

## Template del Rejunte

```

1 #include <bits/stdc++.h>
2 #define sqr(a) ((a)*(a))
3 #define rsz resize
4 #define forr(i,a,b) for(int i=(a);i<(b);i++)
5 #define forn(i,n) forr(i,0,n)
6 #define dforn(i,n) for(int i=n-1;i>=0;i--)
7 #define forall(it,v) for(auto it=v.begin();it!=v.end();it++)
8 #define foreach(i, v) for(auto i:v)
9 #define sz(c) ((int)c.size())
10 #define zero(v) memset(v, 0, sizeof(v))
11 #define pb push_back
12 #define mp make_pair
13 #define lb lower_bound
14 #define ub upper_bound
15 #define fst first
16 #define snd second
17 #define PI 3.1415926535897932384626
18
19 using namespace std;
20
21 typedef long long ll;
22 typedef pair<int,int> ii;
23 typedef vector<int> vi;
24 typedef vector<ii> vii;
25
26 int main()
27 {
28     // agregar g++ -DREJUNTE en compilacin
29     #ifdef REJUNTE
30         freopen("input", "r", stdin);
31         // freopen("output","w", stdout);
32     #endif
33     ios::sync_with_stdio(false);
34     cin.tie(NULL);
35     cout.tie(NULL);
36     return 0;
37 }

```

## Estructuras de datos

### Set Mejorado

Esto solo compila en C++11.

```

1 #include <ext/pb_ds/assoc_container.hpp>

```

```

2 #include <ext/pb_ds/tree_policy.hpp>
3 using namespace __gnu_pbds;
4 //<key,mapped type,comparator,...>
5 typedef tree<int,null_type,less<int>,rb_tree_tag,
6     tree_order_statistics_node_update> ordered_set;
7 //find_by_order(i) devuelve iterador al i-esimo elemento
8 //order_of_key(k): devuelve la pos del lower bound de k
9 //Ej: 12, 100, 505, 1000, 10000.
10 //order_of_key(10) == 0, order_of_key(100) == 1,
11 //order_of_key(707) == 3, order_of_key(9999999) == 5

```

### Union Find

```

1 struct UnionFind{
2     vector<int> f,setSize; //the array f contains the parent of each node
3     int cantSets;
4     void init(int n)
5     {
6         f.clear(); setSize.clear();
7         cantSets=n;
8         f.rsz(n,-1);
9         setSize.rsz(n,1);
10    }
11    int comp(int x){return (f[x]==-1?x:f[x]=comp(f[x]));} //O(1)
12    bool join(int i,int j) //devuelve true si ya estaban juntos
13    {
14        bool con=comp(i)==comp(j);
15        if(!con)
16        {
17            cantSets--;
18            setSize[comp(j)]+=setSize[comp(i)];
19            setSize[comp(i)]=setSize[comp(j)]; //no suma, solo asigna
20            f[comp(i)]=comp(j);
21        }
22        return con;
23    }
24 };

```

### Hash Table

```

1 //Compilar: g++ --std=c++11
2 struct Hash{
3     size_t operator()(const ii &a)const
4     {
5         size_t s=hash<int>()(a.fst);
6         return hash<int>()(a.snd)+0x9e3779b9+(s<<6)+(s>>2);
7     }
8     size_t operator()(const vector<int> &v)const

```

```

9   {
10      size_t s=0;
11      for(auto &e : v) s ^= hash<int>()(e)+0x9e3779b9+(s<<6)+(s>>2);
12      return s;
13   }
14 };
15 unordered_set<ii, Hash> s;
16 unordered_map<ii, int, Hash> m; //map<key, value, hasher>

```

## RMQ

### RMQ (static)

Dado un arreglo y una operacion asociativa *idempotente*,  $get(i, j)$  opera sobre el rango  $[i, j]$ .  
 Restriccion:  $LVL \geq \text{ceil}(\log n)$ ; Usar  $[]$  para llenar arreglo y luego  $build()$ .

```

1 struct RMQ{
2     #define LVL 10
3     tipo vec[LVL][1<<(LVL+1)];
4     tipo &operator[](int p){return vec[0][p];}
5     tipo get(int i, int j) { //intervalo [i,j]
6         int p = 31-__builtin_clz(j-i);
7         return min(vec[p][i], vec[p][j-(1<<p)]);
8     }
9     void build(int n) { //O(nlogn)
10        int mp = 31-__builtin_clz(n);
11        forn(p, mp) forn(x, n-(1<<p))
12            vec[p+1][x] = min(vec[p][x], vec[p][x+(1<<p)]);
13    };

```

### RMQ (dynamic)

```

1 //Dado un arreglo y una operacion asociativa con neutro, get(i, j) opera sobre
  //el rango [i, j].
2 #define MAXN 100000
3 #define operacion(x, y) max(x, y)
4 const int neutro=0;
5 struct RMQ{
6     int sz;
7     tipo t[4*MAXN];
8     tipo &operator[](int p){return t[sz+p];}
9     void init(int n){ //O(nlgn)
10        sz = 1 << (32-__builtin_clz(n));
11        forn(i, 2*sz) t[i]=neutro;
12    }
13    void updall(){ //O(n)
14        dforn(i, sz) t[i]=operacion(t[2*i], t[2*i+1]);}
15    tipo get(int i, int j){return get(i,j,1,0,sz);}

```

```

16    tipo get(int i, int j, int n, int a, int b){ //O(lgn)
17        if(j<=a || i>=b) return neutro;
18        if(i<=a && b<=j) return t[n];
19        int c=(a+b)/2;
20        return operacion(get(i, j, 2*n, a, c), get(i, j, 2*n+1, c, b));
21    }
22    void set(int p, tipo val){ //O(lgn)
23        for(p+=sz; p>0 && t[p]!=val;){
24            t[p]=val;
25            p/=2;
26            val=operacion(t[p*2], t[p*2+1]);
27        }
28    }
29 }rmq;
30 //Usage:
31 cin >> n; rmq.init(n); forn(i, n) cin >> rmq[i]; rmq.updall();

```

### RMQ (lazy)

```

1 //Dado un arreglo y una operacion asociativa con neutro, get(i, j) opera sobre
  //el rango [i, j].
2 typedef int Elem; //Elem de los elementos del arreglo
3 typedef int Alt; //Elem de la alteracion
4 #define operacion(x,y) x+y
5 const Elem neutro=0; const Alt neutro2=0;
6 #define MAXN 100000 //Cambiar segun el N del problema
7 struct RMQ{
8     int sz;
9     Elem t[4*MAXN];
10    Alt dirty[4*MAXN]; //las alteraciones pueden ser de distinto Elem
11    Elem &operator[](int p){return t[sz+p];}
12    void init(int n){ //O(nlgn)
13        sz = 1 << (32-__builtin_clz(n));
14        forn(i, 2*sz) t[i]=neutro;
15        forn(i, 2*sz) dirty[i]=neutro2;
16    }
17    void push(int n, int a, int b){ //propaga el dirty a sus hijos
18        if(dirty[n]!=0){
19            t[n]+=dirty[n]*(b-a); //altera el nodo
20            if(n<sz){
21                dirty[2*n]+=dirty[n];
22                dirty[2*n+1]+=dirty[n];
23            }
24            dirty[n]=0;
25        }
26    }
27    Elem get(int i, int j, int n, int a, int b){ //O(lgn)

```

```

28     if(j<=a || i>=b) return neutro;
29     push(n, a, b);//corrige el valor antes de usarlo
30     if(i<=a && b<=j) return t[n];
31     int c=(a+b)/2;
32     return operacion(get(i, j, 2*n, a, c), get(i, j, 2*n+1, c, b));
33 }
34 Elem get(int i, int j){return get(i,j,1,0,sz);}
35 //altera los valores en [i, j) con una alteracion de val
36 void alterar(Alt val, int i, int j, int n, int a, int b){//0(lgn)
37     push(n, a, b);
38     if(j<=a || i>=b) return;
39     if(i<=a && b<=j){
40         dirty[n]+=val;
41         push(n, a, b);
42         return;
43     }
44     int c=(a+b)/2;
45     alterar(val, i, j, 2*n, a, c), alterar(val, i, j, 2*n+1, c, b);
46     t[n]=operacion(t[2*n], t[2*n+1]);//por esto es el push de arriba
47 }
48 void alterar(Alt val, int i, int j){alterar(val,i,j,1,0,sz);}
49 }rmq;

```

### RMQ (persistente)

```

1  typedef int tipo;
2  tipo oper(const tipo &a, const tipo &b){
3      return a+b;
4  }
5  struct node{
6      tipo v; node *l,*r;
7      node(tipo v):v(v), l(NULL), r(NULL) {}
8      node(node *l, node *r) : l(l), r(r){
9          if(!l) v=r->v;
10         else if(!r) v=l->v;
11         else v=oper(l->v, r->v);
12     }
13 };
14 node *build (tipo *a, int tl, int tr) {//modificar para que tome tipo a
15     if (tl+1==tr) return new node(a[tl]);
16     int tm=(tl + tr)>>1;
17     return new node(build(a, tl, tm), build(a, tm, tr));
18 }
19 node *update(int pos, int new_val, node *t, int tl, int tr){
20     if (tl+1==tr) return new node(new_val);
21     int tm=(tl+tr)>>1;
22     if(pos < tm) return new node(update(pos, new_val, t->l, tl, tm), t->r);

```

```

23     else return new node(t->l, update(pos, new_val, t->r, tm, tr));
24 }
25 tipo get(int l, int r, node *t, int tl, int tr){
26     if(l==tl && tr==r) return t->v;
27     int tm=(tl + tr)>>1;
28     if(r<=tm) return get(l, r, t->l, tl, tm);
29     else if(l>=tm) return get(l, r, t->r, tm, tr);
30     return oper(get(l, tm, t->l, tl, tm), get(tm, r, t->r, tm, tr));
31 }

```

### BIGInt

```

1  #define BASEXP 6
2  #define BASE 1000000
3  #define LMAX 1000
4  struct bint{
5      int l;
6      ll n[LMAX];
7      bint(ll x=0){
8          l=1;
9          forn(i, LMAX){
10             if (x) l=i+1;
11             n[i]=x%BASE;
12             x/=BASE;
13         }
14     }
15     bint(string x){
16         l=(x.size()-1)/BASEXP+1;
17         fill(n, n+LMAX, 0);
18         ll r=1;
19         forn(i, sz(x)){
20             n[i / BASEXP] += r * (x[x.size()-1-i]-'0');
21             r*=10; if(r==BASE)r=1;
22         }
23     }
24     void out(){
25         cout << n[l-1];
26         dforn(i, l-1) printf("%6.6llu", n[i]);//6=BASEXP!
27     }
28     void invar(){
29         fill(n+1, n+LMAX, 0);
30         while(l>1 && !n[l-1]) l--;
31     }
32 };
33 bint operator+(const bint&a, const bint&b){
34     bint c;

```

```

36     c.l = max(a.l, b.l);
37     ll q = 0;
38     forn(i, c.l) q += a.n[i]+b.n[i], c.n[i]=q %BASE, q/=BASE;
39     if(q) c.n[c.l++] = q;
40     c.invar();
41     return c;
42 }
43 pair<bint, bool> lresta(const bint& a, const bint& b)    // c = a - b
44 {
45     bint c;
46     c.l = max(a.l, b.l);
47     ll q = 0;
48     forn(i, c.l) q += a.n[i]-b.n[i], c.n[i]=(q+BASE) %BASE, q=(q+BASE)/BASE-1;
49     c.invar();
50     return make_pair(c, !q);
51 }
52 bint& operator-= (bint& a, const bint& b){return a=lresta(a, b).first;}
53 bint operator- (const bint&a, const bint&b){return lresta(a, b).first;}
54 bool operator< (const bint&a, const bint&b){return !lresta(a, b).second;}
55 bool operator<= (const bint&a, const bint&b){return lresta(b, a).second;}
56 bool operator==(const bint&a, const bint&b){return a <= b && b <= a;}
57 bint operator*(const bint&a, ll b){
58     bint c;
59     ll q = 0;
60     forn(i, a.l) q += a.n[i]*b, c.n[i] = q %BASE, q/=BASE;
61     c.l = a.l;
62     while(q) c.n[c.l++] = q %BASE, q/=BASE;
63     c.invar();
64     return c;
65 }
66 bint operator*(const bint&a, const bint&b){
67     bint c;
68     c.l = a.l+b.l;
69     fill(c.n, c.n+b.l, 0);
70     forn(i, a.l){
71         ll q = 0;
72         forn(j, b.l) q += a.n[i]*b.n[j]+c.n[i+j], c.n[i+j] = q %BASE, q/=BASE;
73         c.n[i+b.l] = q;
74     }
75     c.invar();
76     return c;
77 }
78 pair<bint, ll> ldiv(const bint& a, ll b){// c = a / b ; rm = a % b
79     bint c;
80     ll rm = 0;
81     dforn(i, a.l){

```

```

82         rm = rm * BASE + a.n[i];
83         c.n[i] = rm / b;
84         rm %= b;
85     }
86     c.l = a.l;
87     c.invar();
88     return make_pair(c, rm);
89 }
90 bint operator/(const bint&a, ll b){return ldiv(a, b).first;}
91 ll operator%(const bint&a, ll b){return ldiv(a, b).second;}
92 pair<bint, bint> ldiv(const bint& a, const bint& b){
93     bint c;
94     bint rm = 0;
95     dforn(i, a.l){
96         if (rm.l==1 && !rm.n[0])
97             rm.n[0] = a.n[i];
98         else{
99             dforn(j, rm.l) rm.n[j+1] = rm.n[j];
100            rm.n[0] = a.n[i];
101            rm.l++;
102        }
103        ll q = rm.n[b.l] * BASE + rm.n[b.l-1];
104        ll u = q / (b.n[b.l-1] + 1);
105        ll v = q / b.n[b.l-1] + 1;
106        while (u < v-1){
107            ll m = (u+v)/2;
108            if (b*m <= rm) u = m;
109            else v = m;
110        }
111        c.n[i]=u;
112        rm-=b*u;
113    }
114    c.l=a.l;
115    c.invar();
116    return make_pair(c, rm);
117 }
118 bint operator/(const bint&a, const bint&b){return ldiv(a, b).first;}
119 bint operator%(const bint&a, const bint&b){return ldiv(a, b).second;}

```

## Algoritmos

### Longest Increasing Subsequence

```

1 //Para non-increasing, cambiar comparaciones y revisar busq binaria
2 //Given an array, paint it in the least number of colors so that each color
  turns to a non-increasing subsequence.

```

```

3 //Solution:Min number of colors=Length of the longest increasing subsequence
4 int N, a[MAXN]; //secuencia y su longitud
5 ii d[MAXN+1]; //d[i]=ultimo valor de la subsecuencia de tamaño i
6 int p[MAXN]; //padres
7 vector<int> R; //respuesta
8 void rec(int i){
9     if(i== -1) return;
10    R.push_back(a[i]);
11    rec(p[i]);
12 }
13 int lis(){ //O(nlogn)
14     d[0] = ii(-INF, -1); forn(i, N) d[i+1]=ii(INF, -1);
15     forn(i, N){
16         int j = upper_bound(d, d+N+1, ii(a[i], INF))-d;
17         if (d[j-1].first < a[i]&& a[i] < d[j].first){
18             p[i]=d[j-1].second;
19             d[j] = ii(a[i], i);
20         }
21     }
22     R.clear();
23     dforn(i, N+1) if(d[i].first!=INF){
24         rec(d[i].second); //reconstruir
25         reverse(R.begin(), R.end());
26         return i; //longitud
27     }
28     return 0;
29 }

```

## Mo's

 $O(q * \sqrt{n})$ 

```

1 int n, sq;
2 struct Qu{ //queries [l, r]
3     //intervalos cerrado abiertos !!! importante!!
4     int l, r, id;
5 }qs[MAXN];
6 int ans[MAXN], curans; //ans[i]=ans to ith query
7 bool bymos(const Qu &a, const Qu &b){
8     if(a.l/sq!=b.l/sq) return a.l<b.l;
9     return (a.l/sq)&1? a.r<b.r : a.r>b.r;
10 }
11 void mos(){
12     forn(i, t) qs[i].id=i;
13     sort(qs, qs+t, bymos);
14     int cl=0, cr=0;
15     sq=sqrt(n);
16     curans=0;

```

```

17     forn(i, t){ //intervalos cerrado abiertos !!! importante!!
18         Qu &q=qs[i];
19         while(cl>q.l) add(--cl);
20         while(cr<q.r) add(cr++);
21         while(cl<q.l) remove(cl++);
22         while(cr>q.r) remove(--cr);
23         ans[q.id]=curans;
24     }
25 }

```

## Strings

### KMP

```

1 vector<int> b; //back table b[i] maximo borde de [0..i]
2 void kmppre(string &P) //by gabina with love
3 {
4     b.clear();
5     b.rsz(P.size());
6     int i =0, j=-1; b[0]=-1;
7     while(i<sz(P))
8     {
9         while(j>=0 && P[i] != P[j]) j=b[j];
10        i++, j++;
11        b[i] = j;
12    }
13 }
14 void kmp(string &T, string &P) //Text, Pattern -- O(|T|+|P|)
15 {
16     kmppre(P);
17     int i=0, j=0;
18     while(i<sz(T))
19     {
20         while(j>=0 && T[i]!=P[j]) j=b[j];
21         i++, j++;
22         if(j==sz(P))
23         {
24             //P encontrado en T empezando en [i-j,i)
25             j=b[j];
26         }
27     }
28 }

```

### Z function

```

1 //z[i]=length of longest substring starting from s[i] that is prefix of s

```

```

2 vector<int> z;
3 void zFunction(string &s)
4 {
5     int n=s.size();
6     for(int i=1,l=0,r=0;i<n;i++)
7     {
8         if(i<=r)
9             z[i]=min(r-i+1,z[i-1]);
10        while(i+z[i]<n && s[z[i]]==s[i+z[i]])
11            z[i]++;
12        if(i+z[i]-1>r)
13            l=i, r=i+z[i]-1;
14    }
15 }
16 void match(string &T,string &P) //Text, Pattern -- O(|T|+|P|)
17 {
18     string s=P;
19     s+='$'; //here append a character that is not present in T
20     s.append(T);
21     z.clear();
22     z.rsz(s.size(),0);
23     zFunction(s);
24     forr(i,P.size()+1,s.size())
25         if(z[i]==P.size()); //match found, idx = i-P.size()-1
26 }

```

## Trie

```

1 struct trie{
2     map<char, trie> m;
3     void add(const string &s, int p=0)
4     {
5         if(s[p]) m[s[p]].add(s, p+1);
6     }
7     void dfs()
8     {
9         //Do stuff
10        forall(it, m)
11            it->second.dfs();
12    }
13 };

```

## Manacher

```

1 string s;
2 int d1[MAXN]; //d1[i]=long del maximo palindromo impar con centro en i
3 int d2[MAXN]; //d2[i]=analogo pero para longitud par
4 //0 1 2 3 4

```

```

5 //a a b a a <--d1[2]=3
6 //a a a a <--d2[2]=2 (estan uno antes)
7 void manacher() // O(|S|) - find longest palindromic substring
8 {
9     int l=0, r=-1, n=sz(s);
10    forn(i, n)
11    {
12        int k=(i>r? 1 : min(d1[l+r-i], r-i));
13        while(i+k<n && i-k>=0 && s[i+k]==s[i-k]) ++k;
14        d1[i] = k--;
15        if(i+k > r) l=i-k, r=i+k;
16    }
17    l=0, r=-1;
18    forn(i, n)
19    {
20        int k=(i>r? 0 : min(d2[l+r-i+1], r-i+1))+1;
21        while(i+k-1<n && i-k>=0 && s[i+k-1]==s[i-k]) k++;
22        d2[i] = --k;
23        if(i+k-1 > r) l=i-k, r=i+k-1;
24    }
25 }

```

## Aho Corasick

```

1 struct Trie{
2     map<char, Trie> next;
3     Trie* tran[256]; //transiciones del automata
4     int idhoja, szhoja; //id de la hoja o 0 si no lo es
5     //link lleva al sufijo mas largo, nxthoja lleva al mas largo pero que es hoja
6     Trie *padre, *link, *nxthoja;
7     char pch; //caracter que conecta con padre
8     //Trie(): tran(), idhoja(), padre(), link() {}
9     //coment linea de arriba porque me daba errores usarla.
10    void insert(const string &s, int id=1, int p=0) //id>0!!!
11    {
12        if(p<sz(s))
13        {
14            Trie &ch=next[s[p]];
15            tran[(int)s[p]]=&ch;
16            ch.padre=this, ch.pch=s[p];
17            ch.insert(s, id, p+1);
18        }
19        else idhoja=id, szhoja=sz(s);
20    }
21    Trie* get_link()
22    {
23        if(!link)

```

```

24 {
25     if(!padre) link=this;//es la raiz
26     else if(!padre->padre) link=padre;//hijo de la raiz
27     else link=padre->get_link()->get_tran(pch);
28 }
29 return link;
30 }
31 Trie* get_tran(int c)
32 {
33     if(!tran[c]) tran[c] = !padre? this : this->get_link()->get_tran(c);
34     return tran[c];
35 }
36 Trie *get_nxthoja()
37 {
38     if(!nxthoja) nxthoja = get_link()->idhoja? link : link->nxthoja;
39     return nxthoja;
40 }
41 void print(int p)
42 {
43     if(idhoja) cout << "found_" << idhoja << "_at_position_" << p-szhoja <<
44         endl;
45     if(get_nxthoja()) get_nxthoja()->print(p);
46 }
47 void matching(const string &s, int p=0) //0(|s| + tamao palabras)
48 {
49     print(p); if(p<sz(s)) get_tran(s[p])->matching(s, p+1);
50 }

```

## Geometría

### Punto

```

1 struct pto{
2     double x, y;
3     pto(double x=0, double y=0):x(x),y(y){}
4     pto operator+(pto a){return pto(x+a.x, y+a.y);}
5     pto operator-(pto a){return pto(x-a.x, y-a.y);}
6     pto operator+(double a){return pto(x+a, y+a);}
7     pto operator*(double a){return pto(x*a, y*a);}
8     pto operator/(double a){return pto(x/a, y/a);}
9     //dot product, producto interno:
10    double operator*(pto a){return x*a.x+y*a.y;}
11    //module of the cross product or vectorial product:
12    //if a is less than 180 clockwise from b, a^b>0
13    double operator^(pto a){return x*a.y-y*a.x;}

```

```

14 //returns true if this is at the left side of line qr
15 bool left(pto q, pto r){return ((q-*this)^(r-*this))>0;}
16 bool operator<(const pto &a) const{return x<a.x-EPS || (abs(x-a.x)<EPS && y<a
    .y-EPS);}
17 bool operator==(pto a){return abs(x-a.x)<EPS && abs(y-a.y)<EPS;}
18 double norm(){return sqrt(x*x+y*y);}
19 double norm_sq(){return x*x+y*y;}
20 };
21 double dist(pto a, pto b){return (b-a).norm();}
22 typedef pto vec;
23
24 double angle(pto a, pto o, pto b){
25     pto oa=a-o, ob=b-o;
26     return atan2(oa^ob, oa*ob);}
27
28 //rotate p by theta rads CCW w.r.t. origin (0,0)
29 pto rotate(pto p, double theta){
30     return pto(p.x*cos(theta)-p.y*sin(theta),
31         p.x*sin(theta)+p.y*cos(theta));
32 }

```

### Orden Radial de Puntos

```

1 struct Cmp{//orden total de puntos alrededor de un punto r
2     pto r;
3     Cmp(pto r):r(r) {}
4     int cuad(const pto &a) const{
5         if(a.x > 0 && a.y >= 0)return 0;
6         if(a.x <= 0 && a.y > 0)return 1;
7         if(a.x < 0 && a.y <= 0)return 2;
8         if(a.x >= 0 && a.y < 0)return 3;
9         assert(a.x ==0 && a.y==0);
10        return -1;
11    }
12    bool cmp(const pto&p1, const pto&p2)const{
13        int c1 = cuad(p1), c2 = cuad(p2);
14        if(c1==c2) return p1.y*p2.x<p1.x*p2.y;
15        else return c1 < c2;
16    }
17    bool operator()(const pto&p1, const pto&p2) const{
18        return cmp(pto(p1.x-r.x,p1.y-r.y),pto(p2.x-r.x,p2.y-r.y));
19    }
20 };

```

### Linea

```

1 int sgn(ll x){return x<0? -1 : !!x;}
2 struct line{

```



```

3   line() {}
4   double a,b,c;//Ax+By=C
5   //pto MUST store float coordinates!
6   line(double a, double b, double c):a(a),b(b),c(c){}
7   // TO DO chequear porque paso problema metiendo negativo el C (-(todo el
   calculo como esta))
8   line(pto p, pto q): a(q.y-p.y), b(p.x-q.x), c(a*p.x+b*p.y) {}
9   int side(pto p){return sgn(l1(a) * p.x + l1(b) * p.y - c);}
10 };
11 bool parallels(line l1, line l2){return abs(l1.a*l2.b-l2.a*l1.b)<EPS;}
12 pto inter(line l1, line l2){//intersection
13     double det=l1.a*l2.b-l2.a*l1.b;
14     if(abs(det)<EPS) return pto(INF, INF);//parallels
15     return pto(l2.b*l1.c-l1.b*l2.c, l1.a*l2.c-l2.a*l1.c)/det;
16 }

```

### Segmento

```

1 struct segm{
2     pto s,f;
3     segm(pto s, pto f):s(s), f(f) {}
4     pto closest(pto p) {//use for dist to point
5         double l2 = dist_sq(s, f);
6         if(l2==0.) return s;
7         double t=((p-s)*(f-s))/l2;
8         if (t<0.) return s;//not write if is a line
9         else if(t>1.)return f;//not write if is a line
10        return s+((f-s)*t);
11    }
12    bool inside(pto p){return abs(dist(s, p)+dist(p, f)-dist(s, f))<EPS;}
13 };
14
15 //NOTA: Si los segmentos son colineales solo devuelve un punto de interseccion
16 pto inter(segm s1, segm s2){
17     if(s1.inside(s2.s)) return s2.s; //Fix cuando son colineales
18     if(s1.inside(s2.f)) return s2.f; //Fix cuando son colineales
19     pto r=inter(line(s1.s, s1.f), line(s2.s, s2.f));
20     if(s1.inside(r) && s2.inside(r)) return r;
21     return pto(INF, INF);
22 }

```

### Rectangulo

```

1 struct rect{
2     //lower-left and upper-right corners
3     pto lw, up;
4 };
5 //returns if there's an intersection and stores it in r

```

```

6 bool inter(rect a, rect b, rect &r){
7     r.lw=pto(max(a.lw.x, b.lw.x), max(a.lw.y, b.lw.y));
8     r.up=pto(min(a.up.x, b.up.x), min(a.up.y, b.up.y));
9     //check case when only a edge is common
10    return r.lw.x<r.up.x && r.lw.y<r.up.y;
11 }

```

### Circulo

```

1 vec perp(vec v){return vec(-v.y, v.x);}
2 line bisector(pto x, pto y){
3     line l=line(x, y); pto m=(x+y)/2;
4     return line(-l.b, l.a, -l.b*m.x+l.a*m.y);
5 }
6 struct Circle{
7     pto o;
8     double r;
9     Circle(pto x, pto y, pto z){
10         o=inter(bisector(x, y), bisector(y, z));
11         r=dist(o, x);
12     }
13     pair<pto, pto> ptosTang(pto p){
14         pto m=(p+o)/2;
15         tipo d=dist(o, m);
16         tipo a=r*r/(2*d);
17         tipo h=sqrt(r*r-a*a);
18         pto m2=o+(m-o)*a/d;
19         vec per=perp(m-o)/d;
20         return make_pair(m2-per*h, m2+per*h);
21     }
22 };
23 //finds the center of the circle containing p1 and p2 with radius r
24 //as there may be two solutions swap p1, p2 to get the other
25 bool circle2PtsRad(pto p1, pto p2, double r, pto &c){
26     double d2=(p1-p2).norm_sq(), det=r*r/d2-0.25;
27     if(det<0) return false;
28     c=(p1+p2)/2+perp(p2-p1)*sqrt(det);
29     return true;
30 }
31 #define sqr(a) ((a)*(a))
32 #define feq(a,b) (fabs((a)-(b))<EPS)
33 pair<tipo, tipo> ecCuad(tipo a, tipo b, tipo c){//a*x*x+b*x+c=0
34     tipo dx = sqrt(b*b-4.0*a*c);
35     return make_pair((-b + dx)/(2.0*a),(-b - dx)/(2.0*a));
36 }
37 pair<pto, pto> interCL(Circle c, line l){
38     bool sw=false;

```



```

39  if((sw=freq(0,1.b))) {
40  swap(1.a, 1.b);
41  swap(c.o.x, c.o.y);
42  }
43  pair<tipo, tipo> rc = ecCuad(
44  sqr(1.a)+sqr(1.b),
45  2.0*1.a*1.b*c.o.y-2.0*(sqr(1.b)*c.o.x+1.c*1.a),
46  sqr(1.b)*(sqr(c.o.x)+sqr(c.o.y)-sqr(c.r))+sqr(1.c)-2.0*1.c*1.b*c.o.y
47  );
48  pair<pto, pto> p( pto(rc.first, (1.c - 1.a * rc.first) / 1.b),
49  pto(rc.second, (1.c - 1.a * rc.second) / 1.b) );
50  if(sw){
51  swap(p.first.x, p.first.y);
52  swap(p.second.x, p.second.y);
53  }
54  return p;
55 }
56 pair<pto, pto> interCC(Circle c1, Circle c2){
57  line l;
58  l.a = c1.o.x-c2.o.x;
59  l.b = c1.o.y-c2.o.y;
60  l.c = (sqr(c2.r)-sqr(c1.r)+sqr(c1.o.x)-sqr(c2.o.x)+sqr(c1.o.y)
61  -sqr(c2.o.y))/2.0;
62  return interCL(c1, l);
63 }

```

### Area de poligono

```

1  double area(vector<pto> &p){//0(sz(p))
2  double area=0;
3  forn(i, sz(p)) area+=p[i]^p[(i+1)%sz(p)];
4  //if points are in clockwise order then area is negative
5  return abs(area)/2;
6  }
7  //Area ellipse = M_PI*a*b where a and b are the semi axis lengths
8  //Area triangle = sqrt(s*(s-a)(s-b)(s-c)) where s=(a+b+c)/2

```

### Punto en poligono

```

1  //checks if v is inside of P, using ray casting
2  //works with convex and concave.
3  //excludes boundaries, handle it separately using segment.inside()
4  bool inPolygon(pto v, vector<pto>& P) {
5  bool c = false;
6  forn(i, sz(P)){
7  int j=(i+1)%sz(P);
8  if((P[j].y>v.y) != (P[i].y > v.y) &&
9  (v.x < (P[i].x - P[j].x) * (v.y-P[j].y) / (P[i].y - P[j].y) + P[j].x))

```

```

10  c = !c;
11  }
12  return c;
13 }

```

### Punto en Poligono Convexo

$O(\log n)$

```

1  void normalize(vector<pto> &pt) //delete collinear points first!
2  {
3  //this makes it clockwise:
4  if(pt[2].left(pt[0], pt[1])) reverse(pt.begin(), pt.end());
5  int n=sz(pt), pi=0;
6  forn(i, n)
7  if(pt[i].x<pt[pi].x || (pt[i].x==pt[pi].x && pt[i].y<pt[pi].y))
8  pi=i;
9  vector<pto> shift(n);//puts pi as first point
10  forn(i, n) shift[i]=pt[(pi+i)%n];
11  pt.swap(shift);
12  }
13  bool inPolygon(pto p, const vector<pto> &pt)
14  {
15  //call normalize first!
16  if(p.left(pt[0], pt[1]) || p.left(pt[sz(pt)-1], pt[0])) return false;
17  int a=1, b=sz(pt)-1;
18  while(b-a>1)
19  {
20  int c=(a+b)/2;
21  if(!p.left(pt[0], pt[c])) a=c;
22  else b=c;
23  }
24  return !p.left(pt[a], pt[a+1]);
25  }

```

### Chequeo de Convex

```

1
2  bool isConvex(vector<int> &p){//0(N), delete collinear points!
3  int N=sz(p);
4  if(N<3) return false;
5  bool isLeft=p[0].left(p[1], p[2]);
6  forr(i, 1, N)
7  if(p[i].left(p[(i+1)%N], p[(i+2)%N])!=isLeft)
8  return false;
9  return true; }

```

### Convex Hull

```

1  //stores convex hull of P in S, CCW order

```

```

27 //left must return >=0 to delete collinear points!
28 void CH(vector<pto>& P, vector<pto> &S){
29     S.clear();
30     sort(P.begin(), P.end()); //first x, then y
31     forn(i, sz(P)){ //lower hull
32         while(sz(S)>= 2 && S[sz(S)-1].left(S[sz(S)-2], P[i])) S.pop_back();
33         S.pb(P[i]);
34     }
35     S.pop_back();
36     int k=sz(S);
37     dforn(i, sz(P)){ //upper hull
38         while(sz(S) >= k+2 && S[sz(S)-1].left(S[sz(S)-2], P[i])) S.pop_back();
39         S.pb(P[i]);
40     }
41     S.pop_back();
42 }

```

### Convex Hull Trick

```

1 struct Line{tipo m,h;};
2 tipo inter(Line a, Line b){
3     tipo x=b.h-a.h, y=a.m-b.m;
4     return x/y+(x%y?!((x>0)^(y>0)):0); //==ceil(x/y)
5 }
6 struct CHT {
7     vector<Line> c;
8     bool mx;
9     int pos;
10    CHT(bool mx=0):mx(mx),pos(0){} //mx=1 si las query devuelven el max
11    inline Line acc(int i){return c[c[0].m>c.back().m? i : sz(c)-1-i];}
12    inline bool irre(Line x, Line y, Line z){
13        return c[0].m>z.m? inter(y, z) <= inter(x, y)
14            : inter(y, z) >= inter(x, y);
15    }
16    void add(tipo m, tipo h) { //O(1), los m tienen que entrar ordenados
17        if(mx) m*=-1, h*=-1;
18        Line l=(Line){m, h};
19        if(sz(c) && m==c.back().m) { l.h=min(h, c.back().h), c.pop_back(); if(
20            pos) pos--; }
21        while(sz(c)>=2 && irre(c[sz(c)-2], c[sz(c)-1], l)) { c.pop_back(); if(
22            pos) pos--; }
23        c.pb(l);
24    }
25    inline bool fbin(tipo x, int m) {return inter(acc(m), acc(m+1))>x;}
26    tipo eval(tipo x){
27        int n = sz(c);
28        //query con x no ordenados O(lgn)

```

```

27 int a=-1, b=n-1;
28 while(b-a>1) { int m = (a+b)/2;
29     if(fbin(x, m)) b=m;
30     else a=m;
31 }
32 return (acc(b).m*x+acc(b).h)*(mx?-1:1);
33 //query O(1)
34 while(pos>0 && fbin(x, pos-1)) pos--;
35 while(pos<n-1 && !fbin(x, pos)) pos++;
36 return (acc(pos).m*x+acc(pos).h)*(mx?-1:1);
37 }
38 } ch;

```

### Convex Hull Trick Dinamico

```

1 const ll is_query = -(1LL<<62);
2 struct Line {
3     ll m, b;
4     mutable multiset<Line>::iterator it;
5     const Line *succ(multiset<Line>::iterator it) const;
6     bool operator<(const Line& rhs) const {
7         if (rhs.b != is_query) return m < rhs.m;
8         const Line *s=succ(it);
9         if(!s) return 0;
10        ll x = rhs.m;
11        return b - s->b < (s->m - m) * x;
12    }
13 };
14 struct HullDynamic : public multiset<Line>{ // will maintain upper hull for
15     maximum
16     bool bad(iterator y) {
17         iterator z = next(y);
18         if (y == begin()) {
19             if (z == end()) return 0;
20             return y->m == z->m && y->b <= z->b;
21         }
22         iterator x = prev(y);
23         if (z == end()) return y->m == x->m && y->b <= x->b;
24         return (x->b - y->b)*(z->m - y->m) >= (y->b - z->b)*(y->m - x->m);
25     }
26     iterator next(iterator y){return ++y;}
27     iterator prev(iterator y){return --y;}
28     void insert_line(ll m, ll b) {
29         iterator y = insert((Line) { m, b });
30         y->it=y;
31         if (bad(y)) { erase(y); return; }
32         while (next(y) != end() && bad(next(y))) erase(next(y));

```

```

32     while (y != begin() && bad(prev(y))) erase(prev(y));
33 }
34 ll eval(ll x) {
35     Line l = *lower_bound((Line) { x, is_query });
36     return l.m * x + l.b;
37 }
38 }h;
39 const Line *Line::succ(multiset<Line>::iterator it) const{
40     return (++it==h.end())? NULL : &*it;};

```

### Cortar poligono

```

1 //cuts polygon Q along the line ab
2 //stores the left side (swap a, b for the right one) in P
3 void cutPolygon(pto a, pto b, vector<pto> Q, vector<pto> &P){
4     P.clear();
5     forn(i, sz(Q)){
6         double left1=(b-a)^(Q[i]-a), left2=(b-a)^(Q[(i+1)%sz(Q)]-a);
7         if(left1>=0) P.pb(Q[i]);
8         if(left1*left2<0)
9             P.pb(inter(line(Q[i], Q[(i+1)%sz(Q)]), line(a, b)));
10    }
11 }

```

### Intersección de Circulos

```

1 struct event {
2     double x; int t;
3     event(double xx, int tt) : x(xx), t(tt) {}
4     bool operator <(const event &o) const { return x < o.x; }
5 };
6 typedef vector<Circle> VC;
7 typedef vector<event> VE;
8 int n;
9 double cuenta(VE &v, double A,double B)
10 {
11     sort(v.begin(), v.end());
12     double res = 0.0, lx = ((v.empty())?0.0:v[0].x);
13     int contador = 0;
14     forn(i,sz(v))
15     { //interseccion de todos (contador == n), union de todos (contador > 0)
16         //conjunto de puntos cubierto por exacta k Circulos (contador == k)
17         if (contador == n) res += v[i].x - lx;
18         contador += v[i].t, lx = v[i].x;
19     }
20     return res;
21 }
22 // Primitiva de sqrt(r*r - x*x) como funcion double de una variable x.

```

```

23 inline double primitiva(double x,double r)
24 {
25     if (x >= r) return r*r*M_PI/4.0;
26     if (x <= -r) return -r*r*M_PI/4.0;
27     double raiz = sqrt(r*r-x*x);
28     return 0.5 * (x * raiz + r*r*atan(x/raiz));
29 }
30 double interCircle(VC &v)
31 {
32     vector<double> p; p.reserve(v.size() * (v.size() + 2));
33     forn(i,sz(v)) p.push_back(v[i].c.x + v[i].r), p.push_back(v[i].c.x - v[i].r)
34     ;
35     forn(i,sz(v)) forn(j,i)
36     {
37         Circle &a = v[i], b = v[j];
38         double d = (a.c - b.c).norm();
39         if (fabs(a.r - b.r) < d && d < a.r + b.r)
40         {
41             double alfa = acos((sqr(a.r) + sqr(d) - sqr(b.r)) / (2.0 * d * a.r));
42             pto vec = (b.c - a.c) * (a.r / d);
43             p.pb((a.c + rotate(vec, alfa)).x), p.pb((a.c + rotate(vec, -alfa)).x);
44         }
45     }
46     sort(p.begin(), p.end());
47     double res = 0.0;
48     forn(i,sz(p)-1)
49     {
50         const double A = p[i], B = p[i+1];
51         VE ve; ve.reserve(2 * v.size());
52         forn(j,sz(v))
53         {
54             const Circle &c = v[j];
55             double arco = primitiva(B-c.c.x,c.r) - primitiva(A-c.c.x,c.r);
56             double base = c.c.y * (B-A);
57             ve.push_back(event(base + arco,-1));
58             ve.push_back(event(base - arco, 1));
59         }
60         res += cuenta(ve,A,B);
61     }
62     return res;
63 }

```

### Rotar Matriz

```

1 //rotates matrix t 90 degrees clockwise
2 //using auxiliary matrix t2(faster)
3 void rotate()

```

```

4 | {
5 |     forn(x, n) forn(y, n)
6 |         t2[n-y-1][x]=t[x][y];
7 |     memcpy(t, t2, sizeof(t));
8 | }

```

## Matemática

### Identidades

$$\sum_{i=0}^n \binom{n}{i} = 2^n$$

$$\sum_{i=0}^n i \binom{n}{i} = n * 2^{n-1}$$

$$\sum_{i=m}^n i = \frac{n(n+1)}{2} - \frac{m(m-1)}{2} = \frac{(n+1-m)(n+m)}{2}$$

$$\sum_{i=0}^n i = \sum_{i=1}^n i = \frac{n(n+1)}{2}$$

$$\sum_{i=0}^n i^2 = \frac{n(n+1)(2n+1)}{6} = \frac{n^3}{3} + \frac{n^2}{2} + \frac{n}{6}$$

$$\sum_{i=0}^n i(i-1) = \frac{8}{6} \left(\frac{n}{2} + 1\right) (n+1) \text{ (doubles)} \rightarrow \text{Sino ver caso impar y par}$$

$$\sum_{i=0}^n i^3 = \left(\frac{n(n+1)}{2}\right)^2 = \frac{n^4}{4} + \frac{n^3}{2} + \frac{n^2}{4} = \left[\sum_{i=1}^n i\right]^2$$

$$\sum_{i=0}^n i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} = \frac{n^5}{5} + \frac{n^4}{2} + \frac{n^3}{3} - \frac{n}{30}$$

$$\sum_{i=0}^n i^p = \frac{(n+1)^{p+1}}{p+1} + \sum_{k=1}^p \frac{B_k}{p-k+1} \binom{p}{k} (n+1)^{p-k+1}$$

$$r = e - v + k + 1$$

Teorema de Pick: (Area, puntos interiores y puntos en el borde)

$$A = I + \frac{B}{2} - 1$$

### Ec. Caracteristica

$$a_0 T(n) + a_1 T(n-1) + \dots + a_k T(n-k) = 0$$

$$p(x) = a_0 x^k + a_1 x^{k-1} + \dots + a_k$$

Sean  $r_1, r_2, \dots, r_q$  las raíces distintas, de mult.  $m_1, m_2, \dots, m_q$

$$T(n) = \sum_{i=1}^q \sum_{j=0}^{m_i-1} c_{ij} n^j r_i^n$$

Las constantes  $c_{ij}$  se determinan por los casos base.

### Teorema Chino del Resto

$$y = \sum_{j=1}^n (x_j * (\prod_{i=1, i \neq j}^n m_i)^{-1} * \prod_{i=1, i \neq j}^n m_i)$$

```

1 | //Chinese remainder theorem (special case): find z such that
2 | //z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = lcm(m1, m2).
3 | //Return (z, M). On failure, M = -1.
4 | ii chinese_remainder_theorem(int m1, int r1, int m2, int r2)
5 | { //{xx,yy,d} son variables globales usadas en extendedEuclid
6 |     extendedEuclid(m1, m2);
7 |     if (r1%d != r2%d) return make_pair(0,-1);
8 |     return mp(sumMod(xx*r2*m1, yy*r1*m2, m1*m2) / d, m1*m2 / d);
9 | }
10 | //Chinese remainder theorem: find z such that z % m[i] = r[i] for all i.
11 | //Note that the solution is unique modulo M = lcm_i (m[i]).
12 | //Return (z, M). On failure, M = -1.

```

```

13 | //Note that we do not require the a[i]'s to be relatively prime.
14 | ii chinese_remainder_theorem(const vector<int> &m, const vector<int> &r)
15 | {
16 |     ii ret=mp(r[0], m[0]);
17 |     forr(i,1,m.size())
18 |     {
19 |         ret=chinese_remainder_theorem(ret.snd, ret.fst, m[i], r[i]);
20 |         if (ret.snd==-1) break;
21 |     }
22 |     return ret;
23 | }

```

### GCD & LCM

```

1 | int gcd(int a, int b) {return b? gcd(b,a%b) : a;}
2 | int lcm(int a, int b) {return a*(b/gcd(a,b));}

```

### Euclides Extendido

```

1 | //ecuacin diofntica lineal
2 | //sea d=gcd(a,b); la ecuacin a * x + b * y = c tiene soluciones enteras si
3 | //d|c. La siguiente funcin nos sirve para esto. De forma general ser :
4 | //x = x0 + (b/d)n      x0 = xx*c/d
5 | //y = y0 - (a/d)n      y0 = yy*c/d
6 | ll xx,yy,d;
7 | void extendedEuclid(ll a, ll b)  //a * xx + b * yy = d
8 | {
9 |     if (!b) {xx=1; yy=0; d=a; return;}
10 |    extendedEuclid (b,a%b);
11 |    ll x1=yy;
12 |    ll y1=xx-(a/b)*yy;
13 |    xx=x1; yy=y1;
14 | }

```

### Combinatoria

```

1 | void cargarComb()//O(MAXN^2)
2 | {
3 |     forn(i, MAXN+1) //comb[i][k]=i tomados de a k = i!/(k*(i-k)!)
4 |     {
5 |         comb[0][i]=0;
6 |         comb[i][0]=comb[i][i]=1;
7 |         forr(k, 1, i) comb[i][k]=(comb[i-1][k-1]+comb[i-1][k]) %MOD;
8 |     }
9 | }
10 | ll lucas (ll n, ll k, int p)
11 | { //Calcula (n,k)%p teniendo comb[p][p] precalculado.
12 |     ll aux = 1;
13 |     while (n + k)

```

```

14 {
15     aux = (aux * comb[n %p][k %p]) %p;
16     n/=p, k/=p;
17 }
18 return aux;
19 }

```

## Exponenciación de Matrices y Fibonacci

```

1 #define SIZE 350
2 int NN;
3 void mul(double a[SIZE][SIZE], double b[SIZE][SIZE])
4 {
5     double res[SIZE][SIZE] = {{0}};
6     forn(i, NN) forn(j, NN) forn(k, NN) res[i][j] += a[i][k] * b[k][j];
7     forn(i, NN) forn(j, NN) a[i][j] = res[i][j];
8 }
9 void powmat(double a[SIZE][SIZE], int n, double res[SIZE][SIZE])
10 {
11     forn(i, NN) forn(j, NN) res[i][j] = (i == j);
12     while(n)
13     {
14         if(n & 1) mul(res, a), n--;
15         else mul(a, a), n /= 2;
16     }
17 }
18
19 struct M22{    // |a b|
20     tipo a,b,c,d; // |c d| -- TIPO
21     M22 operator*(const M22 &p) const {
22         return (M22){a*p.a+b*p.c, a*p.b+b*p.d, c*p.a+d*p.c, c*p.b+d*p.d};}
23 };
24 M22 operator^(const M22 &p, int n)
25 { //VER COMO SE PUEDE PONER DENTRO DEL STRUCT
26     if(!n) return (M22){1, 0, 0, 1}; //identidad
27     M22 q = p^(n/2); q = q*q;
28     return n % 2 ? p * q : q;
29 }
30
31 ll fibo(ll n) //calcula el fibonacci enesimo en O(logN)
32 {
33     M22 mat = (M22){0, 1, 1, 1}^n;
34     return mat.a*f0+mat.b*f1; //f0 y f1 son los valores iniciales
35 }

```

## Operaciones Modulares

```

1 ll mulMod(ll a, ll b, ll m=MOD) //O(log b)

```

```

2 { //returns (a*b) %c, and minimize overflow
3     ll x=0, y=a%m;
4     while(b>0)
5     {
6         if(b%2==1) x=(x+y)%m;
7         y=(y*2)%m;
8         b/=2;
9     }
10    return x%m;
11 }
12 ll expMod(ll b, ll e, ll m=MOD) //O(log b)
13 {
14     if(!e) return 1;
15     ll q = expMod(b, e/2, m);
16     q = mulMod(q, q, m);
17     return e%2 ? mulMod(b, q, m) : q;
18 }
19 ll sumMod(ll a, ll b, ll m=MOD)
20 {
21     a %= m;
22     b %= m;
23     if(a < 0) a += m;
24     if(b < 0) b += m;
25     return (a+b)%m;
26 }
27 ll difMod(ll a, ll b, ll m=MOD)
28 {
29     a %= m;
30     b %= m;
31     if(a < 0) a += m;
32     if(b < 0) b += m;
33     ll ret = a - b;
34     if(ret < 0) ret += m;
35     return ret;
36 }
37 ll divMod(ll a, ll b, ll m=MOD)
38 {
39     return mulMod(a, inverso(b), m);
40 }

```

## Funciones de Primos

Sea  $n = \prod p_i^{k_i}$ , fact(n) genera un map donde a cada  $p_i$  le asocia su  $k_i$

```

1 #define MAXP 100000 //no necesariamente primo
2 int criba[MAXP+1];
3 void crearCriba()
4 {

```

```

5   int w[] = {4,2,4,2,4,6,2,6};
6   for(int p=25;p<=MAXP;p+=10) criba[p]=5;
7   for(int p=9;p<=MAXP;p+=6) criba[p]=3;
8   for(int p=4;p<=MAXP;p+=2) criba[p]=2;
9   for(int p=7,cur=0;p*p<=MAXP;p+=w[cur++&7]) if (!criba[p])
10  for(int j=p*p;j<=MAXP;j+=(p<<1)) if(!criba[j]) criba[j]=p;
11 }
12 vector<int> primos;
13 void buscarPrimos()
14 {
15     crearCriba();
16     forr (i,2,MAXP+1) if (!criba[i]) primos.push_back(i);
17 }
18
19 //factoriza bien numeros hasta MAXP^2
20 void fact(ll n,map<ll,ll> &f) //0 (cant primos)
21 { //llamar a buscarPrimos antes
22     forall(p, primos){
23         while(!(n %*p))
24         {
25             f[*p]++;//divisor found
26             n/=*p;
27         }
28     }
29     if(n>1) f[n]++;
30 }
31
32 //factoriza bien numeros hasta MAXP
33 void fact2(ll n,map<ll,ll> &f) //0 (lg n)
34 { //llamar a crearCriba antes
35     while (criba[n])
36     {
37         f[criba[n]]++;
38         n/=criba[n];
39     }
40     if(n>1) f[n]++;
41 }
42
43 //Usar asi: divisores(fac, divs, fac.begin()); NO ESTA ORDENADO
44 void divisores(map<ll,ll> &f,vector<ll> &divs,map<ll,ll>::iterator it,ll n=1)
45 {
46     if(it==f.begin()) divs.clear();
47     if(it==f.end())
48     {
49         divs.pb(n);
50         return;

```

```

51     }
52     ll p=it->fst, k=it->snd; ++it;
53     forn(_, k+1) divisores(f, divs, it, n), n*=p;
54 }
55 ll cantDivs(map<ll,ll> &f)
56 {
57     ll ret=1;
58     forall(it, f) ret*=(it->second+1);
59     return ret;
60 }
61 ll sumDivs(map<ll,ll> &f)
62 {
63     ll ret=1;
64     forall(it, f)
65     {
66         ll pot=1, aux=0;
67         forn(i, it->snd+1) aux+=pot, pot*=it->fst;
68         ret*=aux;
69     }
70     return ret;
71 }
72
73 ll eulerPhi(ll n) // con criba: O(lg n)
74 {
75     map<ll,ll> f;
76     fact(n,f);
77     ll ret=n;
78     forall(it, f) ret-=ret/it->first;
79     return ret;
80 }
81 ll eulerPhi2(ll n) // 0 (sqrt n)
82 {
83     ll r = n;
84     forr(i,2,n+1)
85     {
86         if((ll)i*i>n) break;
87         if(n%i==0)
88         {
89             while(n%i==0) n/=i;
90             r -= r/i;
91         }
92     }
93     if (n != 1) r-= r/n;
94     return r;
95 }

```

## Phollard's Rho

```

1 bool es_primo_prob(ll n, int a)
2 {
3     if(n==a) return true;
4     ll s=0,d=n-1;
5     while(d%2==0) s++,d/=2;
6     ll x=expMod(a,d,n);
7     if((x==1) || (x+1==n)) return true;
8     forn(i,s-1)
9     {
10         x=mulMod(x, x, n);
11         if(x==1) return false;
12         if(x+1==n) return true;
13     }
14     return false;
15 }
16 bool rabin (ll n) //devuelve true si n es primo
17 {
18     if(n==1) return false;
19     const int ar[]={2,3,5,7,11,13,17,19,23};
20     forn(j,9) if(!es_primo_prob(n,ar[j])) return false;
21     return true;
22 }
23 ll rho(ll n)
24 {
25     if((n&1)==0) return 2;
26     ll x=2,y=2,d=1;
27     ll c=rand()%n+1;
28     while(d==1)
29     {
30         x=(mulMod(x,x,n)+c)%n;
31         y=(mulMod(y,y,n)+c)%n;
32         y=(mulMod(y,y,n)+c)%n;
33         if(x-y>=0) d=gcd(n,x-y);
34         else d=gcd(n,y-x);
35     }
36     return d==n? rho(n):d;
37 }
38 void factRho (ll n,map<ll,ll> &f) //O (lg n)^3 un solo numero
39 {
40     if (n == 1) return;
41     if (rabin(n))
42     {
43         f[n]++;
44         return;

```

```

45     }
46     ll factor = rho(n);
47     factRho(factor,f);
48     factRho(n/factor,f);
49 }

```

## Inversos

```

1 #define MAXMOD 15485867
2 ll inv[MAXMOD]; //inv[i]*i=1 mod MOD
3 void calc(int p) //O(p)
4 {
5     inv[1]=1;
6     forr(i,2,p) inv[i]=p-((p/i)*inv[p%i])%p;
7 }
8 int inverso(int x) //O(log x)
9 {
10     return expMod(x, eulerPhi(MOD)-2); //si mod no es primo(sacar a mano)
11     return expMod(x, MOD-2); //si mod es primo
12 }

```

## Fracciones

```

1 struct frac{
2     int p,q;
3     frac(int p=0,int q=1):p(p),q(q) {norm();}
4     void norm()
5     {
6         int a=gcd(q,p);
7         if(a) p/=a, q/=a;
8         else q=1;
9         if (q<0) q=-q, p=-p;
10    }
11    frac operator+(const frac& o)
12    {
13        int a=gcd(o.q,q);
14        return frac(p*(o.q/a)+o.p*(q/a),q*(o.q/a));
15    }
16    frac operator-(const frac& o)
17    {
18        int a=gcd(o.q,q);
19        return frac(p*(o.q/a)-o.p*(q/a),q*(o.q/a));
20    }
21    frac operator*(frac o)
22    {
23        int a=gcd(o.p,q), b=gcd(p,o.q);
24        return frac((p/b)*(o.p/a),(q/a)*(o.q/b));
25    }

```



```
26 | frac operator/(frac o)
27 | {
28 |     int a=gcd(o.q,q), b=gcd(p,o.p);
29 |     return frac((p/b)*(o.q/a),(q/a)*(o.p/b));
30 | }
31 | bool operator<(const frac &o) const{return p*o.q < o.p*q;}
32 | bool operator==(frac o){return p==o.p&&q==o.q;}
33 |};
```

Simpson

```
1 | double integral(double a, double b, int n=10000) //0(n), n=cantdiv
2 | {
3 |     double area=0, h=(b-a)/n, fa=f(a), fb;
4 |     forn(i, n)
5 |     {
6 |         fb=f(a+h*(i+1));
7 |         area+=fa+ 4*f(a+h*(i+0.5)) +fb, fa=fb;
8 |     }
9 |     return area*h/6.;
10 | }
```

Tablas y cotas (Primos, Divisores, Factoriales, etc)

Factoriales	
0! = 1	11! = 39.916.800
1! = 1	12! = 479.001.600 (∈ int)
2! = 2	13! = 6.227.020.800
3! = 6	14! = 87.178.291.200
4! = 24	15! = 1.307.674.368.000
5! = 120	16! = 20.922.789.888.000
6! = 720	17! = 355.687.428.096.000
7! = 5.040	18! = 6.402.373.705.728.000
8! = 40.320	19! = 121.645.100.408.832.000
9! = 362.880	20! = 2.432.902.008.176.640.000 (∈ tint)
10! = 3.628.800	21! = 51.090.942.171.709.400.000
max signed tint = 9.223.372.036.854.775.807	
max unsigned tint = 18.446.744.073.709.551.615	

Primos

2	3	5	7	11	13	17	19	23	29	31	37	41	43	47	53	59	61	67	71	73	79	83	89	97	101	103	107	109	113
127	131	137	139	149	151	157	163	167	173	179	181	191	193	197	199	211	223	227	229	233	239								
241	251	257	263	269	271	277	281	283	293	307	311	313	317	331	337	347	349	353	359	367	373								
379	383	389	397	401	409	419	421	431	433	439	443	449	457	461	463	467	479	487	491	499	503								
509	521	523	541	547	557	563	569	571	577	587	593	599	601	607	613	617	619	631	641	643	647								
653	659	661	673	677	683	691	701	709	719	727	733	739	743	751	757	761	769	773	787	797	809								
811	821	823	827	829	839	853	857	859	863	877	881	883	887	907	911	919	929	937	941	947	953								
967	971	977	983	991	997	1009	1013	1019	1021	1031	1033	1039	1049	1051	1061	1063	1069	1087	1	#define	INF	1e9							

Primos cercanos a 10<sup>n</sup>

9941	9949	9967	9973	10007	10009	10037	10039	10061	10067	10069	10079
99961	99971	99989	99991	100003	100019	100043	100049	100057	100069		
999959	999961	999979	999983	1000003	1000033	1000037	1000039				
9999943	9999971	9999973	9999991	10000019	10000079	10000103	10000121				
99999941	99999959	99999971	99999989	100000007	100000037	100000039	100000049				
999999893	999999929	999999937	1000000007	1000000009	1000000021	1000000033					

Cantidad de primos menores que 10<sup>n</sup>

$\pi(10^1) = 4$	$;$	$\pi(10^2) = 25$	$;$	$\pi(10^3) = 168$	$;$	$\pi(10^4) = 1229$	$;$	$\pi(10^5) = 9592$
$\pi(10^6) = 78.498$	$;$	$\pi(10^7) = 664.579$	$;$	$\pi(10^8) = 5.761.455$	$;$	$\pi(10^9) = 50.847.534$		
$\pi(10^{10}) = 455.052,511$	$;$	$\pi(10^{11}) = 4.118.054.813$	$;$	$\pi(10^{12}) = 37.607.912.018$				

Números Catalanés

Utiles para problemas de Combinatoria

$Cat(n) = \frac{\binom{2n}{n}}{n+1} = \frac{(2n)!}{n!(n+1)!}$

Con  $Cat(0) = 1$ .

Diferentes aplicaciones:

- 1. Contar la cantidad de diferentes arboles binarios con  $n$  nodos que se pueden armar.
- 2. Contar las formas en que un polígono convexo de  $n + 2$  lados puede ser triangulado.
- 3. Contar la cantidad de caminos monotonos a lo largo de los lados de una grilla  $n * n$ , que no cruzan la diagonal.
- 4. Contar el número de expresiones que contienen  $n$  pares de paréntesis correctamente colocados

Primeros 25 Catalanés

1	1	2	5	14	42	132	429	1430	4862	16796	58786	208012	742900	2674440	9694845	35357670
129644790	477638700	1767263190	6564120420	24466267020	91482563640	343059613650										
1289904147324	4861946401452															

Grafos

Dijkstra

```

2 int N;
3 #define MAX_V 250001
4 vector<ii> G[MAX_V];
5 //To add an edge use
6 #define add(a, b, w) G[a].pb(make_pair(w, b))
7 ll dijkstra(int s, int t){//O(|E| log |V|)
8     priority_queue<ii, vector<ii>, greater<ii> > Q;
9     vector<ll> dist(N, INF); vector<int> dad(N, -1);
10    Q.push(make_pair(0, s)); dist[s] = 0;
11    while(sz(Q)){
12        ii p = Q.top(); Q.pop();
13        if(p.snd == t) break;
14        forall(it, G[p.snd])
15            if(dist[p.snd]+it->first < dist[it->snd]){
16                dist[it->snd] = dist[p.snd] + it->fst;
17                dad[it->snd] = p.snd;
18                Q.push(make_pair(dist[it->snd], it->snd)); }
19    }
20    return dist[t];
21    if(dist[t]<INF)//path generator
22        for(int i=t; i!=-1; i=dad[i])
23            printf("%d%c", i, (i==s?'\\n':'_'));}

```

### Bellman-Ford

```

1 //Mas lento que Dijkstra, pero maneja arcos con peso negativo
2 vector<ii> G[MAX_N]; //ady. list with pairs (weight, dst)
3 int dist[MAX_N];
4 void bford(int src){//O(VE)
5     dist[src]=0;
6     forn(i, N-1) forn(j, N) if(dist[j]!=INF) forall(it, G[j])
7         dist[it->snd]=min(dist[it->snd], dist[j]+it->fst);
8 }
9
10 bool hasNegCycle(){
11     forn(j, N) if(dist[j]!=INF) forall(it, G[j])
12         if(dist[it->snd]>dist[j]+it->fst) return true;
13     //inside if: all points reachable from it->snd will have -INF distance(do bfs
14     )
15     return false;
16 }

```

### Floyd-Warshall

```

1 // Camino minimo en grafos dirigidos ponderados, en todas las parejas de nodos.
2 //G[i][j] contains weight of edge (i, j) or INF
3 //G[i][i]=0
4 int G[MAX_N][MAX_N];

```

```

5 void floyd(){//O(N^3)
6     forn(k, N) forn(i, N) if(G[i][k]!=INF) forn(j, N) if(G[k][j]!=INF)
7         G[i][j]=min(G[i][j], G[i][k]+G[k][j]);
8     }
9     bool inNegCycle(int v){
10         return G[v][v]<0;}
11     //checks if there's a neg. cycle in path from a to b
12     bool hasNegCycle(int a, int b){
13         forn(i, N) if(G[a][i]!=INF && G[i][i]<0 && G[i][b]!=INF)
14             return true;
15         return false;
16     }

```

### Kruskal

```

1 // Minimum Spanning Tree in O(e log e)
2 bool operator<(const Ar& a, const Ar &b){return a.w<b.w;}
3 vector<Ar> E;
4 ll kruskal(){
5     ll cost=0;
6     sort(E.begin(), E.end()); //ordenar aristas de menor a mayor
7     uf.init(n);
8     forall(it, E){
9         if(uf.comp(it->a)!=uf.comp(it->b)){//si no estan conectados
10             uf.unir(it->a, it->b); //conectar
11             cost+=it->w;
12         }
13     }
14     return cost;
15 }

```

### Prim

```

1 vector<ii> G[MAXN];
2 bool taken[MAXN];
3 priority_queue<ii, vector<ii>, greater<ii> > pq; //min heap
4 void process(int v){
5     taken[v]=true;
6     forall(e, G[v])
7         if(!taken[e->second]) pq.push(*e);
8 }
9 // Minimum Spanning Tree in O(n^2)
10 ll prim(){
11     zero(taken);
12     process(0);
13     ll cost=0;
14     while(sz(pq)){
15         ii e=pq.top(); pq.pop();

```

```

16         if(!taken[e.second]) cost+=e.first, process(e.second);
17     }
18     return cost;
19 }

```

## Kosaraju SCC

Componente Fuertemente Conexa

```

1 #define MAXN 1000000
2 vector<int> G[MAXN],gt[MAXN]; //Limpiar si se corre mas de una vez
3 //nodos 0...N-1 ; componentes 0...cantcomp-1
4 int comp[MAXN],N,cantcomp,used[MAXN];
5 stack<int> pila;
6 void add(int a, int b){ G[a].pb(b);gt[b].pb(a);}
7 void dfs1(int nodo)
8 {
9     used[nodo]=1;
10    forall(it,G[nodo]) if(!used[*it]) dfs1(*it);
11    pila.push(nodo);
12 }
13 void dfs2(int nodo)
14 {
15     used[nodo]=2;
16     comp[nodo]=cantcomp-1;
17     forall(it,gt[nodo]) if(used[*it]!=2) dfs2(*it);
18 }
19 void kosaraju()
20 {
21     cantcomp=0;
22     memset(used,0,sizeof(used));
23     forn(i,N) if(!used[i]) dfs1(i);
24     while(!pila.empty())
25     {
26         if(used[pila.top()]!=2)
27         {
28             cantcomp++;
29             dfs2(pila.top());
30         }
31         pila.pop();
32     }
33 }

```

## 2-SAT + Tarjan SCC

```

1 //We have a vertex representing a var and other for his negation.
2 //Every edge stored in G represents an implication. To add an equation of the
   form a||b, use addor(a, b)
3 //MAX=max cant var, n=cant var

```

```

4 #define addor(a, b) (G[neg(a)].pb(b), G[neg(b)].pb(a))
5 vector<int> G[MAX*2];
6 //idx[i]=index assigned in the dfs
7 //lw[i]=lowest index(closer from the root) reachable from i
8 int lw[MAX*2], idx[MAX*2], qidx;
9 stack<int> q;
10 int qcmp, cmp[MAX*2];
11 //verdad[cmp[i]]=valor de la variable i
12 bool verdad[MAX*2+1];
13
14 int neg(int x) { return x>=n? x-n : x+n;}
15 void tjn(int v){
16     lw[v]=idx[v]==qidx;
17     q.push(v), cmp[v]=-2;
18     forall(it, G[v]){
19         if(!idx[*it] || cmp[*it]==-2){
20             if(!idx[*it]) tjn(*it);
21             lw[v]=min(lw[v], lw[*it]);
22         }
23     }
24     if(lw[v]==idx[v]){
25         int x;
26         do{x=q.top(); q.pop(); cmp[x]=qcmp;}while(x!=v);
27         verdad[qcmp]=(cmp[neg(v)]<0);
28         qcmp++;
29     }
30 }
31 //remember to CLEAR G!!!
32 bool satisf(){//O(n)
33     memset(idx, 0, sizeof(idx)), qidx=0;
34     memset(cmp, -1, sizeof(cmp)), qcmp=0;
35     forn(i, n){
36         if(!idx[i]) tjn(i);
37         if(!idx[neg(i)]) tjn(neg(i));
38     }
39     forn(i, n) if(cmp[i]==cmp[neg(i)]) return false;
40     return true;
41 }

```

## Puntos de Articulación

```

1 int N;
2 vector<int> G[1000000];
3 //V[i]=node number(if visited), L[i]= lowest V[i] reachable from i
4 int qV, V[1000000], L[1000000], P[1000000];
5 void dfs(int v, int f){
6     L[v]=V[v]=++qV;

```

```

7   forall(it, G[v])
8       if(!V[*it]){
9           dfs(*it, v);
10          L[v] = min(L[v], L[*it]);
11          P[v] += L[*it] >= V[v];
12      }
13      else if(*it != f)
14          L[v] = min(L[v], V[*it]);
15  }
16  int cantart(){ //0(n)
17      qV=0;
18      zero(V), zero(P);
19      dfs(1, 0); P[1]--;
20      int q=0;
21      forn(i, N) if(P[i]) q++;
22  return q;
23  }

```

### Least Common Ancestor + Climb

```

1  const int MAXN=100001, LOGN=20;
2  //f[v][k] holds the 2^k father of v
3  //L[v] holds the level of v
4  int N, f[MAXN][LOGN], L[MAXN];
5  //call before build:
6  void dfs(int v, int fa=-1, int lvl=0){//generate required data
7      f[v][0]=fa, L[v]=lvl;
8      forall(it, G[v])if(*it!=fa) dfs(*it, v, lvl+1); }
9  void build(){//f[i][0] must be filled previously, 0(nlgn)
10     forn(k, LOGN-1) forn(i, N) f[i][k+1]=f[f[i][k]][k];}
11 #define lg(x) (31-__builtin_clz(x))//=floor(log2(x))
12 int climb(int a, int d){//0(lgn)
13     if(!d) return a;
14     dforn(i, lg(L[a])+1) if(1<<i<=d) a=f[a][i], d-=1<<i;
15     return a;}
16 int lca(int a, int b){//0(lgn)
17     if(L[a]<L[b]) swap(a, b);
18     a=climb(a, L[a]-L[b]);
19     if(a==b) return a;
20     dforn(i, lg(L[a])+1) if(f[a][i]!=f[b][i]) a=f[a][i], b=f[b][i];
21     return f[a][0]; }
22 int dist(int a, int b) { //returns distance between nodes
23     return L[a]+L[b]-2*L[lca(a, b)];}

```

### Heavy Light Decomposition

```

1  vector<int> G[MAXN];
2  int treesz[MAXN]; //cantidad de nodos en el subarbol del nodo v

```

```

3  int dad[MAXN]; //dad[v]=padre del nodo v
4  void dfs1(int v, int p=-1){ //pre-dfs
5      dad[v]=p;
6      treesz[v]=1;
7      forall(it, G[v]) if(*it!=p){
8          dfs1(*it, v);
9          treesz[v] += treesz[*it];
10     }
11 }
12 //PONER Q EN 0 !!!!
13 int pos[MAXN], q; //pos[v]=posicion del nodo v en el recorrido de la dfs
14 //Las cadenas aparecen continuas en el recorrido!
15 int cantcad;
16 int homecad[MAXN]; //dada una cadena devuelve su nodo inicial
17 int cad[MAXN]; //cad[v]=cadena a la que pertenece el nodo
18 void heavylight(int v, int cur=-1){
19     if(cur==-1) homecad[cur=cantcad++] = v;
20     pos[v]=q++;
21     cad[v]=cur;
22     int mx=-1;
23     forn(i, sz(G[v])) if(G[v][i] != dad[v])
24         if(mx==-1 || treesz[G[v][mx]] < treesz[G[v][i]]) mx=i;
25     if(mx!=-1) heavylight(G[v][mx], cur);
26     forn(i, sz(G[v])) if(i!=mx && G[v][i] != dad[v])
27         heavylight(G[v][i], -1);
28 }
29 //ejemplo de obtener el maximo numero en el camino entre dos nodos
30 //RTA: max(query(low, u), query(low, v)), con low=lca(u, v)
31 //esta funcion va trepando por las cadenas
32 int query(int an, int v){ //0(logn)
33     //si estan en la misma cadena:
34     if(cad[an]==cad[v]) return rmq.get(pos[an], pos[v]+1);
35     return max(query(an, dad[homecad[cad[v]]]),
36               rmq.get(pos[homecad[cad[v]]], pos[v]+1));
37 }

```

### Centroid Decomposition

```

1  vector<int> G[MAXN];
2  bool taken[MAXN]; //poner todos en FALSE al principio!!
3  int padre[MAXN]; //padre de cada nodo en el centroid tree
4
5  int szt[MAXN];
6  void calcsz(int v, int p) {
7      szt[v] = 1;
8      forall(it, G[v]) if (*it!=p && !taken[*it])
9          calcsz(*it, v), szt[v] += szt[*it];

```

```

10 }
11 void centroid(int v=0, int f=-1, int lvl=0, int tam=-1) { //O(nlogn)
12     if(tam===-1) calcsz(v, -1), tam=szt[v];
13     forall(it, G[v]) if(!taken[*it] && sz[*it]>=tam/2)
14         {szt[v]=0; centroid(*it, f, lvl, tam); return;}
15     taken[v]=true;
16     padre[v]=f;
17     forall(it, G[v]) if(!taken[*it])
18         centroid(*it, v, lvl+1, -1);
19 }

```

### Ciclo Euleriano

```

1 int n,m,ars[MAXE], eq;
2 vector<int> G[MAXN]; //fill G,n,m,ars,eq
3 list<int> path;
4 int used[MAXN];
5 bool usede[MAXE];
6 queue<list<int>::iterator> q;
7 int get(int v){
8     while(used[v]<sz(G[v]) && usede[ G[v][used[v]] ]) used[v]++;
9     return used[v];
10 }
11 void explore(int v, int r, list<int>::iterator it){
12     int ar=G[v][get(v)]; int u=v^ars[ar];
13     usede[ar]=true;
14     list<int>::iterator it2=path.insert(it, u);
15     if(u!=r) explore(u, r, it2);
16     if(get(v)<sz(G[v])) q.push(it);
17 }
18 void euler(){
19     zero(used), zero(usede);
20     path.clear();
21     q=queue<list<int>::iterator>();
22     path.push_back(0); q.push(path.begin());
23     while(sz(q)){
24         list<int>::iterator it=q.front(); q.pop();
25         if(used[*it]<sz(G[*it])) explore(*it, *it, it);
26     }
27     reverse(path.begin(), path.end());
28 }
29 void addEdge(int u, int v){
30     G[u].pb(eq), G[v].pb(eq);
31     ars[eq++]=u^v;
32 }

```

### Diametro Árbol

```

1 vector<int> G[MAXN]; int n,m,p[MAXN],d[MAXN],d2[MAXN];
2 int bfs(int r, int *d) {
3     queue<int> q;
4     d[r]=0; q.push(r);
5     int v;
6     while(sz(q)) { v=q.front(); q.pop();
7         forall(it,G[v]) if (d[*it]==-1)
8             d[*it]=d[v]+1, p[*it]=v, q.push(*it);
9     }
10    return v; //ultimo nodo visitado
11 }
12 vector<int> diams; vector<ii> centros;
13 void diametros(){
14     memset(d,-1,sizeof(d));
15     memset(d2,-1,sizeof(d2));
16     diams.clear(), centros.clear();
17     forn(i, n) if(d[i]==-1){
18         int v,c;
19         c=v=bfs(bfs(i, d2), d);
20         forn(_,d[v]/2) c=p[c];
21         diams.pb(d[v]);
22         if(d[v]&1) centros.pb(ii(c, p[c]));
23         else centros.pb(ii(c, c));
24     }
25 }

```

### Componentes Biconexas y Puentes

```

1 vector<int> G[MAXN];
2
3 struct edge{
4     int u,v, comp;
5     bool bridge;
6 };
7 vector<edge> e;
8 void addEdge(int u, int v)
9 {
10     G[u].pb(sz(e)), G[v].pb(sz(e));
11     e.pb((edge){u,v,-1,false});
12 }
13 //d[i]=id de la dfs
14 //b[i]=lowest id reachable from i
15 int d[MAXN], b[MAXN], t;
16 int nbc; //cant componentes
17 int comp[MAXN]; //comp[i]=cant comp biconexas a la cual pertenece i
18 void initDfs(int n)
19 {

```

```

20 zero(G), zero(comp);
21 e.clear();
22 forn(i,n) d[i]=-1;
23 nbc = t = 0;
24 }
25 stack<int> st;
26 void dfs(int u,int pe) //O(n + m)
27 {
28     b[u]=d[u]=t++;
29     comp[u]=(pe!=-1);
30     forall(ne,G[u]) if(*ne!=pe)
31     {
32         int v=e[*ne].u ^ e[*ne].v ^ u;
33         if(d[v]==-1)
34         {
35             st.push(*ne);
36             dfs(v,*ne);
37             if(b[v]>d[u]) e[*ne].bridge=true; // bridge
38             if(b[v]>=d[u]) // art
39             {
40                 int last;
41                 do
42                 {
43                     las=st.top(); st.pop();
44                     e[last].comp=nbc;
45                 }while(last!=*ne);
46                 nbc++;
47                 comp[u]++;
48             }
49             b[u]=min(b[u],b[v]);
50         }
51         else if(d[v]<d[u]) // back edge
52         {
53             st.push(*ne);
54             b[u]=min(b[u], d[v]);
55         }
56     }
57 }

```

### Hungarian

```

1 //Dado un grafo bipartito completo con costos no negativos, encuentra el
  matching perfecto de minimo costo.
2 #define tipo double
3 tipo cost[N][N], lx[N], ly[N], slack[N]; //llenar: cost=matriz de adyacencia
4 int n, max_match, xy[N], yx[N], slackx[N],prev2[N]; //n=cantidad de nodos
5 bool S[N], T[N]; //sets S and T in algorithm

```

```

6 void add_to_tree(int x, int prevx) {
7     S[x] = true, prev2[x] = prevx;
8     forn(y, n) if (lx[x] + ly[y] - cost[x][y] < slack[y] - EPS)
9         slack[y] = lx[x] + ly[y] - cost[x][y], slackx[y] = x;
10 }
11 void update_labels(){
12     tipo delta = INF;
13     forn (y, n) if (!T[y]) delta = min(delta, slack[y]);
14     forn (x, n) if (S[x]) lx[x] -= delta;
15     forn (y, n) if (T[y]) ly[y] += delta; else slack[y] -= delta;
16 }
17 void init_labels(){
18     zero(lx), zero(ly);
19     forn (x,n) forn(y,n) lx[x] = max(lx[x], cost[x][y]);
20 }
21 void augment() {
22     if (max_match == n) return;
23     int x, y, root, q[N], wr = 0, rd = 0;
24     memset(S, false, sizeof(S)), memset(T, false, sizeof(T));
25     memset(prev2, -1, sizeof(prev2));
26     forn (x, n) if (xy[x] == -1){
27         q[wr++] = root = x, prev2[x] = -2;
28         S[x] = true; break; }
29     forn (y, n) slack[y] = lx[root] + ly[y] - cost[root][y], slackx[y] = root;
30     while (true){
31         while (rd < wr){
32             x = q[rd++];
33             for (y = 0; y < n; y++) if (cost[x][y] == lx[x] + ly[y] && !T[y]){
34                 if (yx[y] == -1) break; T[y] = true;
35                 q[wr++] = yx[y], add_to_tree(yx[y], x); }
36             if (y < n) break; }
37         if (y < n) break;
38         update_labels(), wr = rd = 0;
39         for (y = 0; y < n; y++) if (!T[y] && slack[y] == 0){
40             if (yx[y] == -1){x = slackx[y]; break;}
41             else{
42                 T[y] = true;
43                 if (!S[yx[y]]) q[wr++] = yx[y], add_to_tree(yx[y], slackx[y]);
44             }
45             if (y < n) break; }
46         if (y < n){
47             max_match++;
48             for (int cx = x, cy = y, ty; cx != -2; cx = prev2[cx], cy = ty)
49                 ty = xy[cx], yx[cy] = cx, xy[cx] = cy;
50             augment(); }
51 }

```

```

52 | tipo hungarian(){
53 |     tipo ret = 0; max_match = 0, memset(xy, -1, sizeof(xy));
54 |     memset(yx, -1, sizeof(yx)), init_labels(), augment(); //steps 1-3
55 |     forn (x,n) ret += cost[x][xy[x]]; return ret;
56 | }

```

## Dynamic Connectivity

```

1 | struct UnionFind {
2 |     int n, comp;
3 |     vector<int> pre,si,c;
4 |     UnionFind(int n=0):n(n), comp(n), pre(n), si(n, 1) {
5 |         forn(i,n) pre[i] = i; }
6 |     int find(int u){return u==pre[u]?u:find(pre[u]);}
7 |     bool merge(int u, int v)
8 |     {
9 |         if((u=find(u))==v=find(v)) return false;
10 |        if(si[u]<si[v]) swap(u, v);
11 |        si[u]+=si[v], pre[v]=u, comp--, c.pb(v);
12 |        return true;
13 |    }
14 |    int snap(){return sz(c);}
15 |    void rollback(int snap)
16 |    {
17 |        while(sz(c)>snap)
18 |        {
19 |            int v = c.back(); c.pop_back();
20 |            si[pre[v]] -= si[v], pre[v] = v, comp++;
21 |        }
22 |    }
23 |};
24 |enum {ADD,DEL,QUERY};
25 |struct Query {int type,u,v;};
26 |struct DynCon{//bidirectional graphs; create vble as DynCon name(cant_nodos)
27 |    vector<Query> q;
28 |    UnionFind dsu;
29 |    vector<int> match,res;
30 |    map<ii,int> last;//se puede no usar cuando hay identificador para cada arista
31 |        (mejora poco)
32 |    DynCon(int n=0):dsu(n){}
33 |    void add(int u, int v) //to add an edge
34 |    {
35 |        if(u>v) swap(u,v);
36 |        q.pb((Query){ADD, u, v}), match.pb(-1);
37 |        last[ii(u,v)] = sz(q)-1;
38 |    }
39 |    void remove(int u, int v) //to remove an edge

```

```

39 |    {
40 |        if(u>v) swap(u,v);
41 |        q.pb((Query){DEL, u, v});
42 |        int prev = last[ii(u,v)];
43 |        match[prev] = sz(q)-1;
44 |        match.pb(prev);
45 |    }
46 |    void query() //to add a question (query) type of query
47 |    {
48 |        q.pb((Query){QUERY, -1, -1}), match.pb(-1);
49 |    }
50 |    void process() //call this to process queries in the order of q
51 |    {
52 |        forn(i,sz(q)) if (q[i].type == ADD && match[i] == -1) match[i] = sz(q);
53 |        go(0,sz(q));
54 |    }
55 |    void go(int l, int r)
56 |    {
57 |        if(l+1==r)
58 |        {
59 |            if (q[l].type == QUERY)//Aqui responder la query usando el dsu!
60 |                res.pb(dsu.comp);//aqui query=cantidad de componentes conexas
61 |            return;
62 |        }
63 |        int s=dsu.snap(), m = (l+r) / 2;
64 |        forr(i,m,r) if(match[i]!=-1 && match[i]<l) dsu.merge(q[i].u, q[i].v);
65 |        go(l,m);
66 |        dsu.rollback(s);
67 |        s = dsu.snap();
68 |        forr(i,l,m) if(match[i]!=-1 && match[i]>=r) dsu.merge(q[i].u, q[i].v);
69 |        go(m,r);
70 |        dsu.rollback(s);
71 |    }
72 |};

```

## Flow

### Edmond Karp

```

1 | #define MAX_V 1000
2 | #define INF 1e9
3 | //special nodes
4 | #define SRC 0
5 | #define SNK 1
6 | map<int, int> G[MAX_V];//limpiar esto -- unordered_map mejora
7 | //To add an edge use

```



```

8  #define add(a, b, w) G[a][b]=w
9  int f, p[MAX_V];
10 void augment(int v, int minE)
11 {
12     if(v==SRC) f=minE;
13     else if(p[v]!=-1)
14     {
15         augment(p[v], min(minE, G[p[v]][v]));
16         G[p[v]][v]-=f, G[v][p[v]]+=f;
17     }
18 }
19 ll maxflow()//O(min(VE^2,Mf*E))
20 {
21     ll Mf=0;
22     do
23     {
24         f=0;
25         char used[MAX_V]; queue<int> q; q.push(SRC);
26         zero(used), memset(p, -1, sizeof(p));
27         while(sz(q))
28         {
29             int u=q.front(); q.pop();
30             if(u==SNK) break;
31             forall(it, G[u])
32                 if(it->snd>0 && !used[it->fst])
33                     used[it->fst]=true, q.push(it->fst), p[it->fst]=u;
34         }
35         augment(SNK, INF);
36         Mf+=f;
37     }while(f);
38     return Mf;
39 }

```

### Min Cut

```

1  //Suponemos un grafo con el formato definido en Edmond Karp o Push relabel
2  bitset<MAX_V> type,used; //reset this
3  void dfs1(int nodo)
4  {
5      type.set(nodo);
6      forall(it,G[nodo]) if(!type[it->fst] && it->snd>0) dfs1(it->fst);
7  }
8  void dfs2(int nodo)
9  {
10     used.set(nodo);
11     forall(it,G[nodo])
12     {

```

```

13         if(!type[it->fst])
14         {
15             //edge nodo -> (it->fst) pertenece al min_cut
16             //y su peso original era: it->snd + G[it->fst][nodo]
17             //si no existia arista original al revs
18         }
19         else if(!used[it->fst]) dfs2(it->fst);
20     }
21 }
22 void minCut() //antes correr algn maxflow()
23 {
24     dfs1(SRC);
25     dfs2(SRC);
26     return;
27 }

```

### Push Relabel

```

1  #define MAX_V 1000
2  int N;//valid nodes are [0...N-1]
3  #define INF 1e9
4  //special nodes
5  #define SRC 0
6  #define SNK 1
7  map<int, int> G[MAX_V]; //limpiar esto -- unordered_map mejora
8  //To add an edge use
9  #define add(a, b, w) G[a][b]=w
10 ll excess[MAX_V];
11 int height[MAX_V], active[MAX_V], cuenta[2*MAX_V+1];
12 queue<int> Q;
13
14 void enqueue(int v)
15 {
16     if (!active[v] && excess[v] > 0) active[v]=true, Q.push(v);
17 }
18 void push(int a, int b)
19 {
20     int amt = min(excess[a], ll(G[a][b]));
21     if(height[a] <= height[b] || amt == 0) return;
22     G[a][b]-=amt, G[b][a]+=amt;
23     excess[b] += amt, excess[a] -= amt;
24     enqueue(b);
25 }
26 void gap(int k)
27 {
28     forn(v, N)
29     {

```

```

30     if (height[v] < k) continue;
31     cuenta[height[v]]--;
32     height[v] = max(height[v], N+1);
33     cuenta[height[v]]++;
34     enqueue(v);
35 }
36 }
37 void relabel(int v)
38 {
39     cuenta[height[v]]--;
40     height[v] = 2*N;
41     forall(it, G[v])
42     if(it->snd) height[v] = min(height[v], height[it->fst] + 1);
43     cuenta[height[v]]++;
44     enqueue(v);
45 }
46 ll maxflow() //O(V^3)
47 {
48     zero(height), zero(active), zero(cuenta), zero(excess);
49     cuenta[0]=N-1; cuenta[N]=1;
50     height[SRC] = N;
51     active[SRC] = active[SNK] = true;
52     forall(it, G[SRC])
53     {
54         excess[SRC] += it->snd;
55         push(SRC, it->fst);
56     }
57     while(sz(Q))
58     {
59         int v = Q.front(); Q.pop();
60         active[v]=false;
61         forall(it, G[v]) push(v, it->fst);
62         if(excess[v] > 0)
63             cuenta[height[v]] == 1? gap(height[v]):relabel(v);
64     }
65     ll mf=0;
66     forall(it, G[SRC]) mf+=G[it->fst][SRC];
67     return mf;
68 }

```

## Dinic

```

1 struct Edge {
2     int u, v;
3     ll cap, flow;
4     Edge() {}
5     Edge(int u, int v, ll cap): u(u), v(v), cap(cap), flow(0) {}

```

```

6 };
7 struct Dinic {
8     int N;
9     vector<Edge> E;
10    vector<vector<int>>> g;
11    vector<int> d, pt;
12    Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {} //clear and init
13    void addEdge(int u, int v, ll cap)
14    {
15        if (u != v)
16        {
17            E.emplace_back(Edge(u, v, cap));
18            g[u].emplace_back(E.size() - 1);
19            E.emplace_back(Edge(v, u, 0));
20            g[v].emplace_back(E.size() - 1);
21        }
22    }
23    bool BFS(int S, int T)
24    {
25        queue<int> q({S});
26        fill(d.begin(), d.end(), N + 1);
27        d[S] = 0;
28        while(!q.empty())
29        {
30            int u = q.front(); q.pop();
31            if (u == T) break;
32            for (int k: g[u])
33            {
34                Edge &e = E[k];
35                if (e.flow < e.cap && d[e.v] > d[e.u] + 1)
36                {
37                    d[e.v] = d[e.u] + 1;
38                    q.emplace(e.v);
39                }
40            }
41        }
42        return d[T] != N + 1;
43    }
44    ll DFS(int u, int T, ll flow = -1)
45    {
46        if (u == T || flow == 0) return flow;
47        for (int &i = pt[u]; i < g[u].size(); ++i)
48        {
49            Edge &e = E[g[u][i]];
50            Edge &oe = E[g[u][i]^1];
51            if (d[e.v] == d[e.u] + 1)

```

```

52     {
53         ll amt = e.cap - e.flow;
54         if (flow != -1 && amt > flow) amt = flow;
55         if (ll pushed = DFS(e.v, T, amt))
56         {
57             e.flow += pushed;
58             oe.flow -= pushed;
59             return pushed;
60         }
61     }
62 }
63 return 0;
64 }
65 ll maxFlow(int S,int T)
66 {
67     ll total = 0;
68     while(BFS(S, T))
69     {
70         fill(pt.begin(), pt.end(), 0);
71         while (ll flow = DFS(S, T)) total += flow;
72     }
73     return total;
74 }
75 };

```

### Min cost - Max flow

```

1  const int MAXN=10000;
2  typedef ll tf;
3  typedef ll tc;
4  const tf INFFLUJO = 1e14;
5  const tc INFCOSTO = 1e14;
6  struct edge {
7      int u, v;
8      tf cap, flow;
9      tc cost;
10     tf rem() { return cap - flow; }
11 };
12 int nodes; //numero de nodos
13 vector<int> G[MAXN]; // limpiar!
14 vector<edge> e; // limpiar!
15 void addEdge(int u, int v, tf cap, tc cost)
16 {
17     G[u].pb(sz(e)); e.pb((edge){u,v,cap,0,cost});
18     G[v].pb(sz(e)); e.pb((edge){v,u,0,0,-cost});
19 }
20 tc dist[MAXN], mnCost;

```

```

21 int pre[MAXN];
22 tf cap[MAXN], mxFlow;
23 bool in_queue[MAXN];
24 void flow(int s, int t)
25 {
26     zero(in_queue);
27     mxFlow=mnCost=0;
28     while(1)
29     {
30         fill(dist, dist+nodes, INFCOSTO); dist[s] = 0;
31         memset(pre, -1, sizeof(pre)); pre[s]=0;
32         zero(cap); cap[s] = INFFLUJO;
33         queue<int> q; q.push(s); in_queue[s]=1;
34         while(sz(q))
35         {
36             int u=q.front(); q.pop(); in_queue[u]=0;
37             for(auto it:G[u])
38             {
39                 edge &E = e[it];
40                 if(E.rem() && dist[E.v] > dist[u] + E.cost + 1e-9) // ojo EPS
41                 {
42                     dist[E.v]=dist[u]+E.cost;
43                     pre[E.v] = it;
44                     cap[E.v] = min(cap[u], E.rem());
45                     if(!in_queue[E.v]) q.push(E.v), in_queue[E.v]=1;
46                 }
47             }
48         }
49         if (pre[t] == -1) break;
50         mxFlow +=cap[t];
51         mnCost +=cap[t]*dist[t];
52         for (int v = t; v != s; v = e[pre[v]].u)
53         {
54             e[pre[v]].flow += cap[t];
55             e[pre[v]^1].flow -= cap[t];
56         }
57     }
58 }

```

### Utils

#### Convertir string a num e viceversa

```

1  #include <sstream>
2  string num_to_str(int x){
3      ostringstream convert;

```

```
4 |   convert << x;
5 |   return convert.str();
6 | }
7 |
8 | int str_to_num(string x){
9 |     int ret;
10 |    istringstream (x) >> ret;
11 |    return ret;
12 | }
```

Truquitos para entradas/salidas

```
1 | //Cantidad de decimales
2 | cout << setprecision(2) << fixed;
3 | //Rellenar con espacios(para justificar)
4 | cout << setfill('␣') << setw(3) << 2 << endl;
5 | //Leer hasta fin de linea
6 | //  hacer cin.ignore() antes de getline()
7 | while(getline(cin, line)){
8 |     istringstream is(line);
9 |     while(is >> X)
10 |         cout << X << "␣";
11 |     cout << endl;
12 | }
```

Comparación de Double

```
1 | const double EPS = 1e-9;
2 | x == y <=> fabs(x-y) < EPS
3 | x > y <=> x > y + EPS
4 | x >= y <=> x > y - EPS
```

Iterar subconjuntos

```
1 | for(int sbm=bm; sbm; sbm=(sbm-1)&bm)
```

Limites

```
1 | #include <limits>
2 | numeric_limits<T>
3 |     ::max()
4 |     ::min()
5 |     ::epsilon()
```

Mejorar Lectura de Enteros

```
1 | //Solo para enteros positivos
2 | inline void Scanf(int& a)
3 | {
4 |     char c = 0;
```

```
5 | while(c<33) c = getc(stdin);
6 | a = 0;
7 | while(c>33) a = a*10 + c - '0', c = getc(stdin);
8 | }
```

Compilar C++11 con g++

Dos opciones, útil en Linux.

```
1 | g++ -std=c++11 {file} -o {filename}
2 |
3 | g++ -std=c++0x {file} -o {filename}
```

Funciones Utiles

Algo	Params	Función
fill, fill_n	f, l / n, elem	void llena [f, l) o [f,f+n) con elem
lower_bound, upper_bound	f, l, elem	it al primer ultimo donde se puede insertar elem para que quede ordenada
copy	f, l, resul	hace resul+i=f+i $\forall i$
find, find_if, find_first_of	f, l, elem / pred / f2, l2	it encuentra i $\in [f,l)$ tq. i=elem, pred(i), i $\in [f2,l2)$
count, count_if	f, l, elem/pred	cuenta elem, pred(i)
search	f, l, f2, l2	busca [f2,l2) $\in [f,l)$
replace, replace_if	f, l, old / pred, new	cambia old / pred(i) por new
lexicographical_compare	f1,l1,f2,l2	bool con [f1,l1];[f2,l2]
accumulate	f,l,i,[op]	$T = \sum / \text{oper de } [f,l)$
inner_product	f1, l1, f2, i	$T = i + [f1, l1) \cdot [f2, \dots)$
partial_sum	f, l, r, [op]	$r+i = \sum / \text{oper de } [f,f+i] \forall i \in [f,l)$
__builtin_ffs	unsigned int	Pos. del primer 1 desde la derecha
__builtin_clz	unsigned int	Cant. de ceros desde la izquierda.
__builtin_ctz	unsigned int	Cant. de ceros desde la derecha.
__builtin_popcount	unsigned int	Cant. de 1's en x.
__builtin_parity	unsigned int	1 si x es par, 0 si es impar.
__builtin_XXXXXXll	unsigned ll	= pero para long long's.