

# Unit - II

## Heaps

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**COURSE CODE: 21AI33**



# Priority Queues

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- ❑ In a priority queue, the element with highest (or lowest) priority is deleted from the queue, while elements with arbitrary priority are inserted.
- ❑ A data structure that supports these operations is called a max(min) priority queue.
- ❑ A priority queue can be implemented by a simple, unordered linked list.
- ❑ Insertions can be performed in  $O(1)$  time. However, a deletion requires a search for the element with the largest priority followed by its removal.
- ❑ The search requires time linear in the length of the linked list.
- ❑ A max heap is used, both of these operations can be performed in  $O(\log n)$  time.

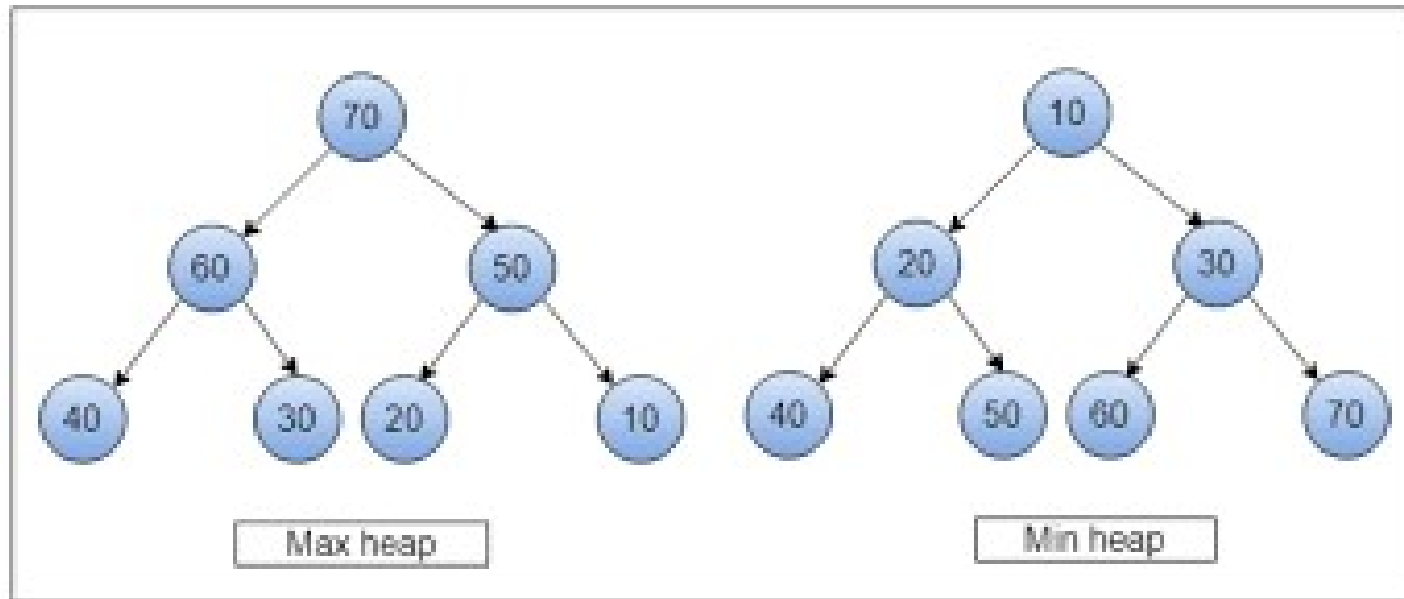
# Definitions

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- **Min-Heap:** The key present at the root node is smaller than or equal to keys of all the nodes present in the children nodes.
  - And this same rule is recursively followed by all the subtrees of the binary tree.
- **Max-Heap:** In this data structure, the key which is present at the root node is greater than or equal to the keys of all the children nodes of the tree.
  - The same property is recursively applicable for all the subtrees of the tree. The maximum key is present at the root of the tree for a Max-Heap.

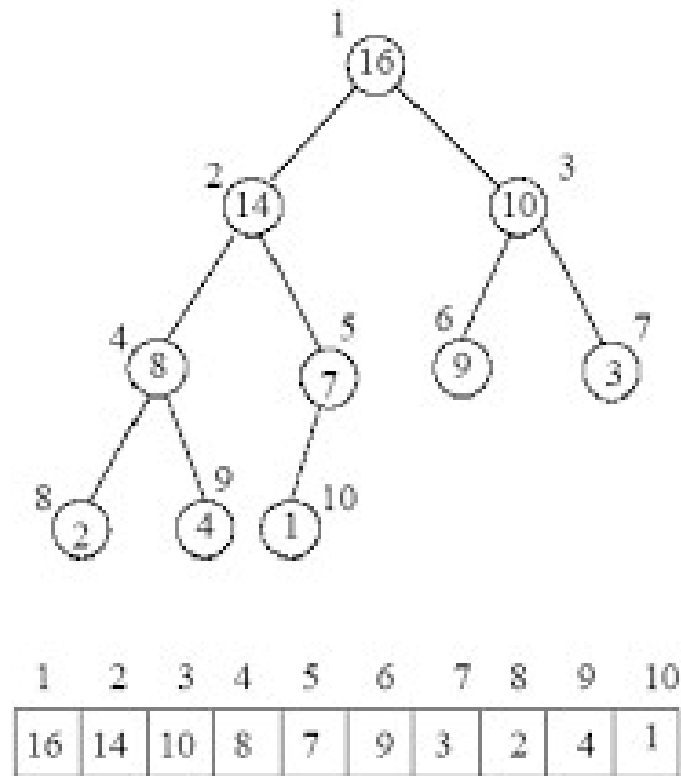
# Max Heap

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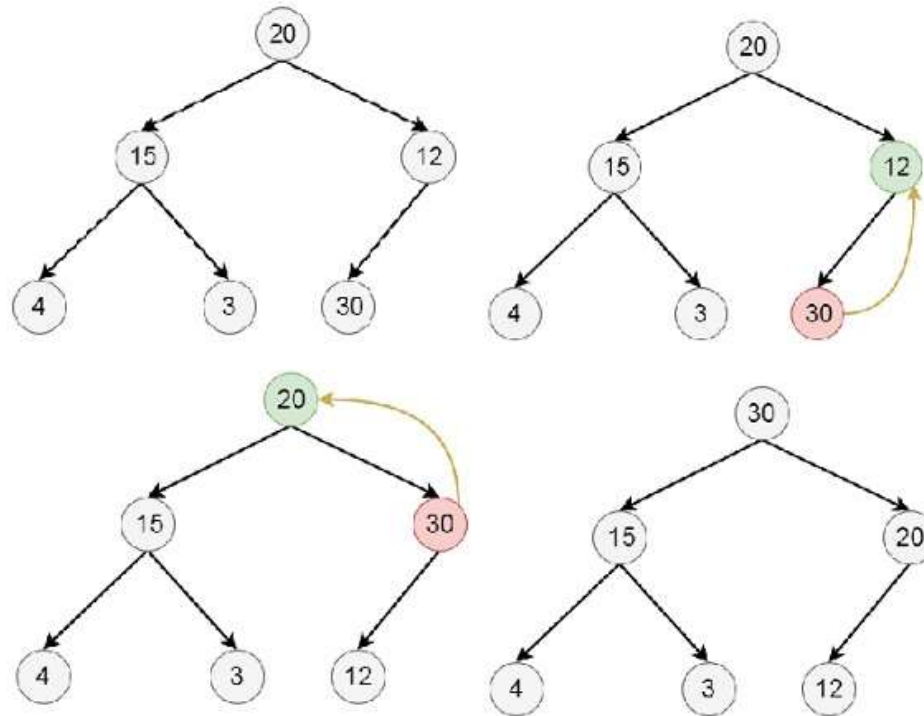
# Array Representation of Heaps

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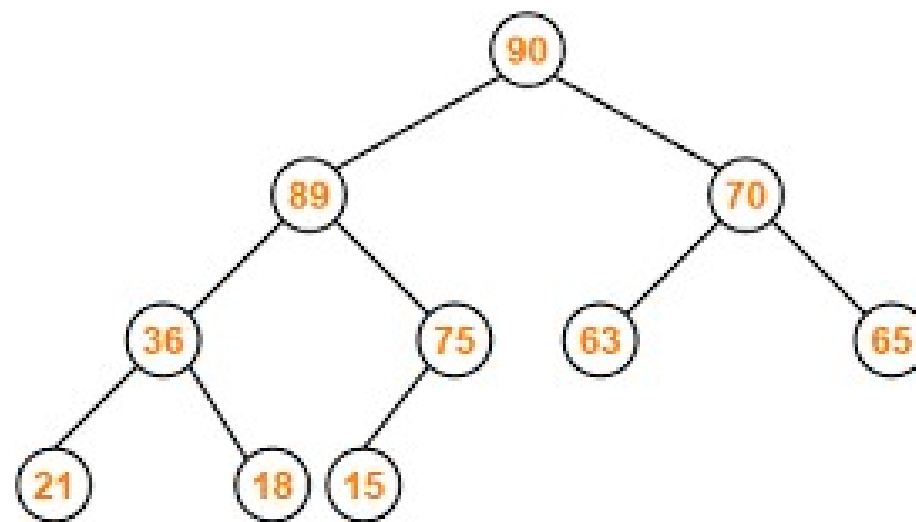
# Max Heap - Insertion

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# Insert 95

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Max Heap Example

# Max Heap Insertion Logic

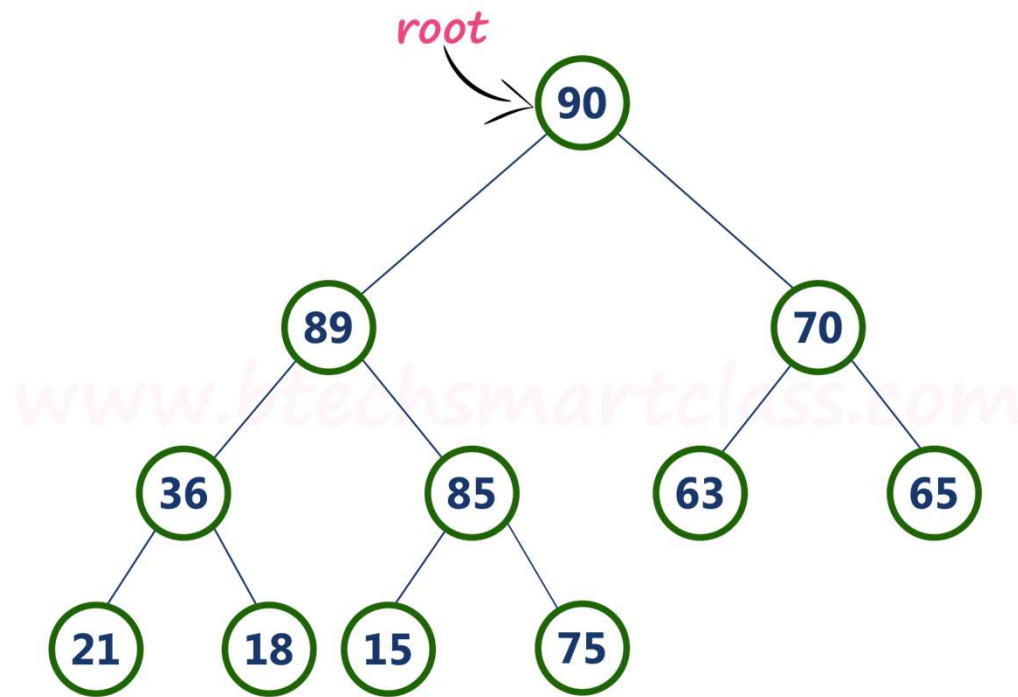
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- Step 1 - Insert the **newNode** as **last leaf** from left to right.
- Step 2 - Compare **newNode value** with its **Parent node**.
- Step 3 - If **newNode value is greater** than its parent, then **swap** both of them.
- Step 4 - Repeat step 2 and step 3 until newNode value is less than its parent node (or) newNode reaches to root.



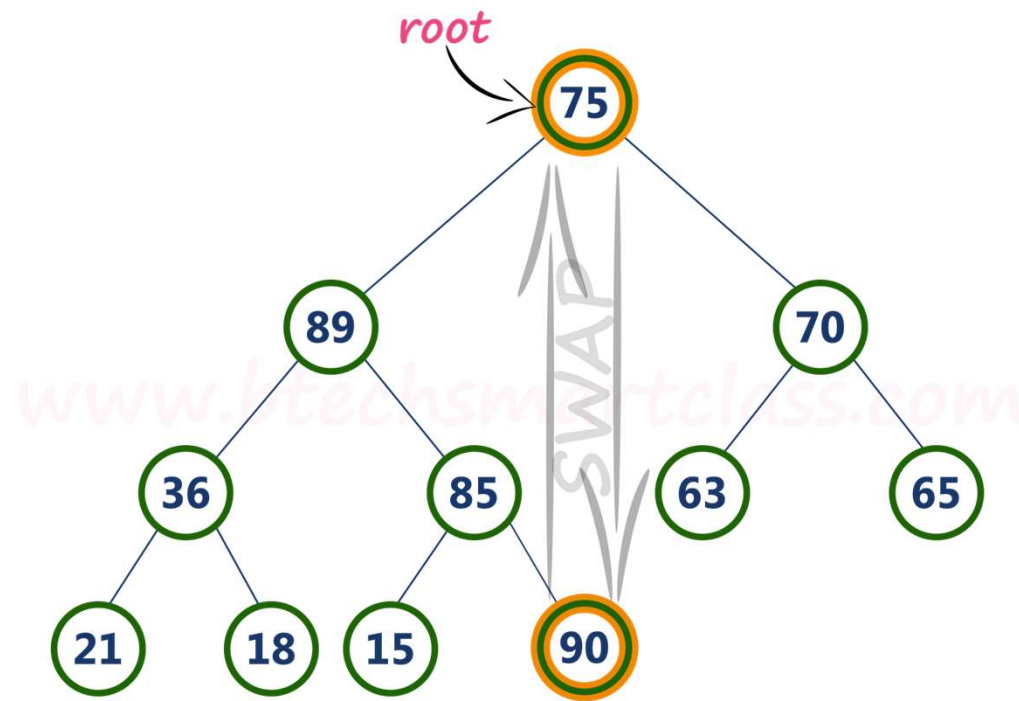
# Max Heap - Deletion

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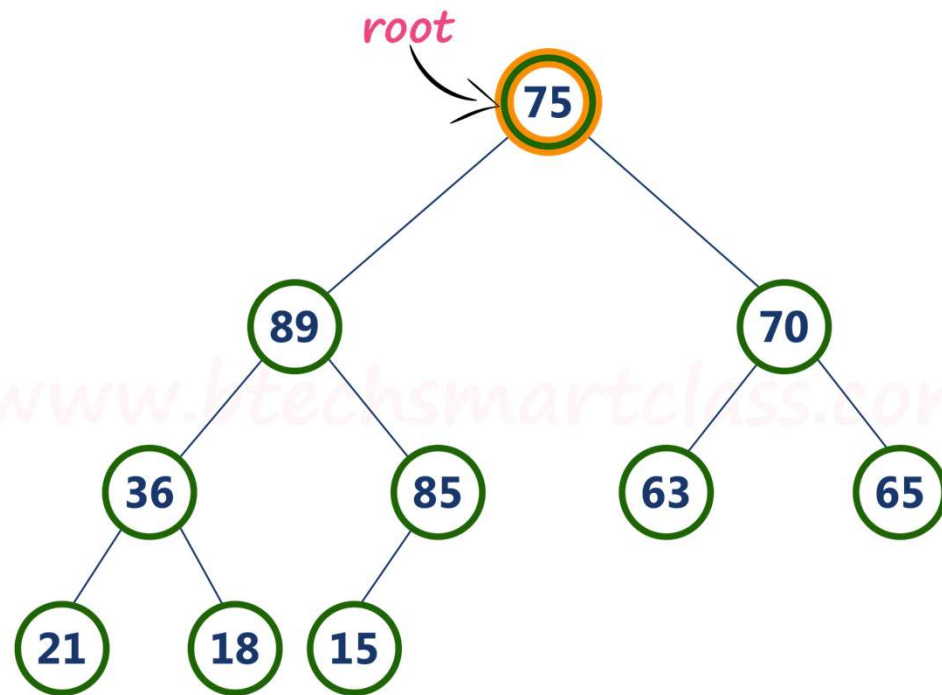
# Max Heap – Deletion – Step 1

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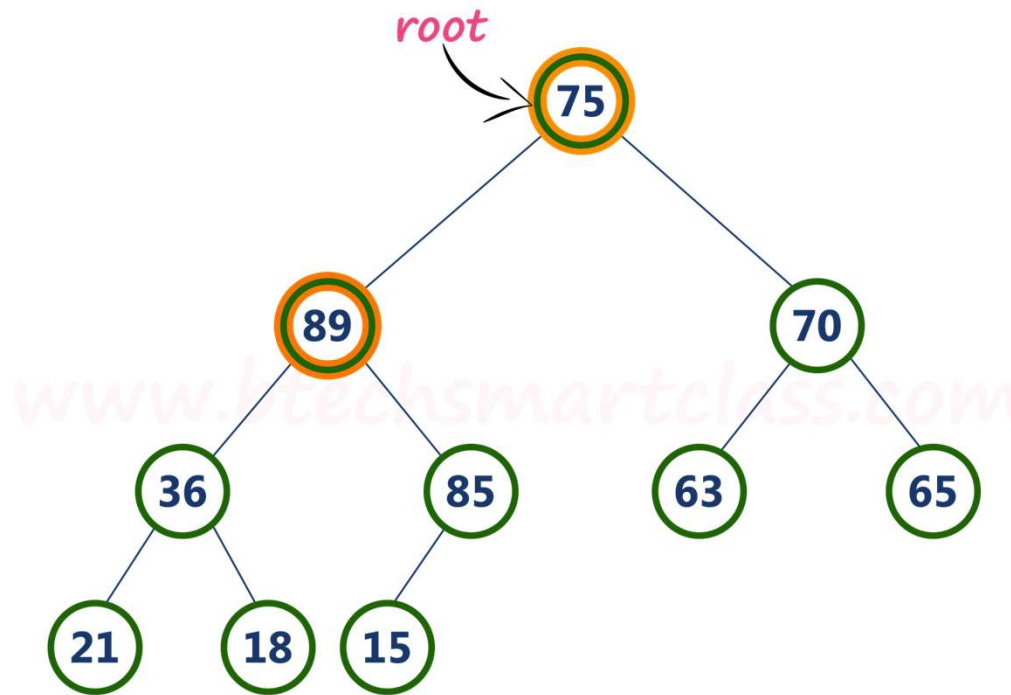
# Max Heap – Deletion – Step 2

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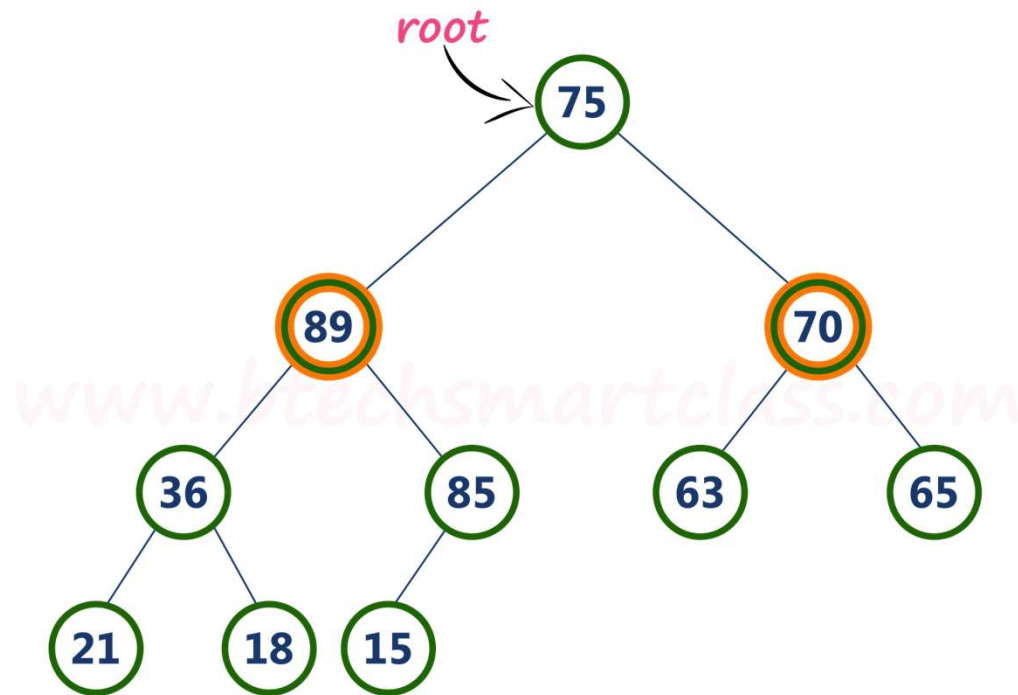
# Max Heap – Deletion – Step 3

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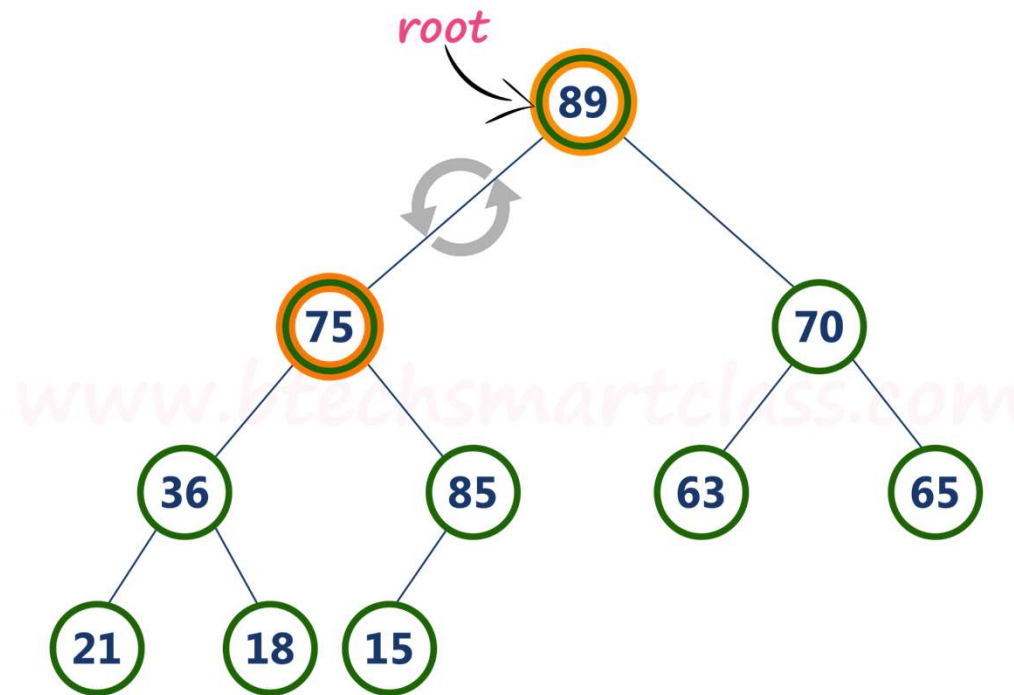
# Max Heap – Deletion – Step 4

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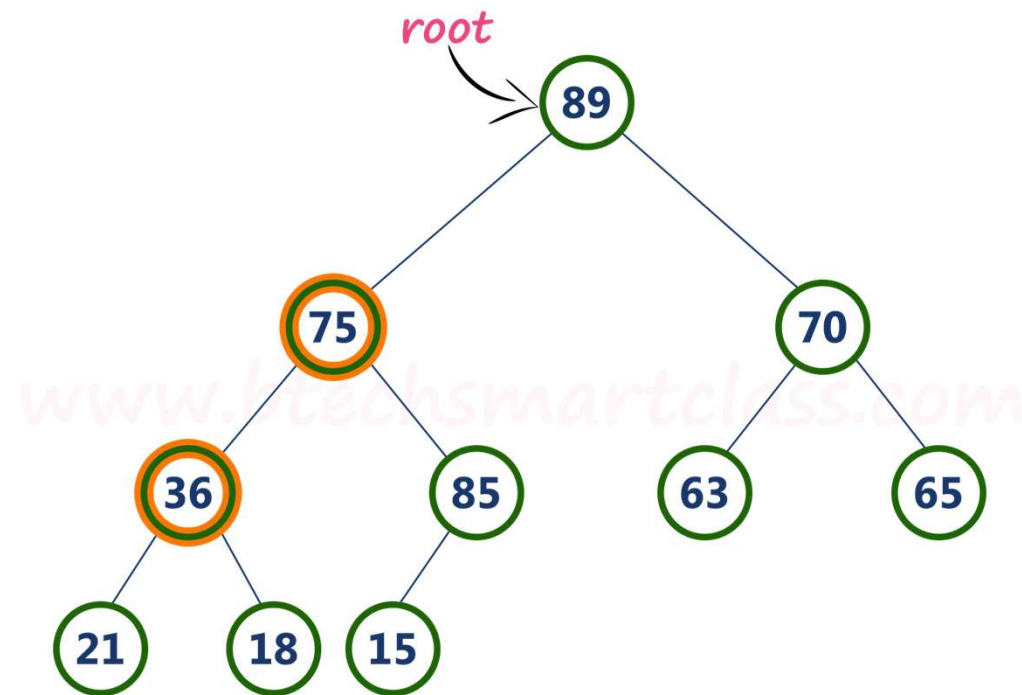
# Max Heap – Deletion – Step 5

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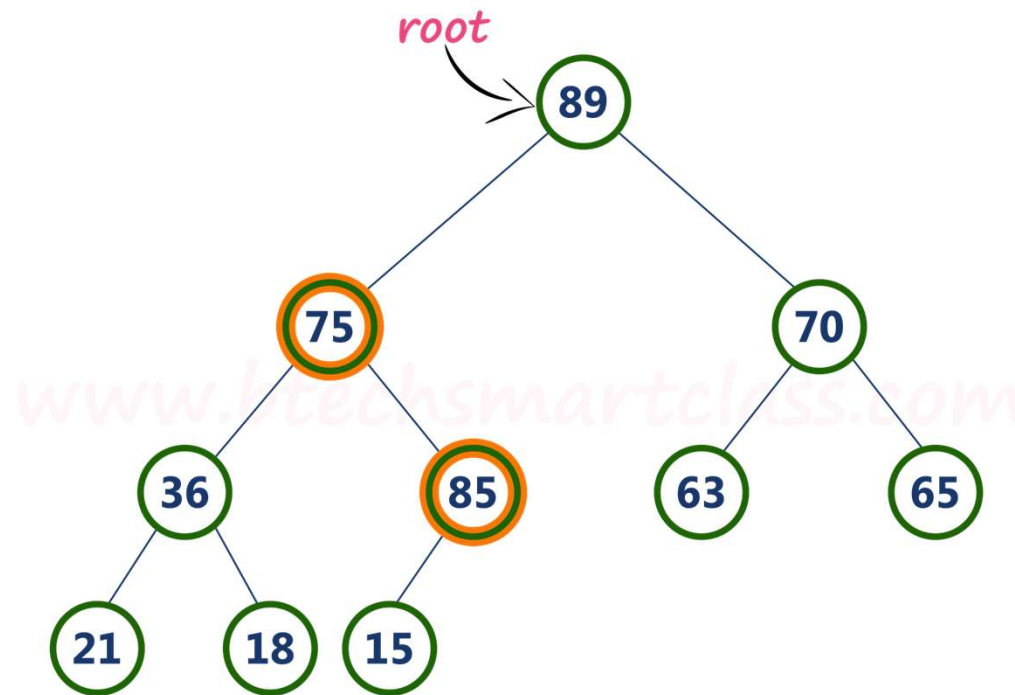
# Max Heap – Deletion – Step 6

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# Max Heap – Deletion – Step 7

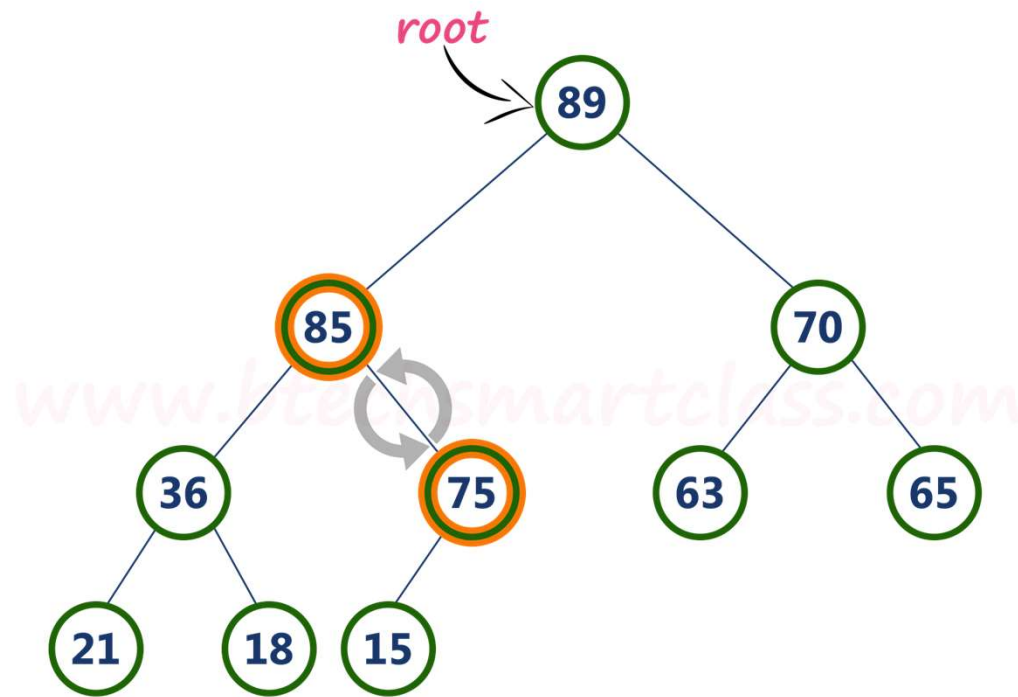
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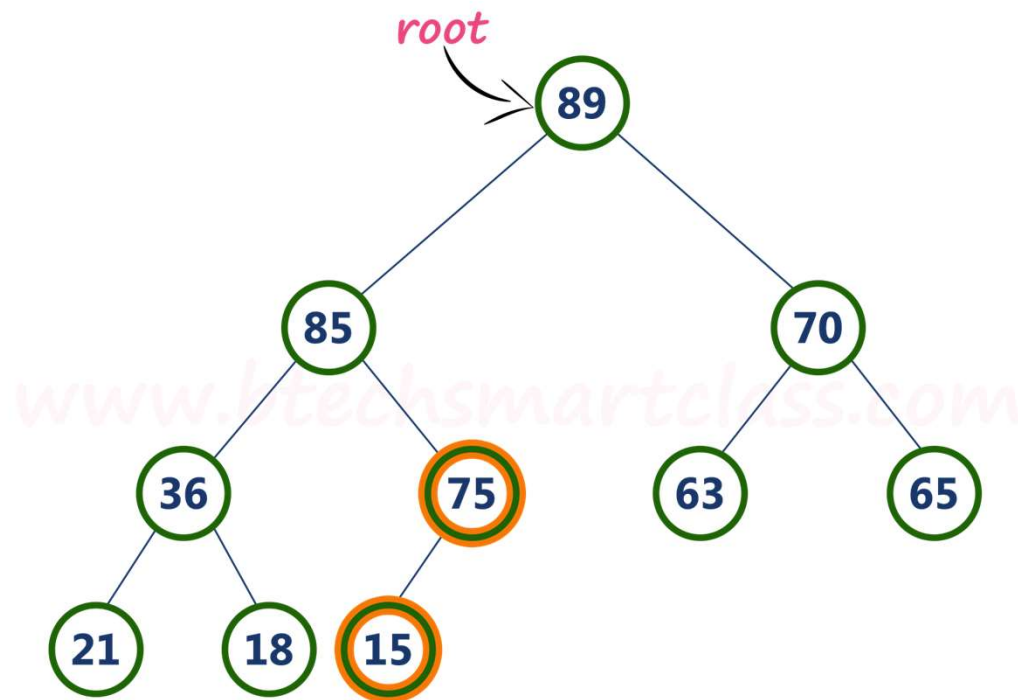
# Max Heap – Deletion – Step 8

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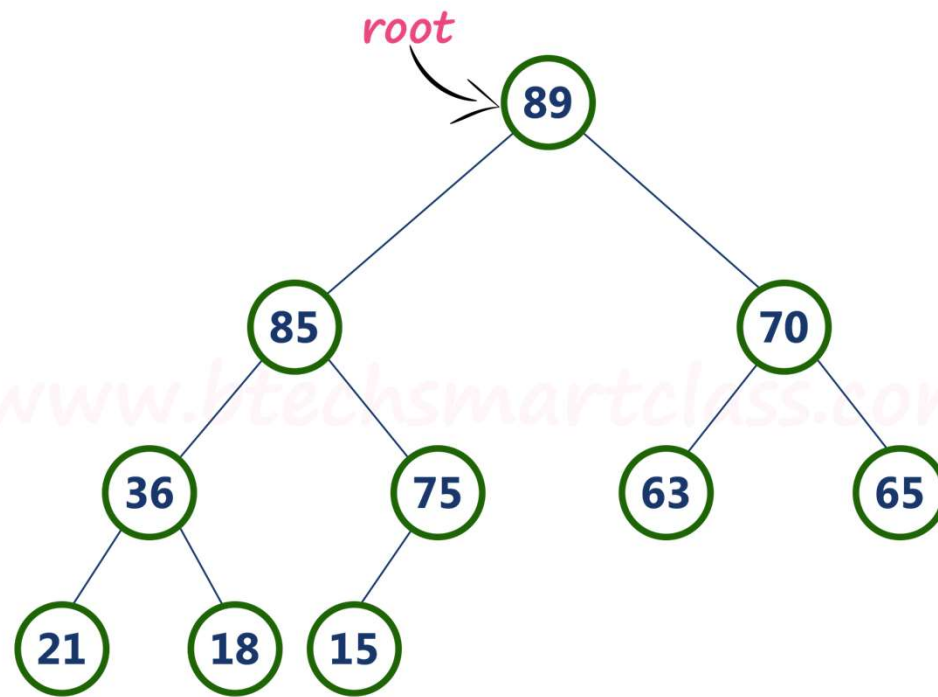
# Max Heap – Deletion – Step 9

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# Max Heap – Deletion – Step 10

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# Max Heap Deletion Logic

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- Step 1 - **Swap** the **root** node with **last** node in max heap
- Step 2 - **Delete** last node.
- Step 3 - Now, compare **root value** with its **left child value**.
- Step 4 - If **root value is smaller** than its left child, then compare **left child** with its **right sibling**. Else goto **Step 6**
- Step 5 - If **left child value is larger** than its **right sibling**, then **swap root** with **left child** otherwise **swap root** with its **right child**.
- Step 6 - If **root value is larger** than its left child, then compare **root value** with its **right child** value.
- Step 7 - If **root value is smaller** than its **right child**, then **swap root** with **right child** otherwise **stop the process**.
- Step 8 - Repeat the same until root node fixes at its exact position.