Recursion

The process in which a function calls itself directly or indirectly is called recursion and the corresponding function is called a recursive function.

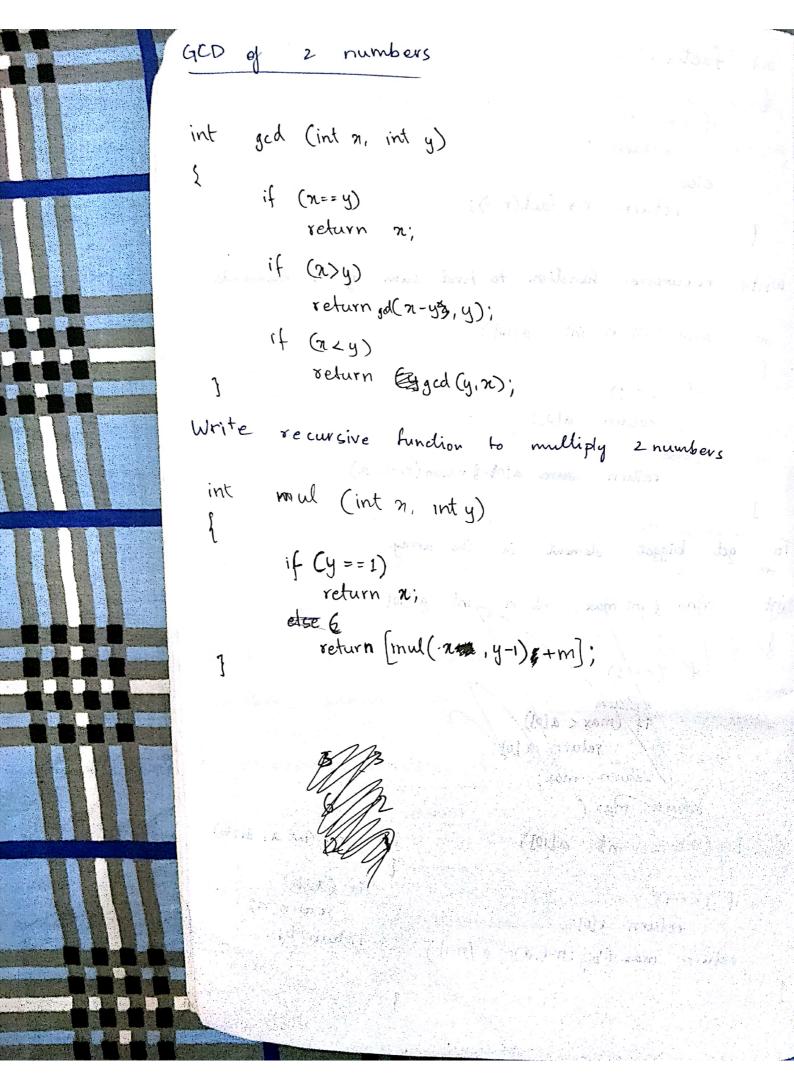
1 1	, ,	
direct recursion.	indirect	recursion
f()(area)	(6/20)	
2 3-(2/1.2)	80	f1()
f()		{
Ì	fic)	FC)
Hollow rome	3	3

In case of a recursive function, all the recursive functions will be pushed onto stack, if there is no base case or base condition, the stack becomes overflow, so, to rainvoid this, every recursive function should have a base case.

A recursive function to find the factorial of a given number

 $fact (n) \begin{cases} 1 & \text{if } h=1 \\ n \neq fact (n-1) & \text{otherwise.} \end{cases}$

```
fact (int n)
 int
 Ę
       if (n == 1)
            return 1;
        else
           return n* fact(n-i);
   3
                   fundion to find sum
      recursive
Write
        sum (int n, int a[10])
 int
  {
         if (n==1)
              return a[0];
          else
                     $$$$$$$ a[n] + sum(n-1, a);
   3
To
     get bigget
                   element
                            in
              (int max, int n, int
                                      a [10]
int
                 /(max < a[p]
          return max (
                                          int max (int a, intb)
                       a [10])
      big Cint no int
int
٤
                                              if (azb)
      if (n==1)
                                                 return a;
            return a [0]:
                                              return bj
      return max (big (n-1,a), a [n-1])
```



Write a recursive tune to perform binary seasoch bs (Marloy, int a [30], int low, int high 3, int key) int Z if (low > high) return -1; int mid = (low+high)/2; orborit [(key = a[mid]) return nid+1; if (key < a [anid]) bs (med Horasa, low, mid-1, key); if (key>a[mid]) bs (a, 4800 mid+1, high, key); રુ Tower of Hanoi solog general, the TOH problem In the tower of Hanoi prob, we have 3 denoted with source (S), destination(D) and temp (T). On the source (s), n discs are stacked one above the other such that the disc with the larger diameter is will be at the bottom. " of rols is alway The objective here is to more all-the discs from the source to destination such that the destination should have the same property of source. To achieve this, the limitations are

- i) only one disc can be moved across the poles.
- ii) at any time, the disc with the largeer dia should always be at the bottom.

Let a denotes the number of disc

I case 1: if n=1,

· me ned only one none is more discovered Sto D.

if n=2, more are as follows.

- · +st (s to)
- · 2nd (5 to D)
- · 1 (T to D).

In general, the TOH problem can be solved recursively in the following manner

step 1: recursively move (n-1) discs from S to T. with the help of D.

step 2: Move nth disce from conce - la D.

step 3'. recursively (n-1) discs from T to D with the help of S.

void tower (int n, char 's', char 'D', char T)

if (n>0)

Tower(n-1, S, T, D);

Pf ("In move 1.d disc from %.c to %c",n,s,D),

Tower (n-1, T, D,s);