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# **EXPERIMENT - IX**

## **FIRST AND FOLLOW**

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**AIM**

To write a program to simulate the FIRST and FOLLOW for any given grammar.

## THEORY

Each time a predictive parser makes a decision, it needs to determine which production rule to apply to the leftmost non-terminal in an intermediate form, based on the next terminal (i.e. the look-ahead symbol). This is the significance of the FIRST sets: they tell you when a non-terminal can produce the look-ahead symbol as the beginning of a statement, so that it can be matched away and reduce the input. FOLLOW covers the possibility that the leftmost non-terminal can disappear, so that the look-ahead symbol is not actually a part of what we're presently expanding, but rather the beginning of the next construct. This is the significance of the FOLLOW sets: they tell you when a non-terminal can hand you the look-ahead symbol at the beginning of a statement by disappearing. Choosing productions that give epsilon doesn't reduce the input string, but you still have to make a rule for when the parser needs to take them, and the appropriate conditions are found from the FOLLOW set of the troublesome non-terminal.

### FIRST(X) for all Grammar Symbols X

Apply the following rules:

1. If X is terminal,  $\text{FIRST}(X)=X$ .
2. If  $X \rightarrow \epsilon$  is a production then add  $\epsilon$  to  $\text{FIRST}(X)$ .
3. If X is a non-terminal, and  $X \rightarrow Y_1 Y_2 Y_k$  is a production, and  $\epsilon$  is in all of  $\text{FIRST}(Y_1)$ ,  $\text{FIRST}(Y_2)$ , ...,  $\text{FIRST}(Y_k)$ , then add  $\epsilon$  to  $\text{FIRST}(X)$ .
4. If X is a non-terminal, and  $X \rightarrow Y_1 Y_2 Y_k$  is a production, then add a to  $\text{FIRST}(X)$  if for some i, a is in  $\text{FIRST}(Y_k)$  and is in all of  $\text{FIRST}(Y_1)$ ,  $\text{FIRST}(Y_2)$ , ...,  $\text{FIRST}(Y_{k-1})$ .

### FOLLOW(A) for all Non-terminal A

Apply the following rules:

1. If \$ is the input end-marker, and S is the start symbol,  $\$ \in \text{FOLLOW}(S)$ .
2. If there is a production,  $A \rightarrow \alpha B \beta$ , then  $\text{FIRST}(\beta) - \epsilon \subseteq \text{FOLLOW}(B)$ .
3. If there is a production,  $A \rightarrow \alpha B$  or  $A \rightarrow \alpha B \beta$ , where  $\epsilon \in \text{FIRST}(\beta)$  then  $\text{FOLLOW}(A) \subseteq \text{FOLLOW}(B)$ .

## ALGORITHM

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**Algorithm 1** Algorithm to find FIRST(X) for all grammar symbols X

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```
1: Start
2: if X is a terminal then
3:   FIRST(X) = X.
4: end if
5: if  $X \rightarrow \epsilon$  is a production then
6:   Add  $\epsilon$  to FIRST(X).
7: end if
8: if X is a non-terminal and  $X \rightarrow Y_1 Y_2 \dots Y_k$  is a production then
9:   if  $\epsilon$  is in all of FIRST( $Y_1$ ), FIRST( $Y_2$ ), ..., FIRST( $Y_k$ ) then
10:    Add  $\epsilon$  to FIRST(X).
11:   end if
12:   if  $\exists i, a \in \text{FIRST}(Y_i)$  and  $\epsilon$  is in all of FIRST( $Y_1$ ), FIRST( $Y_2$ ), ..., FIRST( $Y_k$ ) then
13:    Add a to FIRST(X).
14:   end if
15: end if
16: Stop
```

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**Algorithm 2** Algorithm to find FOLLOW(A) for all non-terminals A

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```
1: Start
2: if $ is the input end-marker and S is the start symbol then
3:   Add $ to FOLLOW(S).
4: end if
5: if  $A \rightarrow \alpha B \beta$  is a production then
6:   FIRST( $\beta$ ) -  $\epsilon \subseteq \text{FOLLOW}(B)$ .
7: end if
8: if there is a production,  $A \rightarrow \alpha B$  or  $A \rightarrow \alpha B \beta$ , where  $\epsilon \in \text{FIRST}(\beta)$  then
9:   FOLLOW(A)  $\subseteq \text{FOLLOW}(B)$ 
10: end if
11: Stop
```

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## SOURCE CODE

```
def FIRST(X):
    global production
    global T
    global N
    first=[]
    if X in T:
        first.append(X)
    if X in N:
        if '#' in production[X]:
            first.append('#')
        for pr in production[X]:
            for p in pr:
                f=FIRST(p)
                for i in f:
                    if i not in first:
                        first.append(i)
                if '#' not in f:
                    break
        flag=0
        for p in pr:
            if '#' not in FIRST(p):
                flag=1
                break
        if flag==0:
            if '#' not in first:
                first.append('#')
    return first

def FOLLOW():
    global production
    global follow
    global first
    global N
```

```

    for X in N:
        for pr in production[X]:
            if pr[-1] in N and '$' not in follow[pr[-1]]:
                follow[pr[-1]].append('$')
            for i in range(len(pr)-1):
                if pr[i] in N:
                    f=first[pr[i+1]]
                    for x in f:
                        if (x not in follow[pr[i]]
                            and x!='#'):
                            follow[pr[i]].append(x)

    for X in N:
        for pr in production[X]:
            for p in pr:
                if ((p==pr[-1]
                    or '#' in first[pr[pr.index(p)+1]])
                    and p in N):
                    f=follow[X]
                    for x in f:
                        if x not in follow[p]:
                            follow[p].append(x)

global production
global T
global N
global follow
global first
production={}
n=int(input("Enter the number of productions:"))
print("Enter the productions:\t\t(Use \# for epsilon)")
T=[]
N=[]
s=[]
first={}
follow={}

```

```

for _ in range(n):
    pr=input().split(" -> ")
    if pr[0] in production.keys():
        production[pr[0]].append(pr[1])
    else:
        production[pr[0]]=[pr[1]]
    if pr[0] not in N:
        N.append(pr[0])
    if pr[0] not in s:
        s.append(pr[0])
    if pr[1]=='#':
        continue
    for i in pr[1]:
        if i not in s:
            s.append(i)

for i in s:
    if i not in N:
        T.append(i)

for i in s:
    first[i]=FIRST(i)
    if i in N:
        follow[i]=[]
follow[N[0]].append('$')
print("Set of Terminals : ",end="")
for i in range(len(T)-1):
    print(T[i],end=",")
print(T[-1])
print("Set of Non-Terminals : ",end="")
for i in range(len(N)-1):
    print(N[i],end=",")
print(N[-1])
print("FIRST")
for i in s:
    print("\t",i,end=" : ")

```

```
        for f in first[i]:
            print(f,end=" ")
        print()
    print("FOLLOW")
    FOLLOW()
    for i in follow.keys():
        print("\t",i,end=" : ")
        for f in follow[i]:
            print(f,end=" ")
        print()
```



## SAMPLE OUTPUT

```
user@adithya-d-rajagopal:~/s7/cd$ python3 p9.py
Enter the number of productions:8
Enter the productions:          (Use \# for epsilon)
E -> TR
F -> (E)
F -> i
R -> #
R -> +TR
T -> FY
Y -> #
Y -> *FY
Set of Terminals : (,),i,+,*
Set of Non-Terminals : E,F,R,T,Y
FIRST
      E : ( i
      T : ( i
      R : # +
      F : ( i
      ( : (
      ) : )
      i : i
      + : +
      Y : # *
      * : *
FOLLOW
      E : $ )
      T : + $ )
      R : $ )
      F : * + $ )
      Y : $ + )
user@adithya-d-rajagopal:~/s7/cd$
```

## **RESULT**

A program to simulate the FIRST and FOLLOW for any given grammar has been implemented using Python and the outputs have been verified.