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### DAA - Tutorial - 3

Q1  
Sol<sup>n</sup>

```
while (low <= high)
{
    mid = (low + high) / 2 ;
    if (arr[mid] == key)
        return true ;
    else if (arr[mid] > key)
        high = mid - 1 ;
    else
        low = mid + 1 ;
}
return false ;
```

Q2

Sol<sup>n</sup>

Iterative insertion sort :

```
for (int i = 1 ; i < n ; i++)
{
    j = i - 1 ;
    n = A[i] ;
    while (j > 0 && A[j] > n)
```

```

{
    A[j+1] = A[j];
    j--;
}
A[j+1] = n;

```

Recursive insertion sort :-

```

void insertion-sort (int arr[], int n)
{
    if (n <= 1)
        return;
    insertion-sort (arr, n-1);
    int last = arr[n-1];
    j = n-2;
    while (j >= 0 && arr[j] > last)
    {
        arr[j+1] = arr[j];
        j--;
    }
    arr[j+1] = last;
}

```

Bubble Sort  $\rightarrow O(n^2)$   
 Insertion Sort  $\rightarrow O(n^2)$   
 Selection Sort  $\rightarrow O(n^2)$   
 Merge sort  $\rightarrow O(n \cdot \log n)$   
 Quick sort  $\rightarrow O(n \log n)$   
 Count sort  $\rightarrow O(n)$   
 Bucket sort  $\rightarrow O(n)$

Q4  
Sol<sup>n</sup>

Online sorting  $\rightarrow$  Insertion Sort  
 Stable sorting  $\rightarrow$  Merge sort, Insertion sort, Bubble sort.  
 Inplace sorting  $\rightarrow$  Bubble sort, Insertion sort, Selection sort.

Q5  
Sol<sup>n</sup>

Iterative Binary Search :  $O(\log n)$

```

while (low <= high)
{
    int mid = (low + high) / 2 ;
    if (arr[mid] == key)
        return true ;
  
```

```

else if (arr[mid] > key)
    high = mid - 1;

```

```

else
{
    low = mid + 1;
}

```

Recursive Binary Search :  $O(\log n)$

```

whileif (low <= high)
{
    int mid = (low + high) / 2;
    if (arr[mid] == key)
        return true;
    else if (arr[mid] > key)
        Binary-search(arr, low, mid - 1);
    else
        Binary-search(arr, mid + 1, high);
    return false;
}

```

Q6  
Sol<sup>n</sup>

$$T(n) = T(n/2) + T(n/2) + c$$



Q7  
Sol<sup>n</sup>

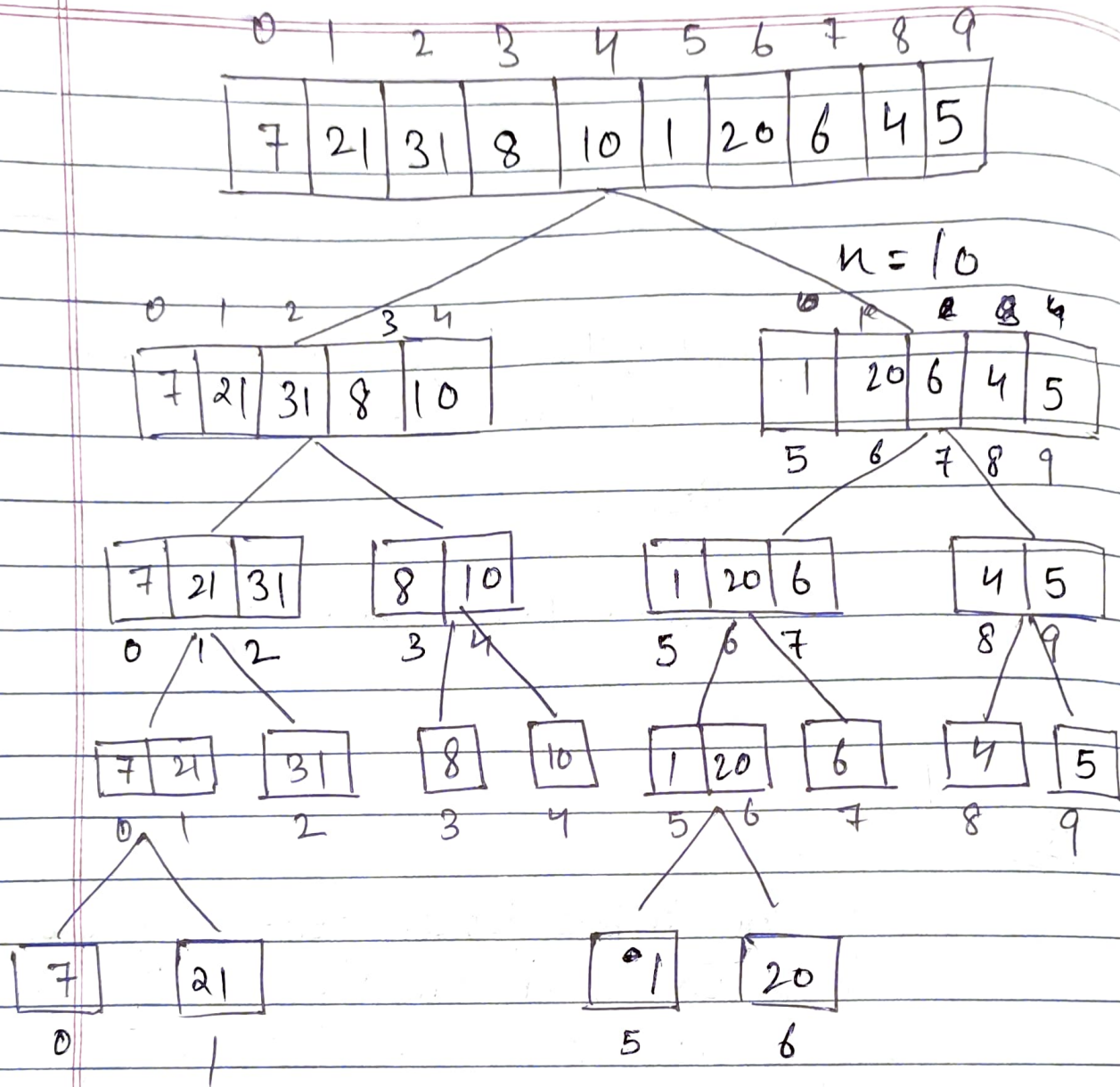
```
map <int, int> m;
for (int i=0; i<arr.size(); i++)
{
    if (m.find(target - arr[i]) != m.end())
        m[arr[i]] = 1;
    else
    {
        cout << i << " " << mp[arr[i]];
    }
}
```

Q8  
Sol<sup>n</sup>

Quick sort is the fastest general purpose sort. In most practical sol<sup>n</sup>, quick sort is the method of choice. If stability is important and space is available, merge sort is the best sort.

Q9  
Sol<sup>n</sup>

Inversion indicators → how far as close the array is from being sorted.



Inversions  $\Rightarrow 31$

Q10

Sol<sup>n</sup>

Worst Case : The worst case occurs when the picked pivot is always an extreme element. This happens when

input array is sorted as reverse sorted and either first or last element is picked as Pivot.

$$O(n^2).$$

Best Case : Best case occurs when Pivot element is the middle element as new to the middle element.

$$O(n \log n)$$

Q11  
Soln

Merge Sort :  $T(n) = 2T\left(\frac{n}{2}\right) + O(n)$

Quick Sort :  $T(n) = 2T(n/2) + n + 1$

Basis	Quick Sort	Merge sort
• Partition	splitting is done in any ratio	array is halved into 2 equal part.
• Works well on	smaller array	fine on any size of array.



Addition of space	less (in-place)	More (Not-in-place)
efficient	works faster on small data set	Has consistent speed for any size array
Stability	Not stable	Stable
Method	Internal sorting method	External sorting method
→	good locality of reference	bad locality of reference

Q12

Sol<sup>n</sup>

void stableselection sort (int a, int n)

{

```

for (int i=0; i<n-1; i++)
{

```

```

    int min = i;

```

```

    for (int j=i+1; j<n; j++)

```

```

        if (a[min] > a[j])

```

```

            min = j;

```



```

int key = a[min];
while (min > i)
{
    a[min] = a[min-1];
    min--;
}
a[i] = key;
}

```

Q13

Sol<sup>n</sup>

```

void bubblesort (int a, int n)
{
    int i, j;
    for (i=0; i<n-1; i++)
        for (j=0; j<n-i-1; j++)
            if (arr[j] > arr[j+1])
            {
                int temp;
                temp = arr[j];
                arr[j] = arr[j+1];
                arr[j+1] = temp;
            }
}

```

Q14Soln

We will be using merge sort algorithm because we can ~~divide~~ ~~the~~ divide the 4 gb data into 4 pockets of 1 gb and sort them seperately and combine them later.

Internal sorting  $\rightarrow$  All the data to sort is stored in memory at all times while sorting is in progress.

External sorting  $\rightarrow$  All the data is stored outside memory and only loaded into memory in small chunks.