

INDIAN INSTITUTE OF TECHNOLOGY GOA

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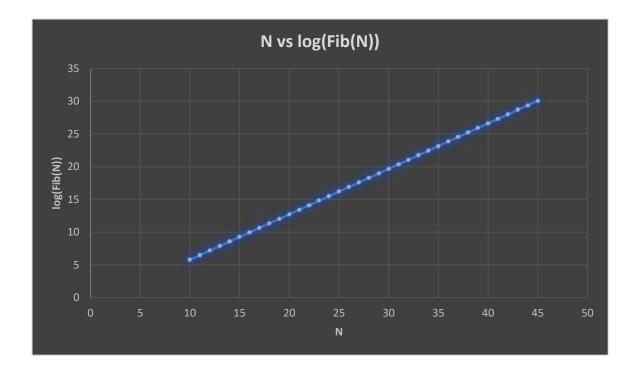
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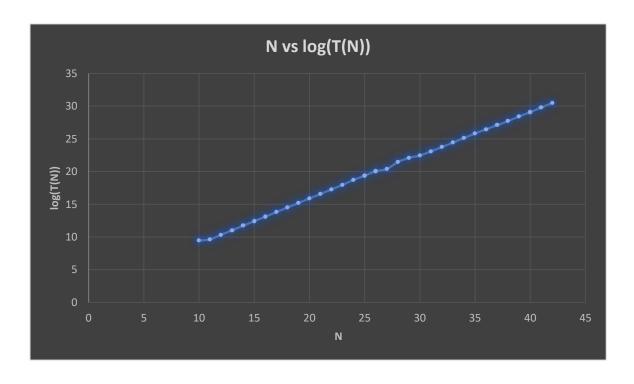
The output of the program written for Q.1]

```
Time taken by the recursive implementation of the 10th fibonacci number is: 500 nanoseconds
Time taken by the recursive implementation of the 11th fibonacci number is: 700 nanoseconds
Time taken by the recursive implementation of the 13th fibonacci number is: 1800 nanoseconds
Time taken by the recursive implementation of the 14th fibonacci number is: 2300 nanoseconds
Time taken by the recursive implementation of the 15th fibonacci number is: 3700 nanoseconds
Time taken by the recursive implementation of the 16th fibonacci number is: 5600 nanoseconds
Time taken by the recursive implementation of the 18th fibonacci number is: 14400 nanoseconds
Time taken by the recursive implementation of the 20th fibonacci number is: 37200 nanoseconds
Time taken by the recursive implementation of the 21th fibonacci number is: 59700 nanoseconds
Time taken by the recursive implementation of the 22th fibonacci number is: 108600 nanoseconds
Time taken by the recursive implementation of the 23th fibonacci number is: 156100 nanoseconds
Time taken by the recursive implementation of the 26th fibonacci number is: 666800 nanoseconds
Time taken by the recursive implementation of the 27th fibonacci number is: 1075800 nanoseconds
Time taken by the recursive implementation of the 28th fibonacci number is: 1737400 nanoseconds
Time taken by the recursive implementation of the 29th fibonacci number is: 2807800 nanoseconds
Time taken by the recursive implementation of the 30th fibonacci number is: 4545800 nanoseconds
Time taken by the recursive implementation of the 31th fibonacci number is: 7356500 nanoseconds
Time taken by the recursive implementation of the 33th fibonacci number is: 23101800 nanoseconds
```

<u>Q.1]</u>

The graphs obtained by storing the data from the recursive implementation of the Fibonacci numbers in a CSV file and plotting them using excel are as follows:





<u>1.1</u>

Each number in the Fibonacci series, which often begins with 0 and 1, is the sum of the two numbers before it. The recurrence relation

$$F(n) = F(n-1) + F(n-2)$$

defines the Fibonacci sequence, where

$$F(0) = 0$$
 and $F(1) = 1$.

Clearly, with the increasing values of n, the growth of the Fibonacci numbers is exponential. It is also evident from the graph of **N vs log(F(N))** which is a straight line that, the Fibonacci numbers grow exponentially with N.

The time required to calculate the n-th Fibonacci number using this recurrence relation increases exponentially with n, as is well known. This is due to the function's ability to call itself twice in a recursive manner, each time with a problem of size n-1 and n-2, which causes exponential growth in the number of recursive calls. When utilizing this recurrence relation to calculate the n-th Fibonacci number, the time required to do so rises exponentially with n since the number of recursive calls grows exponentially with the input size.

1.2

The slope of the first graph can be calculated as follows:

$$\frac{\log(\text{Fib}(31)) - \log(\text{Fib}(28))}{31 - 28} = M1$$

$$M1 = 0.694233$$

The slope of the first graph can be calculated as follow:

$$\frac{\log(T(31)) - \log(T(28))}{31 - 28} = M2$$

$$M2 = 0.5483$$

1.3

It is pretty clear from the previous question that the function F(N) can be written as follows, Since, the fraction log(F(N))/N = 0.694233, thus:

$$F(N) = 2^{0.694233N}$$

1.4

Similar to the previous question the function T(N) can be written as follows, Since, the fraction log(T(N))/N = 0.5483, thus :

$$T(N) = 2^{0.5483N}$$

Q.2]



Because the power function squares the matrix recursively and cuts the size of the problem in half at each step, its time complexity is **O(log n)**.

The time complexity of the multiply function, which multiplies two 2x2 matrices, is **O(1)** because it only requires four n-bit integer multiplications and additions.

The Fibonacci function's overall time complexity is then

$$O(\log n) * O(1) = O. (\log n).$$

The multiply function's time complexity in terms of M(n) is O(M(n)), making the Fibonacci function's overall time complexity in terms of M(n) equal to $O(\log n * M(n))$.

The time complexity of M(n) is dependent on the algorithm used to multiply two n-bit integers, but it is typically $O(n^2)$ or $O(n \log n)$