
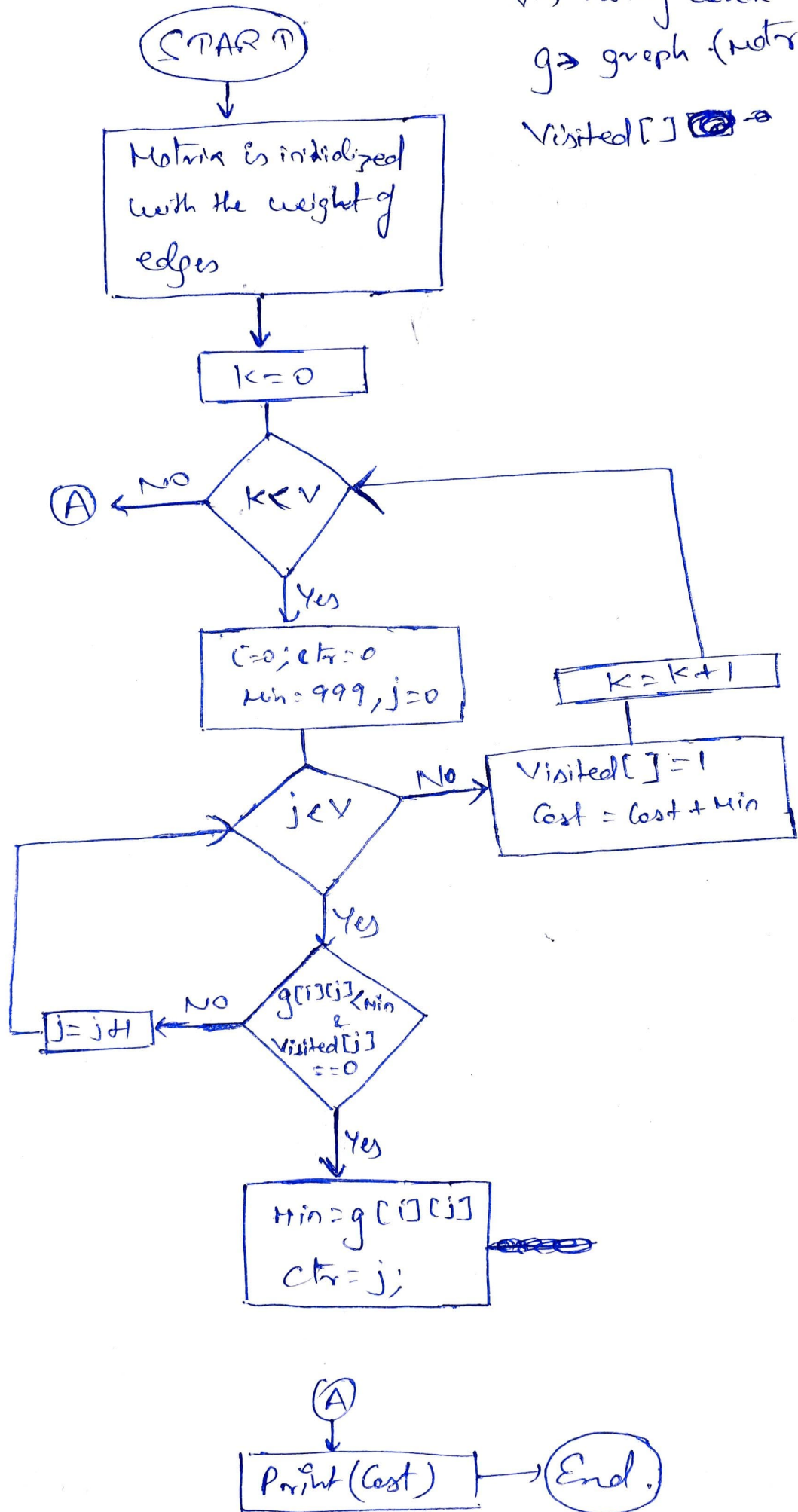


Flow chart (TSP)

$V \rightarrow$ no. of vertices
 $g \rightarrow$ graph (matrix)
 $Visited[]$ 



Time Complexity (Travelling Salesman problem)

Recursive eqⁿ:

$$T(i, S) = \min_{j \in S} (T(i, j) + T(j, S - \{j\})) ; S \neq \phi ;$$

$$= (i, i) ; S = \phi \text{ (base condition)}$$

$T(i, S) \rightarrow$ Vertex $i \rightarrow$ visit each of the element of set S .

$(i, i) \rightarrow$ cost of path i to j .

DP approaches to sub-problem setⁿ approach.

Here, after reaching i^{th} node finding remaining minimum distance to that i^{th} node is the sub problem

On solving eqⁿ: we get $(n-1) \cdot 2^{n-1}$ sub problems
 $\therefore O(n 2^n)$.

Also; each subproblem will take $O(n)$ time

(find the path to remaining not visited nodes)

$$\therefore \text{Total time Complexity} = O(n 2^n) * O(n) \\ = O(n^2 2^n).$$

$$\therefore \text{time Complexity: } \underline{\underline{O(n^2 2^n)}}$$