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# **EXPERIMENT - XI**

## **SHIFT-REDUCE PARSER**

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**AIM**

To construct a shift-reduce parser for a given language.

## **THEORY**

### **Shift-Reduce Parser**

Shift-reduce parsing attempts to construct a parse tree for an input string beginning at the leaves and working up towards the root. In other words, it is a process of reducing (opposite of deriving a symbol using a production rule) a string  $w$  to the start symbol of a grammar. At every (reduction) step, a particular substring matching the RHS of a production rule is replaced by the symbol on the LHS of the production.

A general form of shift-reduce parsing is LR (scanning from Left to right and using Right most derivation in reverse) parsing, which is used in a number of automatic parser generators like Yacc, Bison, etc. A handle of a string is a substring that matches the RHS of a production and whose reduction to the non-terminal (on the LHS of the production) represents one step along the reverse of a rightmost derivation toward reducing to the start symbol. The set of prefixes of right sentential forms that can appear on the stack of a shift-reduce parser are called viable prefixes. It is always possible to add terminal symbols to the end of a viable prefix to obtain a right-sentential form.

## ALGORITHM

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**Algorithm 1** Algorithm to construct a Shift-reduce Parser

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```
1: Loop forever:
2:   for top of stack symbol s and next input symbol a, case action of T[s,a]
3:     Shift x: (X is a STATE number)
4:       push a, then x on the top of the stack and
5:       advance ip to point to the next input symbol
6:     Reduce y: (y is a production number)
7:       Assume that the production is of the form  $A \rightarrow \beta$ 
8:       Pop  $2 * |\beta|$  symbols of the stack.
9:       At this point the top of stack symbol should be a state number say s'.
10:      Push A, then goto of T[s',A] on the top of the stack.
11:      Output the production  $A \rightarrow \beta$ 
12:      accept:
13:        return - - - Successful parse
14:      default:
15:        error - - - Input string is not in language
```

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## SOURCE CODE

```
def printStack():
    global Stack
    for i in Stack:
        print(i, end=" ")
    print("\t\t", end=" ")

def reduce():
    global Stack
    global handle
    global prevhandle
    if Stack[-1]=="i ":
        Stack.pop()
        Stack.append("E")
        prevhandle=handle[0]
        return True
    if len(Stack)>=3:
        if Stack[-1]=="E" and Stack[-3]=="E":
            op=Stack.pop()
            op=Stack.pop()
            if op=="+":
                prevhandle=handle[1]
            elif op=="*":
                prevhandle=handle[2]
            return True
        elif Stack[-1]==")" and Stack[-2]=="E" and Stack[-3]=="(":
            op=Stack.pop()
            op=Stack.pop(-2)
            prevhandle=handle[3]
            return True
    return False

def Shift_Reduce_Parser(str):
    global Stack
```

```

global handle
global prevhandle
T=['+', '*', 'i', '(', ')', '$']
Stack=['$']
ip=0
handle=['i', 'E+E', 'E*E', '(E)']
print("STACK\t\tINPUT\t\tACTION")
print("$\t\t"+str+"\t-")
while ip<len(str):
    Stack.append(str[ip])
    ip=ip+1
    printStack()
    if ip==len(str):
        print("-\t\tShift")
        break
    print(str[ip:], "\t\tShift")
    while(reduce()):
        printStack()
        print(str[ip:], end="\t\t")
        print("Reduce E -> "+prevhandle)
    if Stack[0]=='$' and Stack[1]=='E' and Stack[2]=='$':
        return True
return False

global Stack
global handle
global prevhandle
print("The Grammar is:")
print("E -> E+E | E*E | (E) | i")
s=input("Enter the string to be parsed:")
w=s+'$'
if (Shift_Reduce_Parser(w)):
    print("Successfully parsed")
else:
    print("Error in parsing")

```

## SAMPLE OUTPUT

```
user@adithya-d-rajagopal:~/s7/cd$ python3 p11.py
The Grammar is:
E -> E+E | E*E | (E) | i
Enter the string to be parsed:i+(i*i)
STACK          INPUT          ACTION
$              i+(i*i)$        -
$i             +(i*i)$        Shift
$E             +(i*i)$        Reduce E -> i
$E+            (i*i)$        Shift
$E+(           i*i)$         Shift
$E+(i          *i)$         Shift
$E+(E          *i)$         Reduce E -> i
$E+(E*         i)$         Shift
$E+(E*i        )$         Shift
$E+(E*E        )$         Reduce E -> i
$E+(E          )$         Reduce E -> E*E
$E+(E)         $          Shift
$E+E          $          Reduce E -> (E)
$E            $          Reduce E -> E+E
$E$           -          Shift
Successfully parsed
user@adithya-d-rajagopal:~/s7/cd$
```

## **RESULT**

A program to construct a shift-reduce parser has been implemented using Python and the outputs have been verified.