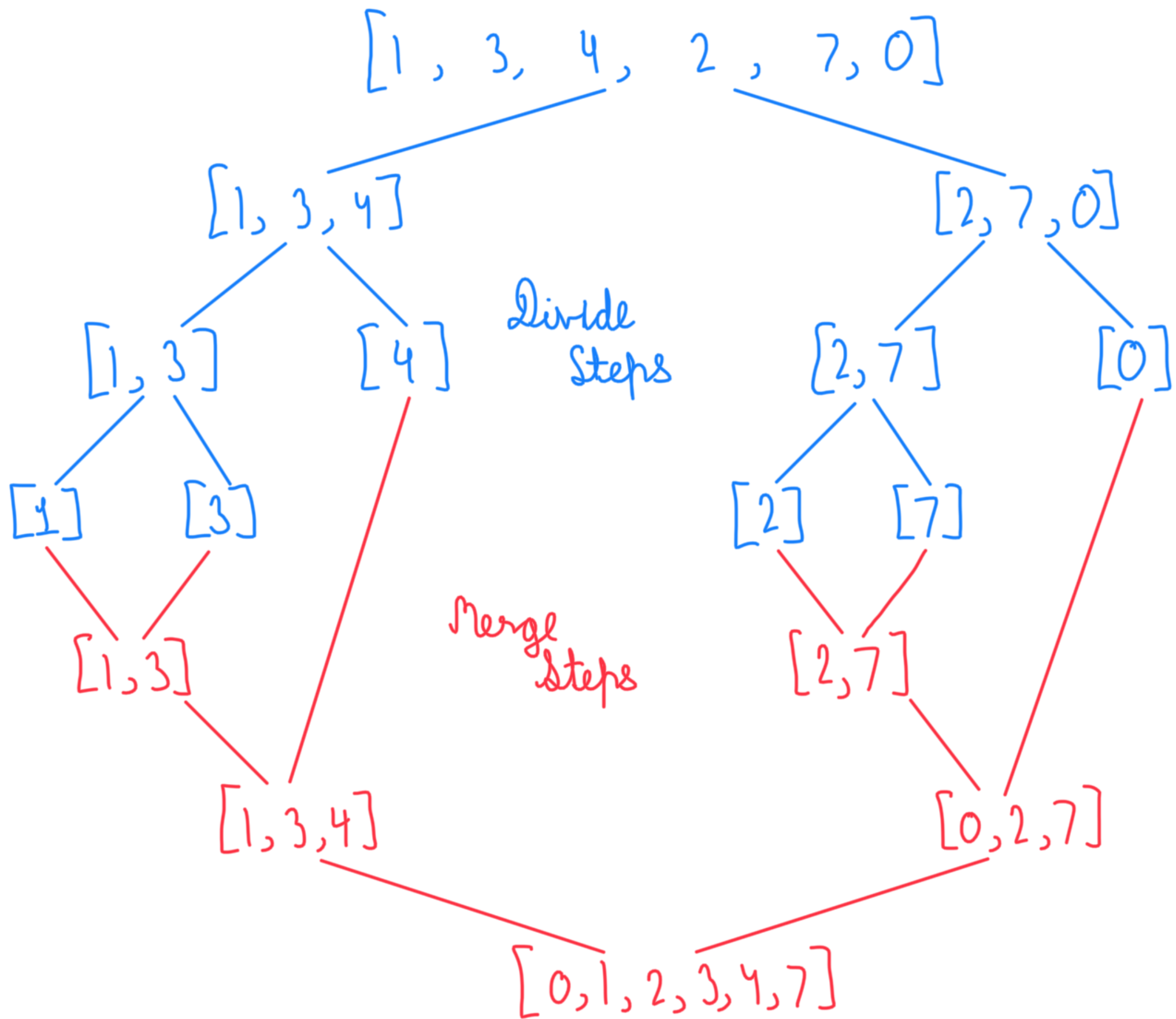
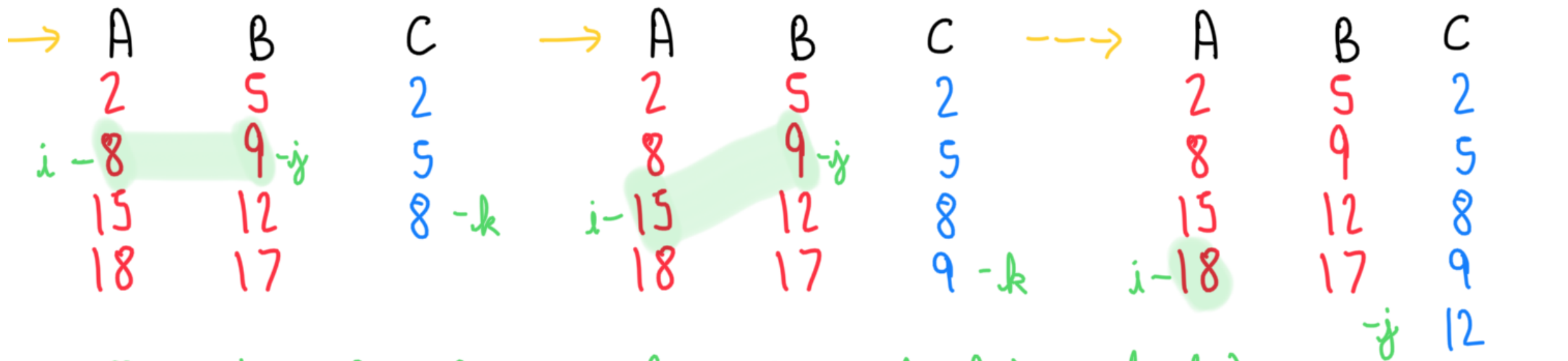
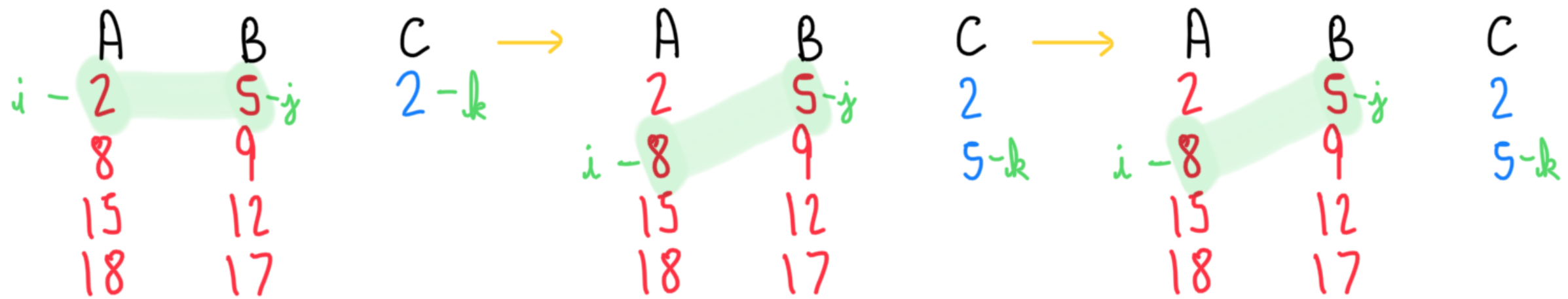


4 Merge Sort



Merge Algorithm of 2 sorted lists:



Merge time complexity: $O(m+n)$ $m = \text{len}(A)$ $n = \text{len}(B)$

```
def merge_sort(arr[lo...hi]):
```

```
    if (len(arr) <= 1):
        return arr
```

$mid = \lfloor (lo + hi) / 2 \rfloor$ — C

$A = \text{merge_sort}(arr[lo \dots mid])$ — $T(n/2)$
 $B = \text{merge_sort}(arr[mid+1 \dots hi])$ — $T(n/2)$
 $\text{return merge}(A, B)$ — $O(n)$

$\text{def merge}(A[i \dots n_1], B[j \dots n_2]):$

$\text{merged} = []$

$\text{while } i \leq n_1 \text{ and } j \leq n_2:$

$\text{if } A[i] \leq B[j]$
 $\text{merged.append}(A[i])$
 $i += 1$

else:
 $\text{merged.append}(B[j])$
 $j += 1$

$\text{return merged} + A[i:] + B[j:]$

Time:

$$T(n) = \begin{cases} b, & \text{if } n=1 \\ 2T(n/2) + n + c, & \text{if } n > 1 \end{cases}$$

$$T(n) = O(n \log n)$$

Best/Worst-Case: $O(n \log n)$

Space: merge list + call stack = $n + \log n = O(n)$

In-place: No, aux space is not $O(1)$ and algo. is recursive.

Stable: Yes, the merge step makes sure that if item in left sub-array is \leq the item in right sub-array, add the item from left sub-array to the merged sorted list. This maintains the relative order of elements with the same value.

Online: No, the merge sort need the entire input array

as it needs to divide it in equal parts.

Code:

```
def mergeSort(arr):
```

```
    if len(arr) <= 1:
```

```
        return arr
```

```
    mid = len(arr)//2
```

```
    A = mergeSort(arr[:mid])
```

```
    B = mergeSort(arr[mid:])
```

```
    return merge(A, B )
```

```
def merge(A, B):
```

```
    i, j = 0, 0
```

```
    merged = []
```

```
    while i < len(A) and j < len(B):
```

```
        if A[ i ] <= B[ j ]:
```

```
            merged.append(A[ i ])
```

```
            i+=1
```

```
        else:
```

```
            merged.append(B[ j ])
```

```
            j+=1
```

```
    return merged + A[i:] + B[j:]
```