

1)

Given that:

U is the orthogonal $m \times m$ matrix with columns $[u_1, u_2, \dots, u_m]$ **V** is the orthogonal $n \times n$ matrix with columns $[v_1, v_2, \dots, v_n]$ with transpose \mathbf{V}^T **Σ** is the diagonal $m \times n$ matrix with values $[\sigma_1, \sigma_2, \dots, \sigma_r]$

We can then say by definition:

$$(1) C = U\Sigma V^T$$

This can then be extrapolated into:

$$C^T = (U\Sigma V^T)^T \Rightarrow C^T = \Sigma^T V U^T$$

$$(2) C^T U = \Sigma^T V U^T U$$

Because U is orthogonal, we can then say:

$$C^T U = \Sigma^T V I$$

$$(3) C^T U = \Sigma^T V$$

Taking the following:

$$U = [u_1, u_2, \dots, u_m], V = [v_1, v_2, \dots, v_n], \Sigma^T = \text{diag}(\sigma_1, \sigma_2, \dots, \sigma_r)^T = [\sigma_1, \sigma_2, \dots, \sigma_r]$$

We derive the general form

$$(4) C^T [u_1, u_2, \dots, u_m] = [\sigma_1, \sigma_2, \dots, \sigma_r] [v_1, v_2, \dots, v_n] = [\sigma_1 v_1, \sigma_2 v_2, \sigma_r v_r, 0, \dots, 0]$$

$$(5) C^T u_i = \sigma_i v_i \text{ for all } i = 1, 2, \dots, r$$

Taking $i = 1$, we observe that:

$$(6) C^T u_1 = \sigma_1 v_1 \Rightarrow \frac{C^T u_1}{\sigma_1} = v_1, \text{ q.e.d.,}$$

2)

Calculations:

Given the directed graph in the problem, we can construct the adjacency matrix of outward edges, A, as such with the following ordering:

AI	Stats	DSci	DMin	DEng	DVis	Prog	PComp	ML	Julia	
0	1	1	0	0	0	0	0	1	0	AI
1	0	1	1	0	0	0	0	0	0	Stats
1	1	0	1	0	1	1	0	0	1	DSci
0	0	0	0	1	0	0	0	0	0	DMin
0	0	0	0	0	1	0	0	0	0	Deng
0	0	0	0	0	0	1	0	0	0	DVis
0	0	0	0	0	0	0	1	0	1	Prog
0	0	0	0	0	0	1	0	0	0	PComp
1	0	0	0	0	0	0	0	0	0	ML
0	0	0	0	0	0	0	0	1	0	Julia

We can then construct the diagonal matrix D by summing the values in each row to get the following:

AI	Stats	DSci	DMin	DEng	DVis	Prog	PComp	ML	Julia	
3										AI
	3									Stats
		6								DSci
			1							DMin
				1						Deng
					1					DVis
						2				Prog
							1			PComp
								1		ML
									1	Julia

We observe that none of the values in the D matrix have a 0 value, so no shifting is required.

To solve the equation $(I - \alpha A^T D^{-1})x = \frac{(1-\alpha)}{n} e$, we find the transpose of the adjacency matrix A and the inverse of the diagonal matrix D. We utilize Julia to perform the numerical calculations as follows, taking values of alpha from 0 to 0.9999 (due to the limitations of the matrix network package) in steps of 0.0001, and observing the unique rankings that occur.

Code and Results:

```
using MatrixNetworks
using SparseArrays
using LinearAlgebra

## define iterable list and ordering list
a_iter = 0:0.0001:0.9999
orderings_man = Set{Vector{Int}}()
orderings_pkg = Set{Vector{Int}}()

## Define the adjacency matrix A
A = [0 1 1 0 0 0 0 0 1 0; ## AI
     1 0 1 1 0 0 0 0 0 0; ## Statistics
     1 1 0 1 0 1 1 0 0 1; ## Data Science
     0 0 0 0 1 0 0 0 0 0; ## Data Mining
     0 0 0 0 0 1 0 0 0 0; ## Data Engineering
     0 0 0 0 0 0 1 0 0 0; ## Data Visualization
     0 0 0 0 0 0 0 1 0 1; ## Programming
     0 0 0 0 0 0 1 0 0 0; ## Parallel Computing
     1 0 0 0 0 0 0 0 0 0; ## Machine Learning
     0 0 0 0 0 0 0 0 1 0]## Julia

sparse_A = sparse(A)

## define the pageRank function
function pageRank(adj_mat::Matrix, alpha::Float64)
    n = size(adj_mat, 1)

    ## Create diagonal matrix
    d_mat = Diagonal(sum(adj_mat, dims=1)[:])
    d_inv = Diagonal(1 ./ diag(d_mat))

    ## Create I - alpha*D^-1*A'
    lhs_mat = I - alpha * adj_mat' * d_inv

    ## Define the right-hand side vector
```

```

rhs_vec = (1- alpha) / n * ones(n)

## Solve the equation
x = lhs_mat \ rhs_vec

return x
end

## Calculate pageRank manually and through MatrixNetworks
for a in a_iter
    p = sortperm(pageRank(A, a), rev=true)
    push!(orderings_man, p)
    p = sortperm(pagerank(sparse_A, a), rev=true)
    push!(orderings_pkg, p)
end

# Print results
for ordering in orderings_pkg
    @printf("Ordering: %s\n", join(ordering, ", "))
end
@printf("MatrixNetworks Ranking Count: %.0f\n\n", length(orderings_pkg))

for ordering in orderings_man
    @printf("Ordering: %s\n", join(ordering, ", "))
end
@printf("Manual Calculation Ranking Count: %.0f\n", length(orderings_man))

```

Ordering: 1, 7, 9, 10, 3, 8, 6, 2, 5, 4
 Ordering: 7, 1, 9, 10, 6, 8, 3, 5, 2, 4
 Ordering: 1, 7, 9, 10, 3, 2, 8, 6, 5, 4
 Ordering: 7, 1, 9, 6, 10, 5, 3, 8, 2, 4
 Ordering: 1, 7, 9, 10, 8, 3, 6, 2, 5, 4
 Ordering: 1, 2, 3, 4, 5, 6, 7, 8, 9, 10
 Ordering: 7, 1, 9, 6, 10, 5, 8, 3, 2, 4
 Ordering: 7, 1, 9, 10, 8, 6, 3, 2, 5, 4
 Ordering: 1, 9, 7, 3, 10, 2, 8, 6, 5, 4
 Ordering: 7, 1, 9, 6, 5, 10, 3, 8, 2, 4

Ordering: 7, 1, 6, 10, 4, 9, 3, 2, 8, 5
 Ordering: 1, 2, 3, 4, 5, 6, 7, 8, 9, 10
 Ordering: 7, 1, 6, 10, 4, 9, 2, 3, 8, 5
 Ordering: 7, 1, 6, 10, 9, 4, 3, 2, 8, 5
 Ordering: 7, 1, 6, 4, 10, 9, 3, 2, 5, 8
 Ordering: 7, 1, 6, 10, 4, 9, 3, 2, 5, 8
 Ordering: 7, 1, 6, 10, 9, 4, 2, 3, 8, 5
 Ordering: 7, 1, 6, 4, 10, 9, 2, 3, 5, 8
 Ordering: 7, 1, 6, 10, 4, 9, 2, 3, 5, 8
 Manual Calculation Ranking Count: 9

Ordering: 7, 1, 9, 6, 10, 8, 3, 5, 2, 4

Ordering: 7, 1, 9, 6, 10, 8, 5, 3, 2, 4

Ordering: 7, 1, 9, 10, 6, 8, 3, 2, 5, 4

Ordering: 7, 1, 9, 10, 8, 3, 6, 2, 5, 4

Ordering: 1, 9, 7, 10, 3, 2, 8, 6, 5, 4

Ordering: 1, 7, 9, 10, 3, 8, 2, 6, 5, 4

MatrixNetworks Ranking Count: 16

Conclusion:

Both a manual calculation using the LinearAlgebra package and the MatrixNetworks package pagerank function were used. From the observations above, we see that there were 16 identifiable rankings using the MatrixNetworks package, and only 9 using the manual calculation of PageRank. It is interesting to note that the manual calculations of page rank do not identify any orderings that begin with node 1 other than the sequential order ranking (1, 2, ..., 10).

3A)

Code and Results

Part 1:

```
using LinearAlgebra
using Printf

function count_darts(N::Int, k::Int)
    num_inside = 0
    for i=1:N
        p = rand(k)
        num_inside += sum(abs.(p)) <= 1
    end
    return num_inside
end

###
N = 1000000
for k=2:20
    n = count_darts(N, k)
    println(k, " -> ", n, " | ", (n/N*100), "%")
end

###
```

Results:

dim	->	count	% volume
2	->	5001839	50.0183900000000004%
3	->	1666466	16.66466%
4	->	417285	4.17285%
5	->	83163	0.8316300000000001%
6	->	13935	0.13935%
7	->	2067	0.02067%
8	->	251	0.00251%
9	->	22	0.00022%
10	->	2	1.9999999999999998e-5%
11	->	0	0.0%
12	->	0	0.0%
13	->	0	0.0%
14	->	0	0.0%
15	->	0	0.0%
16	->	0	0.0%
17	->	0	0.0%
18	->	0	0.0%
19	->	0	0.0%
20	->	0	0.0%

3B)

Based on the results from above, we can see that the volume of the rhomboidal space rapidly decreases with each additional dimension added to the space.

Per expectations, for the 2-dimensional space, we expect that the rhombus, as stated, takes up about 50% of the volume, which we see as 499,395 darts within the rhomboidal space. For the 3-dimensional space, this declines rapidly to only 167009 darts within the space, and so on and so forth for each additional dimension, converging to 0 at the 10th dimension.

Considering the volume of the hypercube at each dimension, $V_{cube} = 1^d$, we can then observe that the rhomboidal region takes up $V_{rhomboid} = \frac{1}{d!}$, which is validated by the volumes observed above.