

Bonus Lecture - 3

Sparse Table

RECAP

Range max/min/gcd/lcm /- --



Introduction

Inefficient

Preprocessing



$O(N^2)$

Efficient

Querying



$O(1)$

No preprocessing

$\hookrightarrow O(1)$

Inefficient querying

$\hookrightarrow O(N)$

Preprocess answers only for subarrays of
length = 1, 2, 4, 8, 16, ...

Intuition behind Sparse Table



Answer for a
few important
subarrays will be
stored

The magic happens when we realize that any number can be split into sum of powers of 2



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More detailed Explanation

Preprocessing

$dp[N][\log_2 N + 1] \Rightarrow dp[i][k]$ will represent
the answer to the subarray
that starts at index i
and has a length of 2^k .

$0, 1, \dots, N-1$
 $2^0, 2^1, 2^2, \dots, 2^{\log_2 N}$

$N \times \log_2 N$ dp states

$$dp[i][0] \Rightarrow arr[i \dots i] \quad len = 1$$

$$dp[i][1] \Rightarrow arr[i \dots i+1] \quad len = 2$$

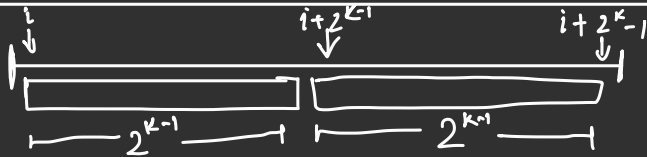
$$dp[i][2] \Rightarrow arr[i \dots i+3] \quad len = 4$$

$$dp[i][3] \Rightarrow arr[i \dots i+7] \quad len = 8$$

Preprocessing for Range Min

// Base Cases

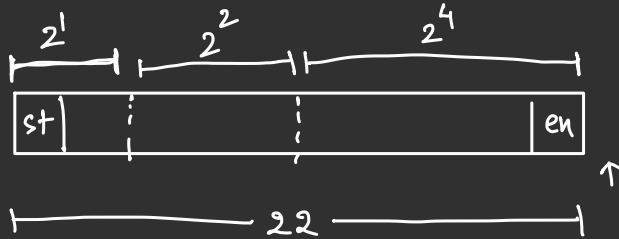
$$dp[i][0] = arr[i];$$



$$dp[i][k] = \min(dp[i][k-1], dp[i + (1 \ll (k-1))][k-1]);$$

Querying:

$$q = [st, en] \quad len = en - st + 1$$



$$ans \Rightarrow a_3, \quad i = st + 2^1 + 2^2 + 2^4$$

$$\Rightarrow 22$$
$$\Downarrow$$
$$2^1, 2^2, 2^4$$

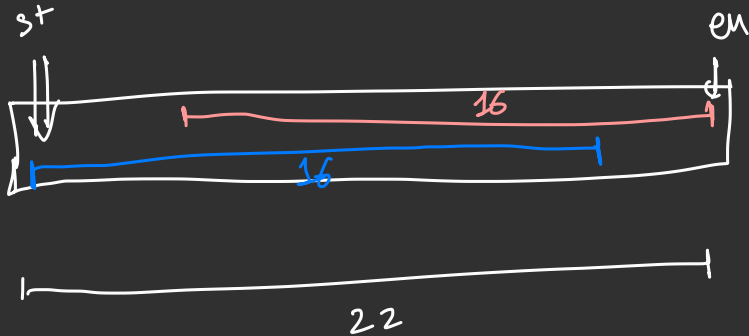
- ~~0~~
- 1 ✓
- 2 ✓
- ~~3~~
- 4 ✓

Let's do some problems!

1. Range Min Queries

Let's implement

Further optimisation for querying



$$N = 5$$

$$\begin{bmatrix} 1, & 2, & 3, & 10, & 50 \end{bmatrix}$$

0 1 2 3 4

2. Sereja and D

$t = 1, d = 1$	\Rightarrow	$i = 0 \Rightarrow 1$
$t = 5, d = 3$	\Rightarrow	$i = 0 \Rightarrow 1$
$t = 11, d = 7$	\Rightarrow	$i = 0 \Rightarrow 1$
$t = 100000, d = 1$	\Rightarrow	$i = 4 \Rightarrow 5$
$t = 10^6, d = 10^6$	\Rightarrow	$i = 0 \Rightarrow 1$
$t = 11, d = 6$	\Rightarrow	$i = 3 \Rightarrow 4$

$$\underline{k \geq i}$$

1.) abs diff for consec.
 $\leq d$

2.) $a[k] \leq t$

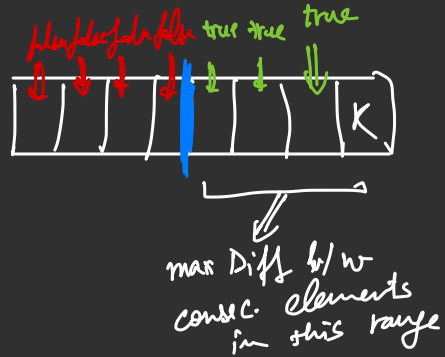
3.) $a[k+1] > t$

Intuition / Solution

Query \Rightarrow Unique K

Smallest i (s.t. $i \leq K$)

$$\hookrightarrow a[i+1] - a[i] \leq d \quad \forall i \in [i, K-1]$$

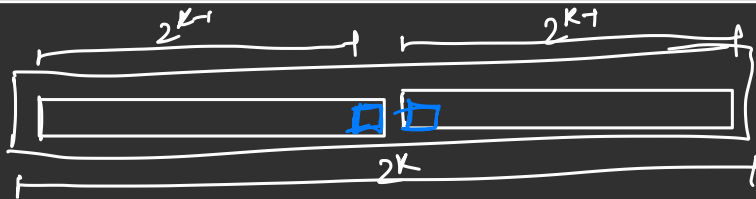


- 1.) Use upper-bound to find accurate K .
- 2.) Use Binary Search to find smallest i s.t. $\text{arr}[i \dots K]$ is *validity*.
- a.) Check \Rightarrow Use sparse Table to find range max Diff & check whether that is $\leq d$ or not?

Pre processing

$$dp[i][0] \Rightarrow 0$$

len = 1

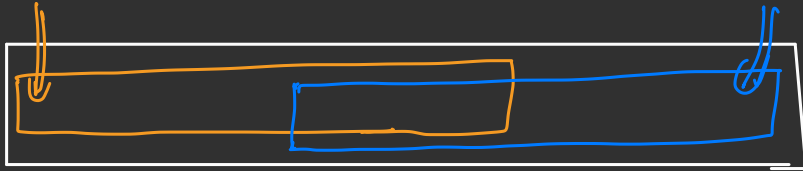


$$dp[i][k] = \max(dp[i][k-1], dp[i + (1 \ll (k-1))][k], arr[i + (1 \ll (k-1))] - arr[i + (1 \ll (k-1)) - 1])$$

Querying

starts at st

ends at en



largest pow of 2 $\geq (x) \leq \text{len} \Rightarrow \log_2 \lceil \text{len} \rceil$

$$2^k > \frac{\text{len}}{2}$$

Let's implement

Thank You!

Reminder: Going to the gym & observing the trainer work out can help you know the right technique, but you'll muscle up only if you lift some weights yourself.

So, PRACTICE, PRACTICE, PRACTICE!