

Bonus Lecture - 1

Digit DP

Introduction

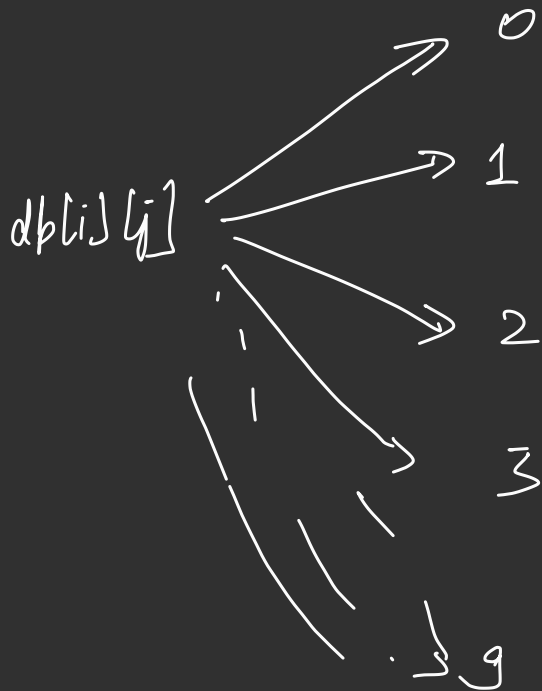
Given a length N , find the number of arrays of size N such that all the array elements are b/w 0 and 9, and the sum of array elements is divisible by K .

$dp[N][K]$

$dp[i][j]$

No. of ways
to fill arr
till index i

$sum(0 \dots i) \% K = j$



i, j



```
dp[i][j] = 0;  
for (d = 2; d <= g; ++d)  
    dp[i][j] += dp[i-1][(j-d+K)%K];
```

Let's say $N = 5$, don't you think that the solution will be equal to number of whole numbers less than or equal to 99999 having digit sum divisible by K ?

0-9 0-9 0-9 0-9 0-9 \Rightarrow sum of digits
has to be a
multiple of
 K .

$\hookrightarrow [0, 0, 5, 1, 2]$

$\hookrightarrow [1, 0, 9, 1, 1] \xrightarrow{512}$

$\Rightarrow 10,911$

What if we are asked to find number of whole numbers less than or equal to NUM, having digit sum divisible by K?

NUM \Rightarrow 50239

numArr \Rightarrow [5, 0, 2, 3, 9]

No. of arrays of size 5 where $a[i] \in [0, 9]$ & sum of array elements div. by K, & are lexicographically \leq numArr.

Let's do some problems!

1. Digit Sum

Intuition / Solution

num = 5 0 2 3 9

$n = \text{numDigits}(\text{num})$
 $K \Rightarrow \text{divisor}$

$dp[N][K][2]$

↓
index until
which digits
have been
filled

↘
digSum
% K
value

↘
The prefix
is for
numPrefix?
or not

$dp[i][j][0] \Rightarrow$ The no. of prefixes of num
(0---i) that they are $< \text{Corresp}$
pref
& $\text{digSum} - K = 2^j$

$dp[i][j][1] \Rightarrow$ The no. of prefixes of num
(0---i) that they are $= 2^j \text{Corresp}$
pref
& $\text{digSum} - K = 2^j$



0 or 1

$(0 \dots n-1)$ ($\text{digSum} + 1 \cdot K \geq 0$) ($< 22 = 2$
both are fine)

$\text{ans} = (\text{dp}[n-1][0][0] + \text{dp}[n-1][0][1]) \% \text{mod};$

$\text{return } (\text{ans} - 1 + \text{mod}) \% \text{mod};$

10 - - - - - i-1 (i, j)

⇒ equal

$dp[i][j][1]$ $\xrightarrow{\text{ } j^{\text{th}} \text{ digit of num} \Rightarrow d_i}$ $dp[i-1][j-d_i+1][1]$

10 - - - - - 11 $\textcircled{i+j}$ $\Rightarrow di \Rightarrow 5 \Rightarrow \text{loss}$

$dp[i][j][0]$

$\downarrow d$
 $d \geq d_i$

$\nearrow d$
 $d < d_i$

$dp[i][j][0] + 2 \cdot dp[i-1][j-d+1][0]$
 $dp[i][j][0] + 2 \cdot dp[i-1][j-d+1][1]$

$$\Leftrightarrow dp[i][j][0] + 1 = dp[i-1][j-d+1][1][0]$$

Let's implement

2. Stepping Number

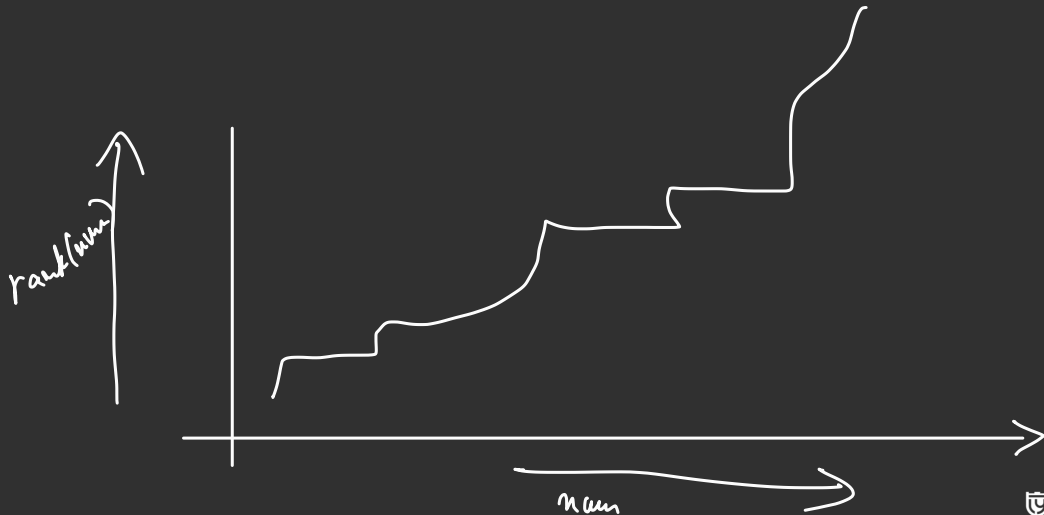
$\text{getRank}(n) \Rightarrow$ Number of whole numbers $\leq n$
that are stepping numbers

$\text{getRank}(\text{ans}) - \text{getRank}(n) \Rightarrow K$

$\text{rankDiff} \Rightarrow K$

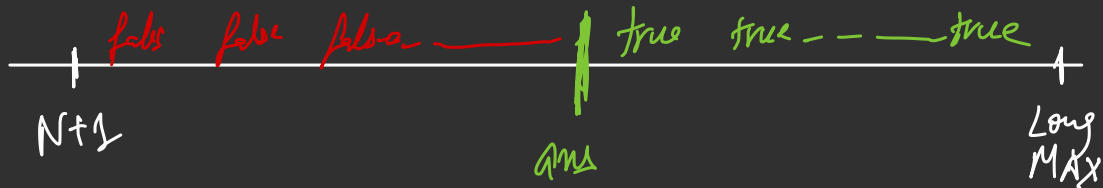
Intuition / Solution

$$N = 7, \quad K = 6$$



find the smallest number

s.t. $\text{getRank}(\text{ans}) - \text{getRank}(n) \geq K$



$$\overbrace{\hspace{10em}} \Rightarrow n \approx \text{numDigits}(num) \\ \leq num$$

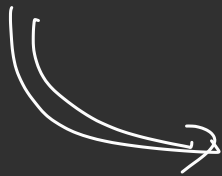
↪ every consec. digits
 $\text{abs}(d_i - d_{i+1}) \leq 1$

$dp[n][10][2]$

equal.

$$dp[i][d][1] \Rightarrow d = z_{d_i}$$

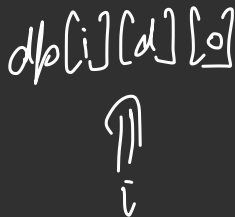
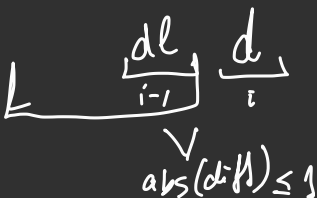
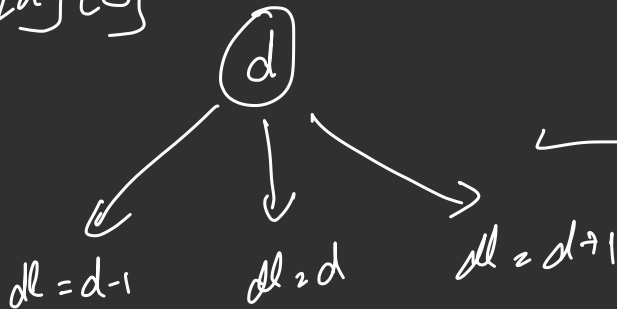
$$\&\& \text{abs}(d_i - d_{e_i}) \leq 1$$



$$dp[i][d][1] = dp[i-1][d_{i-1}][1]$$

$dp[i][d][0]$

less



$d < d_i \Rightarrow i^{\text{th}}$ smaller

$d \geq d_i \Rightarrow$ Should've already
been smaller $(i-1)$

$$dp[i][d][0] += dp[i-1][d][0];$$

if $(d < d_i)$

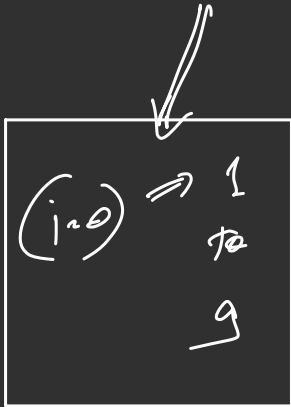
$$dp[i][d][0] += dp[i-1][d][1];$$

5 4 3 2 2 3 5

0 0 5 4 4 5 6 \Rightarrow 5 4 4 5 6 ✓

~~5 4 0 5 6~~

num = 543218 \Rightarrow n = 6



Find all the st—

digits ≤ 5

Let's implement

Thank You!

Reminder: Going to the gym & observing the trainer work out can help you know the right technique, but you'll muscle up only if you lift some weights yourself.

So, PRACTICE, PRACTICE, PRACTICE!