Full selection sort reference →



Sorting Algorithm Cheat Sheet

For coding interviews or computer science classes

A quick reference of the big O (/big-o-notation-time-and-space-complexity) costs and core properties of each sorting algorithm.

| ➤ Expand all () | worst | best | average | space | |
|--|------------------|----------|-------------------------|--------------|--|
| Calastian Cant | worst | best 20 | average | space | |
| ▲ Selection Sort | $O(n^2)$ | $O(n^2)$ | $O(n^2)$ | <i>O</i> (1) | |
| Selection sort works by repeatedly "selecting" | | | 1 2 7 8 4 3 6 | | |
| the next-smallest ele | ement from the u | unsorted | 1 2 7 1 0 1 4 1 3 1 0 1 | | |
| array and moving it to the front. | | | 1 2 3 8 4 7 6 | | |

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Sorting algorithm reference, for coding interviews and computer science...

✓ Expand all () worst best average space

▲ Insertion Sort

worst $O(n^2)$ best

average

space

O(n)

 $O(n^2)$

2

O(1)

Insertion sort works by inserting elements from an unsorted array into a sorted subsection of the array, one item at a time.

3 | 6 | 2

Full insertion sort reference →

▲ Merge Sort

worst

best

average

space

 $O(n \lg n)$

 $O(n \lg n)$

 $O(n \lg n)$

O(n)

Merge sort works by splitting the input in half, recursively sorting each half, and then merging

the sorted halves back together.

4 3 6 8 2 9

Full merge sort reference →

worst

best

average

space

▲ Quicksort

 $O(n^2)$

 $O(n \lg n)$

 $O(n \lg n)$

 $O(\lg n)$

6

space pivot

4

✓ Expand all ()

best

2 3

Quicksort works by recursively dividing the input into two smaller arrays around a pivot item: one half has items smaller than the pivot,

worst

9 < 4 > 4 pivot

9

average

Full quicksort reference →

the other has larger items.

▲ Heapsort

worst $O(n \lg n)$

best

average

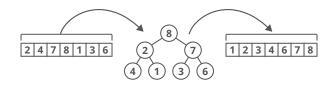
space

O(n)

 $O(n \lg n)$

O(1)

Heapsort is similar to selection sort—we're repeatedly choosing the largest item and moving it to the end of our array. But we use a heap to get the largest item more quickly.



Full heapsort reference →

best worst

average

space

▲ Counting Sort

O(n)

O(n)

O(n)

O(n)

Counting sort works by iterating through the input, counting the number of times each item occurs, and using those counts to compute each item's index in the final, sorted array.

Input: 3 8 3

Counts:

0 3 0 2 0 0 0 0 4's

Full counting sort reference →

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| ∨ Expand all () | worst | best | average | space |
|--|--------------|-------------|---------------------------|-----------------------|
| ▲ Radix Sort | worst $O(n)$ | best $O(n)$ | average $oldsymbol{O(n)}$ | space $O(n)$ |
| Radix sort works by sorting the input numbers one digit at a time. | | | 23 31 81 27 56 | 3 7 5 <u>1</u> |
| Full radix sort refere | nce → | | 31 51 81 23 56 | 27 37 |

Which Sorting Algorithm Should I Use?

It depends. Each algorithm comes with its own set of pros and cons.

- Quicksort is a good default choice. It tends to be fast in practice, and with some small tweaks its dreaded $O(n^2)$ worst-case time complexity becomes very unlikely. A tried and true favorite.
- Heapsort is a good choice if you can't tolerate a worst-case time complexity of O(n²) or need low space costs. The Linux kernel uses heapsort instead of quicksort (https://github.com/torvalds/linux/blob/master/lib/sort.c#L194) for both of those reasons.
- Merge sort is a good choice if you want a stable sorting algorithm (/concept/stable-sort).
 Also, merge sort can easily be extended to handle data sets that can't fit in RAM, where the bottleneck cost is reading and writing the input on disk, not comparing and swapping individual items.
- Radix sort looks fast, with its O(n)worst-case time complexity. But, if you're using it to sort binary numbers, then there's a hidden constant factor that's usually 32 or 64 (depending on how many bits your numbers are). That's often way bigger than $O(\lg(n))$, meaning radix sort

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tends to be slow in practice.

• Counting sort is a good choice in scenarios where there are small number of distinct values to be sorted. This is pretty rare in practice, and counting sort doesn't get much use.

Each sorting algorithm has tradeoffs. You can't have it all.

So you have to know *what*'s *important* in the problem you're working on. How large is your input? How many distinct values are in your input? How much space overhead is acceptable? Can you afford $O(n^2)$ worst-case runtime?

Once you know what's important, you can pick the sorting algorithm that does it best. Being able to compare different algorithms and weigh their pros and cons is the mark of a strong computer programmer and a definite plus when interviewing.

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