Localization in Wireless Sensor Network

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- Location information of nodes in the network is fundamental for a number of reasons:

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- to achieve load balancing in topology control mechanisms; if nodes are densely deployed, geographic information of nodes can be used to selectively shut down some percentage of nodes in each geographic area to conserve energy, and rotate these over time to achieve load balancing

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- to determine the quality of coverage; if node locations are known, the network can keep track of the extent of spatial coverage provided by active sensors at any time

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- If three dimensional coordinates are required, then at least four non-coplanar beacons must be present

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- later provides poor location accuracy subset of nodes have know location a priori and positions of other nodes must be determined using some localization technique

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- the signals may be emitted and measured by the reference nodes, by the unknown nodes, or both

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- Form factor: size of node, multiple sensors, communications requirements, time synchronization

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 - having trouble positioning nodes near the edges
 - fewer range measurements for border nodes
 - sensors outside the main convex body of the network can often prove unlocalizable

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- blind Source Localization: a signal source is localized without any a priori knowledge of the type of signal emitted; done by blind beam-forming which effectively cross-correlates the signals from different receivers
- these techniques generally only work so long as the signals being compared are coherent

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- tradeoff
 - lower resource consumption and equipment cost
 - provide lower accuracy than the detailed information techniques

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- key difference in RFID proximity detection compared with active badges is that the unknown nodes are passive tags, being queried by the reference nodes

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- each node having a simple circular range R in an infinite square mesh of reference nodes spaced a distance d apart
- shown through simulations that, as the overlap ratio R/d is increased from 1 to 4, the average RMS error in localization is reduced from 0.5d to 0.25d

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- if an unknown node hears from several reference nodes, it can determine that it must lie in the geometric region described by the intersection of circles of radius R_{max} centered on these nodes
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- arbitrary shapes can be potentially computed in this manner

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- Identifying codes
- utilizes overlapping coverage regions to provide localization
- referred to as the identifying code construction (ID-CODE) algorithm
- the sensor deployment is planned in such a way as to ensure that each resolvable location is covered by a unique set of sensors

- algorithm runs on a deployment region graph G = (V, E) vertices V represent the different regions, and the edges E represent radio connectivity between regions
- goal is to construct an identifying code for any distinguishable graph, with each vertex in the code corresponding to a region where a reference node must be placed
- Once this is done, by the definition of the identifying code, each location region in the graph will be covered by a unique set of reference nodes
- obtaining a minimal cardinality identifying code is known to be NP-complete

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- triangulation using distance estimates, pattern matching, and sequence decoding used in the large-scale GPS

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- this noise occurs because radio propagation tends to be highly non-uniform in real environments

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- $P_{r,dB}(d)$ is the received power at distance d
- ullet $P(d_0)$ is the received power at some reference distance d_0
- η the path-loss exponent
- $X_{\sigma, {
 m dB}}$ a log-normal random variable with variance σ^2 that accounts for fading effects

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- RSS-based ranging may perform much better in situations where the fading effects can be combated by diversity techniques that take advantage of separate spatio-temporally correlated signal samples

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- the local connectivity information provided by the radio defines an unweighted graph, where the vertices are sensor nodes, and edges represent direct radio links between nodes
- the hop count h_{ij} between sensor nodes s_i and s_j is then defined as the length of the shortest path in the graph between s_i and s_j

- even though RSSI is too inaccurate for many applications, the radio can still be used to assist localization
- hop count to be a useful way to compute inter-node distances
- the local connectivity information provided by the radio defines an unweighted graph, where the vertices are sensor nodes, and edges represent direct radio links between nodes
- the hop count h_{ij} between sensor nodes s_i and s_j is then defined as the length of the shortest path in the graph between s_i and s_j
- if the hop count between s_i and s_j is h_{ij} then the distance between s_i and s_j , d_{ij} , is less than $R*h_{ij}$, where R is again the maximum radio range

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- the speed of the radio signal is much larger than the speed of the acoustic signal, the distance is then simply estimated as $(T_s T_r)$ V_s , where V_s is the speed of the acoustic signal

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- the basic idea is to send a pseudo-random noise sequence as the acoustic signal and use a matched filter for detection
- acoustic TDoA ranging techniques can be very accurate in practical settings

- each node is equipped with a speaker and a microphone
- some systems use ultrasound while others use audible frequencies
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- once they have t_{radio} , t_{sound} , and t_{delay} , the listeners can compute the distance d between themselves and the transmitter using the fact that radio waves travel substantially faster than sound in air

$$d = (s_{radio} - s_{sound}) * (t_{sound} - t_{radio} - t_{delay})$$

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- it can be solved by the use of gradient descent techniques or by iterative successive approximation techniques
- alternative approach provides a numerical solution to an over-determined ($n \ge 3$) linear system

 the over-determined linear system can be obtained as follows: rearranging and squaring terms -

$$x_i^2 + y_i^2 - d_i^2 = 2x_0x_i + 2y_0y_i - (x_0^2 + y_0^2)$$

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- **B** is the (n-1) element column vector whose *i*th term is the expression $x_i^2 + v_i^2 - x_n^2 - v_n^2 - d_i^2 + d_n^2$

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- solving for the above may not directly yield a numerical solution if the matrix A is ill-conditioned
- a recommended approach is to instead use the pseudo-inverse A⁺ of the matrix A

$$\bar{x} = \mathbf{A}^+ \hat{\mathbf{B}}$$

11 June 2010

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- angles can potentially be estimated by using rotating directional beacons, or by using nodes equipped with a phased array of RF or ultrasonic receivers
- involving three rotating reference beacons at the boundary of a sensor network providing localization for all interior nodes
- if the angular information provided to a given reference node can be combined with a good distance estimate to that reference node, then localization can be performed with a single reference using polar coordinate transformation

• in all the localization approaches the network is modeled by a graph

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- in other case, d_{ij} are not given a special value; it is only assumed that $d_{ij} < R R$ is the range of the transmitter of a wireless sensor node
- to describe the positions of the nodes of the network, form a corresponding matrix and store available distance information in matrix $D = \{d_{ij}i, j = 1, 2, \dots, n\}$

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- two types of approaches to distributed localization
- first, beacon-based distributed algorithm: starts with some beacon nodes, few nodes compute their own location using distance measurements and become beacons to help other nodes localize

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- the sub-regions use a peer-to-peer process to merge their local maps into a single global map; this global map approximates the global optimum map

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- only constraints that form convex regions are amenable to representation as an LMI

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- SDP's poor scaling and inability to effectively use range data will likely preclude the algorithm's use in practice

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- Step 2 Run a standard all pairs shortest path algorithm (Dijkstra's, Floyd's) on R to produce a complete matrix of inter-node distances D

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- MDS-MAP estimates improve as ranging improves, it does not use anchor nodes very well, since it effectively ignores their data until stage 4
- its performance lags behind other algorithms as anchor density increases
- poor asymptotic performance, which is $O(n^3)$ on account of stages 2 and 3
- this problem can be partially ameliorated using coordinate system stitching

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- Bounding Box algorithm is a computationally simple method of localizing nodes given their ranges to several beacons

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- otherwise, more mathematically rigorous approaches such as gradient multilateration may be more appropriate

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- after some amount of settling time, each value d_{jk} will be the length of the shortest path between node j and beacon k
- given perfect internode distance estimates, gradient distance estimates will always be longer than (or exactly equal) to corresponding straight line distances
- given imperfect internode distance estimates, gradient based distance estimate can actually be shorter than straight distances

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- the node chooses the centroid of this intersection region as its position estimate

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- the performance of network localization depends very much on the resources and information available within the network

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- other network localization approaches are distributed, often involving the iterative communication of updated location information

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- the objective is then to find the location set that maximizes this likelihood
- the performance of this joint MLE technique has been verified through simulations and experiments to show that localization of the order of 2 meters is possible when there is a high density of unknown nodes, even if there are only a few reference nodes sparsely placed

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- the use of localized unknown nodes as reference nodes can introduce substantial cumulative error

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 approach to tackle the deficiency of iterative multilateration with respect to nodes with insufficient reference neighbors

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- approach to tackle the deficiency of iterative multilateration with respect to nodes with insufficient reference neighbors
- the key insight is to determine collaborative subgraphs within the network that contain reference and unknown nodes in a topology such that
- their positions and inter-node distances can be written as an over-constrained set of quadratic equations with a unique solution for the location of unknown nodes (which can be obtained through gradient descent or local search algorithms)
- used in conjunction with iterative multilateration
- useful in portions of the network where the reference node density is low

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- the average distance per hop depends upon the network density, and is assumed to be known
- the relative performance of these three schemes depends on factors such as the radio range and accuracy of available distance estimates

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- classical MDS is the simplest case of MDS: the proximities of objects are treated as distances in a Euclidean space
- the goal of MDS is to find a configuration of points in a multidimensional space such that the inter-point distances are related to the provided proximities by some transformation (e.g., a linear transformation)

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- for an $n \times n$ P matrix for n points and m dimensions of each point, it can shown that

$$-\frac{1}{2}\left(p_{ij}^2 - \frac{1}{n}\sum_{i=1}^n p_{ij}^2 - \frac{1}{n}\sum_{i=1}^n p_{ij}^2 + \frac{1}{n^2}\sum_{i=1}^n \sum_{j=1}^n p_{ij}^2\right) = \sum_{k=1}^m x_{ik}x_{jk}$$

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- ullet this implies that the summation over k runs from 1 to r instead of m
- this is the best low rank approximation in the least-squares sense
- for example, for a 2D network, we take the first 2 largest eigenvalues and eigenvectors to construct the best 2D approximation

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substituting this approximation back

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$$s \approx \frac{\arg\min}{s} ||As - b||^2$$

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$$A = \begin{bmatrix} \bigtriangledown e(s_0, b_1) \\ \bigtriangledown e(s_0, b_2) \\ \vdots \\ \bigtriangledown e(s_0, b_m) \end{bmatrix}$$

$$b = \begin{bmatrix} -e(s_0, b_1) + \bigtriangledown e(s_0, b_1)s_0) \\ -e(s_0, b_2) + \bigtriangledown e(s_0, b_2)s_0) \\ \vdots \\ -e(s_0, b_m) + \bigtriangledown e(s_0, b_m)s_0) \end{bmatrix}$$

- the right side of equation is in exactly the right form to be solved by an off-the-shelf iterative least squares solver
- the resulting s is a good estimate of the unknown sensor's position, provided b_i and r_i are accurate

① Choose s_0 to be a starting point for the optimization. The choice is somewhat arbitrary, but the centroid \bar{b} is a good one:

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- **1** If $E(s_0) E(s_0') < \epsilon$ then s_0' is the solution, otherwise set $s_0 = s_0'$ and return to Step 2
- many ways to solve the multilateration problem
- the one presented above is equivalent to Newton-Raphson descent on the error function E
- most alternate methods also attempt to minimize squared error using some form of iterative optimization

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- using only ranging data, without anchors or GPS, it can solve for the relative coordinates of a group of sensor nodes
- classical metric MDS: uses only one matrix of dissimilarity or distance information, and metric because the dissimilarity information is quantitative (e.g. distance measurements), as opposed to ordinal

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- the distance measurements d_{ij} must obey the triangular inequality $d_{ij} + d_{ik} \ge d_{jk}$ for all (i, j, k)
- the goal of MDS is to find an assignment of X in low-dimensional space that minimizes a "Stress function" defined as

$$X = \frac{\underset{X}{\operatorname{arg\,min}} \operatorname{Stress}(X)}{X}$$

$$\operatorname{Stress}(X) = \sqrt{\frac{\sum_{i=1}^{n} \sum_{j=1}^{i-1} (d_{ij} - \delta_{ij})^{2}}{\sum_{i=1}^{n} \sum_{j=1}^{i-1} \delta_{ij}^{2}}}$$

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$$d_{jk}^2 = d_{ij}^2 + d_{jk}^2 - 2d_{ij}d_{jk}\cos\theta_{jik}$$

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left side of equation can be rewritten as a dot product

$$(X_j - X_i) \cdot (X_k - X_i) = \frac{1}{2}(d_{ij}^2 + d_{ik}^2 - d_{jk}^2)$$

• if all measurements are perfect, then a good zero-stress way to solve for the positions X is to choose some X_0 from X to be the origin of a coordinate system and construct a matrix $B_{(n-1)\times(n-1)}$ as follows:

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- B is known as the matrix of scalar products; B can be written in terms of X
- $X'_{(n-1)\times m}$ is the matrix X where each of the X_i 's is shifted to have its origin at X_0 : $X'_i = X_i X_0$; then, $X'X'^T = B$
- X' can be solved by taking an eigen decomposition B into an orthonormal matrix of eigenvectors and a diagonal matrix of matching eigen values:

$$B = X'X'^T = UVU^T \quad X' = UV^{1/2}$$

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- above, a single point from our data is chosen to be the origin
- the choice gives X_0 an undue influence on the error of X

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- the squared matrix $D^2 = [d_{ij}^2]$
- \bullet to double center a matrix, subtract the row and column means from each element, then, add the grand mean to each element and finally, multiply by -1/2
- the element-wise formula for double centering is below:

$$b_{ij} = -\frac{1}{2} \left(d_{ij}^2 - \frac{1}{n} \sum_{k=1}^n d_{kj}^2 - \frac{1}{n} \sum_{k=1}^n d_{ik}^2 + \frac{1}{n^2} \sum_{k=1}^n \sum_{l=1}^n d_{kl}^2 \right)$$
$$= \sum_{i=1}^m x_{ia} x_{ja}$$

$$B_{n\times n} = -\frac{1}{2}JD^2J = XX^T$$

$$J_{n\times n} = I_{n\times n} - \frac{1}{n}e^Te$$

$$e_{1\times n} = [1, 1, 1, \dots, 1]$$

- ullet it is an expression for X in terms of D, in m-dimensional space
- if m = n 1, then there is a trivial assignment of $X_1 ... X_n$ that makes Stress(X) = 0
- as m decreases, it turns out that Stress(X) must increase or stay the same; it can not decreases
- the measurements D originates from a 2 or 3 dimensional space
- if the measurements from D are perfect, then there is a zero stress assignment of X when m=2 or 3

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- MDS provides a method of converting a complete matrix of distance measurements to a matching topology in 2-space or 3-space
- this conversion is quite resilient to measurement error, since increased measurement error simply becomes an increase in the stress function

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- classical MDS requires that D contain a distance measurement for all pairs of nodes
- this requirement is impossible to meet with ranging hardware alone in large networks
- implementations of MDS in sensor networks must do pre-processing on measured data to generate D or use coordinate system stitching to distribute the computation

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- Ocompute $X_d = [X_1, X_2, \dots, X_n]^T$ using $X_d = U_d V_d^{1/2}$. $V_d^{1/2}$ can be computed by taking the square root of each of V_d 's diagonal elements

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- **3** (Optional) Transform the X_i 's from X_d into the desired global coordinate space using some coordinate system registration algorithm, these transformed X_i 's are the solution

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- three dimensional version naturally requires four points

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- one caveat: even after a rigid transformation, it is unlikely that the known points will precisely align, since the measurements used to localize the points are likely to have errors
- the best that can be done is a minimization of the misalignment between the two coordinate systems

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 Horn et al approach this problem using squared error; they look for a t, s, and R that meet the following condition

$$(t, s, R) = \arg\min_{t, s, R} \sum_{i=1}^{n} ||e_i||^2$$

$$e_i = x_{r,i} - sRx_{l,i} - t$$



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- Step 1 Compute the centroids of x_l and x_r

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 Step 2 - Shift the points so that they are defined with respect to the centroids:

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the error term can be rewritten as

$$e_i = x'_{r,i} - sRx'_{l,i} - t'$$
 $t' = t - \bar{x}_r + sR\bar{x}_l$

the squared error is minimized when $t' = \mathbf{0}$ independent of s and R

$$t = \bar{x}_r - sR\bar{x}_l$$

$$e_i = \frac{1}{\sqrt{s}}x'_{r,i} - \sqrt{s}Rx'_{l,i}$$

$$(s,R) = \underset{s,R}{\operatorname{arg \, min}} \sum_{i=1}^{n} ||e_{i}||^{2}$$

$$= \underset{s,R}{\operatorname{arg \, min}} \frac{1}{s} \sum_{i=1}^{n} ||x'_{r,i}||^{2} + s \sum_{i=1}^{n} ||r_{l,i}||^{2} - 2 \sum_{i=1}^{n} x'_{r,i} \cdot (Rx'_{l,i})$$

By completing the square in *s* it can be shown that the above equation is minimized when:

$$s = \sqrt{\sum_{i=1}^{n} ||x'_{r,i}||^2 / \sum_{i=1}^{n} ||x'_{l,i}||^2}$$

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Step 3 Use this to compute the optimal scale factor s

$$R = \underset{R}{\operatorname{arg\,min}} 2\left(\sqrt{\left(\sum_{i=1}^{n} ||x'_{r,i}||^{2}\right)\left(\sum_{i=1}^{n} ||x'_{l,i}||^{2}\right)} - \sum_{i=1}^{n} x'_{r,i} \cdot (Rx'_{l,i})\right)$$

• it can be minimized when the following is true:

$$R = \mathop{\mathsf{arg\,max}}_{R} \sum_{i=1}^{n} x'_{r,i} \cdot (Rx'_{l,i})$$

this is the same as:

$$R = \frac{\arg\max}{R} \operatorname{Trace}(R^T M)$$

$$M = \sum_{i=1}^{n} x'_{r,i} (x'_{l,i})^{T}$$

M is a 2 \times 2 or 3 \times 3 matrix, depending on whether the points $x_{I,i}$ and $x_{r,i}$ are two or three dimensional. (Assume M 3 \times 3, the results are similar for the two dimensional case.)

• Step 4 Compute M

• Step 5 Compute the eigen decomposition of M^TM . Find eigenvalues $\lambda_1, \lambda_2, \lambda_3$ and eigenvectors $\hat{u}_1, \hat{u}_2, \hat{u}_3$ such that

$$M^{T}M = \lambda_{1}\hat{u}_{1}\hat{u}_{1}^{T} + \lambda_{2}\hat{u}_{2}\hat{u}_{2}^{T} + \lambda_{3}\hat{u}_{3}\hat{u}_{3}^{T}$$

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• Step 6 Compute $S = (M^T M)^{1/2}$ and $U = MS^{-1}$

$$S = \sqrt{\lambda_1} \hat{u}_1 \hat{u}_1^T + \sqrt{\lambda_2} \hat{u}_2 \hat{u}_2^T + \sqrt{\lambda_3} \hat{u}_3 \hat{u}_3^T$$

$$U = MS^{-1} = M \left(\frac{1}{\sqrt{\lambda_1}} \hat{u}_1 \hat{u}_1^T + \frac{1}{\sqrt{\lambda_2}} \hat{u}_2 \hat{u}_2^T + \frac{1}{\sqrt{\lambda_3}} \hat{u}_3 \hat{u}_3^T \right)$$

M = US and U is orthonormal; $U^TU = I$

• it can be written $Trace(R^TM)$

$$\mathsf{Trace}(R^T US) = \sqrt{\lambda_1} \mathsf{Trace}(R^T U \hat{u}_1 \hat{u}_1^T) + \sqrt{\lambda_2} \mathsf{Trace}(R^T U \hat{u}_2 \hat{u}_2^T) + \sqrt{\lambda_3} \mathsf{Trace}(R^T U \hat{u}_3 \hat{u}_3^T)$$

• Trace($R^T U \hat{u}_i \hat{u}_i^T$) can be rewritten as $(R \hat{u}_i \cdot U \hat{u}_i)$

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$$\mathsf{Trace}(R^T \mathit{US}) \leq \sqrt{\lambda_1} + \sqrt{\lambda_2} + \sqrt{\lambda_3} = \mathsf{Trace}(S)$$

- the maximum value of Trace($R^T US$) occurs when $R^T U = I$ i.e. when R = U
- ullet the rotation R necessary to minimize the error in equation of R

$$R = U = M \left(\frac{1}{\sqrt{\lambda_1}} \hat{u}_1 \hat{u}_1^T + \frac{1}{\sqrt{\lambda_2}} \hat{u}_2 \hat{u}_2^T + \frac{1}{\sqrt{\lambda_3}} \hat{u}_3 \hat{u}_3^T \right)$$

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- \hat{u}_i is a unit vector and U, R are orthonormal transformations, $(R\hat{u}_i \cdot U\hat{u}_i \leq 1 \text{ with equality only when } R\hat{u}_i = U\hat{u}_i$

$$\mathsf{Trace}(R^T \mathit{US}) \leq \sqrt{\lambda_1} + \sqrt{\lambda_2} + \sqrt{\lambda_3} = \mathsf{Trace}(S)$$

- the maximum value of Trace($R^T US$) occurs when $R^T U = I$ i.e. when R = U
- ullet the rotation R necessary to minimize the error in equation of R

$$R = U = M \left(\frac{1}{\sqrt{\lambda_1}} \hat{u}_1 \hat{u}_1^T + \frac{1}{\sqrt{\lambda_2}} \hat{u}_2 \hat{u}_2^T + \frac{1}{\sqrt{\lambda_3}} \hat{u}_3 \hat{u}_3^T \right)$$

• Step 7 Compute R, R is an orthonormal matrix that encapsulates the rotation and possible reflection necessary to transform $x_{l,i}$ into $x_{r,i}$.

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Coordinate System Registration

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- It is straightforward to implement, and gives provably optimal results
- many algorithms depend on coordinate system registration, either to shift a completely localized relative topology into global coordinates, or to stitch together small local topologies into a single consistent coordinate assignment

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$$\left(\begin{array}{cc} l_2R & x_1-x_2 \\ (x_1-x_2)^T & R \end{array}\right) \succeq 0$$

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- all $(i,j) \in E$ where i < j and if j is an anchor are denoted by N_a and unknown is denoted by N_x
- the following constraints must be satisfied:

$$\| a_k - x_j \|^2 = d_{kj}^2 \quad \forall (k, j) \in N_a$$

 $\| x_i - x_j \|^2 = d_{ij}^2 \quad \forall (i, j) \in N_x$

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- when the node distances are known, let $X = [x_1; ...; x_n]^T$ be the matrix in $R^{\dim \times (n)}$ that needs to be determined
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$$(a_k; e_j)^T [I_{\text{dim}}; X^T] [I_{\text{dim}}; X] (a_k; e_j) = d_{kj}^2 \ \forall (k, j) \in N_a$$
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• need to find a symmetric matrix $Y \in R^{\dim \times \dim}$ and X that satisfy the following constraints:

$$(a_k; e_j)^T \begin{pmatrix} I_{\text{dim}} & X \\ X^T & Y \end{pmatrix} (a_k; e_j) = d_{kj}^2 \ \forall (k, j) \in N_a$$

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- this is the SDP formulation of the problem of WSN localization
- the constraint

$$Y = X^T X$$

is relaxed with $Y \succ X^T X$

this condition can be written as

$$Z = \begin{pmatrix} I_{\text{dim}} & X \\ X^T & Y \end{pmatrix} \succeq 0$$



• m known points (anchors) $a_k \in \mathbb{R}^2, k = 1, ..., m$ and n unknown points (sensors) $x_i \in \mathbb{R}^2, j = 1, ..., n$

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- for a pair of two points in N_u , a distance upper bound \bar{r}_{kj} between a_k and x_j or \bar{r}_{ij} between x_i and x_j
- the localization problem is to find x_i s such that

$$\begin{aligned} ||x_{i}-x_{j}||^{2} &= (\hat{d}_{ij})^{2}, ||a_{k}-x_{j}||^{2} &= (\hat{d}_{kj})^{2}, \forall (i,j), (k,j) \in N_{e} \\ ||x_{i}-x_{j}||^{2} &\geq (\underline{r}_{ij})^{2}, ||a_{k}-x_{j}||^{2} \geq (\underline{r}_{kj})^{2}, \forall (i,j), (k,j) \in N_{I} \\ ||x_{i}-x_{j}||^{2} &\leq (\overline{r}_{ij})^{2}, ||a_{k}-x_{j}||^{2} \leq (\overline{r}_{kj})^{2}, \forall (i,j), (k,j) \in N_{U} \end{aligned}$$

 due to measurement error x_js can be chosen to minimize the sum of errors:

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$$\begin{split} & \min \sum_{i,j \in N_{e}, i < j} |\parallel x_{i} - x_{j} \parallel^{2} - (\hat{d}_{ij})^{2} \mid + \sum_{k,j \in N_{e}} |\parallel a_{k} - x_{j} \parallel^{2} - (\hat{d}_{kj})^{2} \mid \\ & + \sum_{i,j \in N_{l}, i < j} (\parallel x_{i} - x_{j} \parallel^{2} - (\underline{r}_{ij})^{2})_{-} + \sum_{k,j \in N_{l}} (\parallel a_{k} - x_{j} \parallel^{2} - (\underline{r}_{kj})^{2})_{-} \\ & + \sum_{i,j \in N_{u}, i < j} (\parallel x_{i} - x_{j} \parallel^{2} - (\overline{r}_{ij})^{2})_{+} + \sum_{k,j \in N_{u}} (\parallel a_{k} - x_{j} \parallel^{2} - (\overline{r}_{kj})^{2})_{+} \end{split}$$

- $(u)_{-}$ and $(u)_{+}$ defined as $(u)_{-} = \max\{0, -u\}$ and $(u)_{+} = \max\{0, u\}$
- by introducing slack variables and relaxing $Y = X^T X$ to $Y \succeq X^T X$, the problem can be rewritten as a standard SDP problem

Thank You