

# Rules for reasoning about code

CSE 331

University of Washington

Michael Ernst

# Review:

## Forward vs. backward reasoning

**Forward reasoning** is more **intuitive** for most people

- Helps you understand what will happen (simulates the code)

- Introduces facts that may be irrelevant to the goal

  - Set of current facts may get large

- Takes longer to realize that the task is hopeless

**Backward reasoning** is usually more **helpful**

- Helps you understand what should happen

- Given a specific goal, indicates how to achieve it

  - Given an error, gives a test case that exposes it

# Reasoning about code statements

Goal: Convert **assertions about programs** into **logic**

Mechanical process; just follow rules:

- Rule for each type of statement

- Rule for combining/eliminating statements

There is a (forward and backward) rule for each statement in the programming language

- Loops have no rule: you have to *guess a loop invariant*

# Hoare triples: A notation for properties about code



Sir Anthony Hoare

A Hoare triple:  $\{ P \} \text{code} \{ Q \}$

$P$  and  $Q$  are logical statements (about program values)

**code** is Java code

“ $\{ P \} \text{code} \{ Q \}$ ” means “if  $P$  is true and you execute **code**, then  $Q$  is true afterward”

“ $\{ P \} \text{code} \{ Q \}$ ” is a logical formula like “ $x + y = z$ ”

Examples:

“ $1 + 2 = 3$ ” is true

“ $2 + 2 = 5$ ” is false

“ $\{ x > 0 \} x++ \{ x > 1 \}$ ” is true

“ $\{ x < 0 \} x++ \{ x < 0 \}$ ” is false

“ $\{ x > 0 \} x++ \{ x > -5 \}$ ” is true

Is this notation good for forward or for backward reasoning?

# Backward reasoning:

## Assignment

// precondition: ??

$x = e;$

// postcondition: Q

Precondition = Q with all (free) occurrences of x replaced by e

Examples:

// assert: ??

$y = x + 1;$

// assert  $y > 0$

Precondition =  $(x+1) > 0$

// assert: ??

$z = z + 1;$

// assert  $z > 0$

Precondition =  $(z+1) > 0$

Notation: wp for “*weakest* precondition”

$\text{wp}(\text{“}x=e\text{”}, Q) = Q$  with x replaced by e

Weakest = most general

Strongest = most specific

## Aside: weaker and stronger

Strength of an assertion corresponds to logical implication

- “x is stronger than y” corresponds to “ $x \Rightarrow y$ ”
- “x is stronger than y” means x guarantees more than y
- “x is stronger than y” means fewer worlds satisfy x

Suppose that all of the following are true:

- $a \Rightarrow b$      $b \Rightarrow c$      $x \Rightarrow y$      $y \Rightarrow z$
- $\{b\}$  **mycode**  $\{y\}$

Then which of these are true?

- $\{a\}$  **mycode**  $\{y\}$  True
- $\{c\}$  **mycode**  $\{y\}$  False
- $\{b\}$  **mycode**  $\{x\}$  False
- $\{b\}$  **mycode**  $\{z\}$  True

# Method calls

```
// precondition: ??  
x = foo() ;  
// postcondition: Q
```

If the method has no side effects: just like ordinary assignment

<pre>// precondition: ?? <b>x</b> = <b>Math.sqrt</b>(<b>y</b>) ; // postcondition: <b>x</b> = 3 Precondition: (<b>y</b> = 9) and (<b>x</b> = anything)</pre>	<pre>// precondition: ?? <b>x</b> = <b>Math.abs</b>(<b>y</b>) ; // postcondition: <b>x</b> = 22 Precondition: (<b>y</b> = 22 or <b>y</b> = -22)</pre>
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If it has side effects: an assignment to every var method may modify

Use the method specification to determine the new value

```
// precondition: ??    z+1 = 22  
incrementZ() ;    // spec: zpost = zpre + 1  
// postcondition: z = 22
```

# Composition (statement sequences; blocks)

```
// precondition: ??  
S1;      // some statement  
S2;      // another statement  
// postcondition: Q
```

Work from back to front

Precondition =  $\text{wp}(\text{"s1; s2;"}, Q) = \text{wp}(\text{"s1;"}, \text{wp}(\text{"s2;"}, Q))$

Example:

```
// precondition: ??  
x = 0;  
y = x+1;  
// postcondition: y > 0
```

Think of this as:

```
// precondition: ??  
x = 0;  
// postcondition: ??  
// precondition: ??  
y = x+1;  
// postcondition: y > 0
```

Same condition



# If statement example

// precondition: ??

```
if (x < 5) {  
    x = x*x;  
} else {  
    x = x+1;  
}
```

// postcondition:  $x \geq 9$

# If statements

// precondition: ??

**if (b) S1 else S2**

// postcondition: Q

Do case analysis:

$\text{wp}(\text{"if (b) S1 else S2"}, Q)$

$= ( b \Rightarrow \text{wp}(\text{"S1"}, Q)$

$\wedge \neg b \Rightarrow \text{wp}(\text{"S2"}, Q) )$

$= ( b \wedge \text{wp}(\text{"S1"}, Q)$

$\vee \neg b \wedge \text{wp}(\text{"S2"}, Q) )$

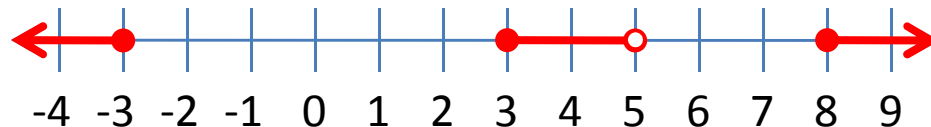
(Why is there no substitution in the condition?)

# If statement example redux

```
// precondition: ??  
if (x < 5) {  
    x = x*x;  
} else {  
    x = x+1;  
}  
// postcondition:  $x \geq 9$ 
```

Precondition

$$\begin{aligned} &= \text{wp}(\text{"if } (x < 5) \{x = x * x;\} \text{ else } \{x = x + 1;\}", x \geq 9) \\ &= (x < 5 \wedge \text{wp}(\text{"x=x*x"}, x \geq 9)) \vee (x \geq 5 \wedge \text{wp}(\text{"x=x+1"}, x \geq 9)) \\ &= (x < 5 \wedge x * x \geq 9) \vee (x \geq 5 \wedge x + 1 \geq 9) \\ &= (x \leq -3) \vee (x \geq 3 \wedge x < 5) \vee (x \geq 8) \end{aligned}$$



# If statements review

## Forward reasoning

```
{P}  
if B  
    {P ∧ B}  
    S1  
    {Q1}  
else  
    {P ∧ !B}  
    S2  
    {Q2}  
{Q1 ∨ Q2}
```

## Backward reasoning

```
{ (B ∧ wp(S1, Q)) ∨ (¬B ∧ wp(S2, Q)) }  
if B  
    {wp(S1, Q)}  
    S1  
    {Q}  
else  
    {wp(S2, Q)}  
    S2  
    {Q}  
{Q}
```

# If statement with one branch empty

// precondition: ??

```
if (x > y) {  
    tmp = x;  
    x = y;  
    y = x;  
}
```

// postcondition:  $x < y$