

Let's start at 9:02 PM

L51

Modular Arithmetic & Euclid's GCD Algorithm

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# RECAP

What is this Modulo Operator? ( $\cdot \%$ )

$a \% b \Rightarrow$  remainder when  $a$  is divided  
by  $b$ .

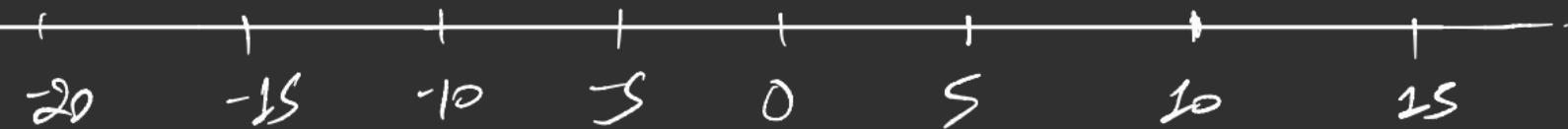
$$a = q * b + r$$

$$r = a - q * b$$

$$q = \left\lfloor \frac{a}{b} \right\rfloor$$

$$\Rightarrow r = a - \left\lfloor \frac{a}{b} \right\rfloor * b$$

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$$10 \div 3 = 1$$

$$13 \div 2 = 1$$

$$18 \div 5 = 3$$

$$5 \div 10 = 5$$

$$-6 \text{ rem } 3 = 0$$

$$-18 \text{ rem } 5 = 2$$

$$-25 \text{ rem } 7 = 3$$

## Weird behaviour on negative numbers

$$\frac{13}{5} = 2 \dots$$
$$-2 \dots$$

$$\frac{13}{5} \Rightarrow 2$$

$$\frac{-13}{5} \Rightarrow -2$$

# What is Modulo Arithmetic?

$$a+b \quad \xrightarrow{\hspace{1cm}} \quad a+b \quad \times$$

$$\xrightarrow{\hspace{1cm}} \quad (a+b) \text{ } \% \text{ } m \quad \checkmark$$

## Why Modular Arithmetic?

$[0, \text{ mod}-1]$

Eg. Given  $N$ , find  $\text{fact}(N)$ .

↑

$$1 \leq N \leq 10^{10}$$

$$\text{mod} = 1000000007$$

$$\text{ans} = 5.6290168 \pi^{10^{30}}$$

↙

$$\text{ans} \% \text{ mod} \Rightarrow 5012531$$

ans % mod.

$$1 * 2 * 3 * \underbrace{n}_{N} \dots$$

$$\text{ans} = 1$$

for( $i=2$ ;  $i < N$ ;  $i++$ )

$$\text{ans} = (\text{ans} * i) \% \text{mod}$$

## Addition

$$(a+b) \cdot m = ((ax.m) + (bx.m)) \cdot m$$

$$a = q_1 \cdot m + r_1 \quad (a \cdot m > r_1)$$

$$b = q_2 \cdot m + r_2 \quad (b \cdot m > r_2)$$

$$q_1 \cdot m + r_1 + q_2 \cdot m + r_2 \Rightarrow (q_1 + q_2) \cdot m + r_1 + r_2$$

$$a = 10^{18} + 5 \quad , \quad b = 10^{30} + 5127$$

$$(a+b) \div 10$$

$$\text{m=10}$$

$$\Rightarrow (r_1 + r_2) \div 10$$

$$\Rightarrow (5+7) \div 10 \Rightarrow 2$$

## Subtraction

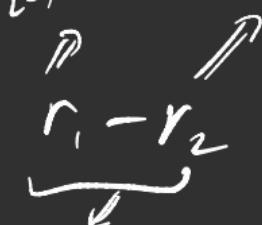
$$(a - b) \times m \Rightarrow ((a \times m) - (b \times m) + m) \times m$$

$$a = q_1 \times m + r_1$$

$$b = q_2 \times m + r_2$$

$$\Rightarrow a - b \Rightarrow (q_1 - q_2) \times m + \underbrace{r_1 - r_2}_{\left[ -(m-1), m-1 \right]}$$

$\left[ 0, m-1 \right]$        $\left[ 0, m-1 \right]$



$$a = 10^{31} + 5 \quad , \quad b = 10^{30} + 5127$$

$$(a - b) \times 10$$

$$\text{m=10}$$

$$\Downarrow (r_1 - r_2 + 10) \times 10$$

$$\Downarrow (5 - 7 + 10) \times 10$$

$$\Rightarrow 8 \times 10 \Rightarrow 8$$

## Multiplication

$$(a * b) \times m \implies ((a \times m) * (b \times m)) \times m$$

$$a = q_1 m + r_1$$

$$b = q_2 m + r_2$$

$$a * b = (q_1 m + r_1) * (q_2 m + r_2)$$

$$= q_1 q_2 m^2 + q_1 r_2 m + q_2 r_1 m + r_1 r_2$$

$$\Rightarrow \underbrace{(q_1 q_2 m + q_1 r_2 + q_2 r_1) * m}_{\text{Irrelevant}} + \underbrace{r_1 * r_2}_{\text{Relevant}}$$

$$a = 10^{31} + 5 \quad , \quad b = 10^{30} + 5127$$

$$(a * b) \times 10$$

$$\text{m} = 10$$

$$\Rightarrow (r_1 * r_2) \times 10$$

$$\Rightarrow (5 * 7) \times 10$$

$$\Rightarrow 35 \times 10 \Rightarrow 5$$

Find  $a^n$  under modulo

Given  $a \geq n$ , find  $a^n$ .

$$1 \leq a \leq 10^9$$
$$1 \leq n \leq 10^{18}$$

Since answer may be large print ans  $\% (10^9 + 7)$

## Introduction to GCD

Also called  
HCF (Highest common  
factor)

↳ stands for  
greatest common  
Divisor

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The largest such numbers which divides  
both A & B.

$$\gcd(2, 8) \Rightarrow 2$$

$$\gcd(24, 32) \Rightarrow 8$$

$$\gcd(15, 15) \Rightarrow 1$$

$$\gcd(0, 10) \Rightarrow 10$$

## Brute Force

```
for ( i = min(a,b); i >= 1; --i )  
    if ( a[i] == 0 && b[i] == 0)  
        return i;
```

Time  $\Rightarrow O(\min(a,b))$

## A better method

$$a = 10500 = 2^2 * 3^1 * 5^3 * 7$$

$$b = 198 = 2^1 * 3^2 * 11$$

$$O(\sqrt{\max(a,b)})$$

$$\gcd = 2^1 * 3^1 \Rightarrow 6$$

# Euclid's GCD Algorithm

$(a > b)$

Observation 1

$(a, b)$



$\{f_1, f_2, f_3\}$

$(a-b, b)$



$\{f_1, f_2, f_3\} \longrightarrow$

If a number is a divisor of both a, b.  
⇒ It will be a divisor of a-b also.

$$\begin{matrix} a & b \\ (25, 8) \end{matrix} \Rightarrow 1, 2, 4,$$

$$\begin{matrix} ab, b \\ (20, 4) \end{matrix}$$

## Observation 2

$$\begin{pmatrix} (a-b) & b \\ \downarrow & \downarrow \end{pmatrix} \longrightarrow \begin{pmatrix} a & b \end{pmatrix}$$

$$\begin{aligned} (a-b) &\Rightarrow q_1 f \\ b &\Rightarrow q_2 f \end{aligned}$$

$$\begin{aligned} a &\Rightarrow (a-b) + b \\ &\quad \Downarrow \qquad \Downarrow \\ a &\Rightarrow q_1 f + q_2 f \\ &\Rightarrow (q_1 + q_2) * f \end{aligned}$$

If a number is a divisor of both  $(a-b)$  &  $b$ , then it will a divisor of  $a$ .

$\Rightarrow$  Common divisors  $a, b$  &  $a-b, b$   
will be same set of numbers.

$$a = 28 \quad b = 8$$

$$(a-b, b) \downarrow \\ (20, 8) \Rightarrow 1, 2, 4$$

$$(28, 8) \Rightarrow 1, 2, 4$$

$$\gcd(a, b) = \gcd(a-b, b)$$

$$\gcd(28, 8) \Rightarrow \gcd(20, 8)$$

$a > b$

### Observation 3

$$\gcd(a, b) = \gcd(a-b, b)$$

$$\gcd(a-b, b) = \gcd(a-2b, b)$$

$$\gcd(a-2b, b) = \gcd(a-3b, b)$$

$$a - \lfloor \frac{a}{b} \rfloor * b$$

$$\gcd(a \mod b, b)$$

Example

$$\text{gcd}(28, 8) \Rightarrow 4$$

$$\hookrightarrow \text{gcd}(4, 8) \rightarrow \underbrace{\text{gcd}(8, 4)}$$

$$\overbrace{\quad}^{\longrightarrow} \text{gcd}(0, 4)$$

$$\text{gcd}(a, b)$$

( $a > b$ )

$$\downarrow$$
$$(b, a \bmod b)$$

Let's implement

```
int gcd (a, b) {  
    if (min(a, b) == 0)  
        return max(a, b);  
    int mx = max(a, b), mn = min(a, b);  
    return gcd (mn, mx % mn);  
}
```

## Time Complexity

cb algorithm euclid's algorithm.

$$\Rightarrow \log(\min(a, b))$$

$\gcd(a, b)$

# *Thank You!*

Reminder: Going to the gym & observing the trainer work out can help you know the right technique, but you'll muscle up only if you lift some weights yourself.

So, PRACTICE, PRACTICE, PRACTICE!