

DMA uge 6

Plan for ugen:

- Mandag: Binære søgetræer
- Tirsdag: Sortering i lineær tid
- Fredag: Nedre grænse for sammenligningsbaseret sortering

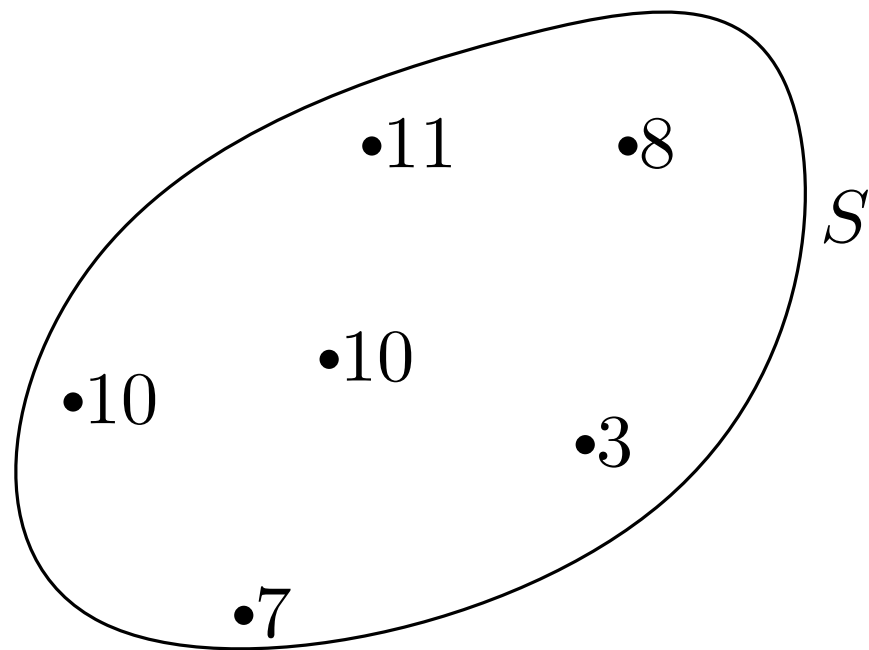
Mikkel Abrahamsen

Ønsket datastruktur

Dynamisk (multi-)mængde S af elementer. Hvert element x har en *nøgle*værdi $x.key$.

Ønskede operationer:

- $\text{Search}(S, k)$: find element x så $k = x.key$.
- $\text{Insert}(S, x)$: tilføj nyt element x .
- $\text{Delete}(S, x)$: fjern x .
- $\text{Predecessor}(S, x)$: find element $y \neq x$ med største nøgle hvor $y.key < x.key$.
- $\text{Successor}(S, x)$: find element $y \neq x$ med mindste nøgle hvor $y.key > x.key$.



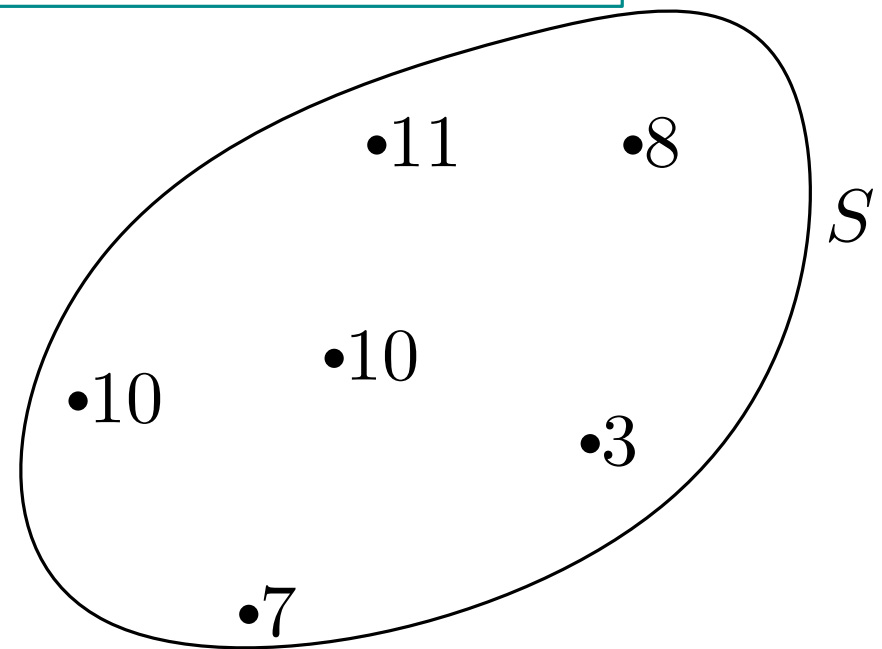
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Som regel i tilfælde hvor alle nøgler er forskellige.



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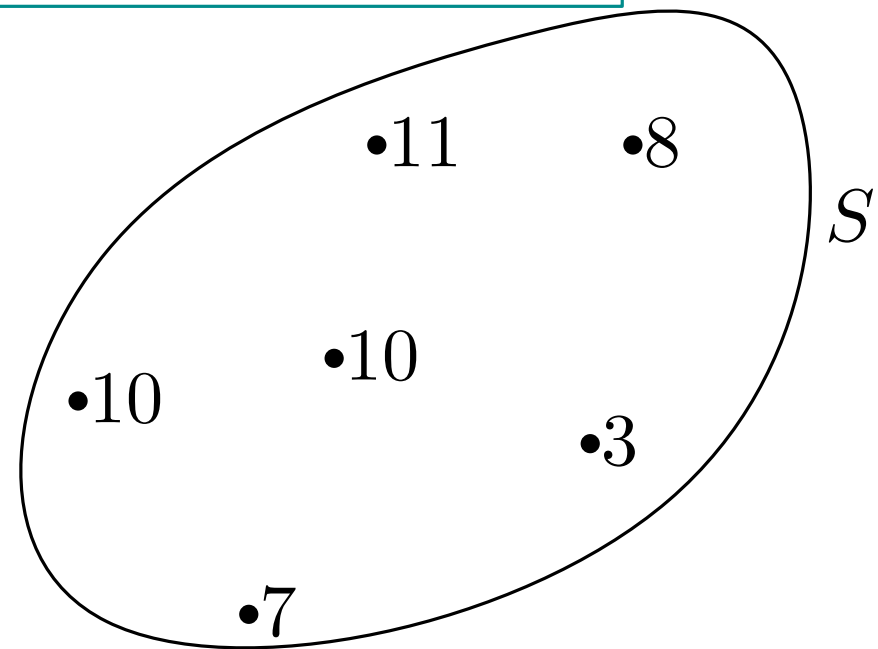
Kendte teknikker?

Array.

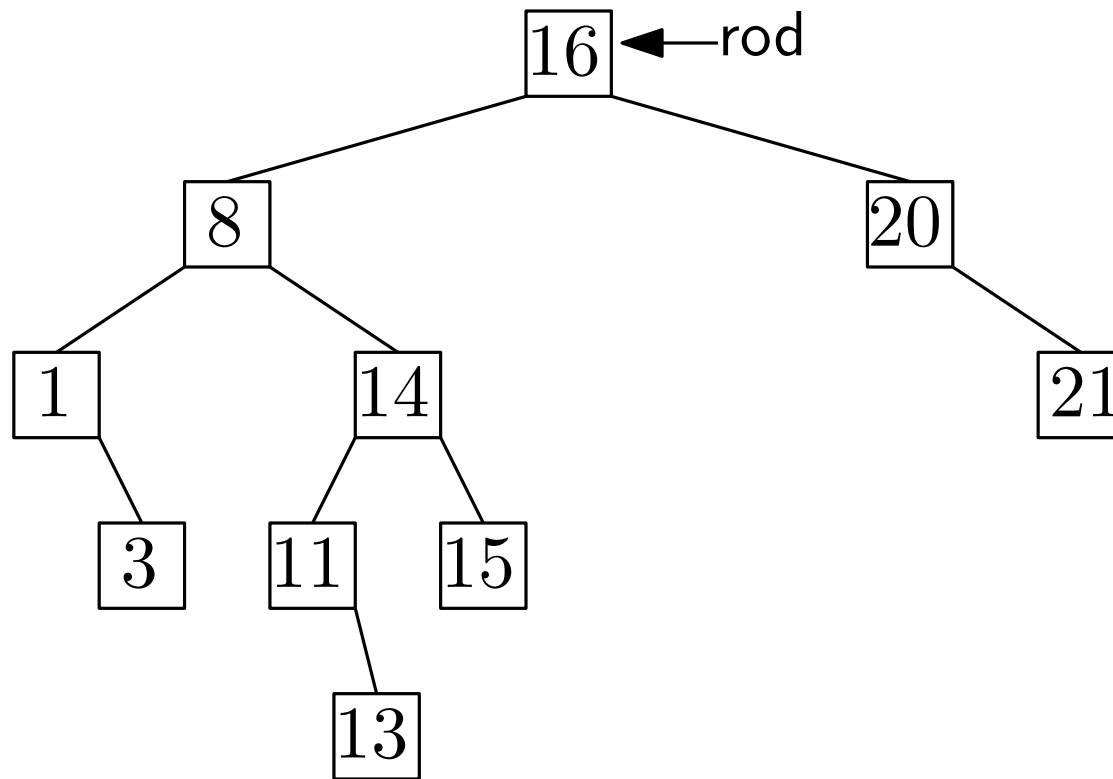
Sorteret array.

Hægtet liste.

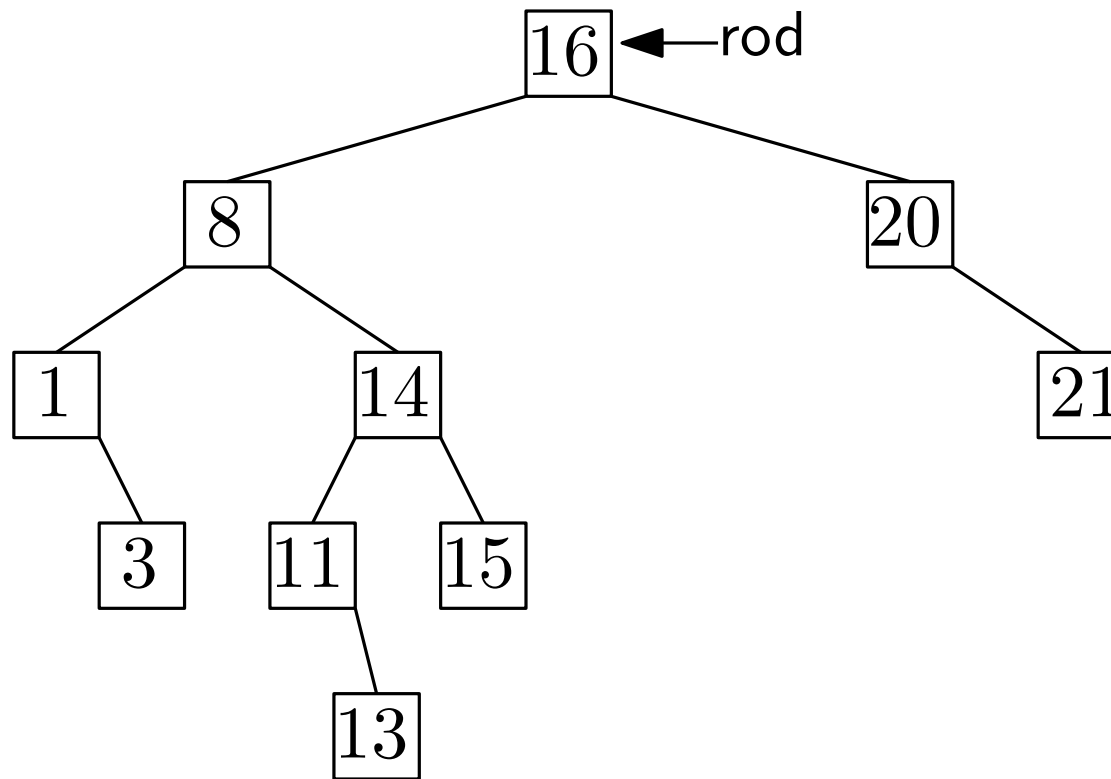
Hob?



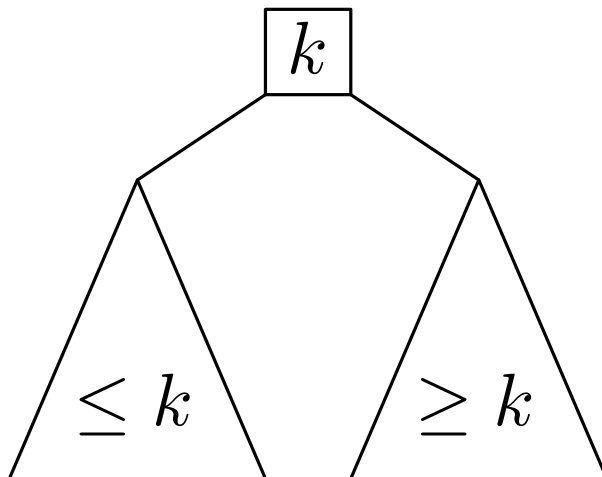
Binære søgetræer



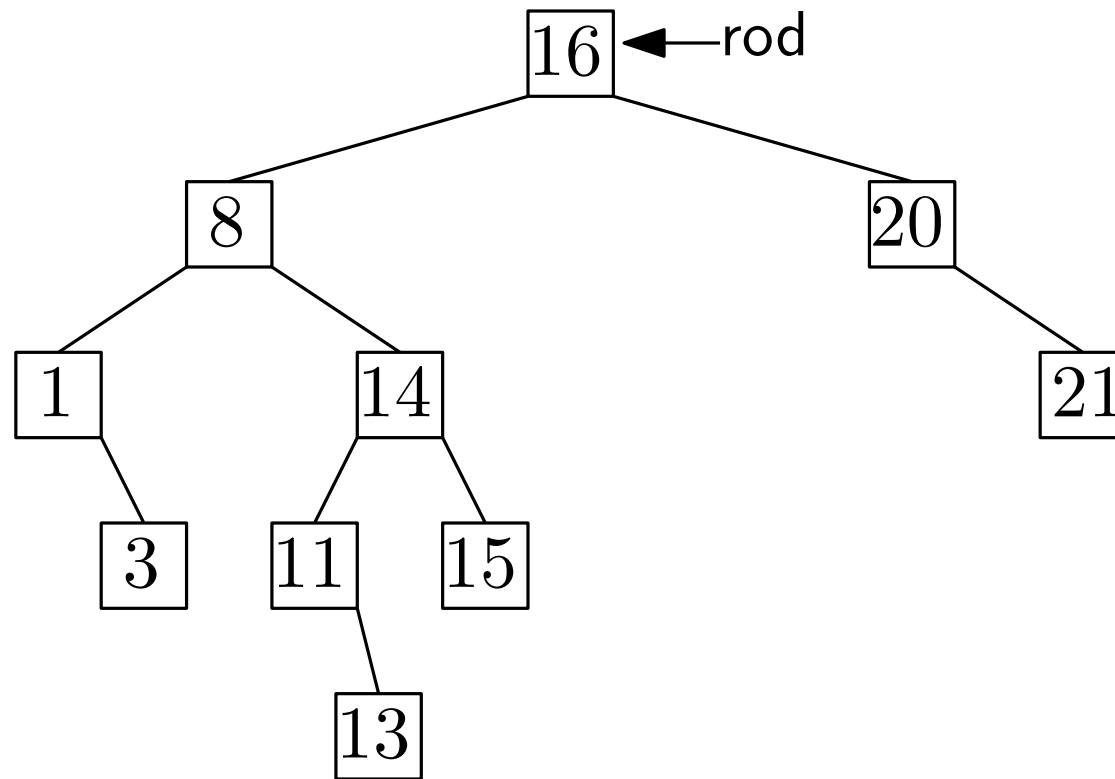
Binære søgetræer



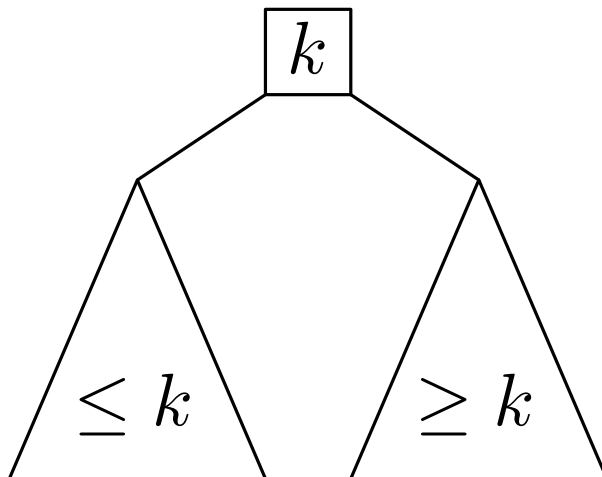
Søgetræsinvariante



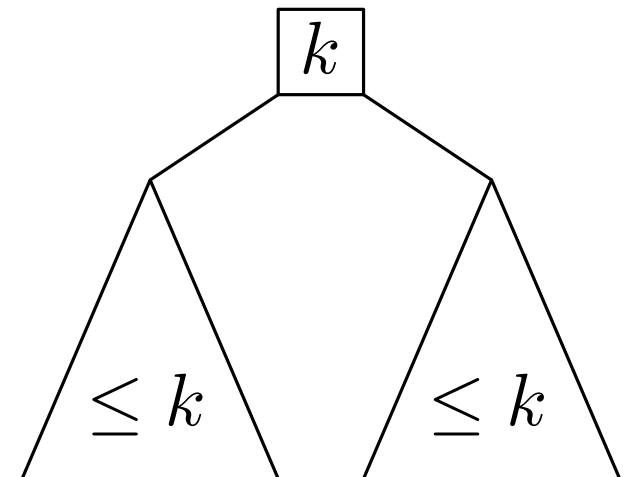
Binære søgetræer



Søgetræsinvariante

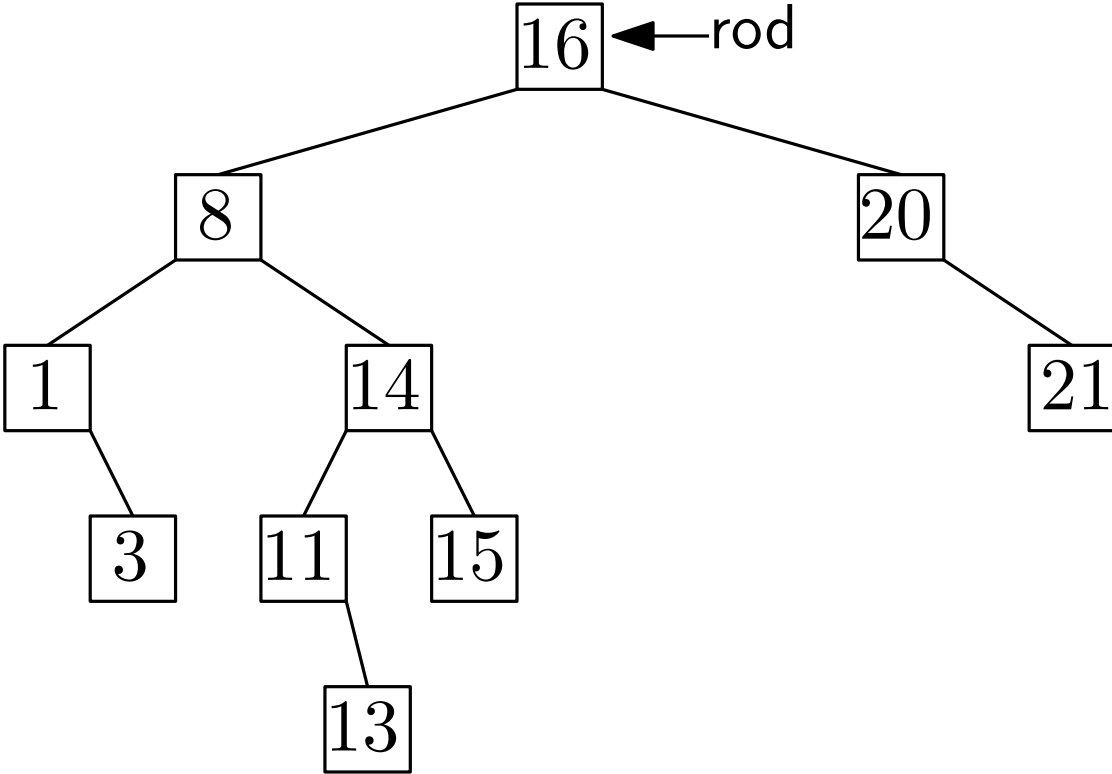


Hobeordenen:

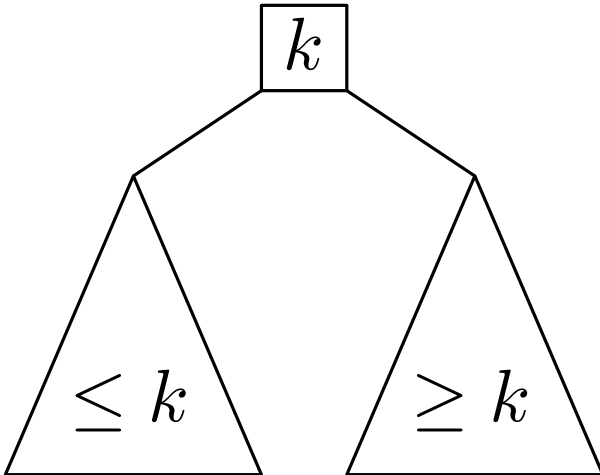


Binære søgetræer

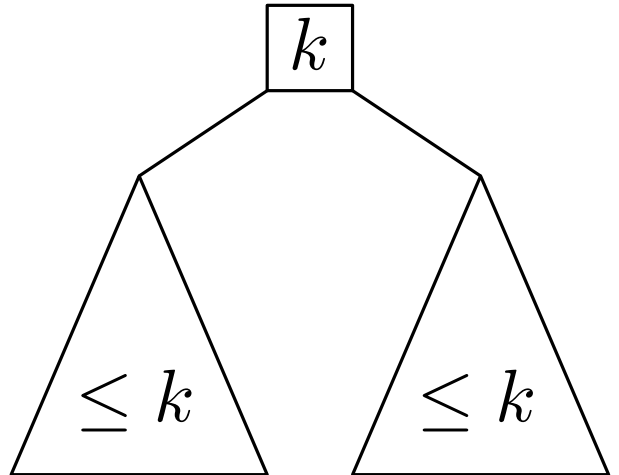
Search(11)



Søgetræsinvarianz



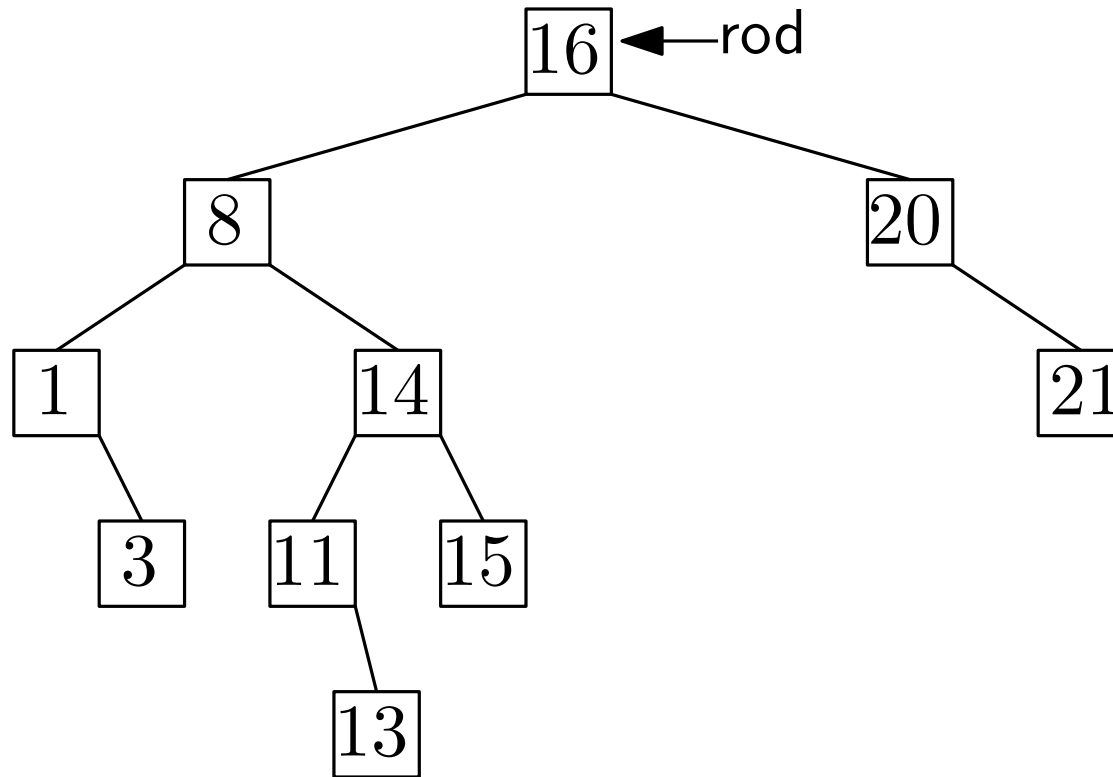
Hobeordenen:



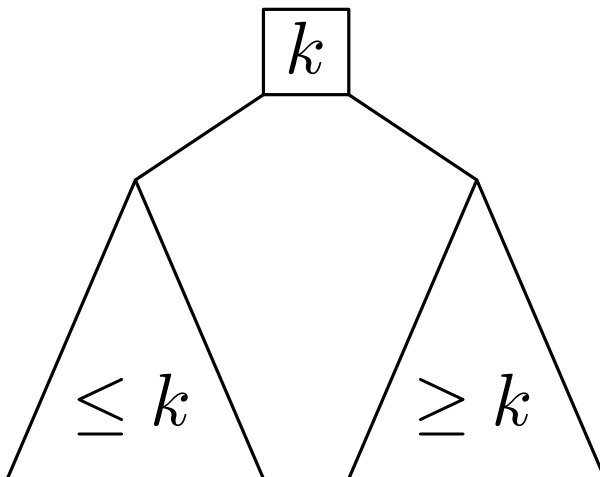
Binære søgetræer

Search(11)

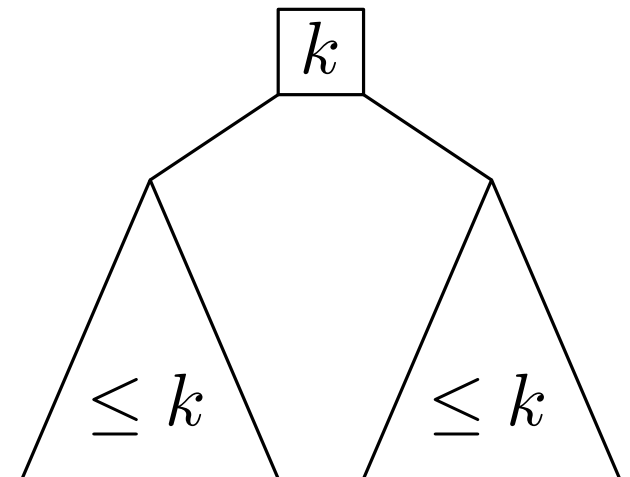
Search(2)



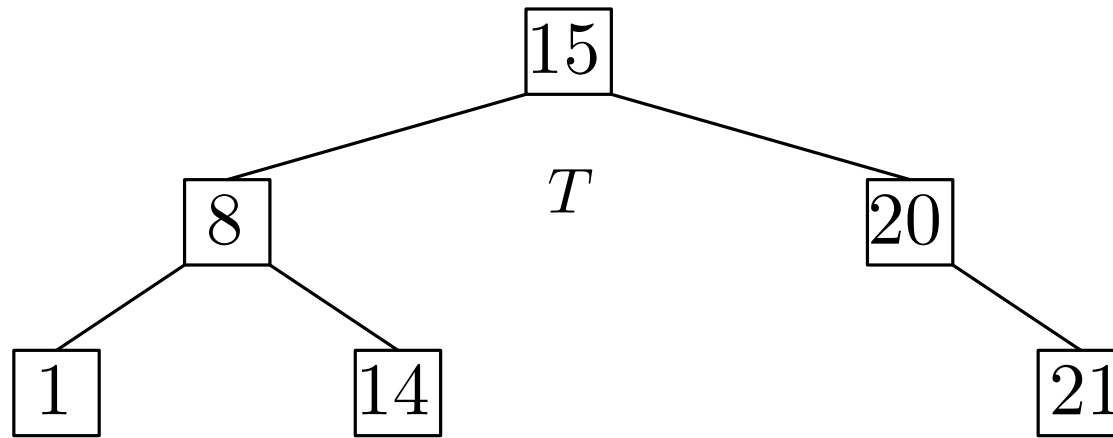
Søgetræsinvariante



Hobeordenen:



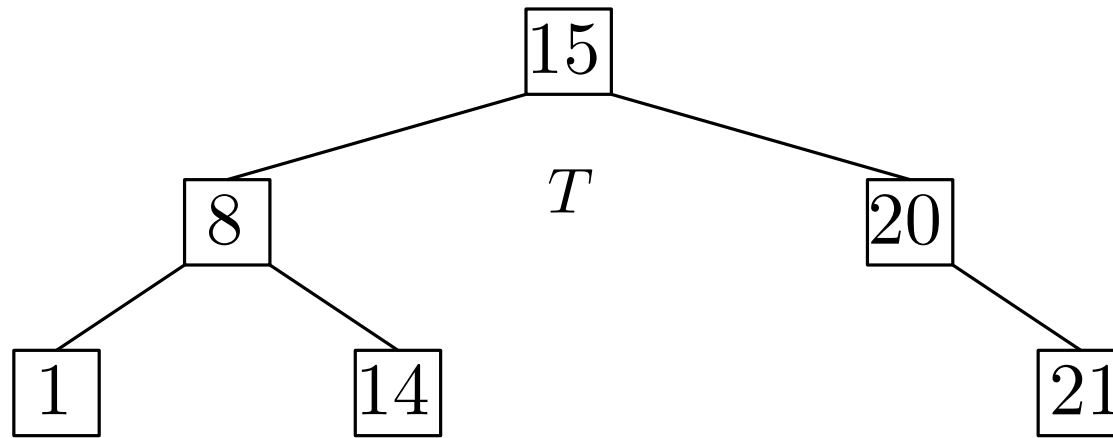
Repræsentation i hukommelsen



Hver knude:

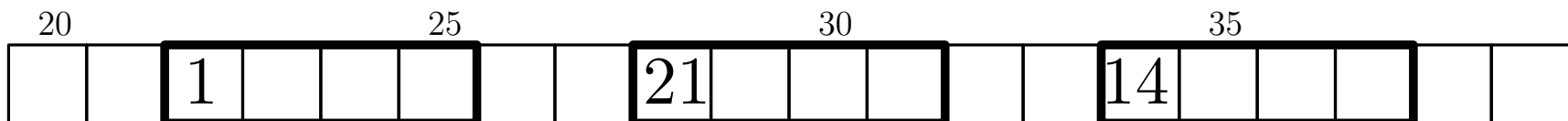
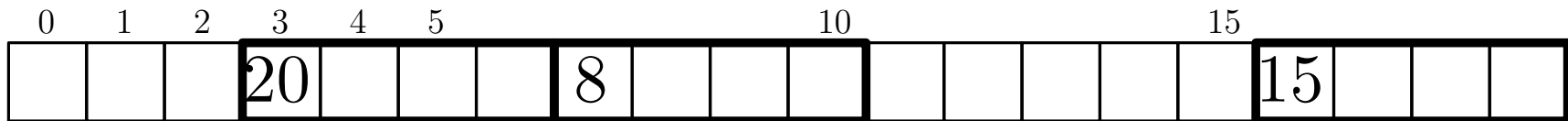
i	$i + 1$	$i + 2$	$i + 3$
<i>key</i>	<i>left</i>	<i>right</i>	<i>parent</i>

Repræsentation i hukommelsen

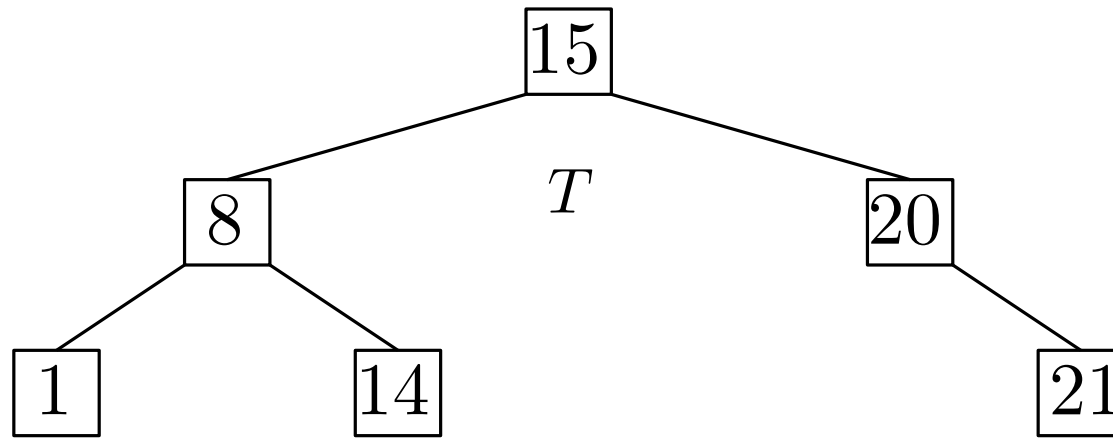


Hver knude:

i	$i + 1$	$i + 2$	$i + 3$
<i>key</i>	<i>left</i>	<i>right</i>	<i>parent</i>

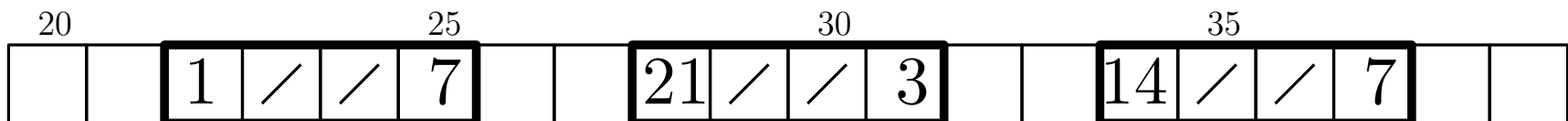
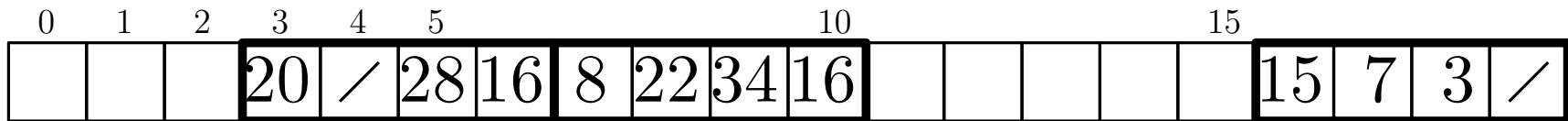


Repræsentation i hukommelsen

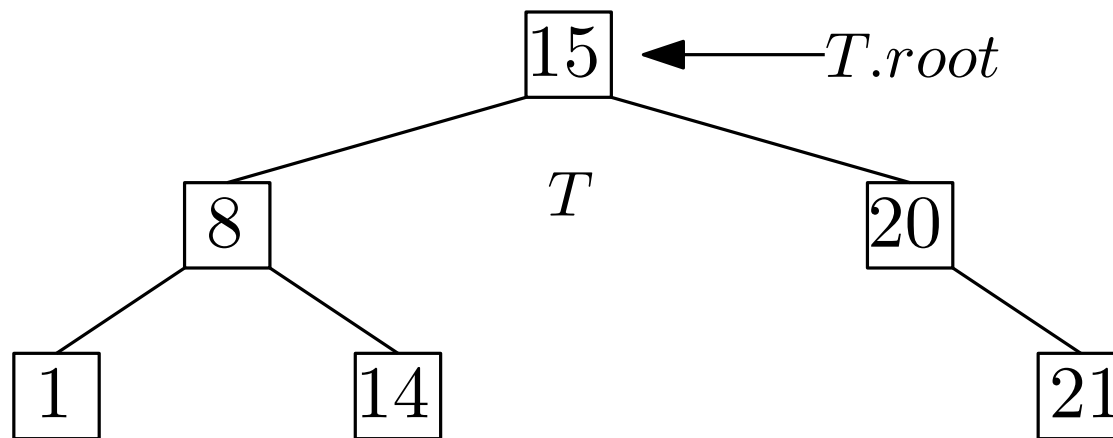


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Repræsentation i hukommelsen

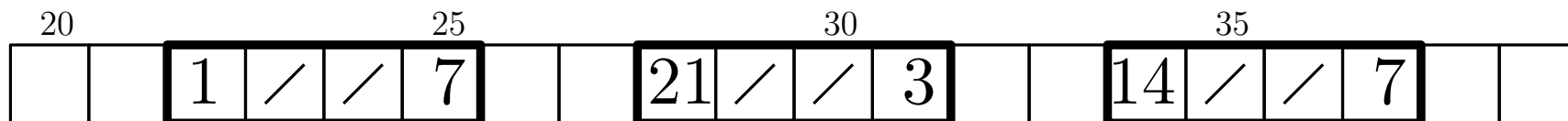
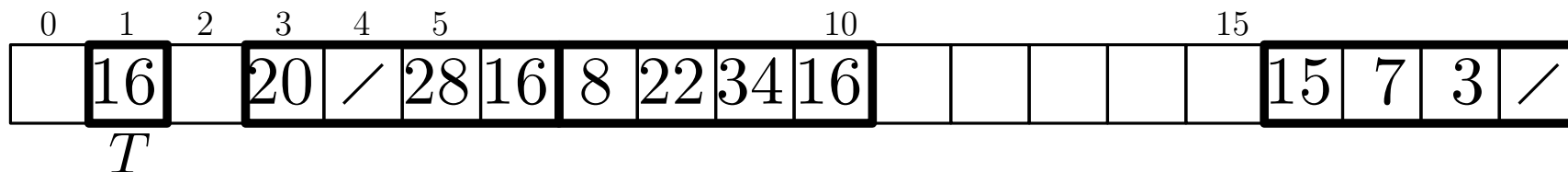


Hver knude:

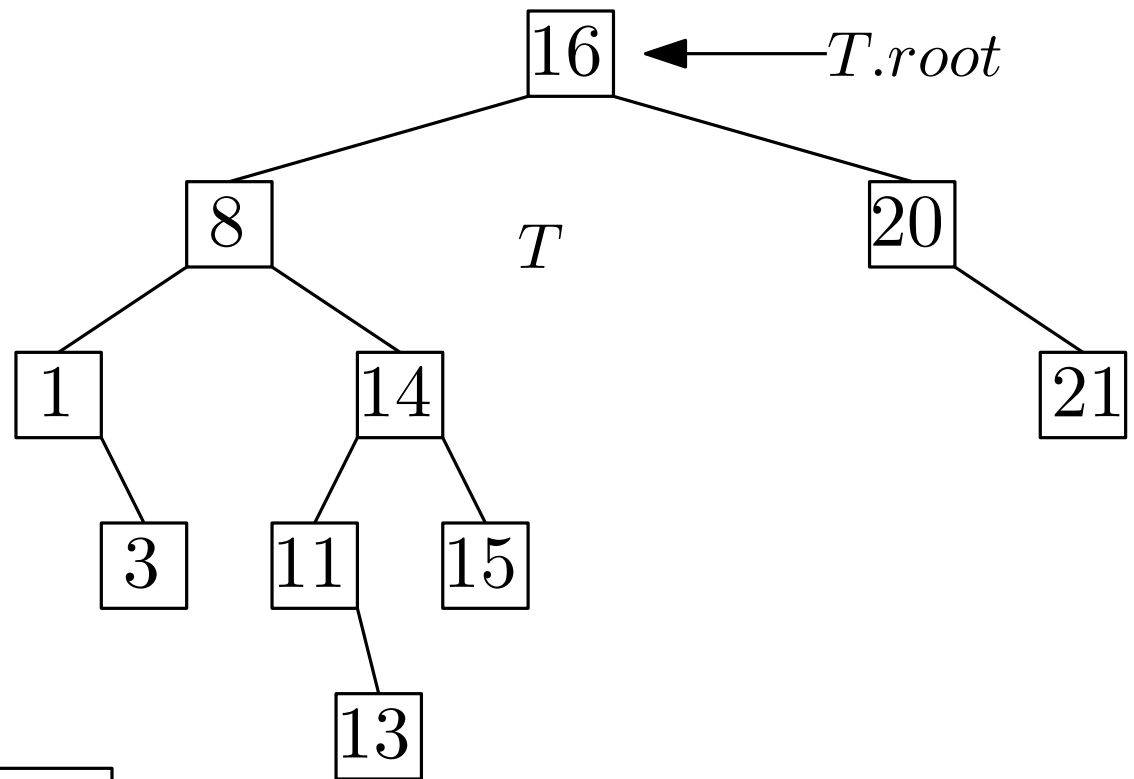
i	$i + 1$	$i + 2$	$i + 3$
<i>key</i>	<i>left</i>	<i>right</i>	<i>parent</i>

Selve træet:

j
<i>root</i>



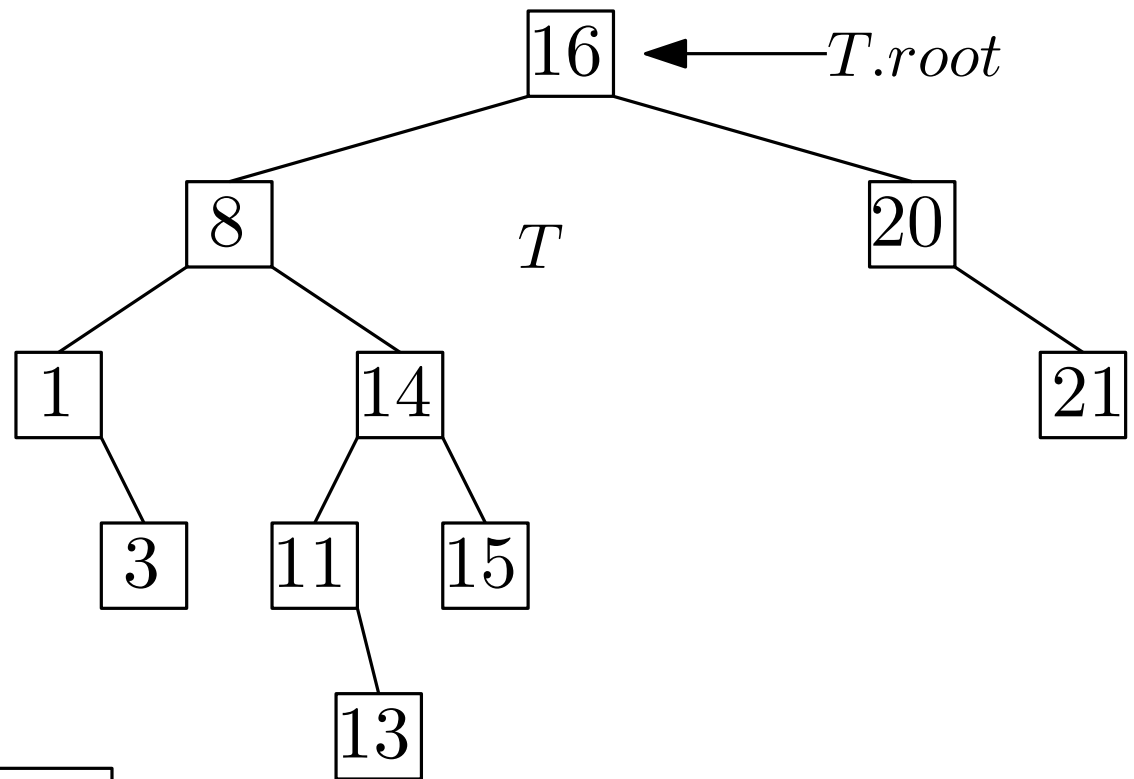
Søgning



```
Tree-Search( $x, k$ )  
  if  $x == \text{NIL}$  or  $k == x.key$   
    return  $x$   
  if  $k < x.key$   
    return Tree-Search( $x.left, k$ )  
  else  
    return Tree-Search( $x.right, k$ )
```

Søgning

Tree-Search($T.root$, 15)

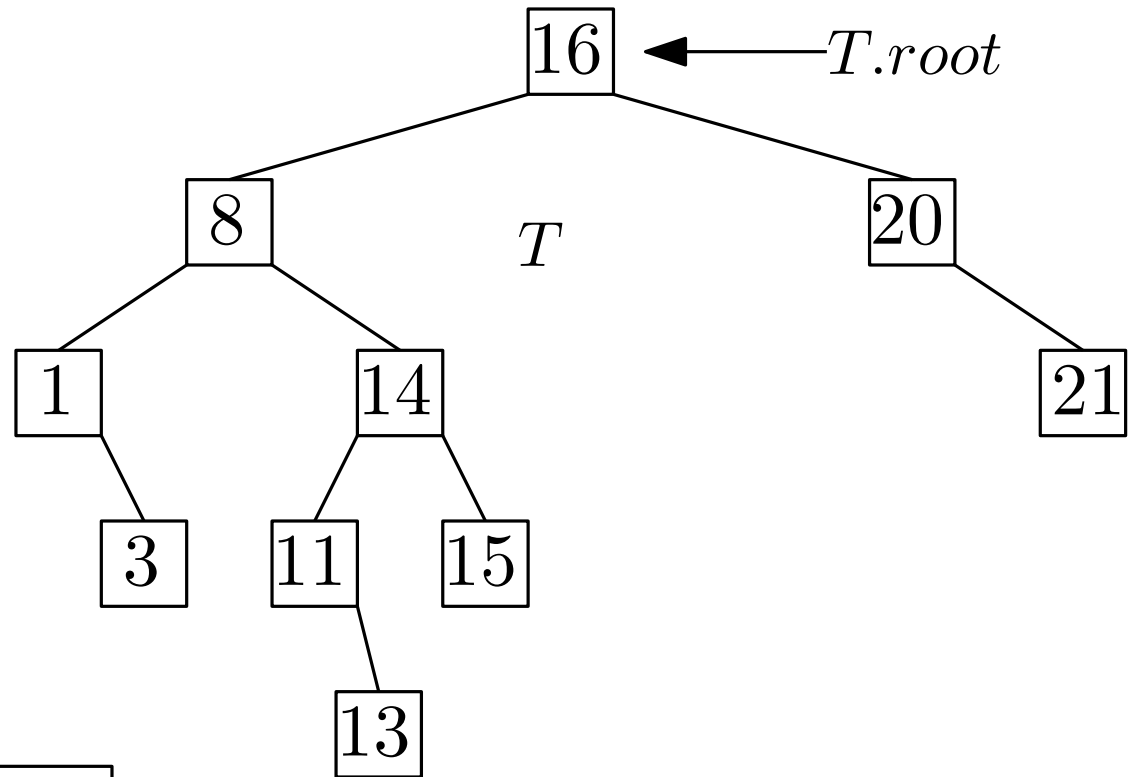


```
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    return  $x$   
  if  $k < x.key$   
    return Tree-Search( $x.left, k$ )  
  else  
    return Tree-Search( $x.right, k$ )
```

Søgning

Tree-Search($T.root$, 15)

Tree-Search($T.root$, 12)



Tree-Search(x, k)

if $x == \text{NIL}$ or $k == x.key$

return x

if $k < x.key$

return Tree-Search($x.left, k$)

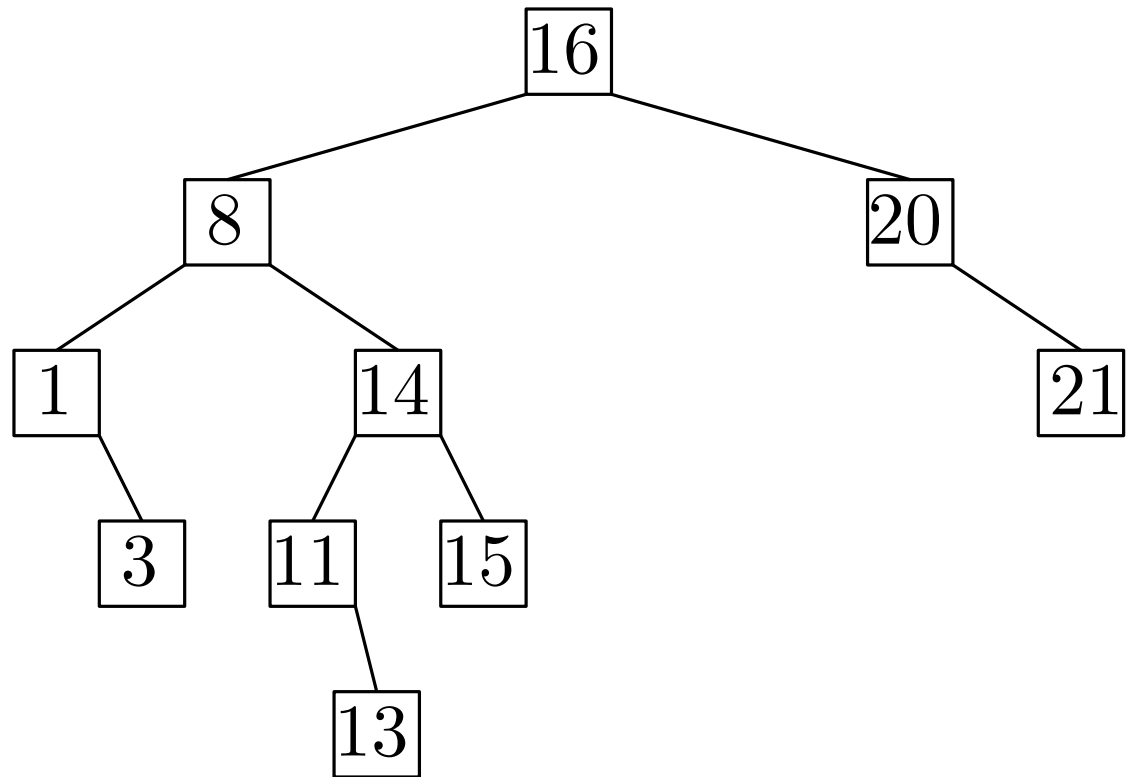
else

return Tree-Search($x.right, k$)

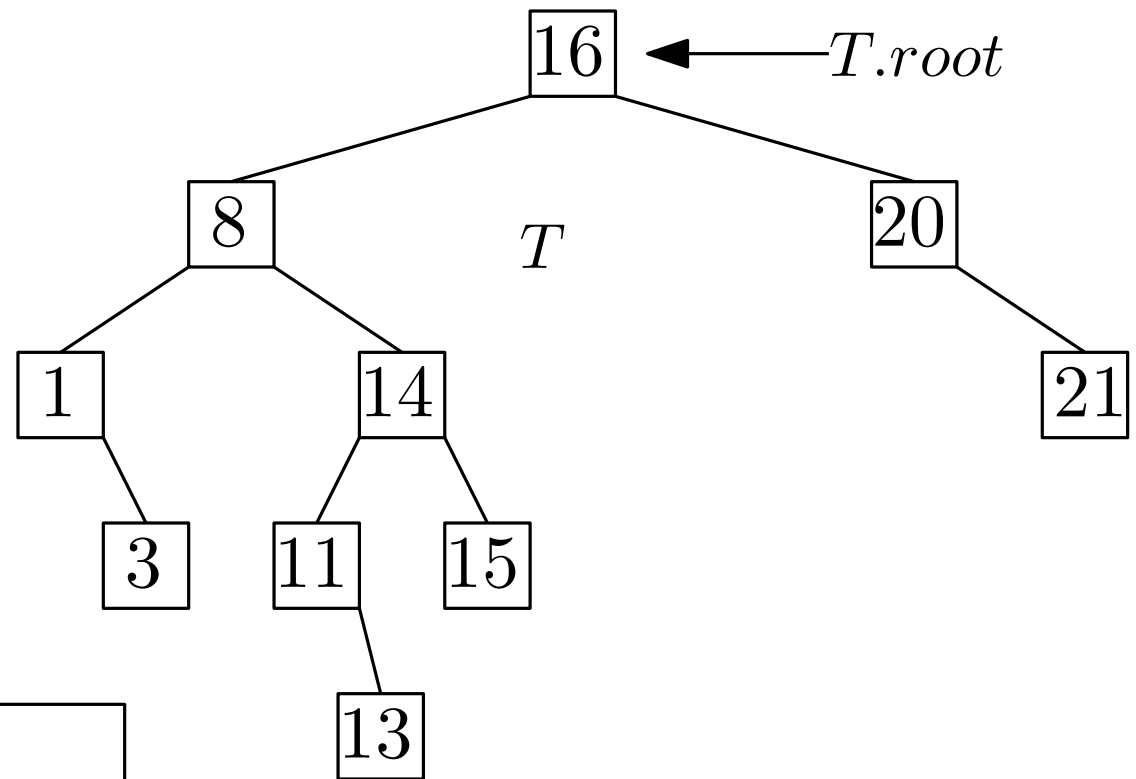
Minimum/maximum

```
Tree-Maximum( $x$ )  
  while  $x.right \neq \text{NIL}$   
     $x = x.right$   
  return  $x$ 
```

```
Tree-Minimum( $x$ )  
  while  $x.left \neq \text{NIL}$   
     $x = x.left$   
  return  $x$ 
```



Efterfølger

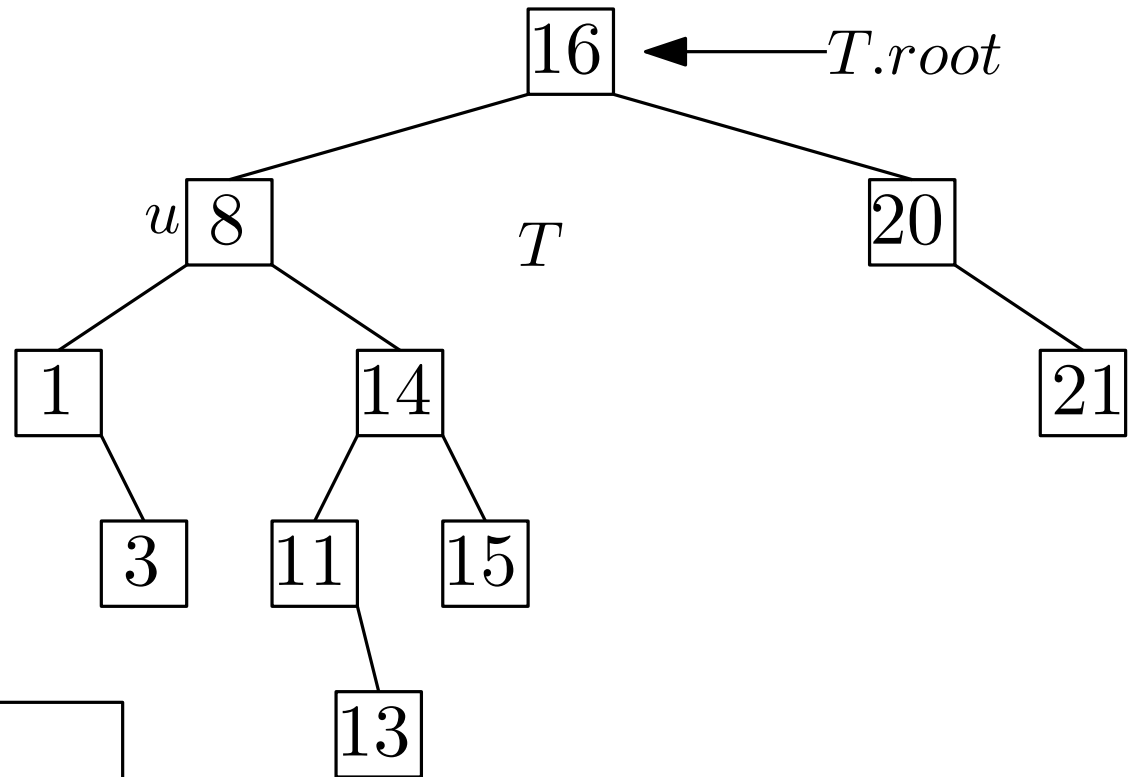


```
Tree-Successor( $x$ )  
  if  $x.right \neq \text{NIL}$   
    return Tree-Minimum( $x.right$ )  
   $y = x.p$   
  while  $y \neq \text{NIL}$  and  $x == y.right$   
     $x = y$   
     $y = y.p$   
  return  $y$ 
```

Antag nøglerne er forskellige!

Efterfølger

Tree-Successor(u)



Tree-Successor(x)

if $x.right \neq \text{NIL}$

return Tree-Minimum($x.right$)

$y = x.p$

while $y \neq \text{NIL}$ and $x == y.right$

$x = y$

$y = y.p$

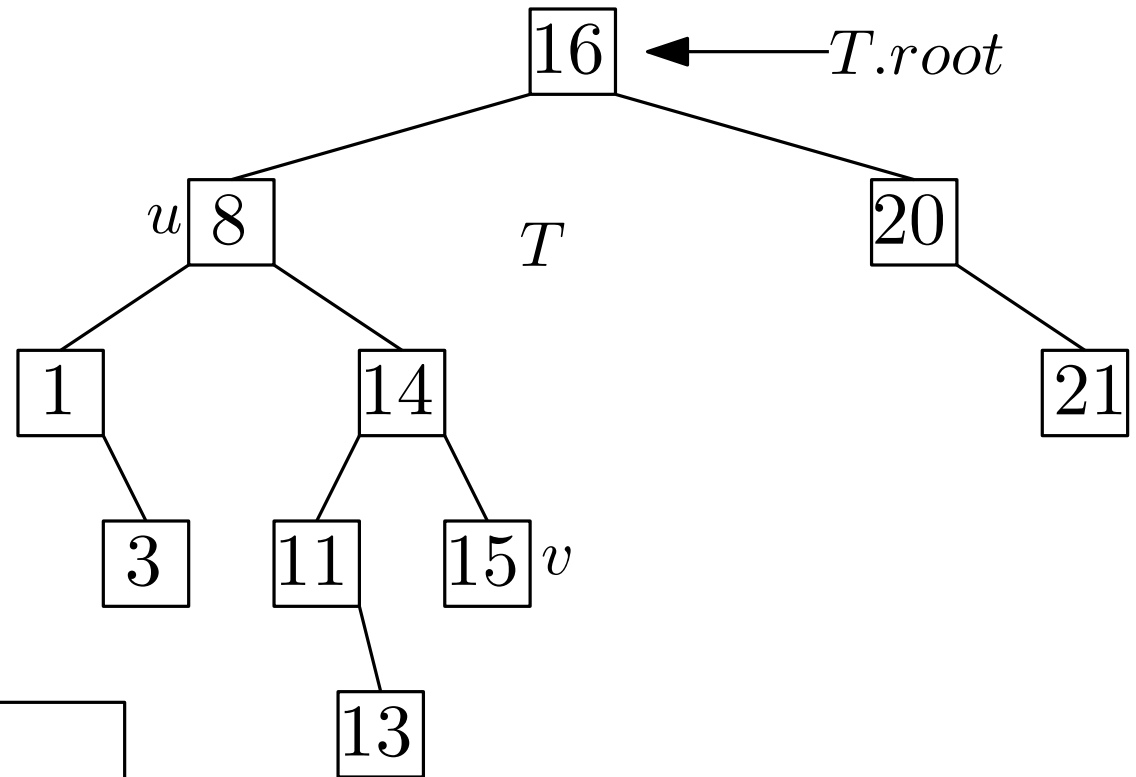
return y

Antag nøglerne er forskellige!

Efterfølger

Tree-Successor(u)

Tree-Successor(v)



Tree-Successor(x)

if $x.right \neq \text{NIL}$

return Tree-Minimum($x.right$)

$y = x.p$

while $y \neq \text{NIL}$ and $x == y.right$

$x = y$

$y = y.p$

return y

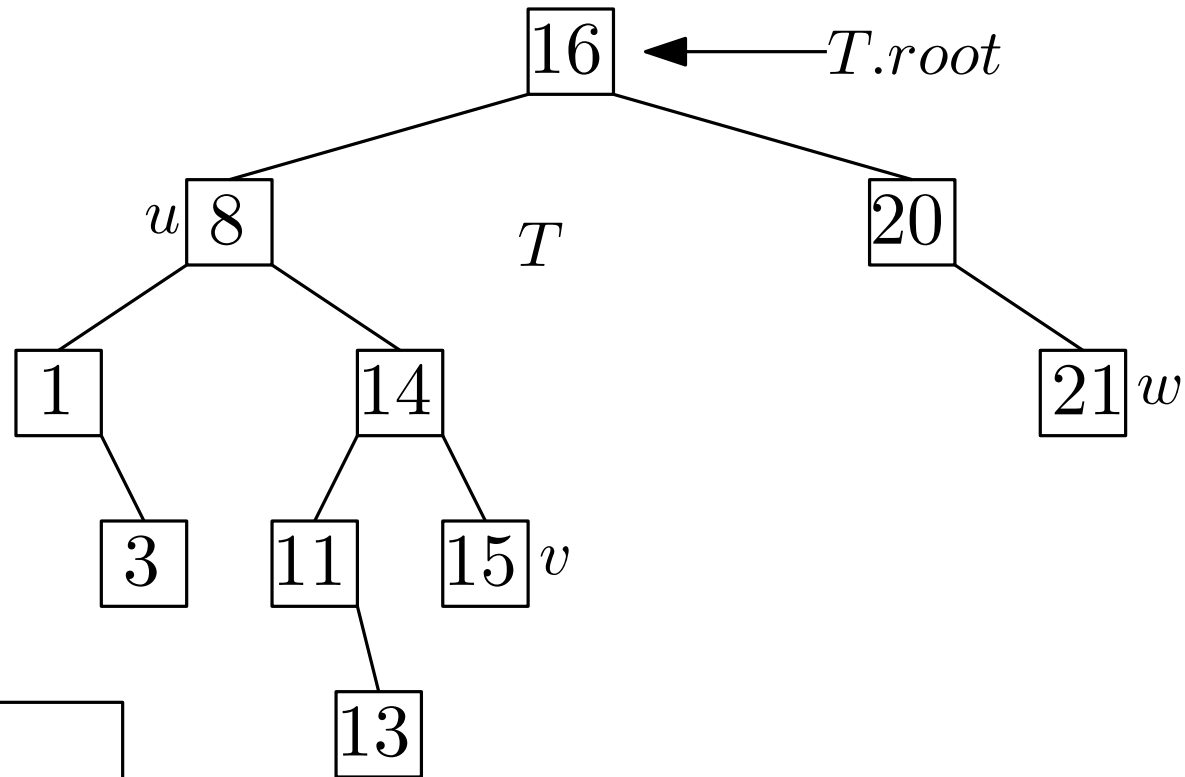
Antag nøglerne er forskellige!

Efterfølger

Tree-Successor(u)

Tree-Successor(v)

Tree-Successor(w)



Tree-Successor(x)

if $x.right \neq \text{NIL}$

return Tree-Minimum($x.right$)

$y = x.p$

while $y \neq \text{NIL}$ and $x == y.right$

$x = y$

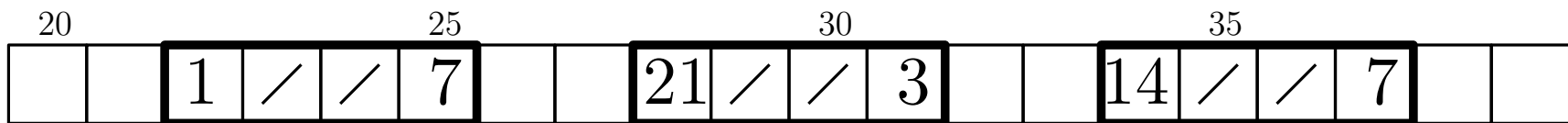
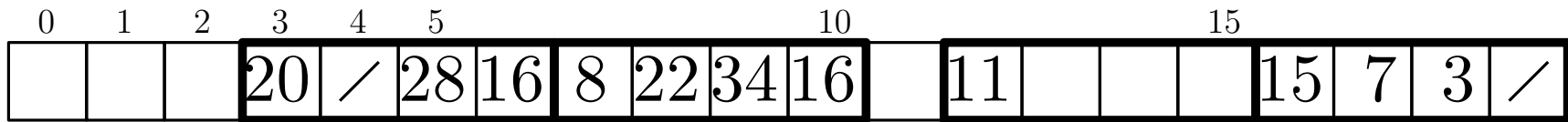
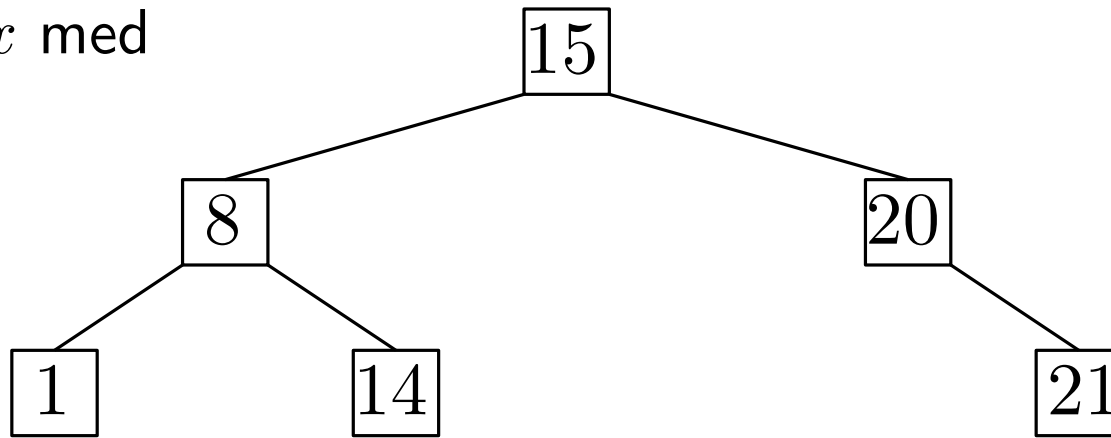
$y = y.p$

return y

Antag nøglerne er forskellige!

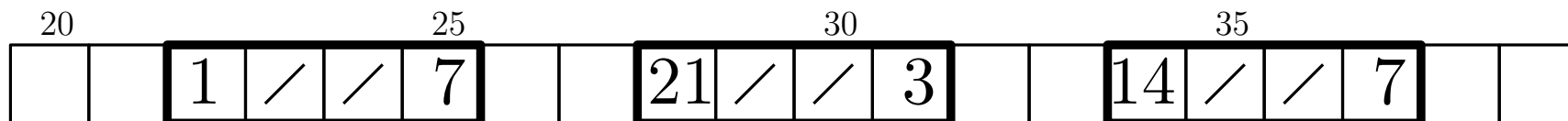
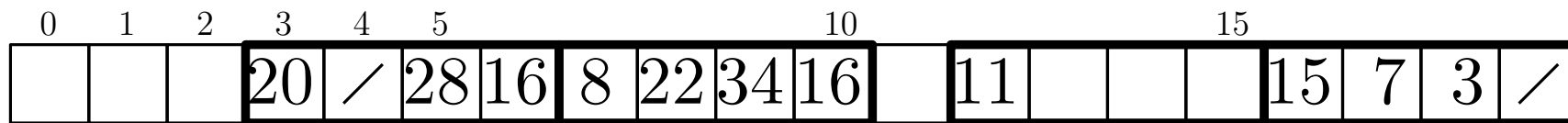
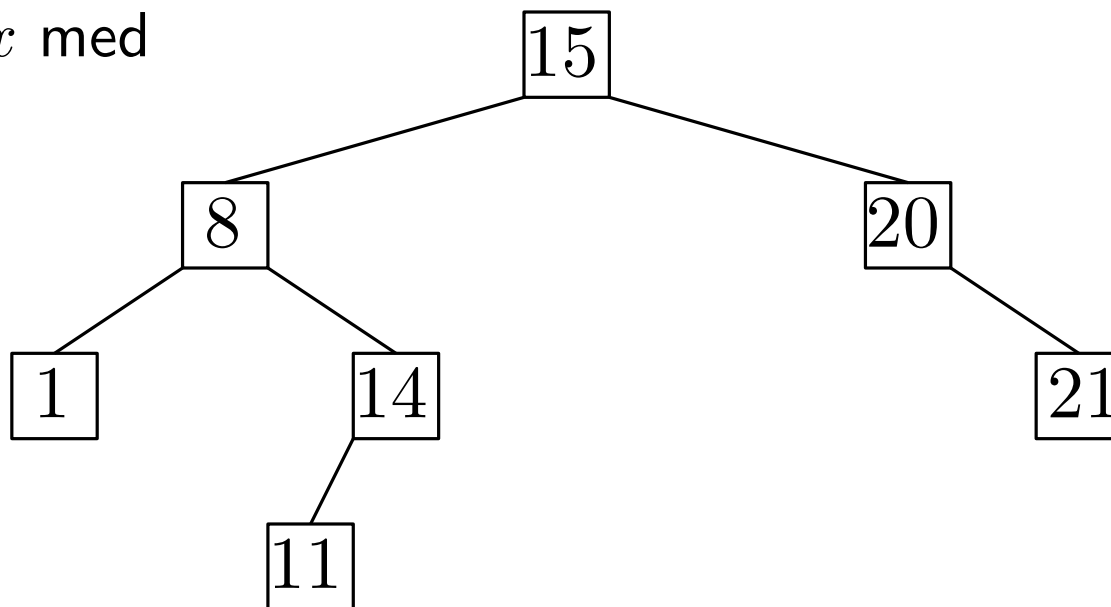
Indsættelse

Indsæt nyt elm. x med
 $x.key = 11$.



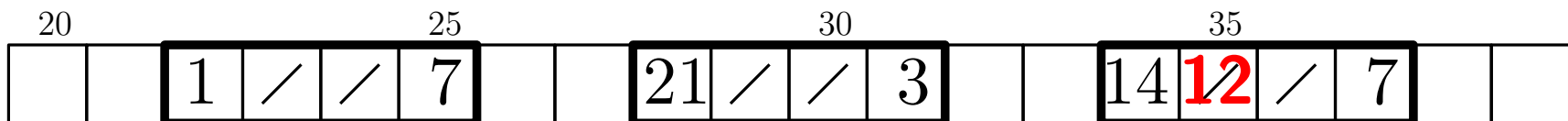
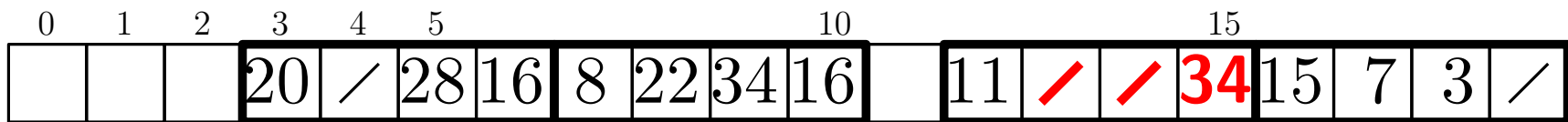
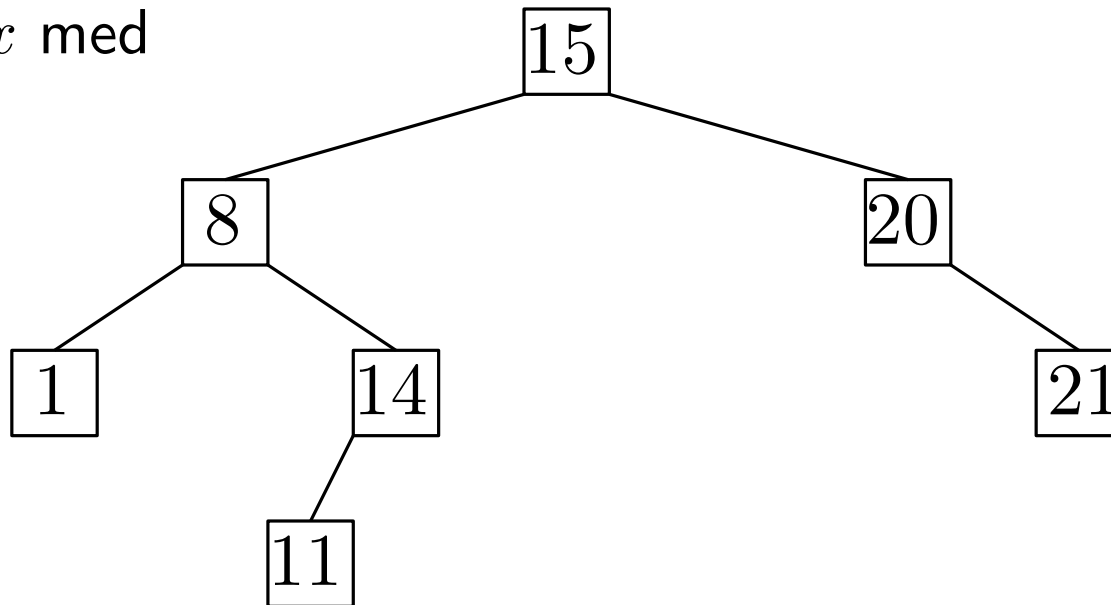
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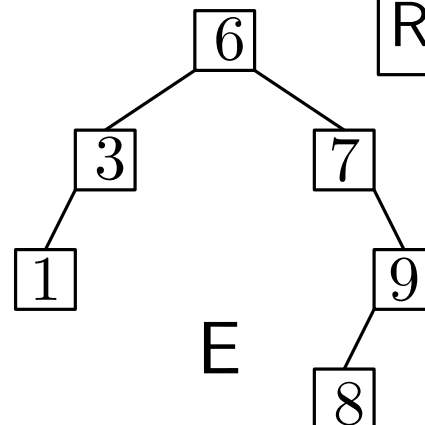
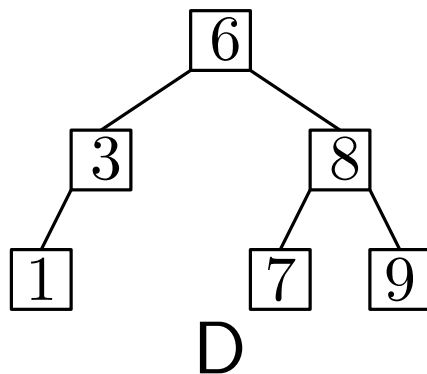
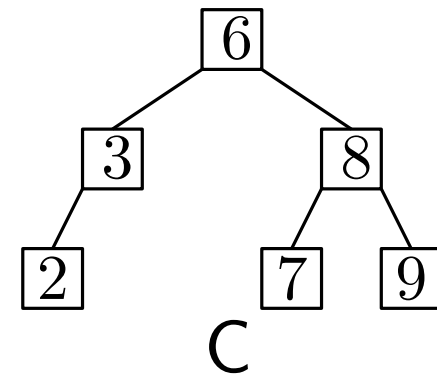
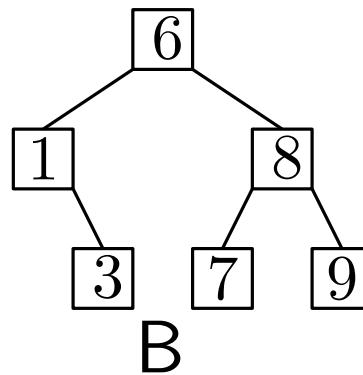
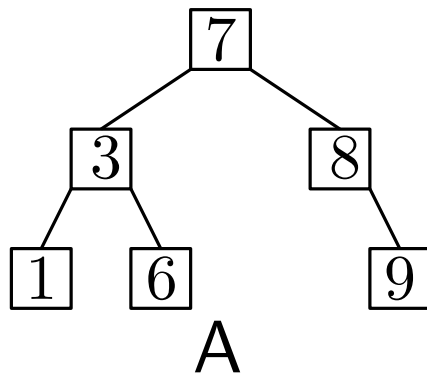


Hvordan ser træet T ud til sidst?

$T.root = NIL$

Indsæt elementer med nøgler

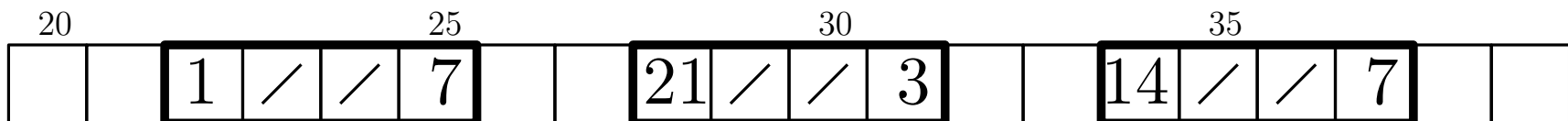
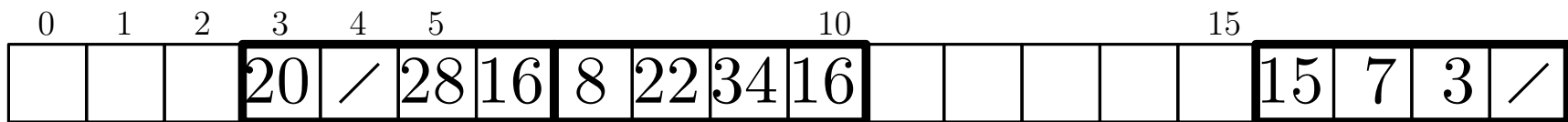
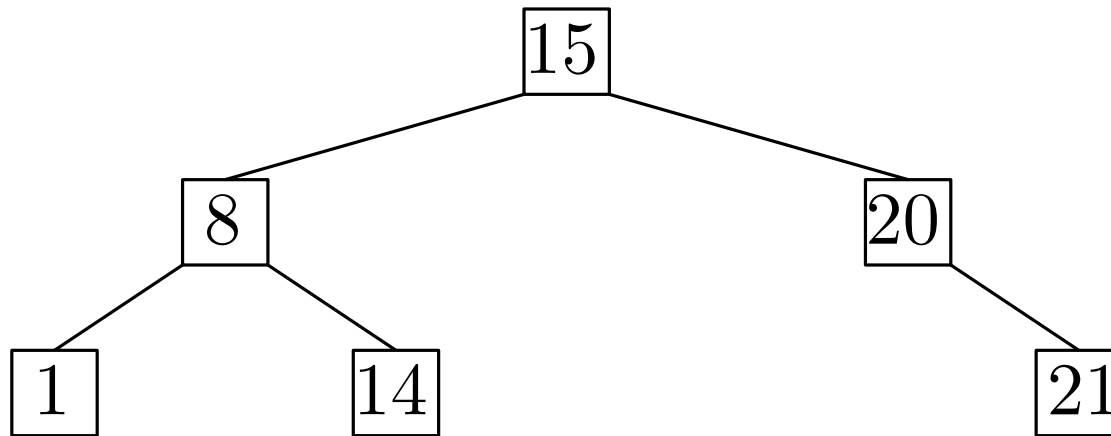
6, 3, 8, 7, 1, 9.



socrative.com → Student login,
Room name: ABRAHAMSEN3464

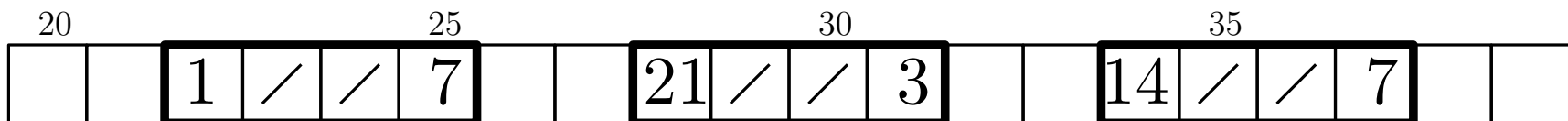
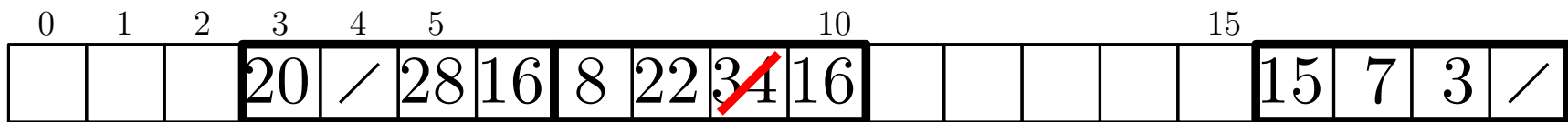
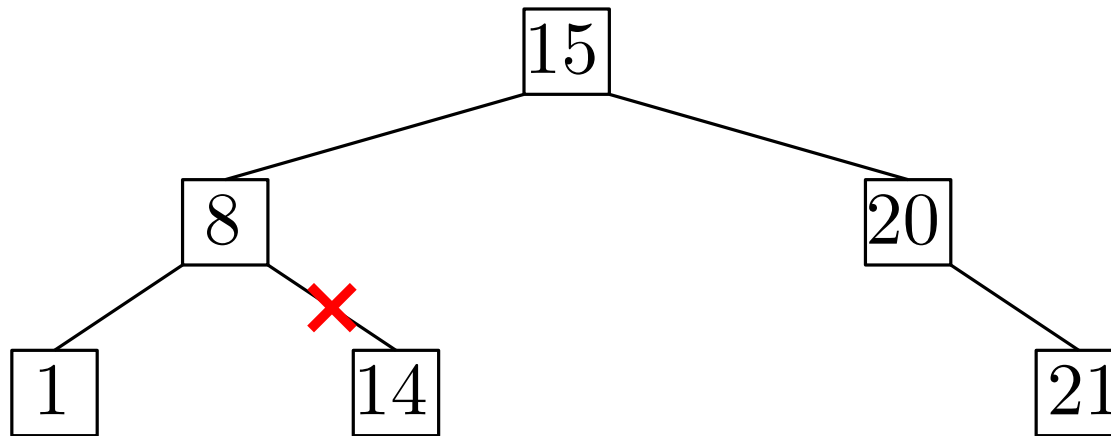
Sletning, 0 børn

Slet 14.



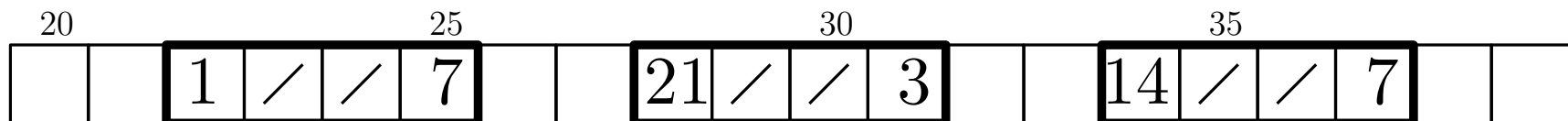
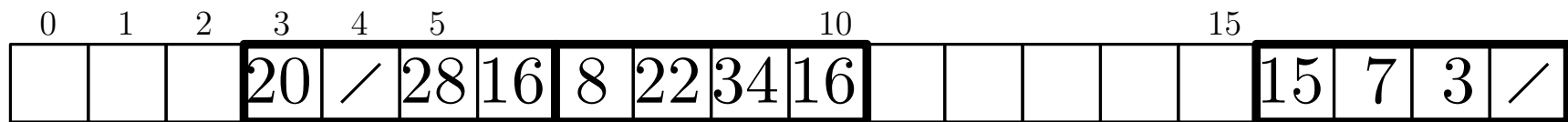
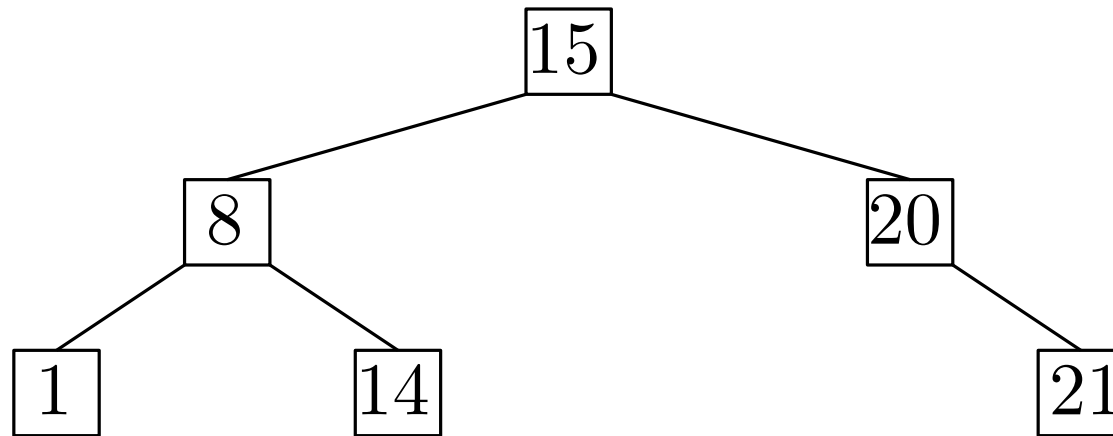
Sletning, 0 børn

Slet 14.



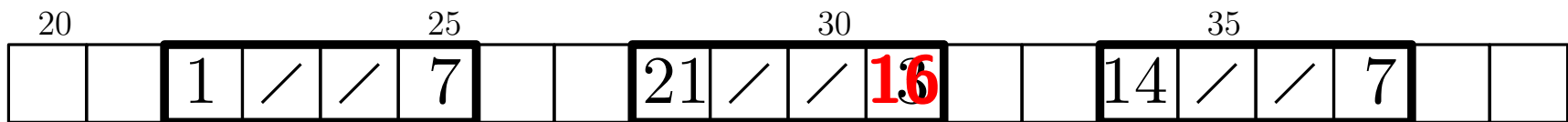
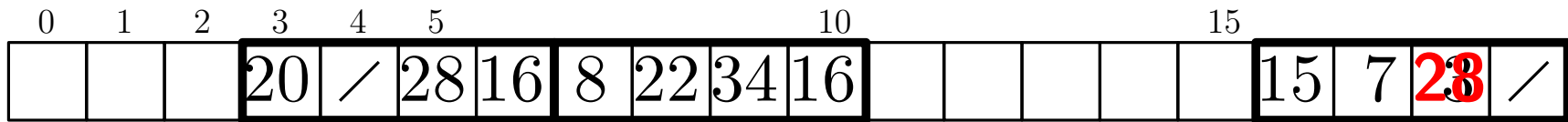
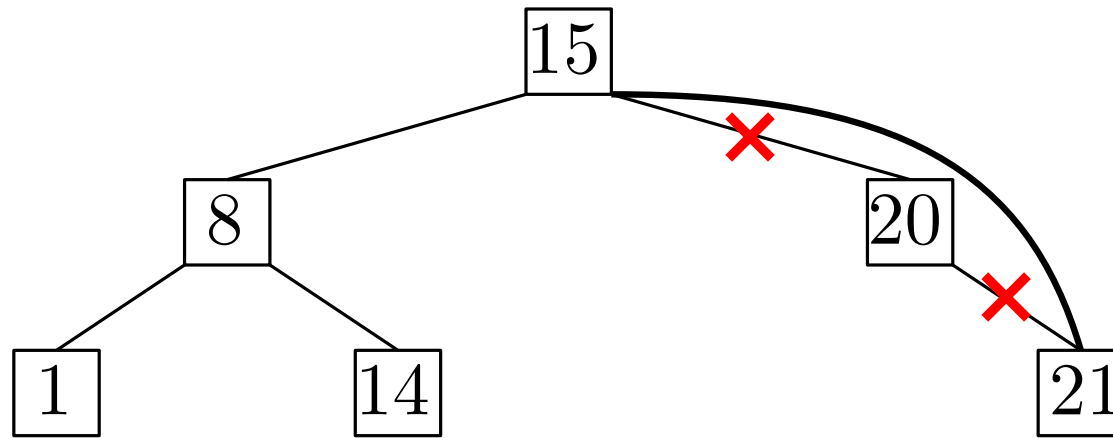
Sletning, 1 barn

Slet 20.



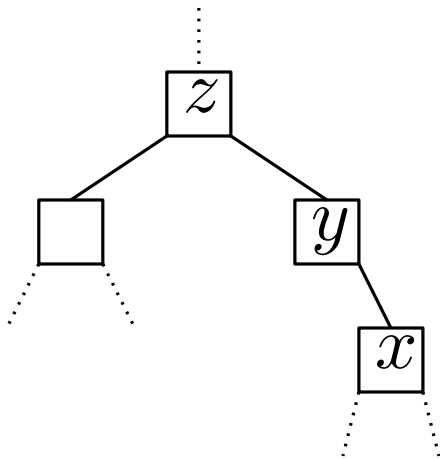
Sletning, 1 barn

Slet 20.



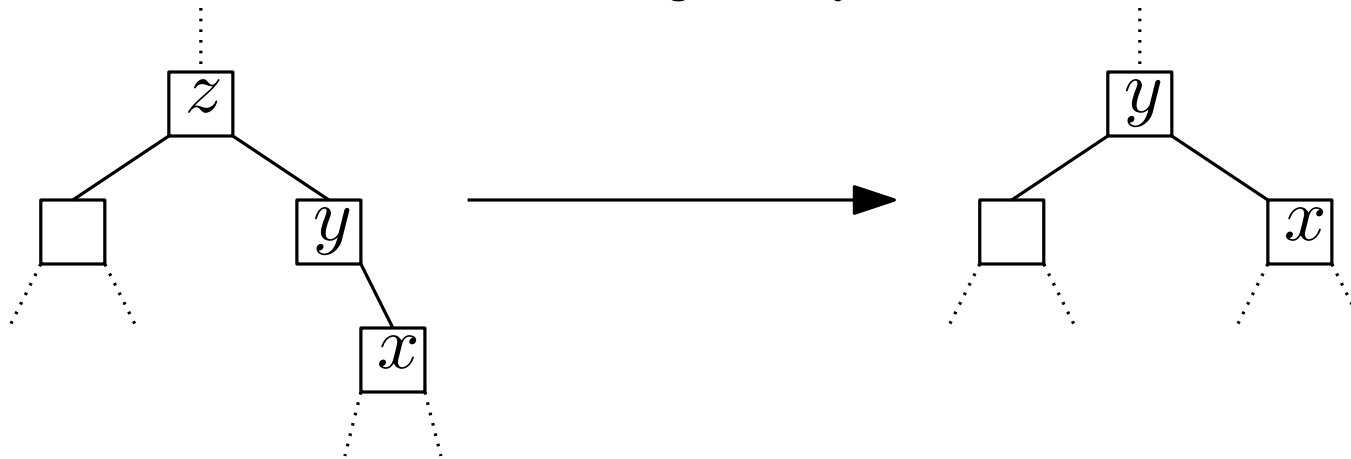
Sletning, 2 børn

Slet z med to børn. Tilfælde 1: $z.right.left == \text{NIL}$



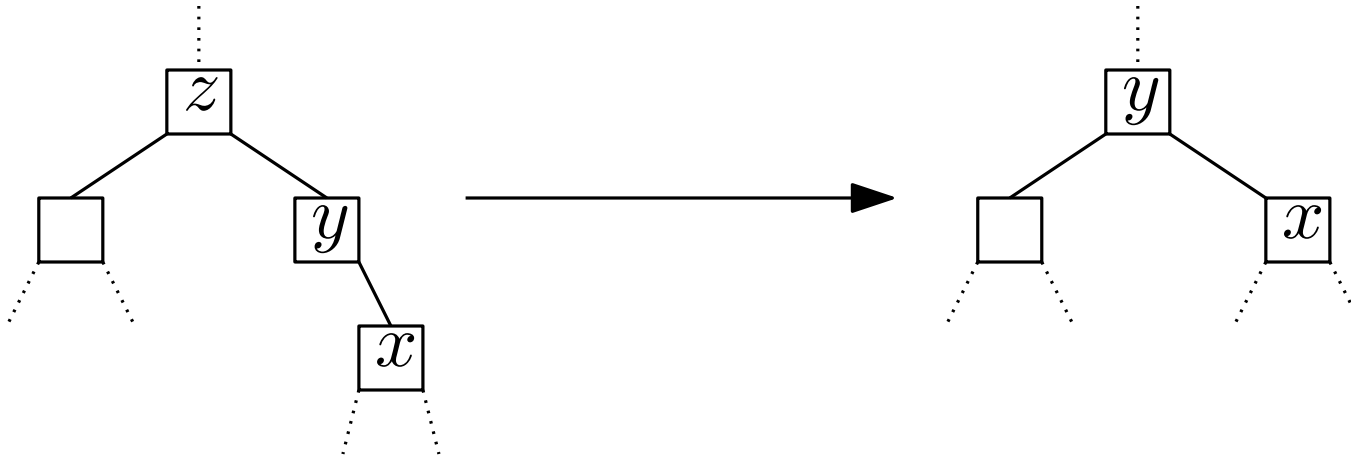
Sletning, 2 børn

Slet z med to børn. Tilfælde 1: $z.right.left == \text{NIL}$

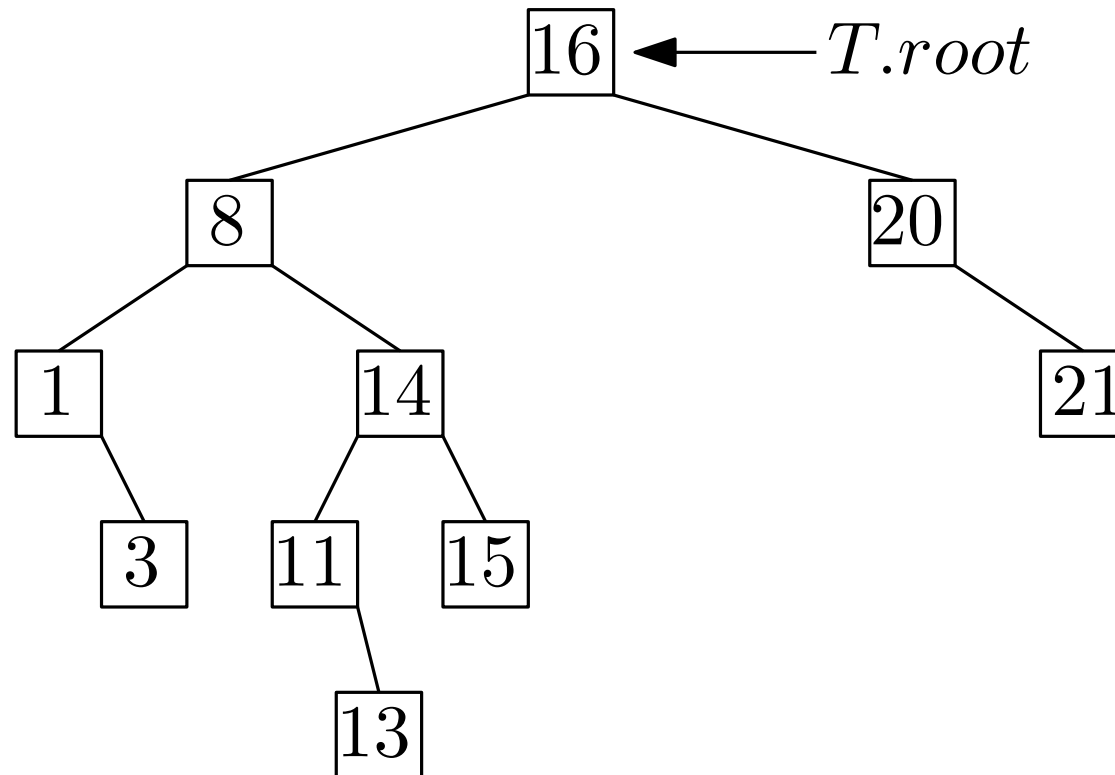


Sletning, 2 børn

Slet z med to børn. Tilfælde 1: $z.right.left == \text{NIL}$

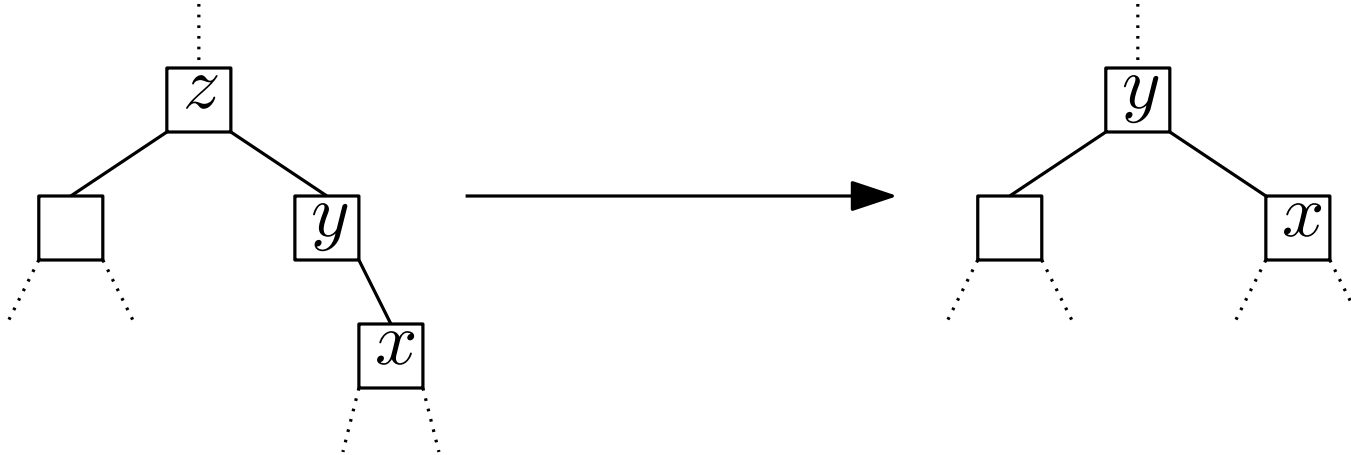


Slet 16.

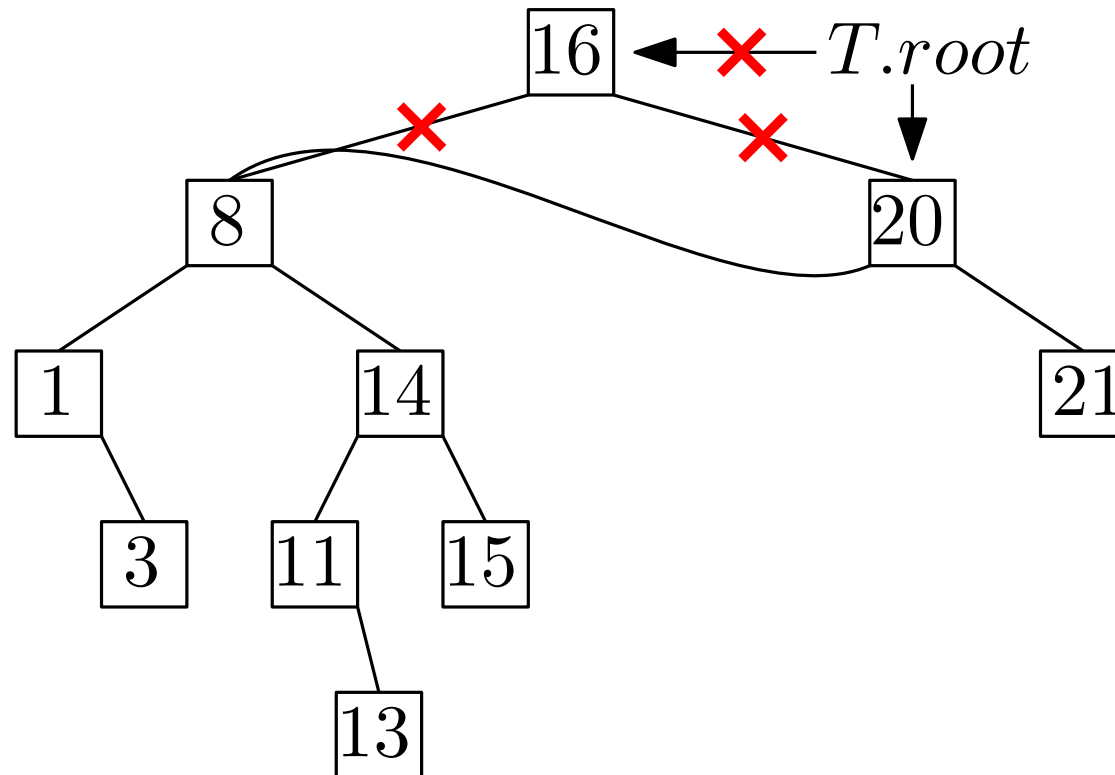


Sletning, 2 børn

Slet z med to børn. Tilfælde 1: $z.right.left == \text{NIL}$

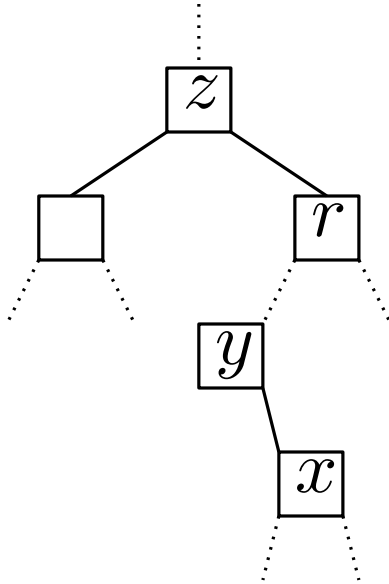


Slet 16.



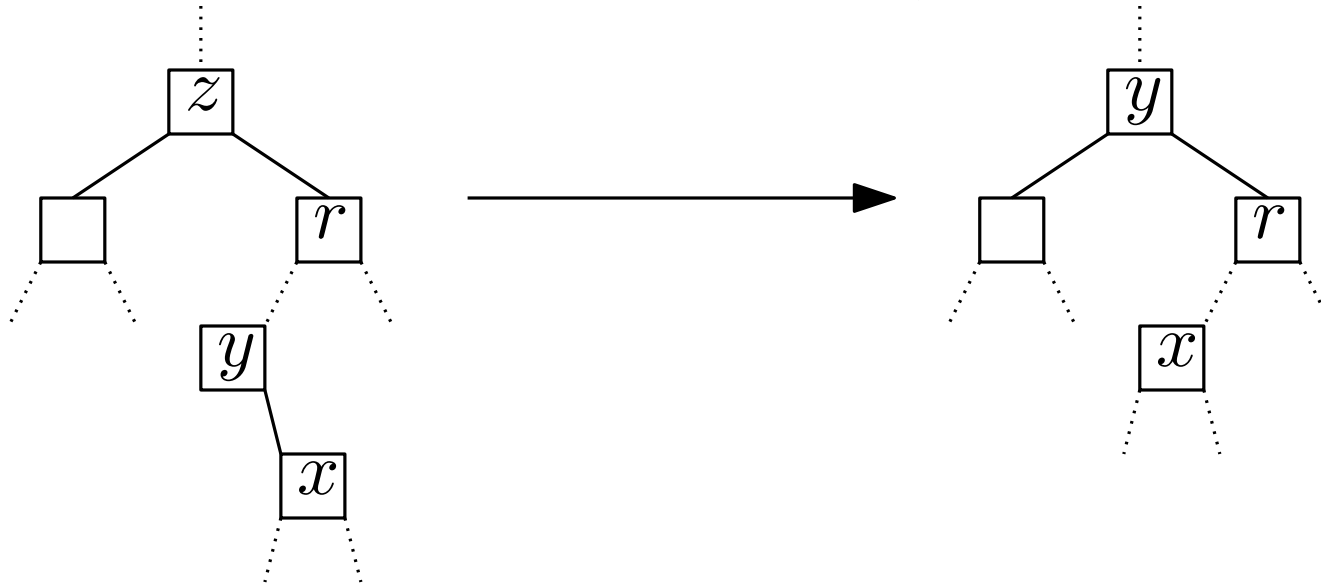
Sletning, 2 børn

Slet z med to børn. Tilfælde 2: $z.right.left \neq \text{NIL}$



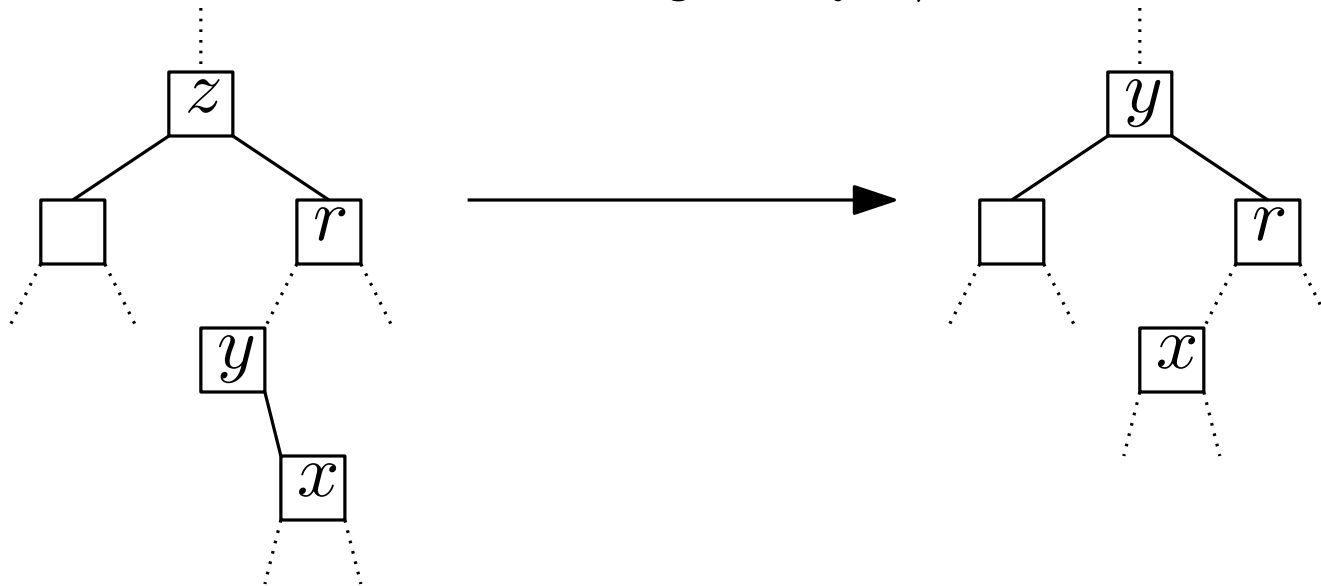
Sletning, 2 børn

Slet z med to børn. Tilfælde 2: $z.right.left \neq \text{NIL}$

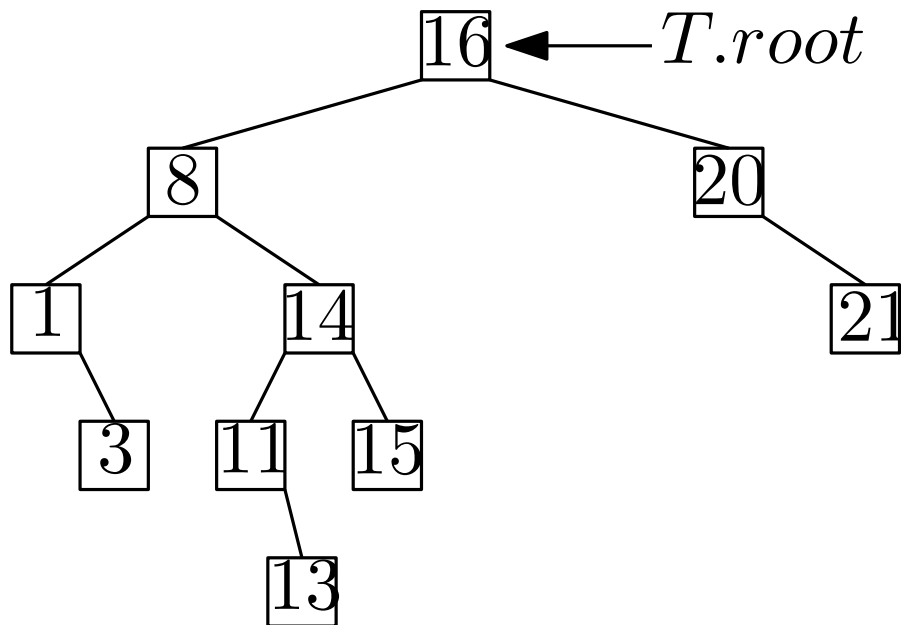


Sletning, 2 børn

Slet z med to børn. Tilfælde 2: $z.right.left \neq \text{NIL}$

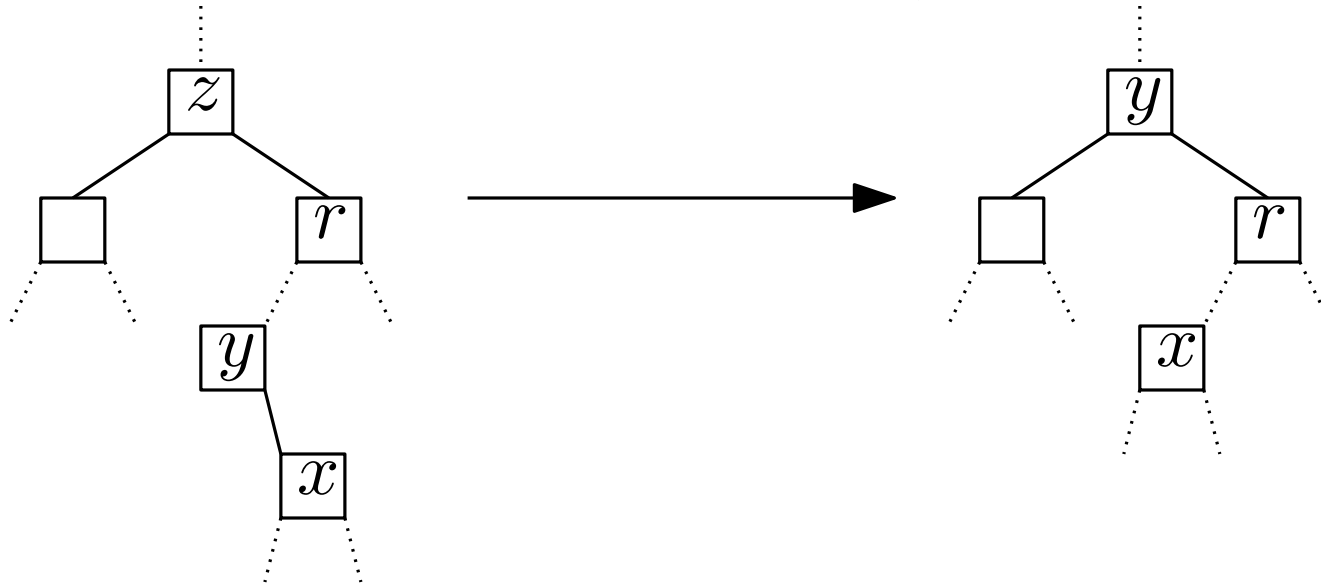


Slet 8.

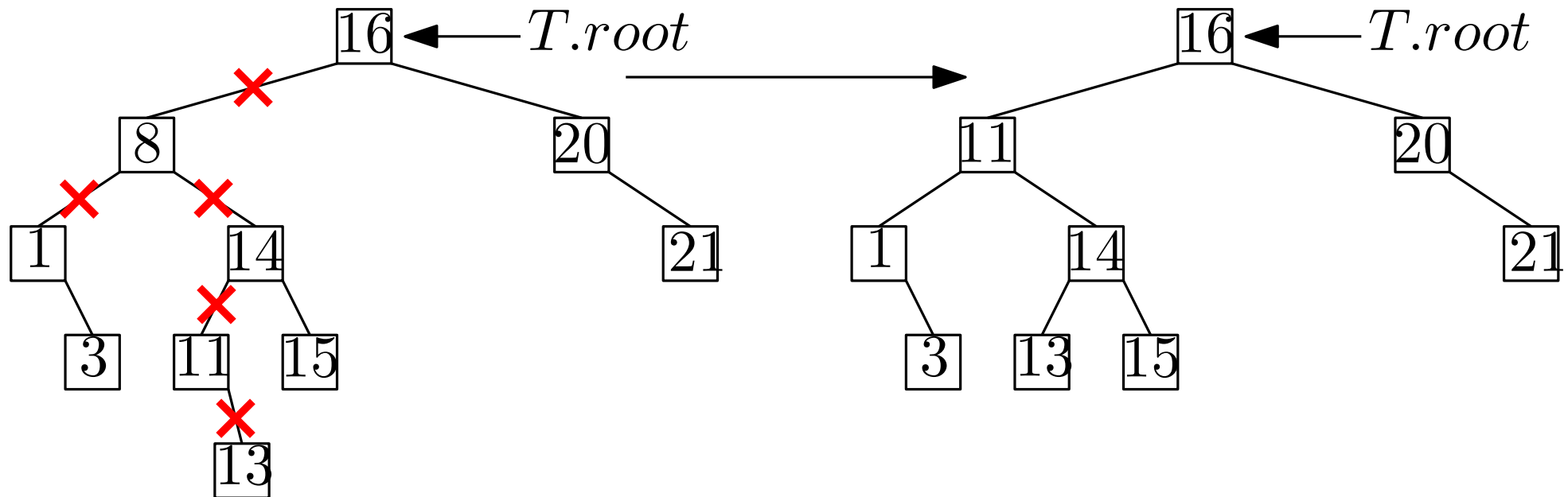


Sletning, 2 børn

Slet z med to børn. Tilfælde 2: $z.right.left \neq \text{NIL}$



Slet 8.



Køretid

Tree-Search(x, k)

if $x == \text{NIL}$ or $k == x.key$

return x

if $k < x.key$

return Tree-Search($x.left, k$)

else

return Tree-Search($x.right, k$)

Tree-Predecessor

Minimum

Maximum

Insert

Delete

Tree-Successor(x)

if $x.right \neq \text{NIL}$

return Tree-Minimum($x.right$)

$y = x.p$

while $y \neq \text{NIL}$ and $x == y.right$

$x = y$

$y = y.p$

return y

Køretid

Tree-Search(x, k)

if $x == \text{NIL}$ or $k == x.key$

return x

if $k < x.key$

return Tree-Search($x.left, k$)

else

return Tree-Search($x.right, k$)

Tree-Predecessor

Minimum

Maximum

Insert

Delete

Tree-Successor(x)

if $x.right \neq \text{NIL}$

return Tree-Minimum($x.right$)

$y = x.p$

while $y \neq \text{NIL}$ and $x == y.right$

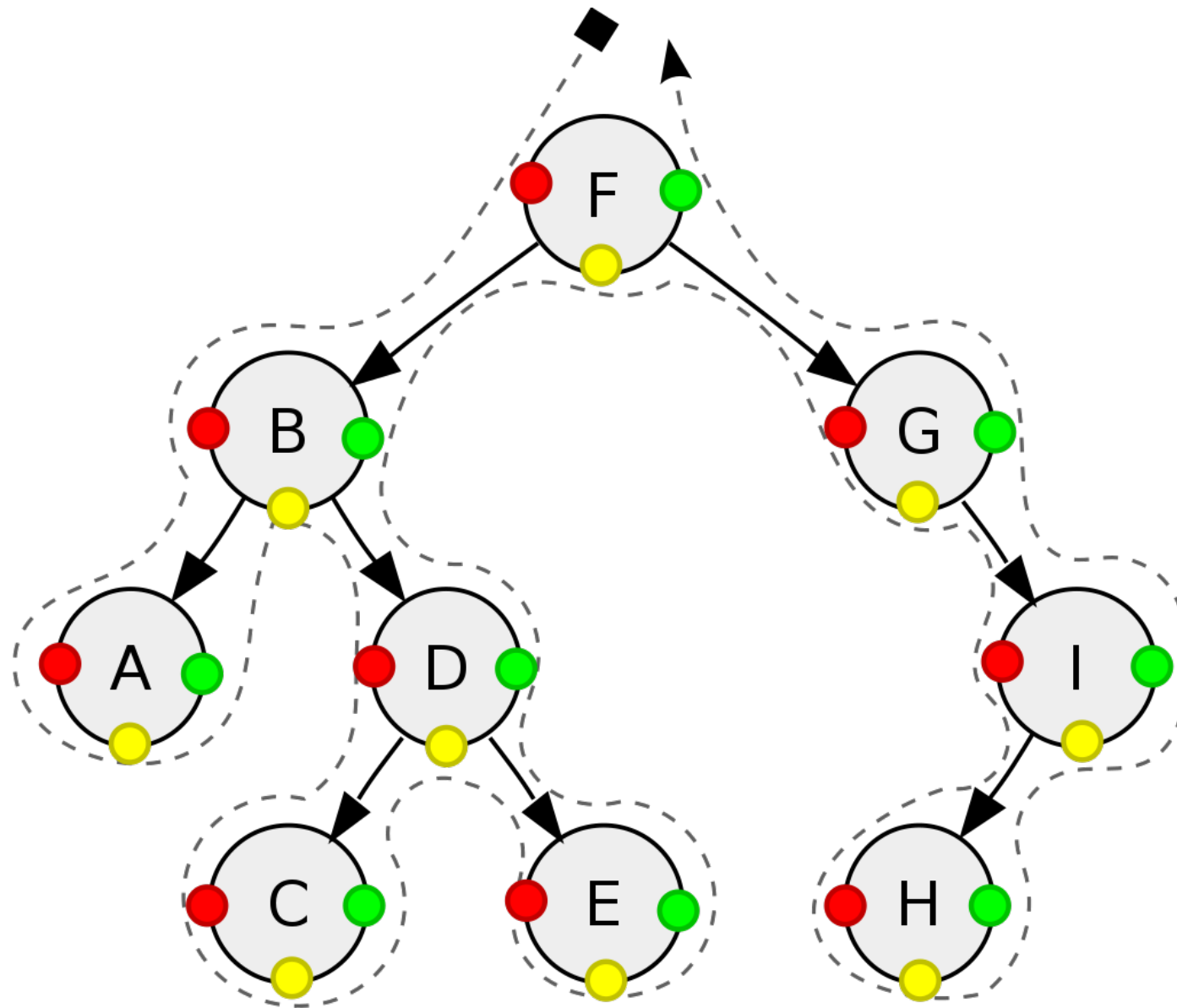
$x = y$

$y = y.p$

return y

Alle tager $\Theta(h)$ tid i værste fald,
hvor h er højden af T .

Pre-, in- og postorder



Pre-, in- og postorder

Preorder-Tree-Walk(x)

if $x \neq \text{NIL}$

print $x.key$

Preorder-Tree-Walk($x.left$)

Preorder-Tree-Walk($x.right$)

Inorder-Tree-Walk(x)

if $x \neq \text{NIL}$

Inorder-Tree-Walk($x.left$)

print $x.key$

Inorder-Tree-Walk($x.right$)

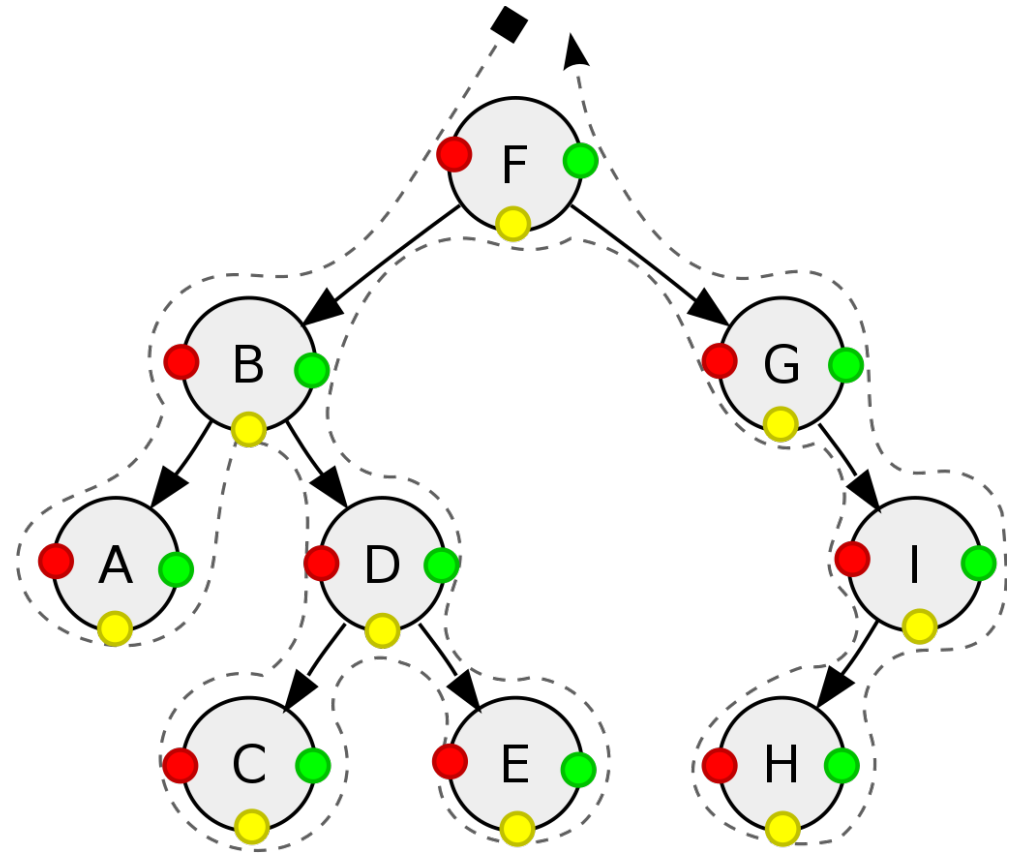
Postorder-Tree-Walk(x)

if $x \neq \text{NIL}$

Postorder-Tree-Walk($x.left$)

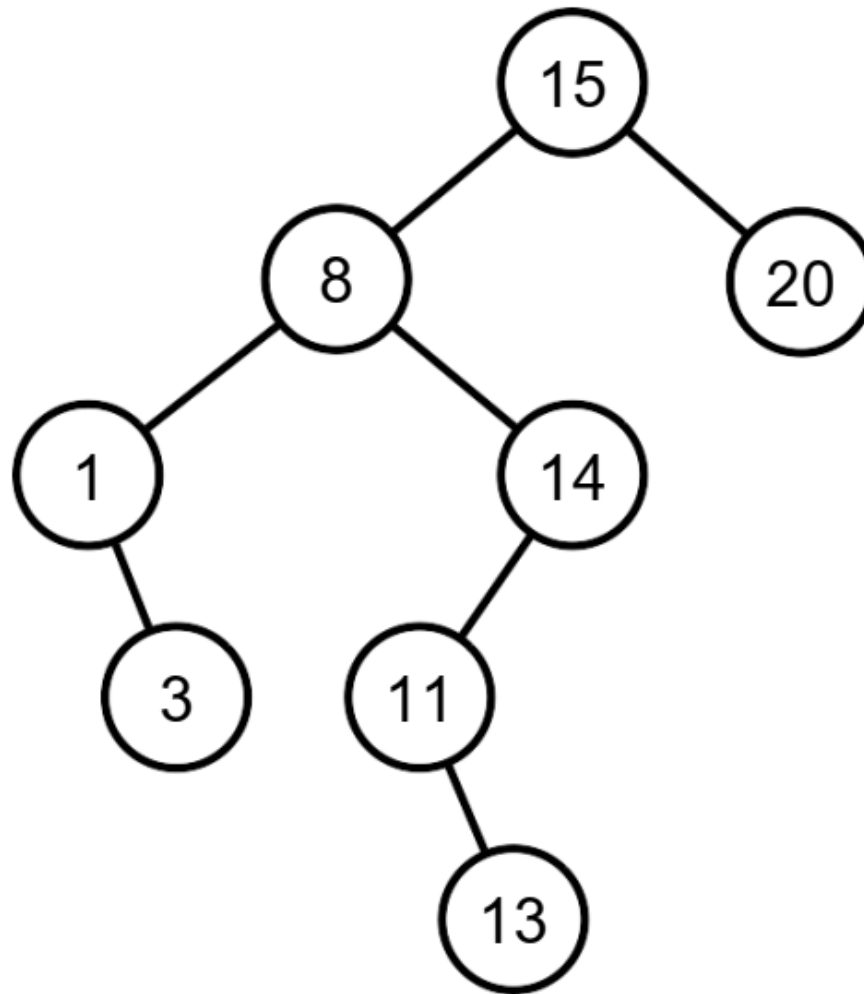
Postorder-Tree-Walk($x.right$)

print $x.key$



Inorder

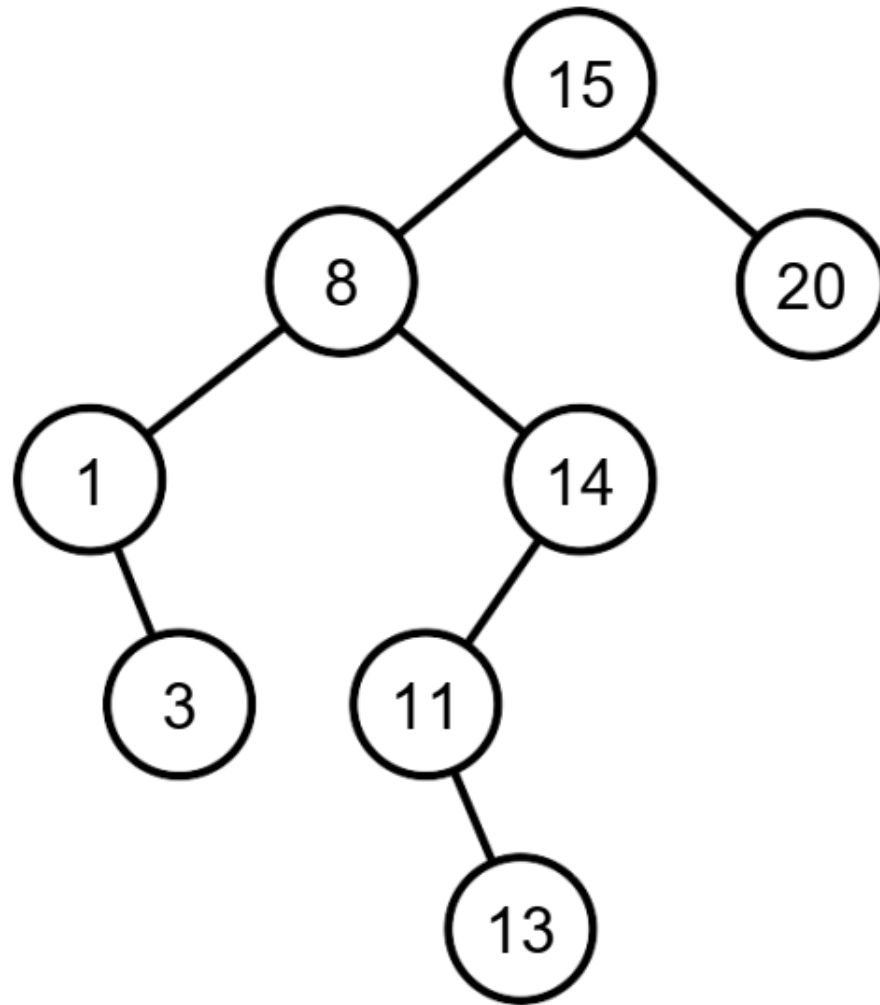
Når man skal løbe knuderne igennem i sorteret rækkefølge.



Inorder: 1, 3, 8, 11, 13, 14, 15, 20

Preorder

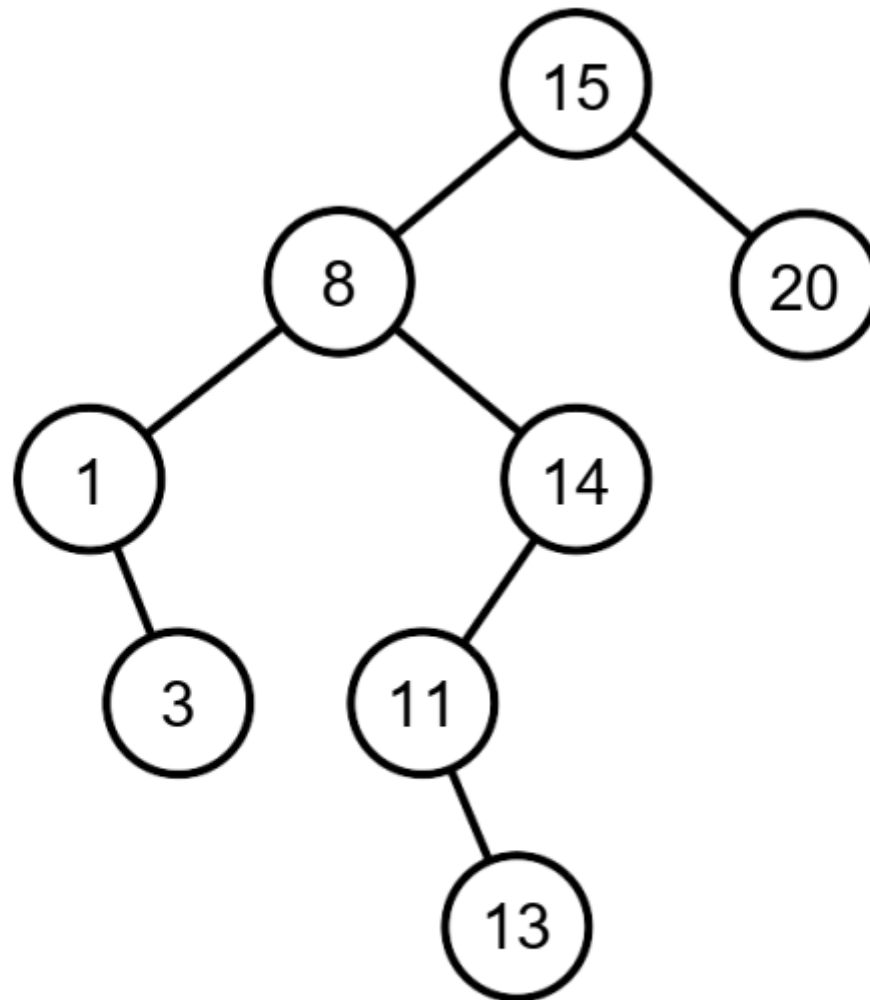
Når man skal lave en kopi.



Preorder: 15, 8, 1, 3, 14, 11, 13, 20

Postorder

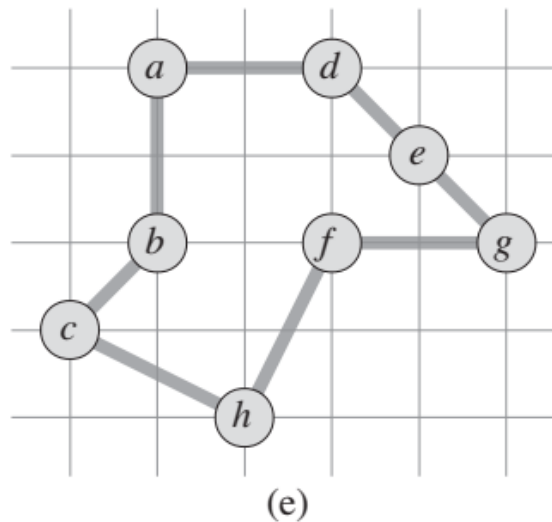
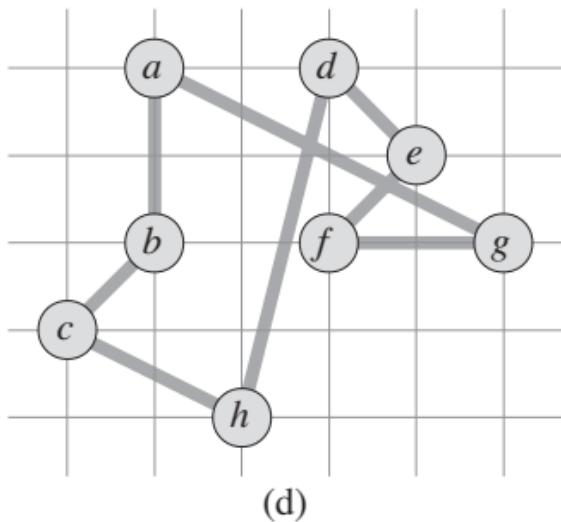
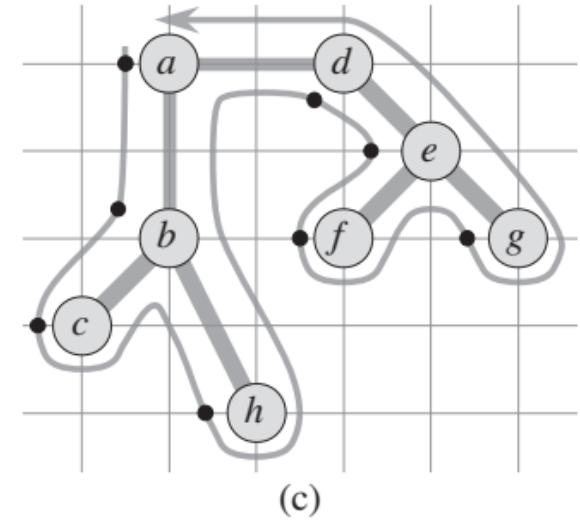
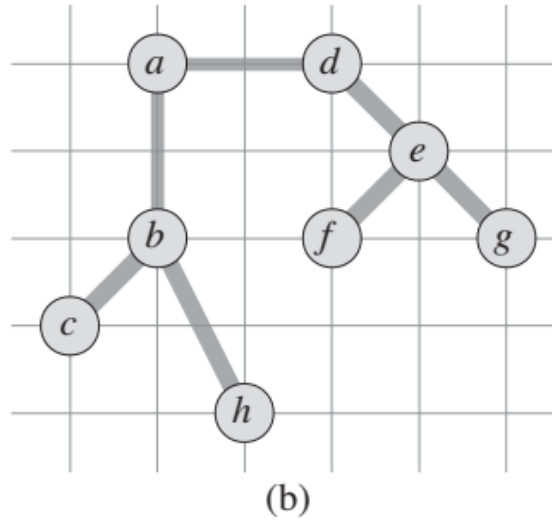
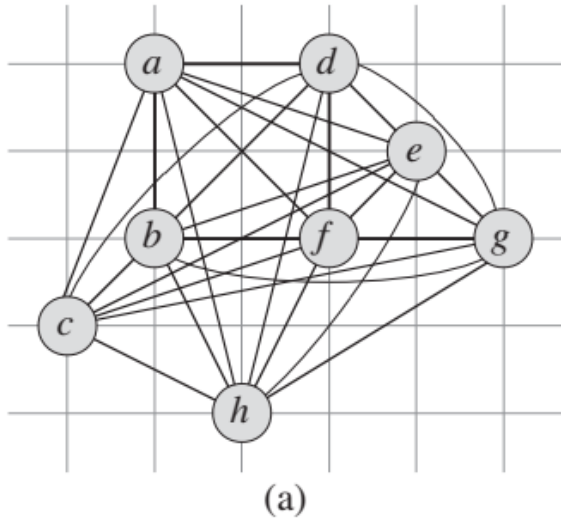
Når man skal løbe knuderne igennem mens man sletter dem.



Postorder: 3, 1, 13, 11, 14, 8, 20, 15

Preorder

Anvendelse i approximationsalgoritme til den handelsrejsendes problem (TSP).



Preorder og postorder

Anvendelse til at repræsentere syntakstræer uden brug af parenteser.

Inorder:

$A*(B-C) + (D+E)$

Preorder:

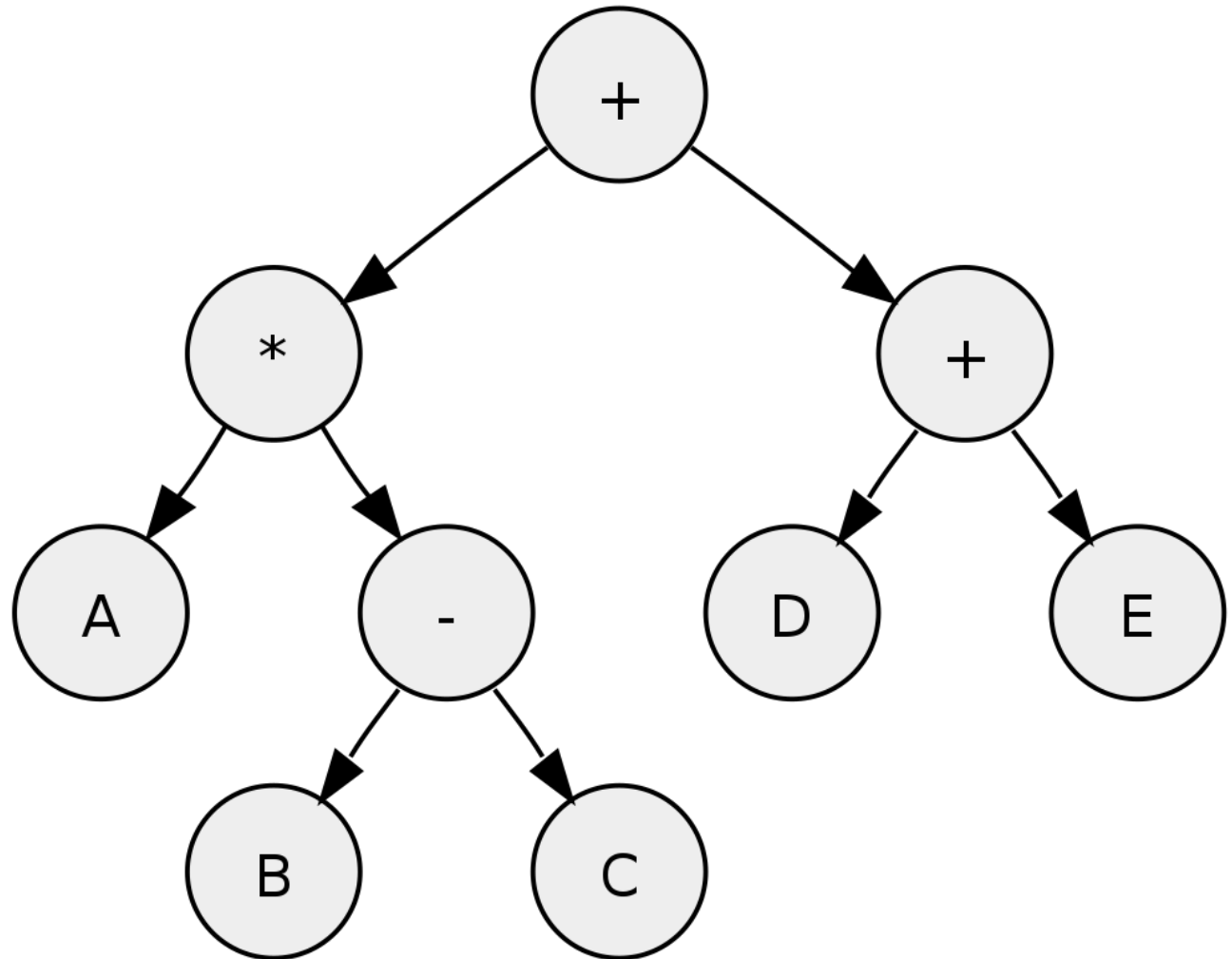
$+ * A - B C + D E$

(polsk notation)

Postorder:

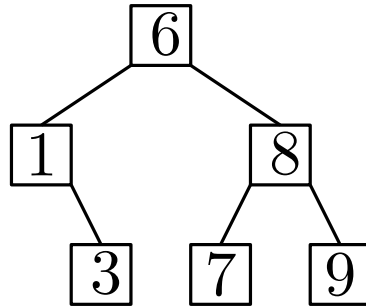
$A B C - * D E + +$

(omvendt polsk notation)



Trægennemløb

Hvad er postorder traversal-rækkefølgen for dette træ?



1,3,6,7,8,9

A

9,7,8,3,1,6

B

6,1,3,8,7,9

C

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9,8,7,6,3,1

D

3,1,7,9,8,6

E