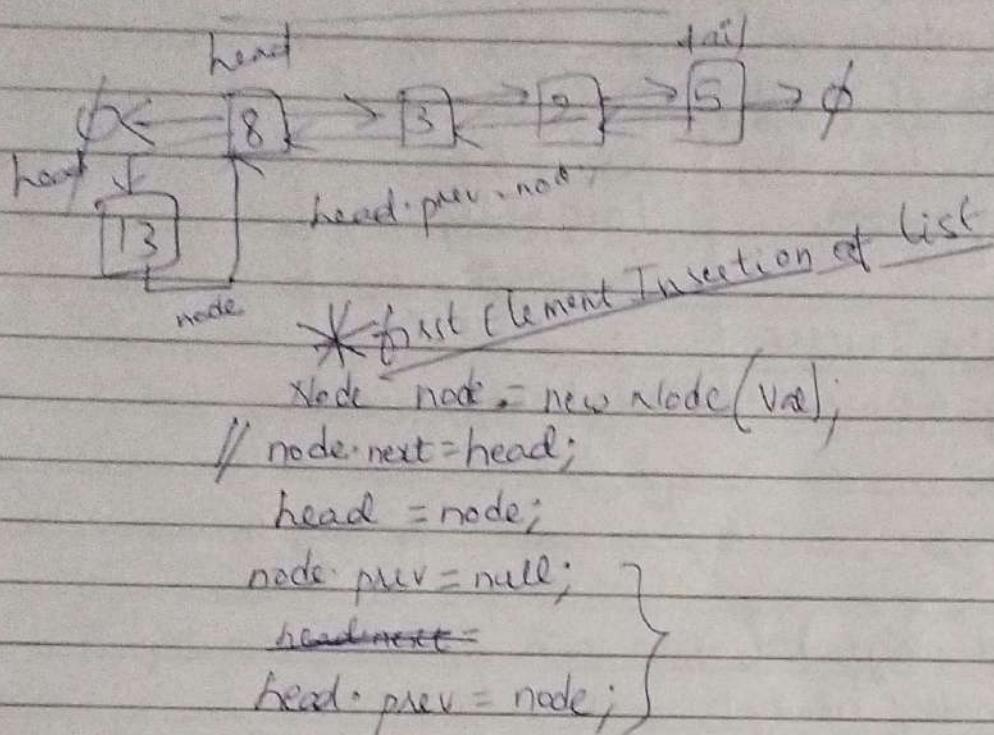
head = node;

y

Here we check for null pointer exception

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Doubly linked list (Backward traversal)



node.next = node;
node.prev = \emptyset ; } causes \emptyset to be assigned
head = node;
head.prev = node;

{ interview
level
questions }

→ causes Null Pointer Exception when checks it.

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// display() Method

```

Node node = head; lastNode (last = null);
while (node != null) {
    cout (node.val + " → ");
    node = node.next;
    last = node;
}
cout ("END");
last = null;
last = node;

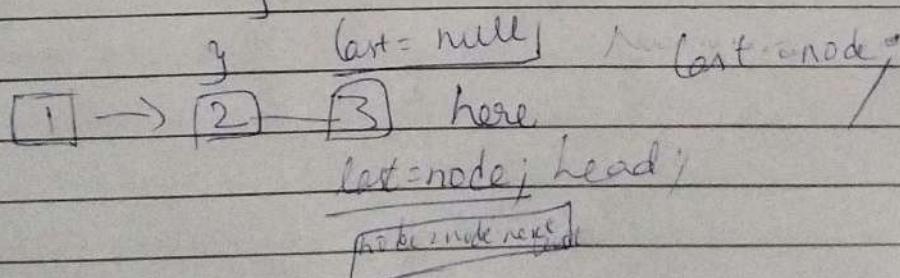
```

also display in Reverse→ firstly we need tail node here;int val

```

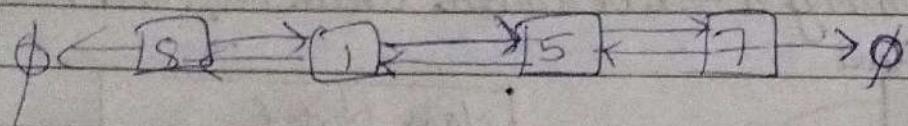
public display() {
    Node node = head;
    Node last = null;
    while (node != null) {
        cout (node.val + " → ");
        node = node.next; // defining lastNode
        last = node; node = node.next;
    }
    cout ("END");
    cout ("In Reverse");
    while (last != null) {
        cout (last.val + " → ");
        last = last.prev;
    }
}

```



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Insert element at last



node = last.next,

node.next = ϕ ;

node.prev = last;

If ($\text{head} == \phi$) {

 head.prev = ϕ ;

 head = node;

}

* for given order in the list :-
traverse order is :-

7 \rightarrow 5 \rightarrow 1 \rightarrow 8 \rightarrow START

* Insertion at middle of the element is also just ok
as same as we did in the Single linked lists.

// Method to find out the value & particularly return that
node then

```
public Node find() {
    Node node = head;
    while (node != null) {
        if (node.val == value) {
            return node;
        }
        node = node.next;
    }
    return null;
}
```

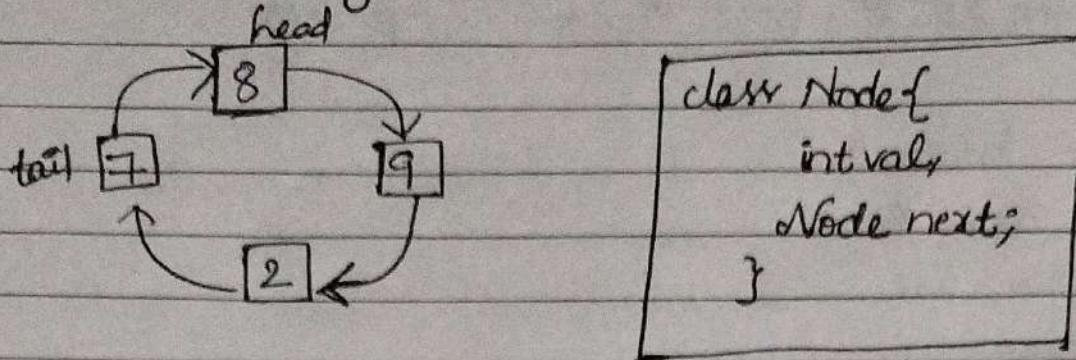
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* Circular linked list :- (only 2 operations :- deletion & Insertion)
 A linked list in a chain.

// here creating a box is called Node i.e., class box,



```

class Node{
    int val;
    Node next;
}
  
```

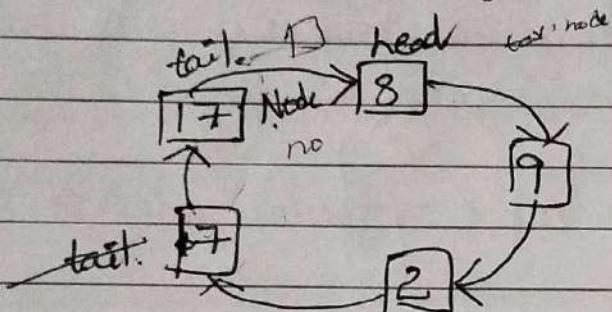
// Just like in the Single linked list.

// Here no things are null. no one is pointing to the null until the list is empty

Insertion :-

i) After the tail :-

insert 17 after tail.



Assumptions :-

// Create a new node first.

$\text{tail} \cdot \text{next} = \text{node}$

$\text{tail} = \text{node}$.

$\text{node} \cdot \text{next} = \text{head};$

if ($\text{head} == \text{null}$) {

$\text{tail} = \text{node};$

$\text{head} = \text{node};$

}

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CLL of

```
private class Node {  
    int val;  
    Node next;  
    Node (int val) {  
        this.val = val;  
    }  
    private Node head;  
    private Node tail;  
    CLL () {  
        this.head = null;  
        this.tail = null;  
    }  
    public void insert (int val) {  
        Node node = new Node (val);  
        if (head == null) {  
            head = node;  
            tail = node;  
            return;  
        }  
        tail.next = node;  
        node.next = head;  
        tail = node;  
    }  
    public void display () {  
        Node node = head;  
        while (node != head) {  
            node = node.next;  
        }  
    }  
}
```

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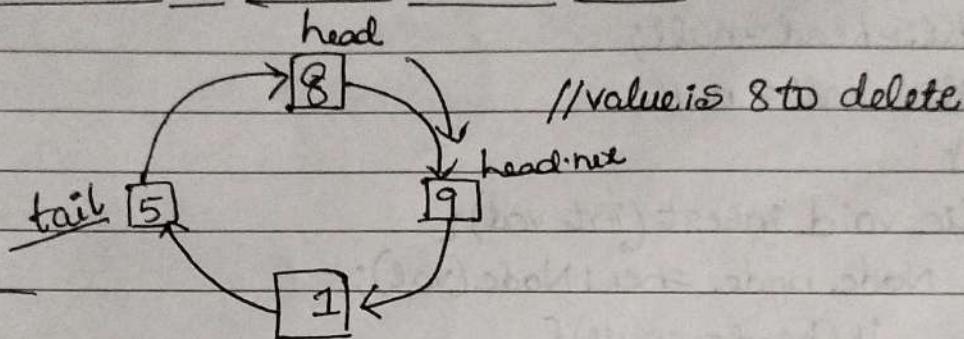
also we can do it by do-while

```

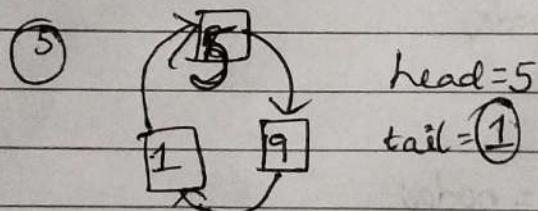
if (head != null) {
    do {
        System.out.println (node.val + " → ");
        node = node.next;
    } while (node != head);
    System.out.println ("HEAD");
}

```

* Deletion in circular linked list

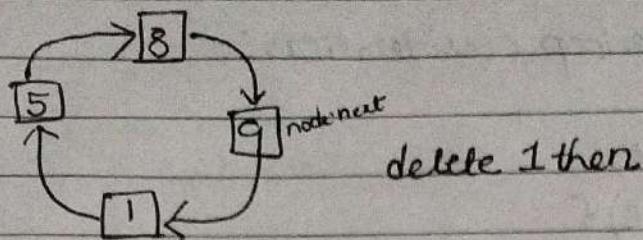


// here you can't reach it to 5 because it is the deleted



* $g = \text{node}$; if we want to delete 9 then
then node starts from 1 to the value we want to
delete.

for example. $\text{node.next} == \text{val}$ (we want to delete)
 $\text{node.next} = n.next$.



① Consider node = 8.

node.next = 9

$(\text{node.next} == \text{val}) X$ then go 2nd pass

② $\text{node} = \text{node.next} \Rightarrow 9$

then $\text{node.next} = 1$

So, here $\underline{\text{node.next} == \text{val}}$

$1 == 1$

then update the criteria to the CL list.
i.e.,

$\text{node} = \text{next} = \text{node.next.next};$

③ Here we are working with 2 pointers.

```
public void delete (int val) {
```

```
    Node node = head;
```

```
    if (node == null) {
```

```
        return;
```

```
}
```

```
    if (node.val == val) {
```

```
        head = head.next;
```

```
        tail = next = head;
```

```
} return;
```

```
}
```

//otherwise I need to check the Item.

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// By using Do-while loop consideration :-

do {

 Node n = node.next;

 if (n.val == val) {

 node.next = n.next;

 break;

}

 node = node.next;

} while (node != head);

// we need to do this until node \neq head.

10 min

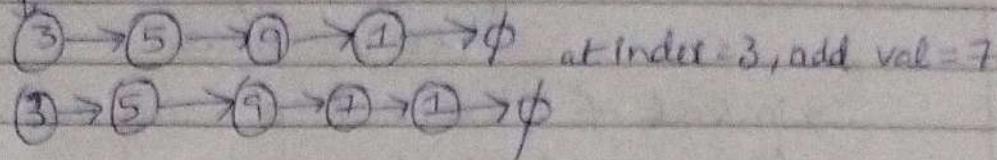


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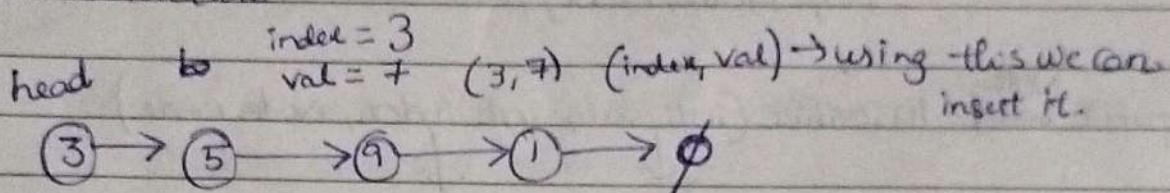
* Recursive Insertion in linked list

Ex:- head



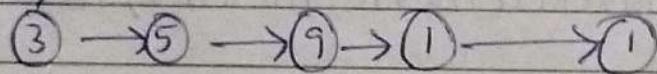
2 things to do in recursive linkedlist:

- 1) void return type & changes made in linked list.
- 2) node return type that return the list node to change the structure.



pass-1:

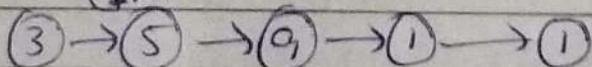
$(3, 3) X$



// then iterate

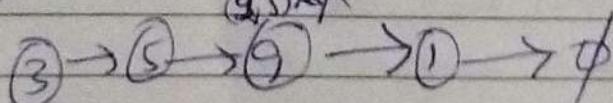
pass-2:

$(4, 5) X$



pass-3:

$(5, 7) X$



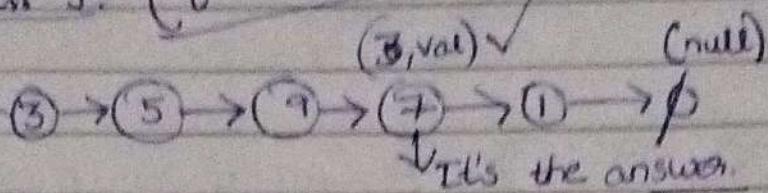
// then here at this index add new Node. & then traverse the operations here.

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//Stack Overflow condition:-
(Sum > (max_val - last) / 10) {
 return 0;
}

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part-3 :- (final result)



→ instead of 7, we need to write the value we want to insert at this place.

Things to do :-

- 1) Visualize what should we do :-
- 2) Try go to get what nodes cause the connections appropriately. (like exactly or without misinterpret)
i.e., node, next, prev etc.

public ^{Node} void insertRec(int val, int index, Node node)

↓ ↓ ↓
what val at ? index reference
is to be add? we need to add?

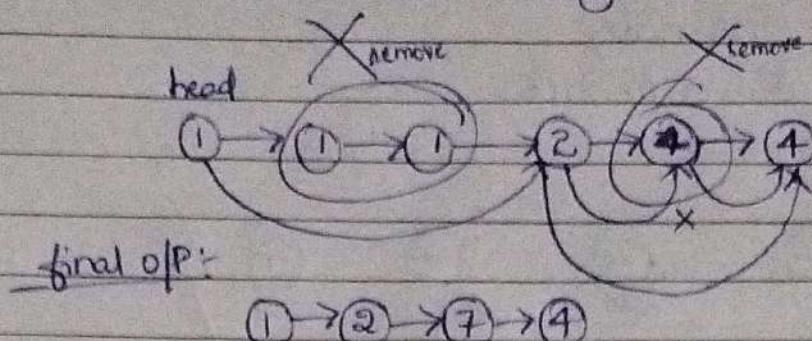
public Node insertRec(int val, int index, Node node) {

Sonal



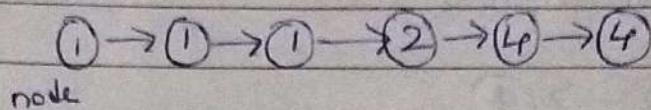
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Remove duplicates from single linked list :-



pass - 1 :-

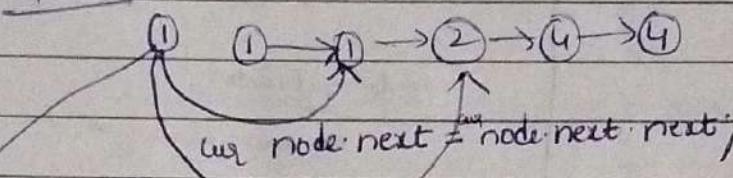
head



node = node.next

skip (or) delete that element.

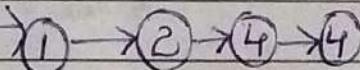
pass - 2 :-



pass - 3 is do check ($node.next == node$)

then also skip it (or) delete here the node.

pass - 4 :-



$curNode.next = curNode.next.next$;

// then $i = i + 1$

then add / move ahead.

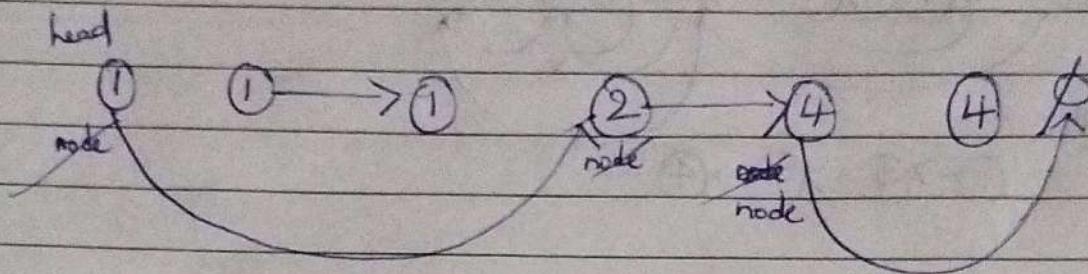
$curNode.next = curNode$; {No it's not}

// then simply $node = node.next$.

pass 5 :- Again $node.next = node$; No (it's not) then move forward.

part - 6 :-

`node.next = curr.next;`
then move forward to null.



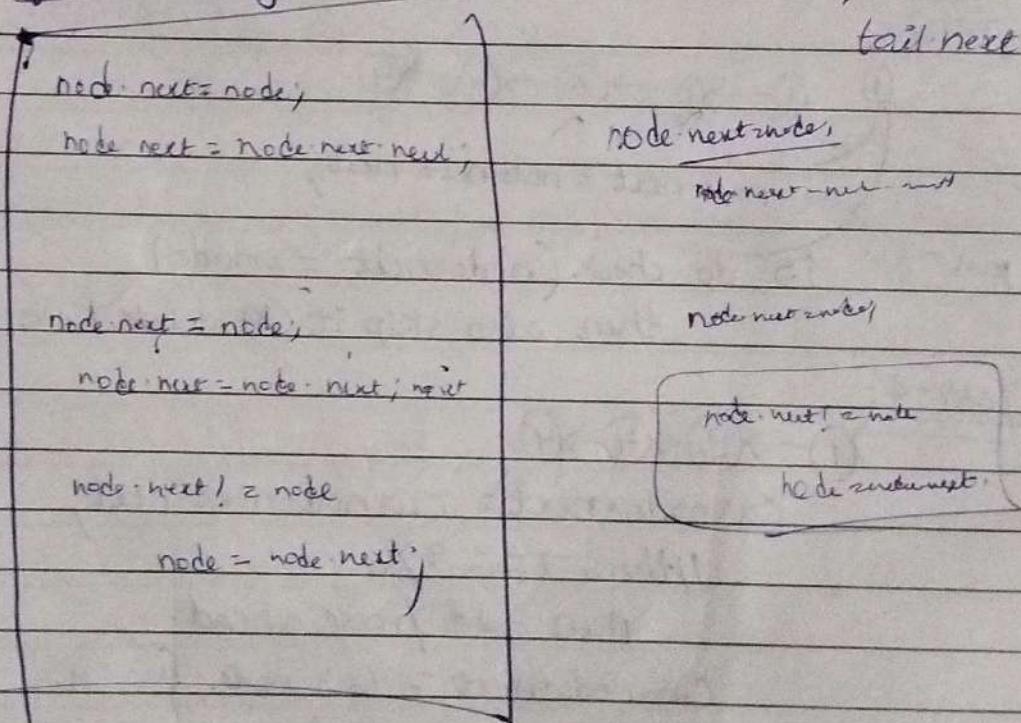
then final answer is :-

$$1 \rightarrow 2 \rightarrow 4$$

Here node starts from head & ends at tail:-

functionality :-

i.e., $\text{tail} = \text{node}$,
 $\text{tail.next} = \emptyset$



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```
public void duplicates() {  
    Node node = head;  
    while (node.next != null) {  
        if (node.val == node.next.value) {  
            node.next = node.next.next;  
            size--;  
        } else {  
            node = node.next;  
        }  
    }  
    tail = node;  
    tail.next = null;  
}
```

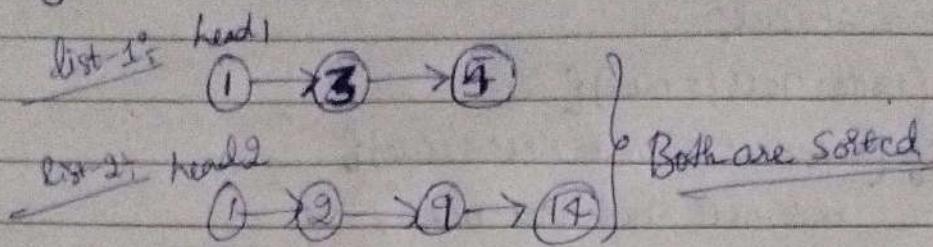
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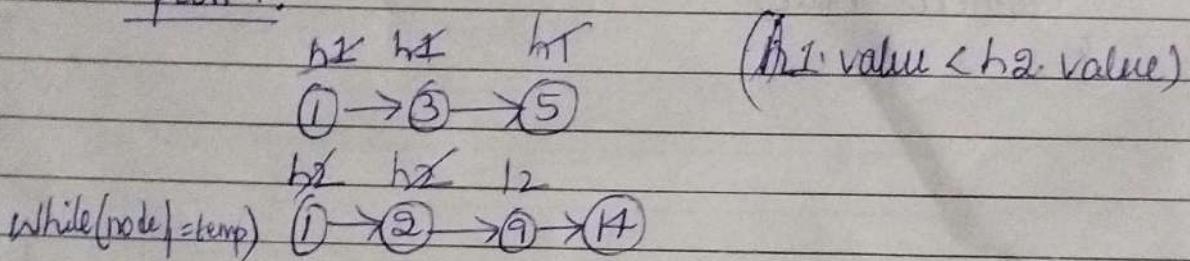
* Merge two sorted lists :-



* final Result: $[1 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 5 \rightarrow 7 \rightarrow 14]$

1/ By comparing smaller elements from both 2 lists & add it to the new lists

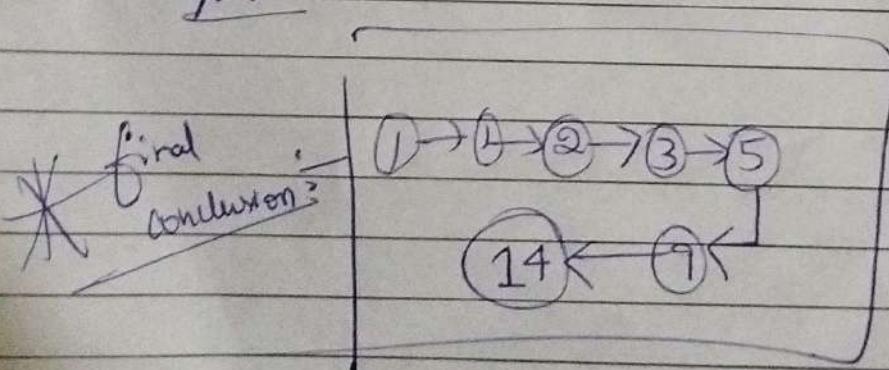
pass-1:-



$h_1 < h_2$ (any 1)

① $h_2 < h_2$ (h_2++)

pass-2:-



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// Merge 2 sorted lists :-

```
listNode dummy = new ListNode(0);
listNode current = dummy;
while (l1 != null && l2 != null) {
    if (l1.val < l2.val) {
        current.next = l1.next;
        l1 = l1.next;
    } else {
        current.next = l2.next;
        l2 = l2.next;
    }
    current = current.next;
}
if (l1 == null) current.next = l2;
if (l2 == null) current.next = l1;
// current.next = (l1 != null) ? l1 : l2;
return dummy.next;
```

listNode dummy = new ListNode(0);

listNode current = dummy;

while (l1 != null && l2 != null) {

if (l1.val < l2.val) {

current.next = l1;

l1 = l1.next;

} else

current.next = l2;

l2 = l2.next;

} current = current.next;

current.next = l1;

if (l1 != null) {

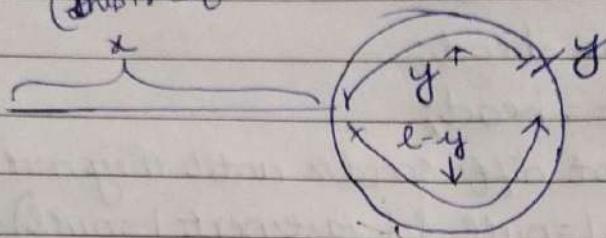
listNode l1 = new Node();
l1 = ...
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Floyd-warshall algorithm (Hare & Tortoise) for cycle detection
at which position in list:
(this is before loop begins)



l = length of the cycle.

* Pointers Distance before meeting the loop :-

$$\text{slow} = x + n_1 \cdot l + y$$

n_1 means no of times by slow pointer rotates the loop

$$\text{fast} = x + n_2 \cdot l + y \quad (\text{same as } n_1)$$

here $\text{fast} = 2 * \text{slow}$ (according to speed)

$$x + n_2 \cdot l + y = 2x + 2 \cdot n_2 \cdot l + y$$

$$n_2 \cdot l = x + y + 2n_2 \cdot l$$

$$x + y = (n_1 \cdot l - 2 \cdot n_2 \cdot l)$$

$$x + y = K \cdot l$$

if $K = 1$ then

$$x + y = l$$

$$\boxed{x = l - y}$$

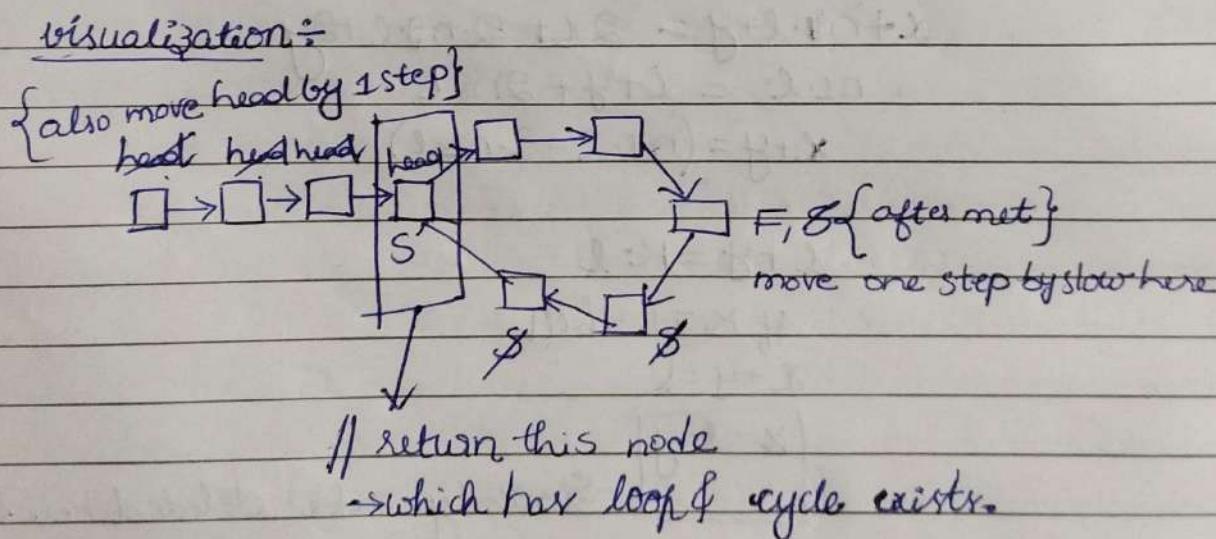
Similarly index detected here. Btw,

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```

public ListNode detectCycle(ListNode head) {
    // start both from head
    ListNode slow = head;
    ListNode fast = head;
    // advance both at diff. speeds until they meet at once.
    while (fast != null && fast.next != null) {
        fast = fast.next.next;
        slow = slow.next;
        if (slow == fast) {
            while (head != slow) {
                head = head.next;
                slow = slow.next;
            }
            return slow;
        }
    }
    return null; // If no ans found in the loop.
}

```



* Happy Number

$$\begin{aligned}
 36 &= 6^2 + 3^2 = 36 + 24 \\
 &= 60 \\
 6^2 + 0 &= 36 \quad \left. \right\} \text{loop continues} \\
 &\text{// return false.}
 \end{aligned}$$

Efficient Approach

- Consider a hashset of all integer values which is been squared & added to the defined variable.
- If again it returns the value, that exactly presents in the hashset then it returns the false value, otherwise it returns true.

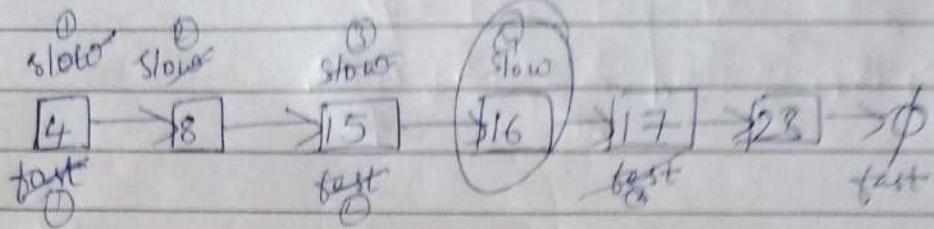
Dry-Run of Code

```

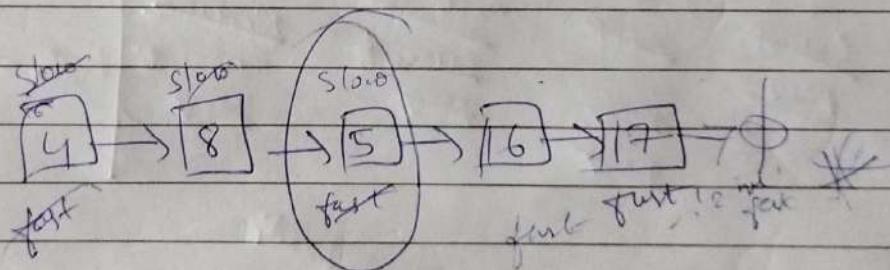
boolean isHappy(int n) {
    Set<Integer> usedIntegers = new HashSet<>();
    while(true) {
        // Find sum of squares
        int sum = 0;
        while(n != 0) {
            sum += Math.pow(n % 10, 2);
            n = n / 10;
        }
        // If sum is 1, then return true;
        if (sum == 1) return true;
        // else the current sum is the new number.
        n = sum;
        // check if we have already encountered that number.
        if (usedIntegers.contains(n))
            return false;
        usedIntegers.add(n);
    }
}
  
```

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Middle of a linked list (876)

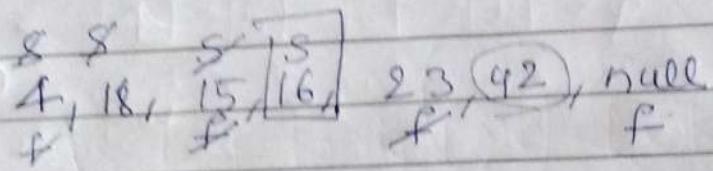


```
ListNode middleNode(ListNode head) {  
    ListNode slowPtr = head;  
    ListNode fastPtr = head;  
    //Travel until the fast pointer reaches until lastnode or null.  
    while(fast != null & & fast.next != null){  
        slowPtr = slowPtr.next;  
        fastPtr = fastPtr.next.next;  
    }  
    return slowPtr;  
}
```



sl. 4, 8, 8, 5, 15
h 6, 5, 16, 17, phi

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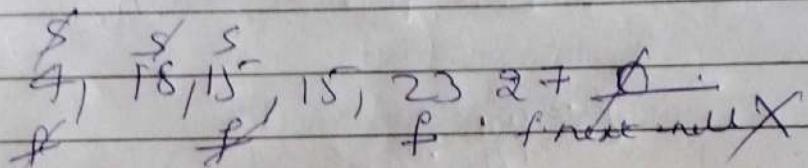


while ($\text{fast} \neq \text{null}$ & $\text{fast} \rightarrow \text{next} \neq \text{null}$) {

23 42

$\text{fast} = \text{fast} \rightarrow \text{next}$; next

$23 \rightarrow \text{null}$; // Kopi vorbere.



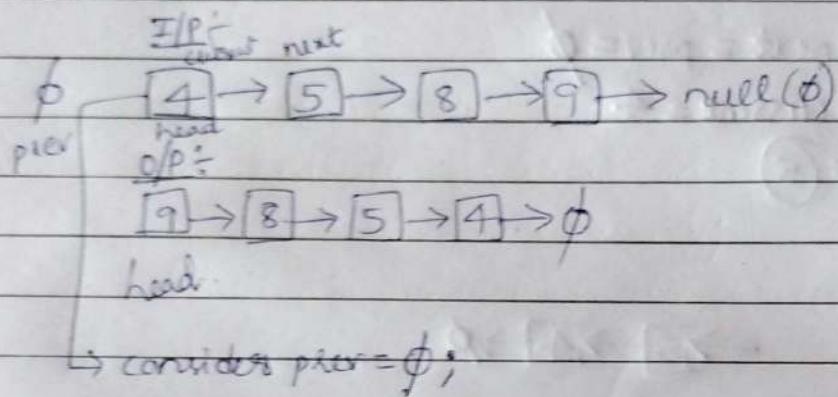
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Reverse a linked list

→ By using Stack Method but not efficient algorithm

Efficient Approach :-

→ using 3 pointers method.

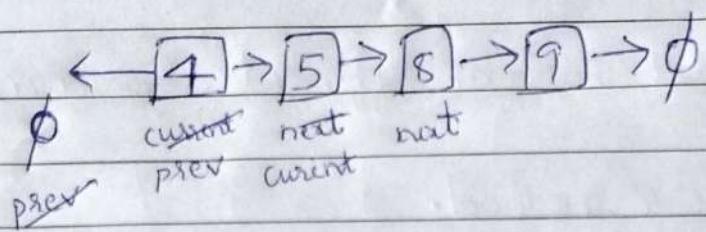


```
current = head;  
while (head != null){  
    next = current.next;  
    current.next = prev;  
    prev = current;  
    current = next;  
}  
head = prev;  
return head;  
Y
```

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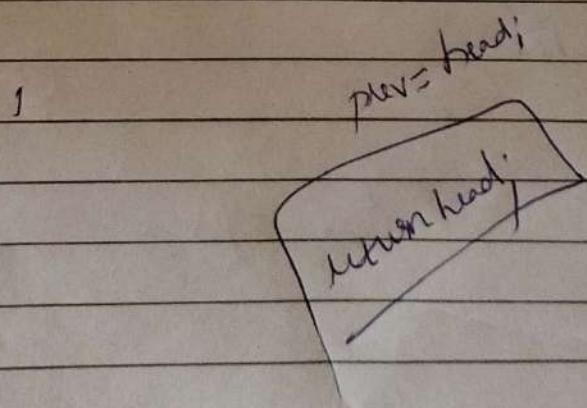
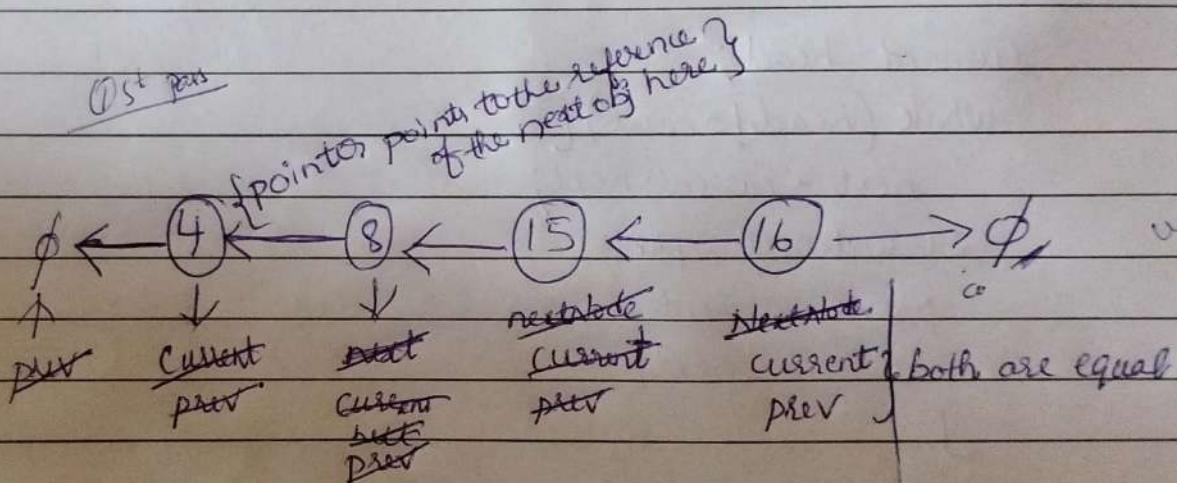
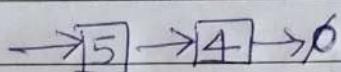
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$$\text{current.next} = \text{prev} = \phi$$

$$\text{prev} = \text{current} = 4$$

$$\text{current} = \text{next} = 5$$



Reverse of a Linkedlist II :-

ListNode reverseBetween(ListNode head, int left, int right) {

 ListNode dummy = new ListNode(0);
 dummy.next = head;

 ListNode letPre = dummy;

 ListNode curNode = head;

 for (int i=0; i<left; i++) {

 letPre = letPre.next;

 curNode = curNode.next;

}

 ListSublistHead = curNode;

 ListNode prev = null;

 for (int i=0; i<=(right-left); i++) {

 ListNode nextNode = curNode.next;

 curNode.next = prev;

 prev = curNode;

 curNode = nextNode;

}

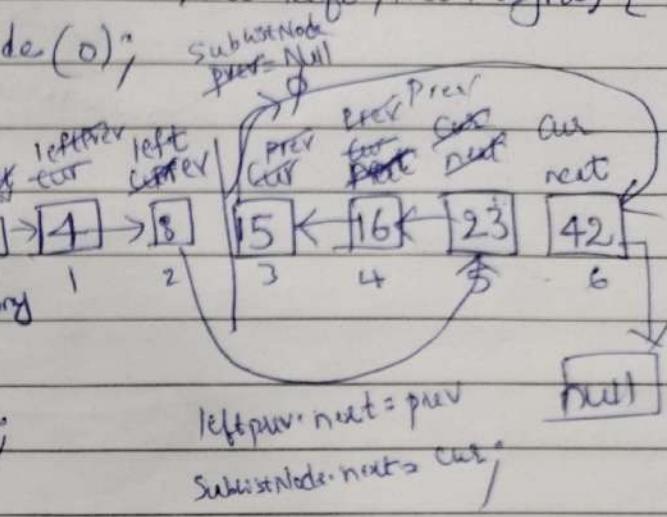
//join the pieces.

 leftPre.next = prev;

 SublistHead.next = curNode;

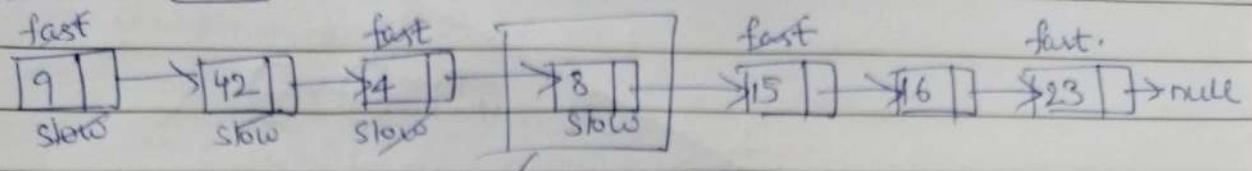
 return dummy.next;

}

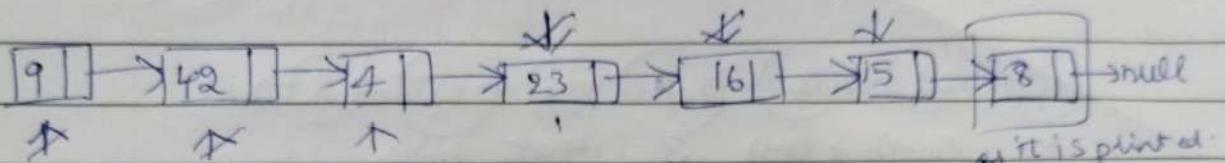


Date

Reorder list



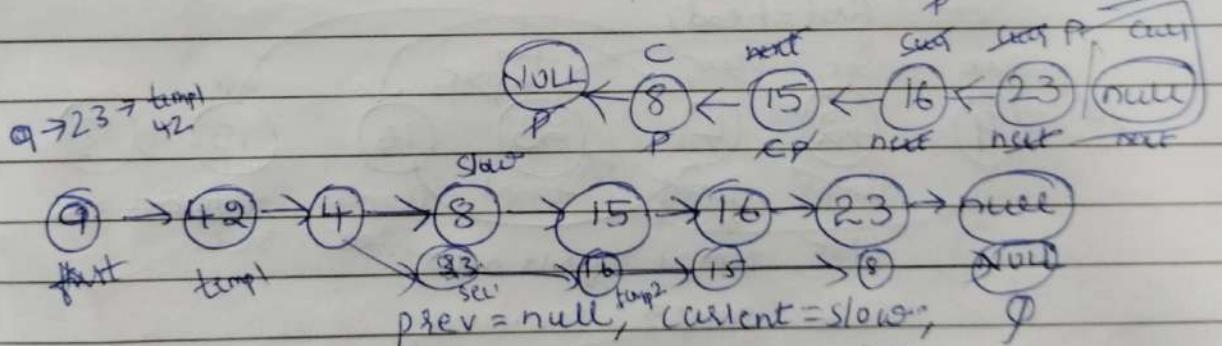
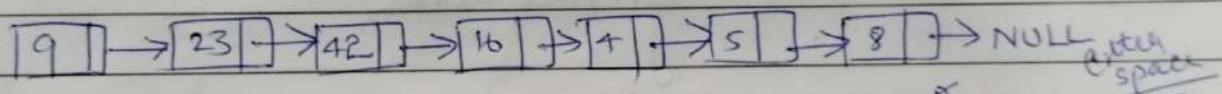
↙ // from here to last position reverse the list



→ Considers 2 pointers from here.

* 1st one at head and 2nd one at from the starting position of the reversed list.

print both the values then,



firsttemp1 = head;

sec temp2 = p->next;

while (sec->next != null) {

temp1 = first->next;

temp2 = sec->next;

first = temp1;

sec = temp2;

}

while (curr->next != null) {

next = curr->next;

curr->next = prev;

prev = curr;

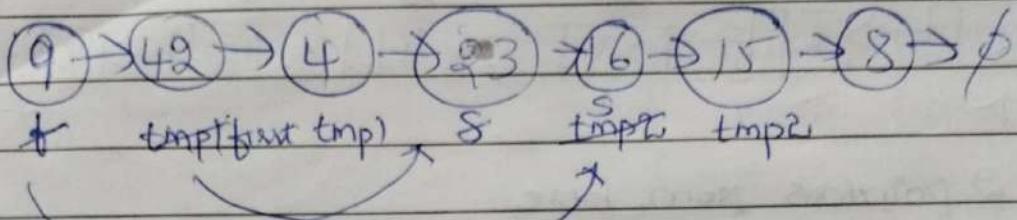
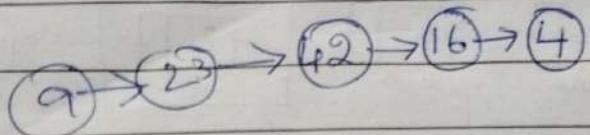
curr = next;

}

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while ($s \neq \text{null}$) {

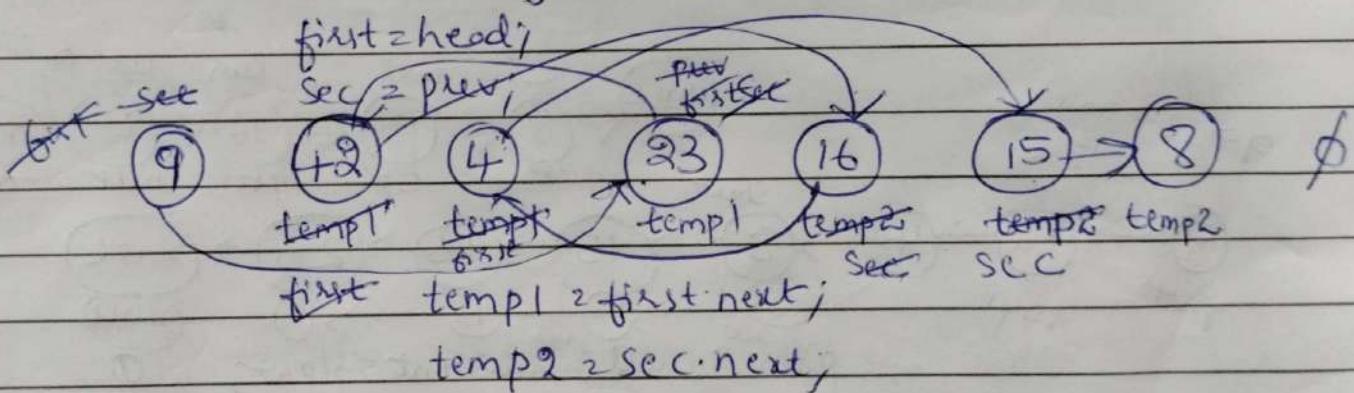
$\text{temp1} = \text{first}.\text{next};$

$\text{temp2} = \text{sec}.\text{next};$

$\text{first}.\text{next} = \text{S};$

$\text{sec}.\text{next} = \text{temp1};$

}



$\text{first}.\text{next} = \text{Sec};$

$\text{sec}.\text{next} = \text{temp2};$

$\text{first} = \text{temp1};$

$\text{sec} = \text{temp2};$

}

9, 12, 4, 23, 16, 15

Date

first = head;

sec = prev;

while (sec != null) {

temp1 = first.next;

temp2 = sec.next; first.next = s;

first = temp1;

s.next = t1;

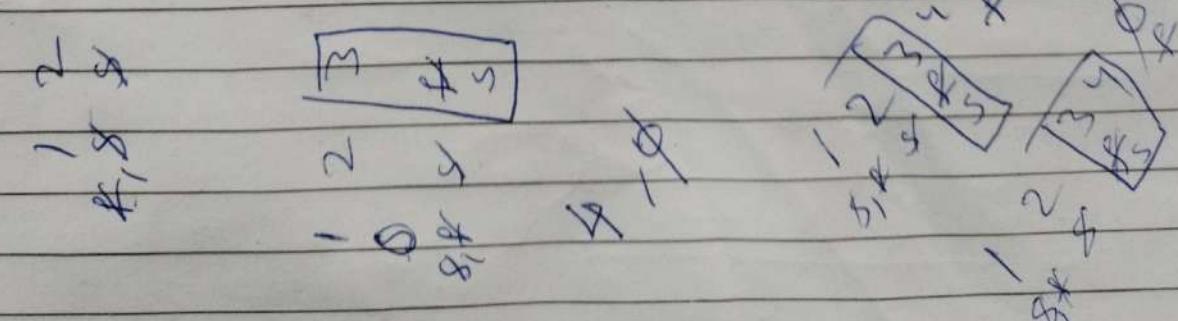
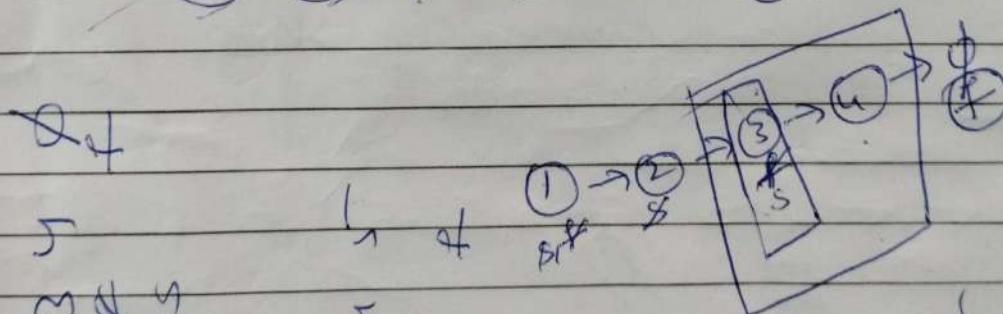
sec = temp2;

y

9 42 23 16
f t1 8 t2
t f

9 42 4 23 16 15 8
t t1 t2 8 8 t2
f t t1 t2 8 8 s

⑨ → ②3 → ④2 → ⑯ → ④ → ⑯ → ②3 →



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CP example-1:-

* Reverse K No. of nodes in a list

I/P :- 1 → 2 → 2 → 4 → 5 → 6 → 7 → 8
 k = 4 [reverse (group) reverse]
O/P :- 4 → 2 → 2 → 1 → 8 → 7 → 6 → 5

I/P :- 1 → 2 → 3 → 4 → 5
 k = 3 :- [reverse] [reverse]
 ↓ ↓
O/P :- 3 → 2 → 1 → 5 → 4

Assumptions :-

* Needn't to print anything / read anything only return the head.

Time complexity :- $O(N)$

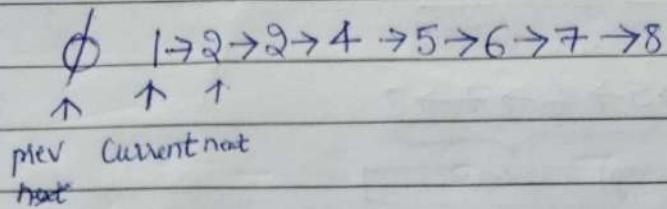
Space complexity :- $O(1)$

$$\left. \begin{array}{l} 1 < N < 10^5 \\ 1 < k < N \end{array} \right\} \text{constraints}$$

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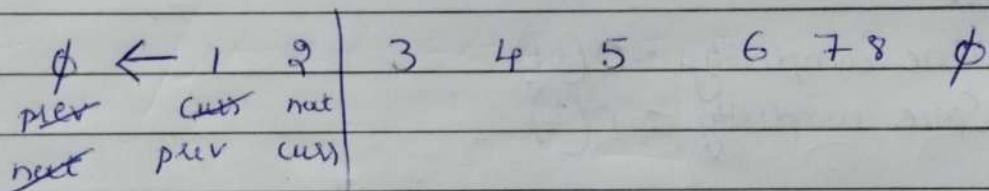
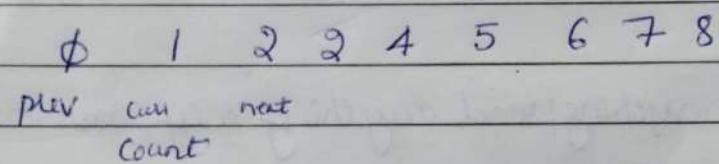
Date

if (`head==null`) {
 return;



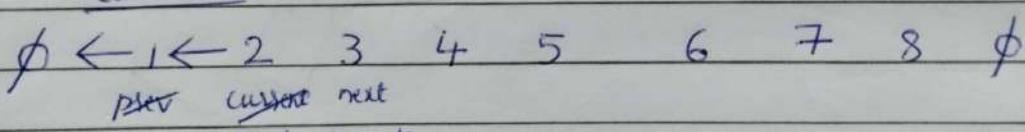
`count = 0;`

pass - 1:



pass - 2:

`count = 1;`



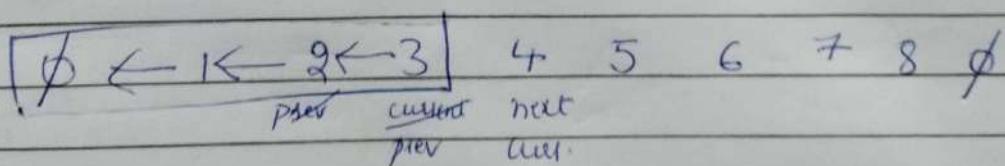
prev current

$3 \rightarrow 2 \rightarrow 1 \rightarrow (\textcircled{1}) \rightarrow (\textcircled{2}) \rightarrow (\textcircled{3}) \rightarrow (\textcircled{4}) \rightarrow 8 \rightarrow \emptyset$

(head \leftarrow curr)

pass - 3:

count = 2



At pass 4: count = 3 then we stop reversing the list.

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1st group: 3 → 2 → 1

Critical check:

head points to 4

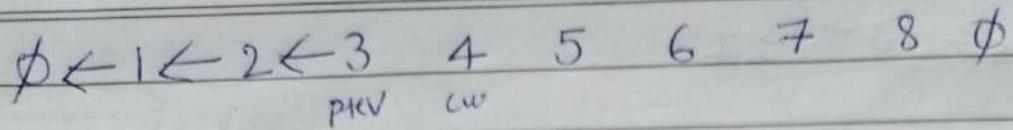
if (next != null) {

Code Snippet :-

```

if (head == null) {
    static Node reverse (Node head, int k) {
        if (head == null) {
            return null;
        }
        Node cur = head; count = 0;
        Node prev = null;
        while (count < k && cur->next != null) {
            Node next = cur->next;
            cur->next = prev;
            prev = cur;
            cur = next;
            count++;
        }
        if (next != null) {
            head->next = reverse(next, k);
        }
    }
    return prev; // after reversing it is the head of the
} // note exactly.

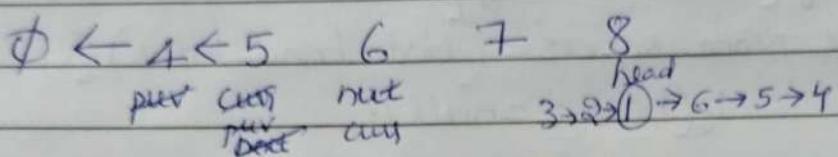
```



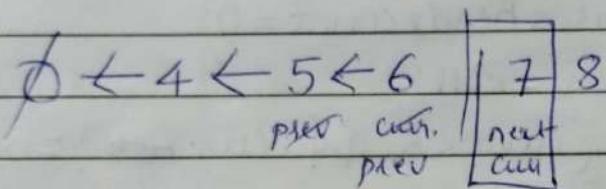
solution 2:
part-1:
 $\emptyset \leftarrow 4 \quad 5 \quad 6 \quad 7 \quad 8$

prev cur
next next
head cur

part-2:



part-3:

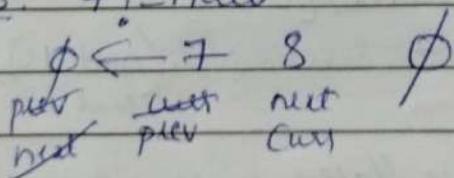


// these Second group is reversed i.e., $4 \rightarrow 5$

$$\boxed{6} \rightarrow 5 \rightarrow 4$$

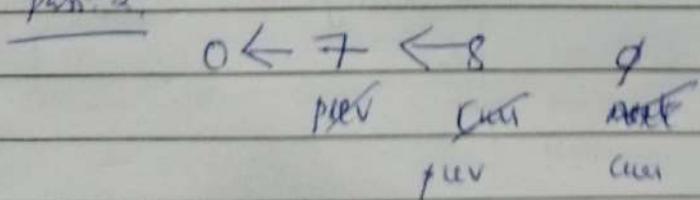
prev

solution-3: $7 = \text{null}$



$3 \rightarrow 2 \rightarrow 1 \rightarrow 6 \rightarrow 5 \rightarrow 4 \rightarrow 8 \rightarrow 7$

part-2:



$8 \rightarrow 7$

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① if (`head == null || k == 1`) {
 return `head`;
}

→ `head` is empty, nothing to reverse here so reverse \emptyset (`null`)

→ also $k = 1$ (reversing 1 node at a time, no need to run a universal logic).

→ // for better efficiency & correctness.
↓

It's important because:

→ prevents nullPointerException, unnecessary computation

 saves time when $k = 1$, need not to traversal.

② `// If K no. of nodes do not exist then return head;`
`ListNode check = head;` otherwise reverse it and
 `int count = 0;` unleash the prev.

`while (check != null) {`
 `check = check.next;`
 `count++;`

}

`if (count < k) return head;`

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Final code:

```

if (head == null || k == 1) {
    return head;
}

// check K nodes exist or not
ListNode check = head, int count = 0;
while (check != null & count < k) {
    check = check.next;
    count++;
}

if (count < k) return head;

// reverse k no. of nodes.
ListNode prev = null;
ListNode next = null, current = head;
while (current != null) {
    next = current.next;
    current.next = prev;
    prev = current;
    current = next;
}

if (next != null) {
    head.next = reverseKGroup(next, k);
}

return prev;
}

```

Date

Rotated Linked List :-

head

$$1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 6 \rightarrow \emptyset$$

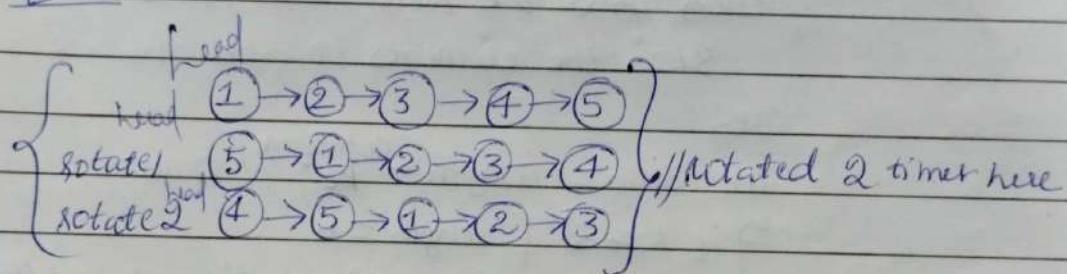
→ rotate 1 time :-

$$2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 6 \rightarrow 1$$

→ rotate 2nd time

$$3 \rightarrow 4 \rightarrow 5 \rightarrow 6 \rightarrow 1 \rightarrow 2 \rightarrow \emptyset$$

Ex:-



$$2 \times 6$$

$$\begin{array}{r} 6 \\ 2 \end{array}$$

$$\begin{array}{r} 12 \\ 12 \end{array} \overline{) 6} \quad \begin{array}{r} 6 \\ 0 \end{array}$$

$$\begin{array}{r} 6 \\ 6 \\ 6 \\ 0 \end{array}$$

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if ($K \leq 0$) || head == null || head.next == null)

{ return head; }

ListNode last = head;

int length = 1;

while (last.next != null) {

last = last.next;

length++;

}

last.next = head;

Now need to make no. of rotations

int rotations = $K \% \text{length}$;

// then find the new node here.

Skip the elements to find the new node

i.e.,

int skip = length - $\frac{rotations}{K}$; // times

\Rightarrow if $K=2$ & $\text{length}=6$

then skip $6-2=4$ times to find the
newhead of the node.

ListNode newlast = head;
// now land at the 4th element Node here.

for (int i=0; i< skip-1; i++) {

newlast = newlast.next;

}

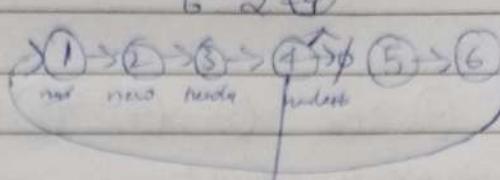
head = newlast.next;

newlast.next = null;

return head;

}

0, 1, 2, 3



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