

Template del Rejunte

```

1 #include <bits/stdc++.h>
2 #define sqr(a) ((a)*(a))
3 #define rsz resize
4 #define forr(i,a,b) for(int i=(a);i<(b);i++)
5 #define forn(i,n) forr(i,0,n)
6 #define dforn(i,n) for(int i=n-1;i>=0;i--)
7 #define forall(it,v) for(auto it=v.begin();it!=v.end();it++)
8 #define foreach(i, v) for(auto i:v)
9 #define sz(c) ((int)c.size())
10 #define zero(v) memset(v, 0, sizeof(v))
11 #define pb push_back
12 #define mp make_pair
13 #define lb lower_bound
14 #define ub upper_bound
15 #define fst first
16 #define snd second
17 #define PI 3.1415926535897932384626
18
19 using namespace std;
20
21 typedef long long ll;
22 typedef pair<int,int> ii;
23 typedef vector<int> vi;
24 typedef vector<ii> vii;
25
26 int main()
27 {
28     // agregar g++ -DREJUNTE en compilacin
29     #ifdef REJUNTE
30         freopen("input", "r", stdin);
31         // freopen("output","w", stdout);
32     #endif
33     ios::sync_with_stdio(false);
34     cin.tie(NULL);
35     cout.tie(NULL);
36     return 0;
37 }
```

Estructuras de datos

Set Mejorado

Esto solo compila en C++11.

```

1 #include <ext/pb_ds/assoc_container.hpp>
```

```

2 #include <ext/pb_ds/tree_policy.hpp>
3 using namespace __gnu_pbds;
4 //<key,mapped type,comparator,...>
5 typedef tree<int,null_type,less<int>,rb_tree_tag,
6     tree_order_statistics_node_update> ordered_set;
7 //find_by_order(i) devuelve iterador al i-esimo elemento
8 //order_of_key(k): devuelve la pos del lower bound de k
9 //Ej: 12, 100, 505, 1000, 10000.
10 //order_of_key(10) == 0, order_of_key(100) == 1,
11 //order_of_key(707) == 3, order_of_key(9999999) == 5
```

Union Find

```

1 struct UnionFind{
2     vector<int> f,setSize; //the array f contains the parent of each node
3     int cantSets;
4     void init(int n)
5     {
6         f.clear(); setSize.clear();
7         cantSets=n;
8         f.rsz(n,-1);
9         setSize.rsz(n,1);
10    }
11    int comp(int x){return (f[x]==-1?x:f[x]=comp(f[x]));} //O(1)
12    bool join(int i,int j) //devuelve true si ya estaban juntos
13    {
14        bool con=comp(i)==comp(j);
15        if(!con)
16        {
17            cantSets--;
18            setSize[comp(j)]+=setSize[comp(i)];
19            setSize[comp(i)]=setSize[comp(j)]; //no suma, solo asigna
20            f[comp(i)]=comp(j);
21        }
22        return con;
23    }
24 };
```

Hash Table

```

1 //Compilar: g++ --std=c++11
2 struct Hash{
3     size_t operator()(const ii &a)const
4     {
5         size_t s=hash<int>()(a.fst);
6         return hash<int>()(a.snd)+0x9e3779b9+(s<<6)+(s>>2);
7     }
8     size_t operator()(const vector<int> &v)const
```

```

9   {
10      size_t s=0;
11      for(auto &e : v) s ^= hash<int>()(e)+0x9e3779b9+(s<<6)+(s>>2);
12      return s;
13   }
14 };
15 unordered_set<ii, Hash> s;
16 unordered_map<ii, int, Hash> m; //map<key, value, hasher>

```

RMQ

RMQ (static)

Dado un arreglo y una operacion asociativa *idempotente*, $get(i, j)$ opera sobre el rango $[i, j]$.
 Restriccion: $LVL \geq \text{ceil}(\log n)$; Usar $[]$ para llenar arreglo y luego $build()$.

```

1 struct RMQ{
2     #define LVL 10
3     tipo vec[LVL][1<<(LVL+1)];
4     tipo &operator[](int p){return vec[0][p];}
5     tipo get(int i, int j) { //intervalo [i,j]
6         int p = 31-__builtin_clz(j-i);
7         return min(vec[p][i], vec[p][j-(1<<p)]);
8     }
9     void build(int n) { //O(n log n)
10        int mp = 31-__builtin_clz(n);
11        forn(p, mp) forn(x, n-(1<<p))
12            vec[p+1][x] = min(vec[p][x], vec[p][x+(1<<p)]);
13    };

```

RMQ (dynamic)

```

1 //Dado un arreglo y una operacion asociativa con neutro, get(i, j) opera sobre
  //el rango [i, j].
2 #define MAXN 100000
3 #define operacion(x, y) max(x, y)
4 const int neutro=0;
5 struct RMQ{
6     int sz;
7     tipo t[4*MAXN];
8     tipo &operator[](int p){return t[sz+p];}
9     void init(int n){ //O(n log n)
10        sz = 1 << (32-__builtin_clz(n));
11        forn(i, 2*sz) t[i]=neutro;
12    }
13    void updall(){ //O(n)
14        dforn(i, sz) t[i]=operacion(t[2*i], t[2*i+1]);}
15    tipo get(int i, int j){return get(i,j,1,0,sz);}

```

```

16    tipo get(int i, int j, int n, int a, int b){ //O(lgn)
17        if(j<=a || i>=b) return neutro;
18        if(i<=a && b<=j) return t[n];
19        int c=(a+b)/2;
20        return operacion(get(i, j, 2*n, a, c), get(i, j, 2*n+1, c, b));
21    }
22    void set(int p, tipo val){ //O(lgn)
23        for(p+=sz; p>0 && t[p]!=val;){
24            t[p]=val;
25            p/=2;
26            val=operacion(t[p*2], t[p*2+1]);
27        }
28    }
29 }rmq;
30 //Usage:
31 cin >> n; rmq.init(n); forn(i, n) cin >> rmq[i]; rmq.updall();

```

RMQ (lazy)

```

1 //Dado un arreglo y una operacion asociativa con neutro, get(i, j) opera sobre
  //el rango [i, j].
2 typedef int Elem; //Elem de los elementos del arreglo
3 typedef int Alt; //Elem de la alteracion
4 #define operacion(x,y) x+y
5 const Elem neutro=0; const Alt neutro2=0;
6 #define MAXN 100000 //Cambiar segun el N del problema
7 struct RMQ{
8     int sz;
9     Elem t[4*MAXN];
10    Alt dirty[4*MAXN]; //las alteraciones pueden ser de distinto Elem
11    Elem &operator[](int p){return t[sz+p];}
12    void init(int n){ //O(n log n)
13        sz = 1 << (32-__builtin_clz(n));
14        forn(i, 2*sz) t[i]=neutro;
15        forn(i, 2*sz) dirty[i]=neutro2;
16    }
17    void push(int n, int a, int b){ //propaga el dirty a sus hijos
18        if(dirty[n]!=0){
19            t[n]+=dirty[n]*(b-a); //altera el nodo
20            if(n<sz){
21                dirty[2*n]+=dirty[n];
22                dirty[2*n+1]+=dirty[n];
23            }
24            dirty[n]=0;
25        }
26    }
27    Elem get(int i, int j, int n, int a, int b){ //O(lgn)

```

```

28     if(j<=a || i>=b) return neutro;
29     push(n, a, b);//corrige el valor antes de usarlo
30     if(i<=a && b<=j) return t[n];
31     int c=(a+b)/2;
32     return operacion(get(i, j, 2*n, a, c), get(i, j, 2*n+1, c, b));
33 }
34 Elem get(int i, int j){return get(i,j,1,0,sz);}
35 //altera los valores en [i, j) con una alteracion de val
36 void alterar(Alt val, int i, int j, int n, int a, int b){//0(lgn)
37     push(n, a, b);
38     if(j<=a || i>=b) return;
39     if(i<=a && b<=j){
40         dirty[n]+=val;
41         push(n, a, b);
42         return;
43     }
44     int c=(a+b)/2;
45     alterar(val, i, j, 2*n, a, c), alterar(val, i, j, 2*n+1, c, b);
46     t[n]=operacion(t[2*n], t[2*n+1]);//por esto es el push de arriba
47 }
48 void alterar(Alt val, int i, int j){alterar(val,i,j,1,0,sz);}
49 }rmq;

```

RMQ (persistente)

```

1  typedef int tipo;
2  tipo oper(const tipo &a, const tipo &b){
3      return a+b;
4  }
5  struct node{
6      tipo v; node *l,*r;
7      node(tipo v):v(v), l(NULL), r(NULL) {}
8      node(node *l, node *r) : l(l), r(r){
9          if(!l) v=r->v;
10         else if(!r) v=l->v;
11         else v=oper(l->v, r->v);
12     }
13 };
14 node *build (tipo *a, int tl, int tr) //modificar para que tome tipo a
15     if (tl+1==tr) return new node(a[tl]);
16     int tm=(tl + tr)>>1;
17     return new node(build(a, tl, tm), build(a, tm, tr));
18 }
19 node *update(int pos, int new_val, node *t, int tl, int tr){
20     if (tl+1==tr) return new node(new_val);
21     int tm=(tl+tr)>>1;
22     if(pos < tm) return new node(update(pos, new_val, t->l, tl, tm), t->r);

```

```

23     else return new node(t->l, update(pos, new_val, t->r, tm, tr));
24 }
25 tipo get(int l, int r, node *t, int tl, int tr){
26     if(l==tl && tr==r) return t->v;
27     int tm=(tl + tr)>>1;
28     if(r<=tm) return get(l, r, t->l, tl, tm);
29     else if(l>=tm) return get(l, r, t->r, tm, tr);
30     return oper(get(l, tm, t->l, tl, tm), get(tm, r, t->r, tm, tr));
31 }

```

BIGInt

```

1  #define BASEXP 6
2  #define BASE 1000000
3  #define LMAX 1000
4  struct bint{
5      int l;
6      ll n[LMAX];
7      bint(ll x=0){
8          l=1;
9          forn(i, LMAX){
10             if (x) l=i+1;
11             n[i]=x%BASE;
12             x/=BASE;
13         }
14     }
15 }
16 bint(string x){
17     l=(x.size()-1)/BASEXP+1;
18     fill(n, n+LMAX, 0);
19     ll r=1;
20     forn(i, sz(x)){
21         n[i / BASEXP] += r * (x[x.size()-1-i]-'0');
22         r*=10; if(r==BASE)r=1;
23     }
24 }
25 void out(){
26     cout << n[l-1];
27     dforn(i, l-1) printf("%6.6llu", n[i]);//6=BASEXP!
28 }
29 void invar(){
30     fill(n+1, n+LMAX, 0);
31     while(l>1 && !n[l-1]) l--;
32 }
33 };
34 bint operator+(const bint&a, const bint&b){
35     bint c;

```

```

36     c.l = max(a.l, b.l);
37     ll q = 0;
38     forn(i, c.l) q += a.n[i]+b.n[i], c.n[i]=q %BASE, q/=BASE;
39     if(q) c.n[c.l++] = q;
40     c.invar();
41     return c;
42 }
43 pair<bint, bool> lresta(const bint& a, const bint& b)    // c = a - b
44 {
45     bint c;
46     c.l = max(a.l, b.l);
47     ll q = 0;
48     forn(i, c.l) q += a.n[i]-b.n[i], c.n[i]=(q+BASE) %BASE, q=(q+BASE)/BASE-1;
49     c.invar();
50     return make_pair(c, !q);
51 }
52 bint& operator-= (bint& a, const bint& b){return a=lresta(a, b).first;}
53 bint operator- (const bint&a, const bint&b){return lresta(a, b).first;}
54 bool operator< (const bint&a, const bint&b){return !lresta(a, b).second;}
55 bool operator<= (const bint&a, const bint&b){return lresta(b, a).second;}
56 bool operator==(const bint&a, const bint&b){return a <= b && b <= a;}
57 bint operator*(const bint&a, ll b){
58     bint c;
59     ll q = 0;
60     forn(i, a.l) q += a.n[i]*b, c.n[i] = q %BASE, q/=BASE;
61     c.l = a.l;
62     while(q) c.n[c.l++] = q %BASE, q/=BASE;
63     c.invar();
64     return c;
65 }
66 bint operator*(const bint&a, const bint&b){
67     bint c;
68     c.l = a.l+b.l;
69     fill(c.n, c.n+b.l, 0);
70     forn(i, a.l){
71         ll q = 0;
72         forn(j, b.l) q += a.n[i]*b.n[j]+c.n[i+j], c.n[i+j] = q %BASE, q/=BASE;
73         c.n[i+b.l] = q;
74     }
75     c.invar();
76     return c;
77 }
78 pair<bint, ll> ldiv(const bint& a, ll b){// c = a / b ; rm = a % b
79     bint c;
80     ll rm = 0;
81     dforn(i, a.l){

```

```

82         rm = rm * BASE + a.n[i];
83         c.n[i] = rm / b;
84         rm %= b;
85     }
86     c.l = a.l;
87     c.invar();
88     return make_pair(c, rm);
89 }
90 bint operator/(const bint&a, ll b){return ldiv(a, b).first;}
91 ll operator%(const bint&a, ll b){return ldiv(a, b).second;}
92 pair<bint, bint> ldiv(const bint& a, const bint& b){
93     bint c;
94     bint rm = 0;
95     dforn(i, a.l){
96         if (rm.l==1 && !rm.n[0])
97             rm.n[0] = a.n[i];
98         else{
99             dforn(j, rm.l) rm.n[j+1] = rm.n[j];
100            rm.n[0] = a.n[i];
101            rm.l++;
102        }
103        ll q = rm.n[b.l] * BASE + rm.n[b.l-1];
104        ll u = q / (b.n[b.l-1] + 1);
105        ll v = q / b.n[b.l-1] + 1;
106        while (u < v-1){
107            ll m = (u+v)/2;
108            if (b*m <= rm) u = m;
109            else v = m;
110        }
111        c.n[i]=u;
112        rm-=b*u;
113    }
114    c.l=a.l;
115    c.invar();
116    return make_pair(c, rm);
117 }
118 bint operator/(const bint&a, const bint&b){return ldiv(a, b).first;}
119 bint operator%(const bint&a, const bint&b){return ldiv(a, b).second;}

```

Algoritmos

Longest Increasing Subsequence

```

1 //Para non-increasing, cambiar comparaciones y revisar busq binaria
2 //Given an array, paint it in the least number of colors so that each color
  turns to a non-increasing subsequence.

```

```

3 //Solution:Min number of colors=Length of the longest increasing subsequence
4 int N, a[MAXN]; //secuencia y su longitud
5 ii d[MAXN+1]; //d[i]=ultimo valor de la subsecuencia de tamaño i
6 int p[MAXN]; //padres
7 vector<int> R; //respuesta
8 void rec(int i){
9     if(i== -1) return;
10    R.push_back(a[i]);
11    rec(p[i]);
12 }
13 int lis(){ //O(nlogn)
14     d[0] = ii(-INF, -1); forn(i, N) d[i+1]=ii(INF, -1);
15     forn(i, N){
16         int j = upper_bound(d, d+N+1, ii(a[i], INF))-d;
17         if (d[j-1].first < a[i]&&a[i] < d[j].first){
18             p[i]=d[j-1].second;
19             d[j] = ii(a[i], i);
20         }
21     }
22     R.clear();
23     dforn(i, N+1) if(d[i].first!=INF){
24         rec(d[i].second); //reconstruir
25         reverse(R.begin(), R.end());
26         return i; //longitud
27     }
28     return 0;
29 }

```

Mo's

 $O(q * \sqrt{n})$

```

1 int n,sq;
2 struct Qu{//queries [l, r]
3     //intervalos cerrado abiertos !!! importante!!
4     int l, r, id;
5 }qs[MAXN];
6 int ans[MAXN], curans; //ans[i]=ans to ith query
7 bool bymos(const Qu &a, const Qu &b){
8     if(a.l/sq!=b.l/sq) return a.l<b.l;
9     return (a.l/sq)&1? a.r<b.r : a.r>b.r;
10 }
11 void mos(){
12     forn(i, t) qs[i].id=i;
13     sort(qs, qs+t, bymos);
14     int cl=0, cr=0;
15     sq=sqrt(n);
16     curans=0;

```

```

17     forn(i, t){ //intervalos cerrado abiertos !!! importante!!
18         Qu &q=qs[i];
19         while(cl>q.l) add(--cl);
20         while(cr<q.r) add(cr++);
21         while(cl<q.l) remove(cl++);
22         while(cr>q.r) remove(--cr);
23         ans[q.id]=curans;
24     }
25 }

```

Strings

KMP

```

1 vector<int> b; //back table b[i] maximo borde de [0..i]
2 void kmppre(string &P) //by gabina with love
3 {
4     b.clear();
5     b.rsz(P.size());
6     int i =0, j=-1; b[0]=-1;
7     while(i<sz(P))
8     {
9         while(j>=0 && P[i] != P[j]) j=b[j];
10        i++, j++;
11        b[i] = j;
12    }
13 }
14 void kmp(string &T,string &P) //Text, Pattern -- O(|T|+|P|)
15 {
16     kmppre(P);
17     int i=0, j=0;
18     while(i<sz(T))
19     {
20         while(j>=0 && T[i]!=P[j]) j=b[j];
21         i++, j++;
22         if(j==sz(P))
23         {
24             //P encontrado en T empezando en [i-j,i)
25             j=b[j];
26         }
27     }
28 }

```

Z function

```

1 //z[i]=length of longest substring starting from s[i] that is prefix of s

```

```

2 vector<int> z;
3 void zFunction(string &s)
4 {
5     int n=s.size();
6     for(int i=1,l=0,r=0;i<n;i++)
7     {
8         if(i<=r)
9             z[i]=min(r-i+1,z[i-1]);
10        while(i+z[i]<n && s[z[i]]==s[i+z[i]])
11            z[i]++;
12        if(i+z[i]-1>r)
13            l=i, r=i+z[i]-1;
14    }
15 }
16 void match(string &T,string &P) //Text, Pattern -- O(|T|+|P|)
17 {
18     string s=P;
19     s+='$'; //here append a character that is not present in T
20     s.append(T);
21     z.clear();
22     z.rsz(s.size(),0);
23     zFunction(s);
24     forr(i,P.size()+1,s.size())
25         if(z[i]==P.size()); //match found, idx = i-P.size()-1
26 }

```

Trie

```

1 struct trie{
2     map<char, trie> m;
3     void add(const string &s, int p=0)
4     {
5         if(s[p]) m[s[p]].add(s, p+1);
6     }
7     void dfs()
8     {
9         //Do stuff
10        forall(it, m)
11            it->second.dfs();
12    }
13 };

```

Manacher

```

1 string s;
2 int d1[MAXN]; //d1[i]=long del maximo palindromo impar con centro en i
3 int d2[MAXN]; //d2[i]=analogo pero para longitud par
4 //0 1 2 3 4

```

```

5 //a a b a a <--d1[2]=3
6 //a a a a <--d2[2]=2 (estan uno antes)
7 void manacher() // O(|S|) - find longest palindromic substring
8 {
9     int l=0, r=-1, n=sz(s);
10    forn(i, n)
11    {
12        int k=(i>r? 1 : min(d1[l+r-i], r-i));
13        while(i+k<n && i-k>=0 && s[i+k]==s[i-k]) ++k;
14        d1[i] = k--;
15        if(i+k > r) l=i-k, r=i+k;
16    }
17    l=0, r=-1;
18    forn(i, n)
19    {
20        int k=(i>r? 0 : min(d2[l+r-i+1], r-i+1))+1;
21        while(i+k-1<n && i-k>=0 && s[i+k-1]==s[i-k]) k++;
22        d2[i] = --k;
23        if(i+k-1 > r) l=i-k, r=i+k-1;
24    }
25 }

```

Aho Corasick

```

1 struct Trie{
2     map<char, Trie> next;
3     Trie* tran[256]; //transiciones del automata
4     int idhoja, szhoja; //id de la hoja o 0 si no lo es
5     //link lleva al sufijo mas largo, nxthoja lleva al mas largo pero que es hoja
6     Trie *padre, *link, *nxthoja;
7     char pch; //caracter que conecta con padre
8     //Trie(): tran(), idhoja(), padre(), link() {}
9     //coment linea de arriba porque me daba errores usarla.
10    void insert(const string &s, int id=1, int p=0) //id>0!!!
11    {
12        if(p<sz(s))
13        {
14            Trie &ch=next[s[p]];
15            tran[(int)s[p]]=&ch;
16            ch.padre=this, ch.pch=s[p];
17            ch.insert(s, id, p+1);
18        }
19        else idhoja=id, szhoja=sz(s);
20    }
21    Trie* get_link()
22    {
23        if(!link)

```

```

24     {
25         if(!padre) link=this;//es la raiz
26         else if(!padre->padre) link=padre;//hijo de la raiz
27         else link=padre->get_link()->get_tran(pch);
28     }
29     return link;
30 }
31 Trie* get_tran(int c)
32 {
33     if(!tran[c]) tran[c] = !padre? this : this->get_link()->get_tran(c);
34     return tran[c];
35 }
36 Trie *get_nxthoja()
37 {
38     if(!nxthoja) nxthoja = get_link()->idhoja? link : link->nxthoja;
39     return nxthoja;
40 }
41 void print(int p)
42 {
43     if(idhoja) cout << "found_" << idhoja << "_at_position_" << p-szhoja <<
44         endl;
45     if(get_nxthoja()) get_nxthoja()->print(p);
46 }
47 void matching(const string &s, int p=0) //0(|s| + tamao palabras)
48 {
49     print(p); if(p<sz(s)) get_tran(s[p])->matching(s, p+1);
50 }

```

Geometría

Punto

```

1 struct pto{
2     double x, y;
3     pto(double x=0, double y=0):x(x),y(y){}
4     pto operator+(pto a){return pto(x+a.x, y+a.y);}
5     pto operator-(pto a){return pto(x-a.x, y-a.y);}
6     pto operator+(double a){return pto(x+a, y+a);}
7     pto operator*(double a){return pto(x*a, y*a);}
8     pto operator/(double a){return pto(x/a, y/a);}
9     //dot product, producto interno:
10    double operator*(pto a){return x*a.x+y*a.y;}
11    //module of the cross product or vectorial product:
12    //if a is less than 180 clockwise from b, a^b>0
13    double operator^(pto a){return x*a.y-y*a.x;}

```

```

14    //returns true if this is at the left side of line qr
15    bool left(pto q, pto r){return ((q-*this)^(r-*this))>0;}
16    bool operator<(const pto &a) const{return x<a.x-EPS || (abs(x-a.x)<EPS && y<a
17        .y-EPS);}
18    bool operator==(pto a){return abs(x-a.x)<EPS && abs(y-a.y)<EPS;}
19    double norm(){return sqrt(x*x+y*y);}
20    double norm_sq(){return x*x+y*y;}
21 };
22 double dist(pto a, pto b){return (b-a).norm();}
23 typedef pto vec;
24 double angle(pto a, pto o, pto b){
25     pto oa=a-o, ob=b-o;
26     return atan2(oa^ob, oa*ob);}
27
28 //rotate p by theta rads CCW w.r.t. origin (0,0)
29 pto rotate(pto p, double theta){
30     return pto(p.x*cos(theta)-p.y*sin(theta),
31         p.x*sin(theta)+p.y*cos(theta));
32 }

```

Orden Radial de Puntos

```

1 struct Cmp{//orden total de puntos alrededor de un punto r
2     pto r;
3     Cmp(pto r):r(r) {}
4     int cuad(const pto &a) const{
5         if(a.x > 0 && a.y >= 0)return 0;
6         if(a.x <= 0 && a.y > 0)return 1;
7         if(a.x < 0 && a.y <= 0)return 2;
8         if(a.x >= 0 && a.y < 0)return 3;
9         assert(a.x ==0 && a.y==0);
10        return -1;
11    }
12    bool cmp(const pto&p1, const pto&p2)const{
13        int c1 = cuad(p1), c2 = cuad(p2);
14        if(c1==c2) return p1.y*p2.x<p1.x*p2.y;
15        else return c1 < c2;
16    }
17    bool operator()(const pto&p1, const pto&p2) const{
18        return cmp(pto(p1.x-r.x,p1.y-r.y),pto(p2.x-r.x,p2.y-r.y));
19    }
20 };

```

Linea

```

1 int sgn(ll x){return x<0? -1 : !!x;}
2 struct line{

```



```

3   line() {}
4   double a,b,c;//Ax+By=C
5   //pto MUST store float coordinates!
6   line(double a, double b, double c):a(a),b(b),c(c){}
7   // TO DO chequear porque paso problema metiendo negativo el C (-(todo el
   calculo como esta))
8   line(pto p, pto q): a(q.y-p.y), b(p.x-q.x), c(a*p.x+b*p.y) {}
9   int side(pto p){return sgn(l1(a) * p.x + l1(b) * p.y - c);}
10 };
11 bool parallels(line l1, line l2){return abs(l1.a*l2.b-l2.a*l1.b)<EPS;}
12 pto inter(line l1, line l2){//intersection
13     double det=l1.a*l2.b-l2.a*l1.b;
14     if(abs(det)<EPS) return pto(INF, INF);//parallels
15     return pto(l2.b*l1.c-l1.b*l2.c, l1.a*l2.c-l2.a*l1.c)/det;
16 }

```

Segmento

```

1 struct segm{
2     pto s,f;
3     segm(pto s, pto f):s(s), f(f) {}
4     pto closest(pto p) {//use for dist to point
5         double l2 = dist_sq(s, f);
6         if(l2==0.) return s;
7         double t=((p-s)*(f-s))/l2;
8         if (t<0.) return s;//not write if is a line
9         else if(t>1.)return f;//not write if is a line
10        return s+((f-s)*t);
11    }
12    bool inside(pto p){return abs(dist(s, p)+dist(p, f)-dist(s, f))<EPS;}
13 };
14
15 //NOTA: Si los segmentos son colineales solo devuelve un punto de interseccion
16 pto inter(segm s1, segm s2){
17     if(s1.inside(s2.s)) return s2.s; //Fix cuando son colineales
18     if(s1.inside(s2.f)) return s2.f; //Fix cuando son colineales
19     pto r=inter(line(s1.s, s1.f), line(s2.s, s2.f));
20     if(s1.inside(r) && s2.inside(r)) return r;
21     return pto(INF, INF);
22 }

```

Rectangulo

```

1 struct rect{
2     //lower-left and upper-right corners
3     pto lw, up;
4 };
5 //returns if there's an intersection and stores it in r

```

```

6 bool inter(rect a, rect b, rect &r){
7     r.lw=pto(max(a.lw.x, b.lw.x), max(a.lw.y, b.lw.y));
8     r.up=pto(min(a.up.x, b.up.x), min(a.up.y, b.up.y));
9     //check case when only a edge is common
10    return r.lw.x<r.up.x && r.lw.y<r.up.y;
11 }

```

Circulo

```

1 vec perp(vec v){return vec(-v.y, v.x);}
2 line bisector(pto x, pto y){
3     line l=line(x, y); pto m=(x+y)/2;
4     return line(-l.b, l.a, -l.b*m.x+l.a*m.y);
5 }
6 struct Circle{
7     pto o;
8     double r;
9     Circle(pto x, pto y, pto z){
10         o=inter(bisector(x, y), bisector(y, z));
11         r=dist(o, x);
12     }
13     pair<pto, pto> ptosTang(pto p){
14         pto m=(p+o)/2;
15         tipo d=dist(o, m);
16         tipo a=r*r/(2*d);
17         tipo h=sqrt(r*r-a*a);
18         pto m2=o+(m-o)*a/d;
19         vec per=perp(m-o)/d;
20         return make_pair(m2-per*h, m2+per*h);
21     }
22 };
23 //finds the center of the circle containing p1 and p2 with radius r
24 //as there may be two solutions swap p1, p2 to get the other
25 bool circle2PtsRad(pto p1, pto p2, double r, pto &c){
26     double d2=(p1-p2).norm_sq(), det=r*r/d2-0.25;
27     if(det<0) return false;
28     c=(p1+p2)/2+perp(p2-p1)*sqrt(det);
29     return true;
30 }
31 #define sqr(a) ((a)*(a))
32 #define feq(a,b) (fabs((a)-(b))<EPS)
33 pair<tipo, tipo> ecCuad(tipo a, tipo b, tipo c){//a*x*x+b*x+c=0
34     tipo dx = sqrt(b*b-4.0*a*c);
35     return make_pair((-b + dx)/(2.0*a),(-b - dx)/(2.0*a));
36 }
37 pair<pto, pto> interCL(Circle c, line l){
38     bool sw=false;

```



```

39  if((sw=freq(0,1.b))) {
40  swap(1.a, 1.b);
41  swap(c.o.x, c.o.y);
42  }
43  pair<tipo, tipo> rc = ecCuad(
44  sqr(1.a)+sqr(1.b),
45  2.0*1.a*1.b*c.o.y-2.0*(sqr(1.b)*c.o.x+1.c*1.a),
46  sqr(1.b)*(sqr(c.o.x)+sqr(c.o.y)-sqr(c.r))+sqr(1.c)-2.0*1.c*1.b*c.o.y
47  );
48  pair<pto, pto> p( pto(rc.first, (1.c - 1.a * rc.first) / 1.b),
49  pto(rc.second, (1.c - 1.a * rc.second) / 1.b) );
50  if(sw){
51  swap(p.first.x, p.first.y);
52  swap(p.second.x, p.second.y);
53  }
54  return p;
55  }
56  pair<pto, pto> interCC(Circle c1, Circle c2){
57  line l;
58  l.a = c1.o.x-c2.o.x;
59  l.b = c1.o.y-c2.o.y;
60  l.c = (sqr(c2.r)-sqr(c1.r)+sqr(c1.o.x)-sqr(c2.o.x)+sqr(c1.o.y)
61  -sqr(c2.o.y))/2.0;
62  return interCL(c1, l);
63  }

```

Area de poligono

```

1  double area(vector<pto> &p){//0(sz(p))
2  double area=0;
3  forn(i, sz(p)) area+=p[i]^p[(i+1)%sz(p)];
4  //if points are in clockwise order then area is negative
5  return abs(area)/2;
6  }
7  //Area ellipse = M_PI*a*b where a and b are the semi axis lengths
8  //Area triangle = sqrt(s*(s-a)(s-b)(s-c)) where s=(a+b+c)/2

```

Punto en poligono

```

1  //checks if v is inside of P, using ray casting
2  //works with convex and concave.
3  //excludes boundaries, handle it separately using segment.inside()
4  bool inPolygon(pto v, vector<pto>& P) {
5  bool c = false;
6  forn(i, sz(P)){
7  int j=(i+1)%sz(P);
8  if((P[j].y>v.y) != (P[i].y > v.y) &&
9  (v.x < (P[i].x - P[j].x) * (v.y-P[j].y) / (P[i].y - P[j].y) + P[j].x))

```

```

10  c = !c;
11  }
12  return c;
13  }

```

Punto en Poligono Convexo

$O(\log n)$

```

1  void normalize(vector<pto> &pt) //delete collinear points first!
2  {
3  //this makes it clockwise:
4  if(pt[2].left(pt[0], pt[1])) reverse(pt.begin(), pt.end());
5  int n=sz(pt), pi=0;
6  forn(i, n)
7  if(pt[i].x<pt[pi].x || (pt[i].x==pt[pi].x && pt[i].y<pt[pi].y))
8  pi=i;
9  vector<pto> shift(n);//puts pi as first point
10  forn(i, n) shift[i]=pt[(pi+i)%n];
11  pt.swap(shift);
12  }
13  bool inPolygon(pto p, const vector<pto> &pt)
14  {
15  //call normalize first!
16  if(p.left(pt[0], pt[1]) || p.left(pt[sz(pt)-1], pt[0])) return false;
17  int a=1, b=sz(pt)-1;
18  while(b-a>1)
19  {
20  int c=(a+b)/2;
21  if(!p.left(pt[0], pt[c])) a=c;
22  else b=c;
23  }
24  return !p.left(pt[a], pt[a+1]);
25  }

```

Chequeo de Convex

```

1
2  bool isConvex(vector<int> &p){//0(N), delete collinear points!
3  int N=sz(p);
4  if(N<3) return false;
5  bool isLeft=p[0].left(p[1], p[2]);
6  forr(i, 1, N)
7  if(p[i].left(p[(i+1)%N], p[(i+2)%N])!=isLeft)
8  return false;
9  return true; }

```

Convex Hull

```

1  //stores convex hull of P in S, CCW order

```

```

2 //left must return >=0 to delete collinear points!
3 void CH(vector<pto>& P, vector<pto> &S){
4     S.clear();
5     sort(P.begin(), P.end()); //first x, then y
6     forn(i, sz(P)){ //lower hull
7         while(sz(S)>= 2 && S[sz(S)-1].left(S[sz(S)-2], P[i])) S.pop_back();
8         S.pb(P[i]);
9     }
10    S.pop_back();
11    int k=sz(S);
12    dforn(i, sz(P)){ //upper hull
13        while(sz(S) >= k+2 && S[sz(S)-1].left(S[sz(S)-2], P[i])) S.pop_back();
14        S.pb(P[i]);
15    }
16    S.pop_back();
17 }

```

Convex Hull Trick

```

1 struct Line{tipo m,h;};
2 tipo inter(Line a, Line b){
3     tipo x=b.h-a.h, y=a.m-b.m;
4     return x/y+(x%y?!((x>0)^(y>0)):0); //==ceil(x/y)
5 }
6 struct CHT {
7     vector<Line> c;
8     bool mx;
9     int pos;
10    CHT(bool mx=0):mx(mx),pos(0){ //mx=1 si las query devuelven el max
11    inline Line acc(int i){return c[c[0].m>c.back().m? i : sz(c)-1-i];}
12    inline bool irre(Line x, Line y, Line z){
13        return c[0].m>z.m? inter(y, z) <= inter(x, y)
14            : inter(y, z) >= inter(x, y);
15    }
16    void add(tipo m, tipo h) { //O(1), los m tienen que entrar ordenados
17        if(mx) m*=-1, h*=-1;
18        Line l=(Line){m, h};
19        if(sz(c) && m==c.back().m) { l.h=min(h, c.back().h), c.pop_back(); if(
20            pos) pos--; }
21        while(sz(c)>=2 && irre(c[sz(c)-2], c[sz(c)-1], l)) { c.pop_back(); if(
22            pos) pos--; }
23        c.pb(l);
24    }
25    inline bool fbin(tipo x, int m) {return inter(acc(m), acc(m+1))>x;}
26    tipo eval(tipo x){
27        int n = sz(c);
28        //query con x no ordenados O(lgn)

```

```

27     int a=-1, b=n-1;
28     while(b-a>1) { int m = (a+b)/2;
29         if(fbin(x, m)) b=m;
30         else a=m;
31     }
32     return (acc(b).m*x+acc(b).h)*(mx?-1:1);
33     //query O(1)
34     while(pos>0 && fbin(x, pos-1)) pos--;
35     while(pos<n-1 && !fbin(x, pos)) pos++;
36     return (acc(pos).m*x+acc(pos).h)*(mx?-1:1);
37 }
38 } ch;

```

Convex Hull Trick Dinamico

```

1 const ll is_query = -(1LL<<62);
2 struct Line {
3     ll m, b;
4     mutable multiset<Line>::iterator it;
5     const Line *succ(multiset<Line>::iterator it) const;
6     bool operator<(const Line& rhs) const {
7         if (rhs.b != is_query) return m < rhs.m;
8         const Line *s=succ(it);
9         if(!s) return 0;
10        ll x = rhs.m;
11        return b - s->b < (s->m - m) * x;
12    }
13 };
14 struct HullDynamic : public multiset<Line>{ // will maintain upper hull for
15     maximum
16     bool bad(iterator y) {
17         iterator z = next(y);
18         if (y == begin()) {
19             if (z == end()) return 0;
20             return y->m == z->m && y->b <= z->b;
21         }
22         iterator x = prev(y);
23         if (z == end()) return y->m == x->m && y->b <= x->b;
24         return (x->b - y->b)*(z->m - y->m) >= (y->b - z->b)*(y->m - x->m);
25     }
26     iterator next(iterator y){return ++y;}
27     iterator prev(iterator y){return --y;}
28     void insert_line(ll m, ll b) {
29         iterator y = insert((Line) { m, b });
30         y->it=y;
31         if (bad(y)) { erase(y); return; }
32         while (next(y) != end() && bad(next(y))) erase(next(y));

```

```

32     while (y != begin() && bad(prev(y))) erase(prev(y));
33 }
34 ll eval(ll x) {
35     Line l = *lower_bound((Line) { x, is_query });
36     return l.m * x + l.b;
37 }
38 }h;
39 const Line *Line::succ(multiset<Line>::iterator it) const{
40     return (++it==h.end())? NULL : &*it;};

```

Cortar poligono

```

1 //cuts polygon Q along the line ab
2 //stores the left side (swap a, b for the right one) in P
3 void cutPolygon(pto a, pto b, vector<pto> Q, vector<pto> &P){
4     P.clear();
5     forn(i, sz(Q)){
6         double left1=(b-a)^(Q[i]-a), left2=(b-a)^(Q[(i+1)%sz(Q)]-a);
7         if(left1>=0) P.pb(Q[i]);
8         if(left1*left2<0)
9             P.pb(inter(line(Q[i], Q[(i+1)%sz(Q)]), line(a, b)));
10    }
11 }

```

Intersección de Circulos

```

1 struct event {
2     double x; int t;
3     event(double xx, int tt) : x(xx), t(tt) {}
4     bool operator <(const event &o) const { return x < o.x; }
5 };
6 typedef vector<Circle> VC;
7 typedef vector<event> VE;
8 int n;
9 double cuenta(VE &v, double A,double B)
10 {
11     sort(v.begin(), v.end());
12     double res = 0.0, lx = ((v.empty())?0.0:v[0].x);
13     int contador = 0;
14     forn(i,sz(v))
15     { //interseccion de todos (contador == n), union de todos (contador > 0)
16         //conjunto de puntos cubierto por exacta k Circulos (contador == k)
17         if (contador == n) res += v[i].x - lx;
18         contador += v[i].t, lx = v[i].x;
19     }
20     return res;
21 }
22 // Primitiva de sqrt(r*r - x*x) como funcion double de una variable x.

```

```

23 inline double primitiva(double x,double r)
24 {
25     if (x >= r) return r*r*M_PI/4.0;
26     if (x <= -r) return -r*r*M_PI/4.0;
27     double raiz = sqrt(r*r-x*x);
28     return 0.5 * (x * raiz + r*r*atan(x/raiz));
29 }
30 double interCircle(VC &v)
31 {
32     vector<double> p; p.reserve(v.size() * (v.size() + 2));
33     forn(i,sz(v)) p.push_back(v[i].c.x + v[i].r), p.push_back(v[i].c.x - v[i].r)
34     ;
35     forn(i,sz(v)) forn(j,i)
36     {
37         Circle &a = v[i], b = v[j];
38         double d = (a.c - b.c).norm();
39         if (fabs(a.r - b.r) < d && d < a.r + b.r)
40         {
41             double alfa = acos((sqr(a.r) + sqr(d) - sqr(b.r)) / (2.0 * d * a.r));
42             pto vec = (b.c - a.c) * (a.r / d);
43             p.pb((a.c + rotate(vec, alfa)).x), p.pb((a.c + rotate(vec, -alfa)).x);
44         }
45     }
46     sort(p.begin(), p.end());
47     double res = 0.0;
48     forn(i,sz(p)-1)
49     {
50         const double A = p[i], B = p[i+1];
51         VE ve; ve.reserve(2 * v.size());
52         forn(j,sz(v))
53         {
54             const Circle &c = v[j];
55             double arco = primitiva(B-c.c.x,c.r) - primitiva(A-c.c.x,c.r);
56             double base = c.c.y * (B-A);
57             ve.push_back(event(base + arco,-1));
58             ve.push_back(event(base - arco, 1));
59         }
60         res += cuenta(ve,A,B);
61     }
62     return res;
63 }

```

Rotar Matriz

```

1 //rotates matrix t 90 degrees clockwise
2 //using auxiliary matrix t2(faster)
3 void rotate()

```

```

4 | {
5 |     forn(x, n) forn(y, n)
6 |         t2[n-y-1][x]=t[x][y];
7 |     memcpy(t, t2, sizeof(t));
8 | }

```

Matemática

Identidades

$$\sum_{i=0}^n \binom{n}{i} = 2^n$$

$$\sum_{i=0}^n i \binom{n}{i} = n * 2^{n-1}$$

$$\sum_{i=m}^n i = \frac{n(n+1)}{2} - \frac{m(m-1)}{2} = \frac{(n+1-m)(n+m)}{2}$$

$$\sum_{i=0}^n i = \sum_{i=1}^n i = \frac{n(n+1)}{2}$$

$$\sum_{i=0}^n i^2 = \frac{n(n+1)(2n+1)}{6} = \frac{n^3}{3} + \frac{n^2}{2} + \frac{n}{6}$$

$$\sum_{i=0}^n i(i-1) = \frac{8}{6} \left(\frac{n}{2} + 1\right) (n+1) \text{ (doubles)} \rightarrow \text{Sino ver caso impar y par}$$

$$\sum_{i=0}^n i^3 = \left(\frac{n(n+1)}{2}\right)^2 = \frac{n^4}{4} + \frac{n^3}{2} + \frac{n^2}{4} = \left[\sum_{i=1}^n i\right]^2$$

$$\sum_{i=0}^n i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} = \frac{n^5}{5} + \frac{n^4}{2} + \frac{n^3}{3} - \frac{n}{30}$$

$$\sum_{i=0}^n i^p = \frac{(n+1)^{p+1}}{p+1} + \sum_{k=1}^p \frac{B_k}{p-k+1} \binom{p}{k} (n+1)^{p-k+1}$$

$$r = e - v + k + 1$$

Teorema de Pick: (Area, puntos interiores y puntos en el borde)

$$A = I + \frac{B}{2} - 1$$

Ec. Caracteristica

$$a_0 T(n) + a_1 T(n-1) + \dots + a_k T(n-k) = 0$$

$$p(x) = a_0 x^k + a_1 x^{k-1} + \dots + a_k$$

Sean r_1, r_2, \dots, r_q las raíces distintas, de mult. m_1, m_2, \dots, m_q

$$T(n) = \sum_{i=1}^q \sum_{j=0}^{m_i-1} c_{ij} n^j r_i^n$$

Las constantes c_{ij} se determinan por los casos base.

Teorema Chino del Resto

$$y = \sum_{j=1}^n (x_j * (\prod_{i=1, i \neq j}^n m_i)^{-1} * \prod_{i=1, i \neq j}^n m_i)$$

```

1 | //Chinese remainder theorem (special case): find z such that
2 | //z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = lcm(m1, m2).
3 | //Return (z, M). On failure, M = -1.
4 | ii chinese_remainder_theorem(int m1, int r1, int m2, int r2)
5 | { //{xx,yy,d} son variables globales usadas en extendedEuclid
6 |     extendedEuclid(m1, m2);
7 |     if (r1%d != r2%d) return make_pair(0,-1);
8 |     return mp(sumMod(xx*r2*m1, yy*r1*m2, m1*m2) / d, m1*m2 / d);
9 | }
10 | //Chinese remainder theorem: find z such that z % m[i] = r[i] for all i.
11 | //Note that the solution is unique modulo M = lcm_i (m[i]).
12 | //Return (z, M). On failure, M = -1.

```

```

13 | //Note that we do not require the a[i]'s to be relatively prime.
14 | ii chinese_remainder_theorem(const vector<int> &m, const vector<int> &r)
15 | {
16 |     ii ret=mp(r[0], m[0]);
17 |     forr(i,1,m.size())
18 |     {
19 |         ret=chinese_remainder_theorem(ret.snd, ret.fst, m[i], r[i]);
20 |         if (ret.snd==-1) break;
21 |     }
22 |     return ret;
23 | }

```

GCD & LCM

```

1 | int gcd(int a, int b) {return b? gcd(b,a%b) : a;}
2 | int lcm(int a, int b) {return a*(b/gcd(a,b));}

```

Euclides Extendido

```

1 | //ecuacin diofntica lineal
2 | //sea d=gcd(a,b); la ecuacin a * x + b * y = c tiene soluciones enteras si
3 | //d|c. La siguiente funcin nos sirve para esto. De forma general ser :
4 | //x = x0 + (b/d)n      x0 = xx*c/d
5 | //y = y0 - (a/d)n      y0 = yy*c/d
6 | ll xx,yy,d;
7 | void extendedEuclid(ll a, ll b)  //a * xx + b * yy = d
8 | {
9 |     if (!b) {xx=1; yy=0; d=a; return;}
10 |    extendedEuclid (b,a%b);
11 |    ll x1=yy;
12 |    ll y1=xx-(a/b)*yy;
13 |    xx=x1; yy=y1;
14 | }

```

Combinatoria

```

1 | void cargarComb()//O(MAXN^2)
2 | {
3 |     forn(i, MAXN+1) //comb[i][k]=i tomados de a k = i!/(k*(i-k)!)
4 |     {
5 |         comb[0][i]=0;
6 |         comb[i][0]=comb[i][i]=1;
7 |         forr(k, 1, i) comb[i][k]=(comb[i-1][k-1]+comb[i-1][k]) %MOD;
8 |     }
9 | }
10 | ll lucas (ll n, ll k, int p)
11 | { //Calcula (n,k)%p teniendo comb[p][p] precalculado.
12 |     ll aux = 1;
13 |     while (n + k)

```

```

14 {
15     aux = (aux * comb[n %p][k %p]) %p;
16     n/=p, k/=p;
17 }
18 return aux;
19 }

```

Exponenciación de Matrices y Fibonacci

```

1 #define SIZE 350
2 int NN;
3 void mul(double a[SIZE][SIZE], double b[SIZE][SIZE])
4 {
5     double res[SIZE][SIZE] = {{0}};
6     forn(i, NN) forn(j, NN) forn(k, NN) res[i][j] += a[i][k] * b[k][j];
7     forn(i, NN) forn(j, NN) a[i][j] = res[i][j];
8 }
9 void powmat(double a[SIZE][SIZE], int n, double res[SIZE][SIZE])
10 {
11     forn(i, NN) forn(j, NN) res[i][j] = (i == j);
12     while(n)
13     {
14         if(n & 1) mul(res, a), n--;
15         else mul(a, a), n /= 2;
16     }
17 }
18
19 struct M22{    // |a b|
20     tipo a,b,c,d; // |c d| -- TIPO
21     M22 operator*(const M22 &p) const {
22         return (M22){a*p.a+b*p.c, a*p.b+b*p.d, c*p.a+d*p.c, c*p.b+d*p.d};}
23 };
24 M22 operator^(const M22 &p, int n)
25 { //VER COMO SE PUEDE PONER DENTRO DEL STRUCT
26     if(!n) return (M22){1, 0, 0, 1}; //identidad
27     M22 q = p^(n/2); q = q*q;
28     return n % 2 ? p * q : q;
29 }
30
31 ll fibo(ll n) //calcula el fibonacci enesimo en O(logN)
32 {
33     M22 mat = (M22){0, 1, 1, 1}^n;
34     return mat.a*f0+mat.b*f1; //f0 y f1 son los valores iniciales
35 }

```

Operaciones Modulares

```

1 ll mulMod(ll a, ll b, ll m=MOD) //O(log b)

```

```

2 { //returns (a*b) %c, and minimize overflow
3     ll x=0, y=a%m;
4     while(b>0)
5     {
6         if(b%2==1) x=(x+y)%m;
7         y=(y*2)%m;
8         b/=2;
9     }
10    return x%m;
11 }
12 ll expMod(ll b, ll e, ll m=MOD) //O(log b)
13 {
14     if(!e) return 1;
15     ll q = expMod(b, e/2, m);
16     q = mulMod(q, q, m);
17     return e%2 ? mulMod(b, q, m) : q;
18 }
19 ll sumMod(ll a, ll b, ll m=MOD)
20 {
21     a %= m;
22     b %= m;
23     if(a < 0) a += m;
24     if(b < 0) b += m;
25     return (a+b)%m;
26 }
27 ll difMod(ll a, ll b, ll m=MOD)
28 {
29     a %= m;
30     b %= m;
31     if(a < 0) a += m;
32     if(b < 0) b += m;
33     ll ret = a - b;
34     if(ret < 0) ret += m;
35     return ret;
36 }
37 ll divMod(ll a, ll b, ll m=MOD)
38 {
39     return mulMod(a, inverso(b), m);
40 }

```

Funciones de Primos

Sea $n = \prod p_i^{k_i}$, fact(n) genera un map donde a cada p_i le asocia su k_i

```

1 #define MAXP 100000 //no necesariamente primo
2 int criba[MAXP+1];
3 void crearCriba()
4 {

```

```

5   int w[] = {4,2,4,2,4,6,2,6};
6   for(int p=25;p<=MAXP;p+=10) criba[p]=5;
7   for(int p=9;p<=MAXP;p+=6) criba[p]=3;
8   for(int p=4;p<=MAXP;p+=2) criba[p]=2;
9   for(int p=7,cur=0;p*p<=MAXP;p+=w[cur++&7]) if (!criba[p])
10  for(int j=p*p;j<=MAXP;j+=(p<<1)) if(!criba[j]) criba[j]=p;
11 }
12 vector<int> primos;
13 void buscarPrimos()
14 {
15     crearCriba();
16     forr (i,2,MAXP+1) if (!criba[i]) primos.push_back(i);
17 }
18
19 //factoriza bien numeros hasta MAXP^2
20 void fact(ll n,map<ll,ll> &f) //0 (cant primos)
21 { //llamar a buscarPrimos antes
22     forall(p, primos){
23         while(!(n %*p))
24             {
25                 f[*p]++;//divisor found
26                 n/=*p;
27             }
28     }
29     if(n>1) f[n]++;
30 }
31
32 //factoriza bien numeros hasta MAXP
33 void fact2(ll n,map<ll,ll> &f) //0 (lg n)
34 { //llamar a crearCriba antes
35     while (criba[n])
36     {
37         f[criba[n]]++;
38         n/=criba[n];
39     }
40     if(n>1) f[n]++;
41 }
42
43 //Usar asi: divisores(fac, divs, fac.begin()); NO ESTA ORDENADO
44 void divisores(map<ll,ll> &f,vector<ll> &divs,map<ll,ll>::iterator it,ll n=1)
45 {
46     if(it==f.begin()) divs.clear();
47     if(it==f.end())
48     {
49         divs.pb(n);
50         return;

```

```

51     }
52     ll p=it->fst, k=it->snd; ++it;
53     forn(_, k+1) divisores(f, divs, it, n), n*=p;
54 }
55 ll cantDivs(map<ll,ll> &f)
56 {
57     ll ret=1;
58     forall(it, f) ret*=(it->second+1);
59     return ret;
60 }
61 ll sumDivs(map<ll,ll> &f)
62 {
63     ll ret=1;
64     forall(it, f)
65     {
66         ll pot=1, aux=0;
67         forn(i, it->snd+1) aux+=pot, pot*=it->fst;
68         ret*=aux;
69     }
70     return ret;
71 }
72
73 ll eulerPhi(ll n) // con criba: O(lg n)
74 {
75     map<ll,ll> f;
76     fact(n,f);
77     ll ret=n;
78     forall(it, f) ret-=ret/it->first;
79     return ret;
80 }
81 ll eulerPhi2(ll n) // 0 (sqrt n)
82 {
83     ll r = n;
84     forr(i,2,n+1)
85     {
86         if((ll)i*i>n) break;
87         if(n%i==0)
88         {
89             while(n%i==0) n/=i;
90             r -= r/i;
91         }
92     }
93     if (n != 1) r-= r/n;
94     return r;
95 }

```

Phollard's Rho

```

1 bool es_primo_prob(ll n, int a)
2 {
3     if(n==a) return true;
4     ll s=0,d=n-1;
5     while(d%2==0) s++,d/=2;
6     ll x=expMod(a,d,n);
7     if((x==1) || (x+1==n)) return true;
8     forn(i,s-1)
9     {
10         x=mulMod(x, x, n);
11         if(x==1) return false;
12         if(x+1==n) return true;
13     }
14     return false;
15 }
16 bool rabin (ll n) //devuelve true si n es primo
17 {
18     if(n==1) return false;
19     const int ar[]={2,3,5,7,11,13,17,19,23};
20     forn(j,9) if(!es_primo_prob(n,ar[j])) return false;
21     return true;
22 }
23 ll rho(ll n)
24 {
25     if((n&1)==0) return 2;
26     ll x=2,y=2,d=1;
27     ll c=rand()%n+1;
28     while(d==1)
29     {
30         x=(mulMod(x,x,n)+c)%n;
31         y=(mulMod(y,y,n)+c)%n;
32         y=(mulMod(y,y,n)+c)%n;
33         if(x-y>=0) d=gcd(n,x-y);
34         else d=gcd(n,y-x);
35     }
36     return d==n? rho(n):d;
37 }
38 void factRho (ll n,map<ll,ll> &f) //0 (lg n)^3 un solo numero
39 {
40     if (n == 1) return;
41     if (rabin(n))
42     {
43         f[n]++;
44         return;

```

```

45     }
46     ll factor = rho(n);
47     factRho(factor,f);
48     factRho(n/factor,f);
49 }

```

Inversos

```

1 #define MAXMOD 15485867
2 ll inv[MAXMOD]; //inv[i]*i=1 mod MOD
3 void calc(int p) //0(p)
4 {
5     inv[1]=1;
6     forr(i,2,p) inv[i]=p-((p/i)*inv[p%i])%p;
7 }
8 int inverso(int x) //0(log x)
9 {
10     return expMod(x, eulerPhi(MOD)-2); //si mod no es primo(sacar a mano)
11     return expMod(x, MOD-2); //si mod es primo
12 }

```

Fracciones

```

1 struct frac{
2     int p,q;
3     frac(int p=0,int q=1):p(p),q(q) {norm();}
4     void norm()
5     {
6         int a=gcd(q,p);
7         if(a) p/=a, q/=a;
8         else q=1;
9         if (q<0) q=-q, p=-p;
10    }
11    frac operator+(const frac& o)
12    {
13        int a=gcd(o.q,q);
14        return frac(p*(o.q/a)+o.p*(q/a),q*(o.q/a));
15    }
16    frac operator-(const frac& o)
17    {
18        int a=gcd(o.q,q);
19        return frac(p*(o.q/a)-o.p*(q/a),q*(o.q/a));
20    }
21    frac operator*(frac o)
22    {
23        int a=gcd(o.p,q), b=gcd(p,o.q);
24        return frac((p/b)*(o.p/a),(q/a)*(o.q/b));
25    }

```



```
26 | frac operator/(frac o)
27 | {
28 |     int a=gcd(o.q,q), b=gcd(p,o.p);
29 |     return frac((p/b)*(o.q/a),(q/a)*(o.p/b));
30 | }
31 | bool operator<(const frac &o) const{return p*o.q < o.p*q;}
32 | bool operator==(frac o){return p==o.p&&q==o.q;}
33 |};
```

Simpson

```
1 | double integral(double a, double b, int n=10000) //0(n), n=cantdiv
2 | {
3 |     double area=0, h=(b-a)/n, fa=f(a), fb;
4 |     forn(i, n)
5 |     {
6 |         fb=f(a+h*(i+1));
7 |         area+=fa+ 4*f(a+h*(i+0.5)) +fb, fa=fb;
8 |     }
9 |     return area*h/6.;
10 | }
```

Tablas y cotas (Primos, Divisores, Factoriales, etc)

Maxs

max signed tint = 9.223.372.036.854.775.807

max unsigned tint = 18.446.744.073.709.551.615

Primos

2 3 5 7 11 13 17 19 23 29 31 37 41 43 47 53 59 61 67 71 73 79 83 89 97 101 103 107 109 113 127 131 137 139 149 151 157 163 167 173

Cantidad de primos menores que 10ⁿ

π(10¹) = 4 ; π(10²) = 25 ; π(10³) = 168 ; π(10⁴) = 1229 ; π(10⁵) = 9592

π(10⁶) = 78.498 ; π(10⁷) = 664.579 ; π(10⁸) = 5.761.455 ; π(10⁹) = 50.847.534

π(10¹⁰) = 455.052,511 ; π(10¹¹) = 4.118.054.813 ; π(10¹²) = 37.607.912.018

Números Catalanés

Utiles para problemas de Combinatoria

$$Cat(n) = \frac{\binom{2n}{n}}{n+1} = \frac{(2n)!}{n!(n+1)!}$$

Con $Cat(0) = 1$.

Diferentes aplicaciones:

1. Contar la cantidad de diferentes arboles binarios con n nodos que se pueden armar.

2. Contar las formas en que un polígono convexo de $n + 2$ lados puede ser triangulado.

3. Contar la cantidad de caminos monotonos a lo largo de los lados de una grilla $n * n$, que no cruzan la diagonal.

4. Contar el número de expresiones que contienen n pares de paréntesis correctamente colocados

Primeros 25 Catalanés

1 1 2 5 14 42 132 429 1430 4862 16796 58786 208012 742900 2674440 9694845 35357670

129644790 477638700 1767263190 6564120420 24466267020 91482563640 343059613650

1289904147324 4861946401452

Grafos

Dijkstra

```
1 | #define INF 1e9
2 | int N;
3 | #define MAX_V 250001
4 | vector<ii> G[MAX_V];
5 | //To add an edge use
6 | #define add(a, b, w) G[a].pb(make_pair(w, b))
7 | ll dijkstra(int s, int t){//0(|E| log |V|)
8 |     priority_queue<ii, vector<ii>, greater<ii> > Q;
9 |     vector<ll> dist(N, INF); vector<int> dad(N, -1);
10 |     Q.push(make_pair(0, s)); dist[s] = 0;
11 |     while(sz(Q)){
12 |         ii p = Q.top(); Q.pop();
13 |         if(p.snd == t) break;
14 |         forall(it, G[p.snd])
15 |             if(dist[p.snd]+it->first < dist[it->snd]){
16 |                 dist[it->snd] = dist[p.snd] + it->fst;
17 |                 dad[it->snd] = p.snd;
18 |                 Q.push(make_pair(dist[it->snd], it->snd)); }
19 |     }
20 |     return dist[t];
21 |     if(dist[t]<INF)//path generator
22 |         for(int i=t; i!=-1; i=dad[i])
23 |             printf("%d%c", i, (i==s?' \n':' \u'));
```

Bellman-Ford

```
1 | //Mas lento que Dijkstra, pero maneja arcos con peso negativo
2 | vector<ii> G[MAX_N]; //ady. list with pairs (weight, dst)
3 | int dist[MAX_N];
4 | void bford(int src){//0(VE)
5 |     dist[src]=0;
6 |     forn(i, N-1) forn(j, N) if(dist[j]!=INF) forall(it, G[j])
7 |         dist[it->snd]=min(dist[it->snd], dist[j]+it->fst);
8 | }
9 |
10 | bool hasNegCycle(){
11 |     forn(j, N) if(dist[j]!=INF) forall(it, G[j])
```

```

12     if(dist[it->snd]>dist[j]+it->fst) return true;
13     //inside if: all points reachable from it->snd will have -INF distance(do bfs
14     )
15     return false;
16 }

```

Floyd-Warshall

```

1 // Camino minimo en grafos dirigidos ponderados, en todas las parejas de nodos.
2 //G[i][j] contains weight of edge (i, j) or INF
3 //G[i][i]=0
4 int G[MAX_N][MAX_N];
5 void floyd(){//O(N^3)
6     forn(k, N) forn(i, N) if(G[i][k]!=INF) forn(j, N) if(G[k][j]!=INF)
7         G[i][j]=min(G[i][j], G[i][k]+G[k][j]);
8 }
9 bool inNegCycle(int v){
10     return G[v][v]<0;}
11 //checks if there's a neg. cycle in path from a to b
12 bool hasNegCycle(int a, int b){
13     forn(i, N) if(G[a][i]!=INF && G[i][i]<0 && G[i][b]!=INF)
14         return true;
15     return false;
16 }

```

Kruskal

```

1 // Minimum Spanning Tree in O(e log e)
2 bool operator<(const Ar& a, const Ar &b){return a.w<b.w;}
3 vector<Ar> E;
4 ll kruskal(){
5     ll cost=0;
6     sort(E.begin(), E.end());//ordenar aristas de menor a mayor
7     uf.init(n);
8     forall(it, E){
9         if(uf.comp(it->a)!=uf.comp(it->b)){//si no estan conectados
10             uf.unir(it->a, it->b);//conectar
11             cost+=it->w;
12         }
13     }
14     return cost;
15 }

```

Prim

```

1 vector<ii> G[MAXN];
2 bool taken[MAXN];
3 priority_queue<ii, vector<ii>, greater<ii> > pq;//min heap
4 void process(int v){

```

```

5     taken[v]=true;
6     forall(e, G[v])
7         if(!taken[e->second]) pq.push(*e);
8 }
9 // Minimum Spanning Tree in O(n^2)
10 ll prim(){
11     zero(taken);
12     process(0);
13     ll cost=0;
14     while(sz(pq)){
15         ii e=pq.top(); pq.pop();
16         if(!taken[e.second]) cost+=e.first, process(e.second);
17     }
18     return cost;
19 }

```

Kosaraju SCC

Componente Fuertemente Conexa

```

1 #define MAXN 1000000
2 vector<int> G[MAXN],gt[MAXN]; //Limpiar si se corre mas de una vez
3 //nodos 0...N-1 ; componentes 0...cantcomp-1
4 int comp[MAXN],N,cantcomp,used[MAXN];
5 stack<int> pila;
6 void add(int a, int b){ G[a].pb(b);gt[b].pb(a);}
7 void dfs1(int nodo)
8 {
9     used[nodo]=1;
10    forall(it,G[nodo]) if(!used[*it]) dfs1(*it);
11    pila.push(nodo);
12 }
13 void dfs2(int nodo)
14 {
15     used[nodo]=2;
16     comp[nodo]=cantcomp-1;
17     forall(it,gt[nodo]) if(used[*it]!=2) dfs2(*it);
18 }
19 void kosaraju()
20 {
21     cantcomp=0;
22     memset(used,0,sizeof(used));
23     forn(i,N) if(!used[i]) dfs1(i);
24     while(!pila.empty())
25     {
26         if(used[pila.top()]!=2)
27         {
28             cantcomp++;

```

```

29     dfs2(pila.top());
30 }
31 pila.pop();
32 }
33 }

```

2-SAT + Tarjan SCC

```

1 //We have a vertex representing a var and other for his negation.
2 //Every edge stored in G represents an implication. To add an equation of the
   form a||b, use addor(a, b)
3 //MAX=max cant var, n=cant var
4 #define addor(a, b) (G[neg(a)].pb(b), G[neg(b)].pb(a))
5 vector<int> G[MAX*2];
6 //idx[i]=index assigned in the dfs
7 //lw[i]=lowest index(closer from the root) reachable from i
8 int lw[MAX*2], idx[MAX*2], qidx;
9 stack<int> q;
10 int qcmp, cmp[MAX*2];
11 //verdad[cmp[i]]=valor de la variable i
12 bool verdad[MAX*2+1];
13
14 int neg(int x) { return x>=n? x-n : x+n;}
15 void tjn(int v){
16     lw[v]=idx[v]=++qidx;
17     q.push(v), cmp[v]=-2;
18     forall(it, G[v]){
19         if(!idx[*it] || cmp[*it]==-2){
20             if(!idx[*it]) tjn(*it);
21             lw[v]=min(lw[v], lw[*it]);
22         }
23     }
24     if(lw[v]==idx[v]){
25         int x;
26         do{x=q.top(); q.pop(); cmp[x]=qcmp;}while(x!=v);
27         verdad[qcmp]=(cmp[neg(v)]<0);
28         qcmp++;
29     }
30 }
31 //remember to CLEAR G!!!
32 bool satisf(){//O(n)
33     memset(idx, 0, sizeof(idx)), qidx=0;
34     memset(cmp, -1, sizeof(cmp)), qcmp=0;
35     forn(i, n){
36         if(!idx[i]) tjn(i);
37         if(!idx[neg(i)]) tjn(neg(i));
38     }

```

```

39     forn(i, n) if(cmp[i]==cmp[neg(i)]) return false;
40     return true;
41 }

```

Puntos de Articulación

```

1 int N;
2 vector<int> G[1000000];
3 //V[i]=node number(if visited), L[i]= lowest V[i] reachable from i
4 int qV, V[1000000], L[1000000], P[1000000];
5 void dfs(int v, int f){
6     L[v]=V[v]=++qV;
7     forall(it, G[v])
8         if(!V[*it]){
9             dfs(*it, v);
10            L[v] = min(L[v], L[*it]);
11            P[v]+= L[*it]>=V[v];
12        }
13        else if(*it!=f)
14            L[v]=min(L[v], V[*it]);
15    }
16    int cantart(){ //O(n)
17        qV=0;
18        zero(V), zero(P);
19        dfs(1, 0); P[1]--;
20        int q=0;
21        forn(i, N) if(P[i]) q++;
22    return q;
23 }

```

Least Common Ancestor + Climb

```

1 const int MAXN=100001, LOGN=20;
2 //f[v][k] holds the 2^k father of v
3 //L[v] holds the level of v
4 int N, f[MAXN][LOGN], L[MAXN];
5 //call before build:
6 void dfs(int v, int fa=-1, int lvl=0){//generate required data
7     f[v][0]=fa, L[v]=lvl;
8     forall(it, G[v])if(*it!=fa) dfs(*it, v, lvl+1); }
9 void build(){//f[i][0] must be filled previously, 0<nlgn
10     forn(k, LOGN-1) forn(i, N) f[i][k+1]=f[f[i][k]][k];}
11 #define lg(x) (31-__builtin_clz(x))//=floor(log2(x))
12 int climb(int a, int d){//O(lgn)
13     if(!d) return a;
14     dforn(i, lg(L[a])+1) if(1<=i<=d) a=f[a][i], d-=1<=i;
15     return a;}
16 int lca(int a, int b){//O(lgn)

```

```

17   if(L[a]<L[b]) swap(a, b);
18   a=climb(a, L[a]-L[b]);
19   if(a==b) return a;
20   dforn(i, lg(L[a])+1) if(f[a][i]!=f[b][i]) a=f[a][i], b=f[b][i];
21   return f[a][0]; }
22 int dist(int a, int b) { //returns distance between nodes
23   return L[a]+L[b]-2*L[lca(a, b)];}

```

Heavy Light Decomposition

```

1 vector<int> G[MAXN];
2 int treesz[MAXN]; //cantidad de nodos en el subarbol del nodo v
3 int dad[MAXN]; //dad[v]=padre del nodo v
4 void dfs1(int v, int p=-1) { //pre-DFS
5   dad[v]=p;
6   treesz[v]=1;
7   forall(it, G[v]) if(*it!=p){
8     dfs1(*it, v);
9     treesz[v]+=treesz[*it];
10  }
11 }
12 //PONER Q EN 0 !!!!
13 int pos[MAXN], q; //pos[v]=posicion del nodo v en el recorrido de la dfs
14 //Las cadenas aparecen continuas en el recorrido!
15 int cantcad;
16 int homecad[MAXN]; //dada una cadena devuelve su nodo inicial
17 int cad[MAXN]; //cad[v]=cadena a la que pertenece el nodo
18 void heavylight(int v, int cur=-1){
19   if(cur==-1) homecad[cur=cantcad++] = v;
20   pos[v]=q++;
21   cad[v]=cur;
22   int mx=-1;
23   forn(i, sz(G[v])) if(G[v][i]!=dad[v])
24     if(mx==-1 || treesz[G[v][mx]]<treesz[G[v][i]]) mx=i;
25   if(mx!=-1) heavylight(G[v][mx], cur);
26   forn(i, sz(G[v])) if(i!=mx && G[v][i]!=dad[v])
27     heavylight(G[v][i], -1);
28 }
29 //ejemplo de obtener el maximo numero en el camino entre dos nodos
30 //RTA: max(query(low, u), query(low, v)), con low=lca(u, v)
31 //esta funcion va trepando por las cadenas
32 int query(int an, int v) { //O(logn)
33   //si estan en la misma cadena:
34   if(cad[an]==cad[v]) return rmq.get(pos[an], pos[v]+1);
35   return max(query(an, dad[homecad[cad[v]]]),
36             rmq.get(pos[homecad[cad[v]]], pos[v]+1));
37 }

```

Centroid Decomposition

```

1 vector<int> G[MAXN];
2 bool taken[MAXN]; //poner todos en FALSE al principio!!
3 int padre[MAXN]; //padre de cada nodo en el centroid tree
4
5 int szt[MAXN];
6 void calcsz(int v, int p) {
7   szt[v] = 1;
8   forall(it, G[v]) if (*it!=p && !taken[*it])
9     calcsz(*it, v), szt[v]+=szt[*it];
10 }
11 void centroid(int v=0, int f=-1, int lvl=0, int tam=-1) { //O(nlogn)
12   if(tam==-1) calcsz(v, -1), tam=szt[v];
13   forall(it, G[v]) if(!taken[*it] && szt[*it]>=tam/2)
14     {szt[v]=0; centroid(*it, f, lvl, tam); return;}
15   taken[v]=true;
16   padre[v]=f;
17   forall(it, G[v]) if(!taken[*it])
18     centroid(*it, v, lvl+1, -1);
19 }

```

Ciclo Euleriano

```

1 int n, m, ars[MAXE], eq;
2 vector<int> G[MAXN]; //fill G, n, m, ars, eq
3 list<int> path;
4 int used[MAXN];
5 bool usede[MAXE];
6 queue<list<int>::iterator> q;
7 int get(int v) {
8   while(used[v]<sz(G[v]) && usede[ G[v][used[v]] ]) used[v]++;
9   return used[v];
10 }
11 void explore(int v, int r, list<int>::iterator it) {
12   int ar=G[v][get(v)]; int u=v^ars[ar];
13   usede[ar]=true;
14   list<int>::iterator it2=path.insert(it, u);
15   if(u!=r) explore(u, r, it2);
16   if(get(v)<sz(G[v])) q.push(it);
17 }
18 void euler() {
19   zero(used), zero(usede);
20   path.clear();
21   q=queue<list<int>::iterator>();
22   path.push_back(0); q.push(path.begin());
23   while(sz(q)){
24     list<int>::iterator it=q.front(); q.pop();

```

```

25     if(used[*it]<sz(G[*it])) explore(*it, *it, it);
26 }
27 reverse(path.begin(), path.end());
28 }
29 void addEdge(int u, int v){
30     G[u].pb(eq), G[v].pb(eq);
31     ars[eq++]=u^v;
32 }

```

Diametro Árbol

```

1 vector<int> G[MAXN]; int n,m,p[MAXN],d[MAXN],d2[MAXN];
2 int bfs(int r, int *d) {
3     queue<int> q;
4     d[r]=0; q.push(r);
5     int v;
6     while(sz(q)) { v=q.front(); q.pop();
7         forall(it,G[v]) if (d[*it]==-1)
8             d[*it]=d[v]+1, p[*it]=v, q.push(*it);
9     }
10    return v;//ultimo nodo visitado
11 }
12 vector<int> diams; vector<ii> centros;
13 void diametros(){
14     memset(d,-1,sizeof(d));
15     memset(d2,-1,sizeof(d2));
16     diams.clear(), centros.clear();
17     forn(i, n) if(d[i]==-1){
18         int v,c;
19         c=v=bfs(bfs(i, d2), d);
20         forn(_,d[v]/2) c=p[c];
21         diams.pb(d[v]);
22         if(d[v]&1) centros.pb(ii(c, p[c]));
23         else centros.pb(ii(c, c));
24     }
25 }

```

Componentes Biconexas y Puentes

```

1 vector<int> G[MAXN];
2
3 struct edge{
4     int u,v, comp;
5     bool bridge;
6 };
7 vector<edge> e;
8 void addEdge(int u, int v)
9 {

```

```

10     G[u].pb(sz(e)), G[v].pb(sz(e));
11     e.pb((edge){u,v,-1,false});
12 }
13 //d[i]=id de la dfs
14 //b[i]=lowest id reachable from i
15 int d[MAXN], b[MAXN], t;
16 int nbc;//cant componentes
17 int comp[MAXN];//comp[i]=cant comp biconexas a la cual pertenece i
18 void initDfs(int n)
19 {
20     zero(G), zero(comp);
21     e.clear();
22     forn(i,n) d[i]=-1;
23     nbc = t = 0;
24 }
25 stack<int> st;
26 void dfs(int u,int pe) //O(n + m)
27 {
28     b[u]=d[u]=t++;
29     comp[u]=(pe!=-1);
30     forall(ne,G[u]) if(*ne!=pe)
31     {
32         int v=e[*ne].u ^ e[*ne].v ^ u;
33         if(d[v]==-1)
34         {
35             st.push(*ne);
36             dfs(v,*ne);
37             if(b[v]>d[u]) e[*ne].bridge=true; // bridge
38             if(b[v]>=d[u]) // art
39             {
40                 int last;
41                 do
42                 {
43                     las=st.top(); st.pop();
44                     e[last].comp=nbc;
45                 }while(last!=*ne);
46                 nbc++;
47                 comp[u]++;
48             }
49             b[u]=min(b[u],b[v]);
50         }
51         else if(d[v]<d[u]) // back edge
52         {
53             st.push(*ne);
54             b[u]=min(b[u], d[v]);
55         }

```

Hungarian

```

56     }
57 }

//Dado un grafo bipartito completo con costos no negativos, encuentra el
// matching perfecto de minimo costo.

#define tipo double
tipo cost[N][N], lx[N], ly[N], slack[N]; //llenar: cost=matriz de adyacencia
int n, max_match, xy[N], yx[N], slackx[N], prev2[N]; //n=cantidad de nodos
bool S[N], T[N]; //sets S and T in algorithm

void add_to_tree(int x, int prevx) {
    S[x] = true, prev2[x] = prevx;
    forn(y, n) if (lx[x] + ly[y] - cost[x][y] < slack[y] - EPS)
        slack[y] = lx[x] + ly[y] - cost[x][y], slackx[y] = x;
}

void update_labels(){
    tipo delta = INF;
    forn(y, n) if (!T[y]) delta = min(delta, slack[y]);
    forn(x, n) if (S[x]) lx[x] -= delta;
    forn(y, n) if (T[y]) ly[y] += delta; else slack[y] -= delta;
}

void init_labels(){
    zero(lx), zero(ly);
    forn(x, n) forn(y, n) lx[x] = max(lx[x], cost[x][y]);
}

void augment() {
    if (max_match == n) return;
    int x, y, root, q[N], wr = 0, rd = 0;
    memset(S, false, sizeof(S)), memset(T, false, sizeof(T));
    memset(prev2, -1, sizeof(prev2));
    forn(x, n) if (xy[x] == -1){
        q[wr++] = root = x, prev2[x] = -2;
        S[x] = true; break; }
    forn(y, n) slack[y] = lx[root] + ly[y] - cost[root][y], slackx[y] = root;
    while (true){
        while (rd < wr){
            x = q[rd++];
            for (y = 0; y < n; y++) if (cost[x][y] == lx[x] + ly[y] && !T[y]){
                if (yx[y] == -1) break; T[y] = true;
                q[wr++] = yx[y], add_to_tree(yx[y], x); }
            if (y < n) break; }
        if (y < n) break;
    }
    update_labels(), wr = rd = 0;
    for (y = 0; y < n; y++) if (!T[y] && slack[y] == 0){
        if (yx[y] == -1){x = slackx[y]; break;}
        else{

```

```

42         T[y] = true;
43         if (!S[yx[y]]) q[wr++] = yx[y], add_to_tree(yx[y], slackx[y]);
44     }}
45     if (y < n) break; }
46 if (y < n){
47     max_match++;
48     for (int cx = x, cy = y, ty; cx != -2; cx = prev2[cx], cy = ty)
49         ty = xy[cx], yx[cy] = cx, xy[cx] = cy;
50     augment(); }
51 }
52 tipo hungarian(){
53     tipo ret = 0; max_match = 0, memset(xy, -1, sizeof(xy));
54     memset(yx, -1, sizeof(yx)), init_labels(), augment(); //steps 1-3
55     forn(x, n) ret += cost[x][xy[x]]; return ret;
56 }

```

Dynamic Connectivity

```

1 struct UnionFind {
2     int n, comp;
3     vector<int> pre, si, c;
4     UnionFind(int n=0):n(n), comp(n), pre(n), si(n, 1) {
5         forn(i, n) pre[i] = i; }
6     int find(int u){return u==pre[u]?u:find(pre[u]);}
7     bool merge(int u, int v)
8     {
9         if((u=find(u))==v) return false;
10        if(si[u]<si[v]) swap(u, v);
11        si[u]+=si[v], pre[v]=u, comp--, c.pb(v);
12        return true;
13    }
14    int snap(){return sz(c);}
15    void rollback(int snap)
16    {
17        while(sz(c)>snap)
18        {
19            int v = c.back(); c.pop_back();
20            si[pre[v]] -= si[v], pre[v] = v, comp++;
21        }
22    }
23 };
24 enum {ADD, DEL, QUERY};
25 struct Query {int type, u, v;};
26 struct DynCon{//bidirectional graphs; create vble as DynCon name(cant_nodos)
27     vector<Query> q;
28     UnionFind dsu;
29     vector<int> match, res;

```

Flow

Edmond Karp

```

30 map<ii,int> last;//se puede no usar cuando hay identificador para cada arista
    (mejora poco)
31 DynCon(int n=0):dsu(n){}
32 void add(int u, int v) //to add an edge
33 {
34     if(u>v) swap(u,v);
35     q.pb((Query){ADD, u, v}), match.pb(-1);
36     last[ii(u,v)] = sz(q)-1;
37 }
38 void remove(int u, int v) //to remove an edge
39 {
40     if(u>v) swap(u,v);
41     q.pb((Query){DEL, u, v});
42     int prev = last[ii(u,v)];
43     match[prev] = sz(q)-1;
44     match.pb(prev);
45 }
46 void query() //to add a question (query) type of query
47 {
48     q.pb((Query){QUERY, -1, -1}), match.pb(-1);
49 }
50 void process() //call this to process queries in the order of q
51 {
52     forn(i,sz(q)) if (q[i].type == ADD && match[i] == -1) match[i] = sz(q);
53     go(0,sz(q));
54 }
55 void go(int l, int r)
56 {
57     if(l+1==r)
58     {
59         if (q[l].type == QUERY)//Aqui responder la query usando el dsu!
60             res.pb(dsu.comp);//aqui query=cantidad de componentes conexas
61         return;
62     }
63     int s=dsu.snap(), m = (l+r) / 2;
64     forr(i,m,r) if(match[i]!=-1 && match[i]<l) dsu.merge(q[i].u, q[i].v);
65     go(l,m);
66     dsu.rollback(s);
67     s = dsu.snap();
68     forr(i,l,m) if(match[i]!=-1 && match[i]>=r) dsu.merge(q[i].u, q[i].v);
69     go(m,r);
70     dsu.rollback(s);
71 }
72 };

```

```

1 #define MAX_V 1000
2 #define INF 1e9
3 //special nodes
4 #define SRC 0
5 #define SNK 1
6 map<int, int> G[MAX_V];//limpiar esto -- unordered_map mejora
7 //To add an edge use
8 #define add(a, b, w) G[a][b]=w
9 int f, p[MAX_V];
10 void augment(int v, int minE)
11 {
12     if(v==SRC) f=minE;
13     else if(p[v]!=-1)
14     {
15         augment(p[v], min(minE, G[p[v]][v]));
16         G[p[v]][v]-=f, G[v][p[v]]+=f;
17     }
18 }
19 ll maxflow()//O(min(VE^2,Mf*E))
20 {
21     ll Mf=0;
22     do
23     {
24         f=0;
25         char used[MAX_V]; queue<int> q; q.push(SRC);
26         zero(used), memset(p, -1, sizeof(p));
27         while(sz(q))
28         {
29             int u=q.front(); q.pop();
30             if(u==SNK) break;
31             forall(it, G[u])
32                 if(it->snd>0 && !used[it->fst])
33                     used[it->fst]=true, q.push(it->fst), p[it->fst]=u;
34         }
35         augment(SNK, INF);
36         Mf+=f;
37     }while(f);
38     return Mf;
39 }

```

Min Cut

```

1 //Suponemos un grafo con el formato definido en Edmond Karp o Push relabel

```



```

2  bitset<MAX_V> type,used; //reset this
3  void dfs1(int nodo)
4  {
5      type.set(nodo);
6      forall(it,G[nodo]) if(!type[it->fst] && it->snd>0) dfs1(it->fst);
7  }
8  void dfs2(int nodo)
9  {
10     used.set(nodo);
11     forall(it,G[nodo])
12     {
13         if(!type[it->fst])
14         {
15             //edge nodo -> (it->fst) pertenece al min_cut
16             //y su peso original era: it->snd + G[it->fst][nodo]
17             //si no existia arista original al revs
18         }
19         else if(!used[it->fst]) dfs2(it->fst);
20     }
21 }
22 void minCut() //antes correr algn maxflow()
23 {
24     dfs1(SRC);
25     dfs2(SRC);
26     return;
27 }

```

Push Relabel

```

1  #define MAX_V 1000
2  int N;//valid nodes are [0...N-1]
3  #define INF 1e9
4  //special nodes
5  #define SRC 0
6  #define SNK 1
7  map<int, int> G[MAX_V]; //limpiar esto -- unordered_map mejora
8  //To add an edge use
9  #define add(a, b, w) G[a][b]=w
10 ll excess[MAX_V];
11 int height[MAX_V], active[MAX_V], cuenta[2*MAX_V+1];
12 queue<int> Q;
13
14 void enqueue(int v)
15 {
16     if (!active[v] && excess[v] > 0) active[v]=true, Q.push(v);
17 }
18 void push(int a, int b)

```

```

19 {
20     int amt = min(excess[a], ll(G[a][b]));
21     if(height[a] <= height[b] || amt == 0) return;
22     G[a][b]-=amt, G[b][a]+=amt;
23     excess[b] += amt, excess[a] -= amt;
24     enqueue(b);
25 }
26 void gap(int k)
27 {
28     forn(v, N)
29     {
30         if (height[v] < k) continue;
31         cuenta[height[v]]--;
32         height[v] = max(height[v], N+1);
33         cuenta[height[v]]++;
34         enqueue(v);
35     }
36 }
37 void relabel(int v)
38 {
39     cuenta[height[v]]--;
40     height[v] = 2*N;
41     forall(it, G[v])
42     if(it->snd) height[v] = min(height[v], height[it->fst] + 1);
43     cuenta[height[v]]++;
44     enqueue(v);
45 }
46 ll maxflow() //O(V^3)
47 {
48     zero(height), zero(active), zero(cuenta), zero(excess);
49     cuenta[0]=N-1; cuenta[N]=1;
50     height[SRC] = N;
51     active[SRC] = active[SNK] = true;
52     forall(it, G[SRC])
53     {
54         excess[SRC] += it->snd;
55         push(SRC, it->fst);
56     }
57     while(sz(Q))
58     {
59         int v = Q.front(); Q.pop();
60         active[v]=false;
61         forall(it, G[v]) push(v, it->fst);
62         if(excess[v] > 0)
63             cuenta[height[v]] == 1? gap(height[v]):relabel(v);
64     }

```

```

65     ll mf=0;
66     forall(it, G[SRC]) mf+=G[it->fst][SRC];
67     return mf;
68 }

```

Dinic

```

1 struct Edge {
2     int u, v;
3     ll cap, flow;
4     Edge() {}
5     Edge(int u, int v, ll cap): u(u), v(v), cap(cap), flow(0) {}
6 };
7 struct Dinic {
8     int N;
9     vector<Edge> E;
10    vector<vector<int>> g;
11    vector<int> d, pt;
12    Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {} //clear and init
13    void addEdge(int u, int v, ll cap)
14    {
15        if (u != v)
16        {
17            E.emplace_back(Edge(u, v, cap));
18            g[u].emplace_back(E.size() - 1);
19            E.emplace_back(Edge(v, u, 0));
20            g[v].emplace_back(E.size() - 1);
21        }
22    }
23    bool BFS(int S, int T)
24    {
25        queue<int> q({S});
26        fill(d.begin(), d.end(), N + 1);
27        d[S] = 0;
28        while(!q.empty())
29        {
30            int u = q.front(); q.pop();
31            if (u == T) break;
32            for (int k: g[u])
33            {
34                Edge &e = E[k];
35                if (e.flow < e.cap && d[e.v] > d[e.u] + 1)
36                {
37                    d[e.v] = d[e.u] + 1;
38                    q.emplace(e.v);
39                }
40            }

```

```

41        }
42        return d[T] != N + 1;
43    }
44    ll DFS(int u, int T, ll flow = -1)
45    {
46        if (u == T || flow == 0) return flow;
47        for (int &i = pt[u]; i < g[u].size(); ++i)
48        {
49            Edge &e = E[g[u][i]];
50            Edge &oe = E[g[u][i]^1];
51            if (d[e.v] == d[e.u] + 1)
52            {
53                ll amt = e.cap - e.flow;
54                if (flow != -1 && amt > flow) amt = flow;
55                if (ll pushed = DFS(e.v, T, amt))
56                {
57                    e.flow += pushed;
58                    oe.flow -= pushed;
59                    return pushed;
60                }
61            }
62        }
63        return 0;
64    }
65    ll maxFlow(int S, int T)
66    {
67        ll total = 0;
68        while(BFS(S, T))
69        {
70            fill(pt.begin(), pt.end(), 0);
71            while (ll flow = DFS(S, T)) total += flow;
72        }
73        return total;
74    }
75 };

```

Min cost - Max flow

```

1 const int MAXN=10000;
2 typedef ll tf;
3 typedef ll tc;
4 const tf INFFLUJO = 1e14;
5 const tc INFCOSTO = 1e14;
6 struct edge {
7     int u, v;
8     tf cap, flow;
9     tc cost;

```

```
10   tf rem() { return cap - flow; }
11 };
12 int nodes; //numero de nodos
13 vector<int> G[MAXN]; // limpiar!
14 vector<edge> e; // limpiar!
15 void addEdge(int u, int v, tf cap, tc cost)
16 {
17     G[u].pb(sz(e)); e.pb((edge){u,v,cap,0,cost});
18     G[v].pb(sz(e)); e.pb((edge){v,u,0,0,-cost});
19 }
20 tc dist[MAXN], mnCost;
21 int pre[MAXN];
22 tf cap[MAXN], mxFlow;
23 bool in_queue[MAXN];
24 void flow(int s, int t)
25 {
26     zero(in_queue);
27     mxFlow=mnCost=0;
28     while(1)
29     {
30         fill(dist, dist+nodes, INFCOSTO); dist[s] = 0;
31         memset(pre, -1, sizeof(pre)); pre[s]=0;
32         zero(cap); cap[s] = INFFLUJO;
33         queue<int> q; q.push(s); in_queue[s]=1;
34         while(sz(q))
35         {
36             int u=q.front(); q.pop(); in_queue[u]=0;
37             for(auto it:G[u])
38             {
39                 edge &E = e[it];
40                 if(E.rem() && dist[E.v] > dist[u] + E.cost + 1e-9) // ojo EPS
41                 {
42                     dist[E.v]=dist[u]+E.cost;
43                     pre[E.v] = it;
44                     cap[E.v] = min(cap[u], E.rem());
45                     if(!in_queue[E.v]) q.push(E.v), in_queue[E.v]=1;
46                 }
47             }
48         }
49         if (pre[t] == -1) break;
50         mxFlow +=cap[t];
51         mnCost +=cap[t]*dist[t];
52         for (int v = t; v != s; v = e[pre[v]].u)
53         {
54             e[pre[v]].flow += cap[t];
55             e[pre[v]^1].flow -= cap[t];
```

```
56     }
57 }
58 }
```

Utils

Truquitos para entradas/salidas

```
1 //Cantidad de decimales
2 cout << setprecision(2) << fixed;
3 //Rellenar con espacios(para justificar)
4 cout << setfill('␣') << setw(3) << 2 << endl;
5 //Leer hasta fin de linea
6 // hacer cin.ignore() antes de getline()
7 while(getline(cin, line)){
8     istringstream is(line);
9     while(is >> X)
10         cout << X << "␣";
11     cout << endl;
12 }
```

Comparación de Double

```
1 const double EPS = 1e-9;
2 x == y <=> fabs(x-y) < EPS
3 x > y <=> x > y + EPS
4 x >= y <=> x > y - EPS
```

Iterar subconjuntos

```
1 for(int sbm=bm; sbm; sbm=(sbm-1)&bm)
```

Funciones Utiles

Algo	Params	Función
count, count_if	f, l, elem/pred	cuenta elem, pred(i)
lexicographical_compare	f1,l1,f2,l2	bool con [f1,l1]i[f2,l2]
inner_product	f1, l1, f2, i	$T = i + [f1, l1) . [f2, \dots)$
partial_sum	f, l, r, [op]	$r+i = \sum / \text{oper de } [f,f+i] \forall i \in [f,l)$
__builtin_ffs	unsigned int	Pos. del primer 1 desde la derecha
__builtin_clz	unsigned int	Cant. de ceros desde la izquierda.
__builtin_ctz	unsigned int	Cant. de ceros desde la derecha.
__builtin_popcount	unsigned int	Cant. de 1's en x.
__builtin_parity	unsigned int	1 si x es par, 0 si es impar.
__builtin_XXXXXXll	unsigned ll	= pero para long long's.