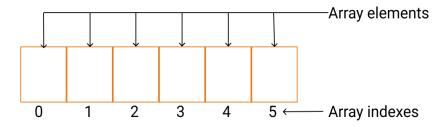
# **Array**

- An Array is a Linear data structure
- It is a collection of data items having similar data types stored in contiguous memory locations.
- By knowing the address of the first item we can easily access all items/elements of an array.
- Arrays and its representation is given below.



## • Array Dsicription:

- **Array Index:** The location of an element in an array has an index, which identifies the element. Array index starts from 0.
- Array element: Items stored in an array is called an element. The elements can be accessed via its index.
- **Array Length:** The length of an array is defined based on the number of elements an array can store. In the above example, array length is 6 which means that it can store 6 elements.
- When an array of size and type is declared, the compiler allocates enough memory to hold all elements of data.
- E.g. an array face [10] will have 10 elements with index starting from 0 to 9 and the memory allocated contiguously will be 20 bytes (integer = 2 bytes).
- The compiler knows the address of the first byte of the array only. Also, the address of the first byte is considered as the memory address for the whole array.

## **Types of Arrays**

The various types of arrays are as follows.

- One dimensional array
- Multi-dimensional array

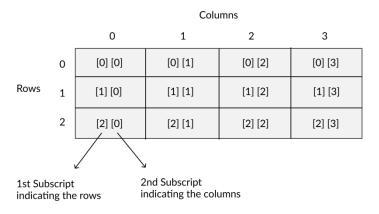
#### One-Dimensional Array

- o A one-dimensional array is also called a single dimensional array
- where the elements will be accessed in sequential order.
- o This type of array will be accessed by the subscript of either a column or row index.

#### Multi-Dimensional Array

- When the number of dimensions specified is more than one, then it is called as a multi-dimensional array.
- Multidimensional arrays include 2D arrays and 3D arrays.

## **Two-dimensional Array**



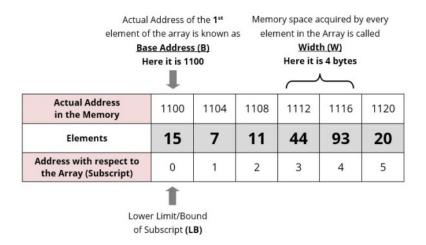
### • E.g. A two-dimensional array

- Accessed with subscript of row and column index
- o For traversal the value of the rows and columns will be considered
- E.g. **face [3] [4],** the first index specifies the number of rows and the second index specifies the number of columns and the array can hold 12 elements (3 \* 4)

#### • A three-dimensional array

o The array **face [5] [10] [15]** can hold 750 elements (5 \* 10 \* 15).

### Address Calculation in single (one) Dimension Array:



• Address of an element A[I] is calculated using the following formula:

- 1100 + (1-0)\*4
- 1100 + (4-0) \* 4
  - $\circ$  Address of A [I] = B + W \* (I LB)
  - o Where.

 $\mathbf{B} = \text{Base address}$ 

**W** = Storage Size of one element stored in the array (in byte)

I = Subscript of element whose address is to be found

**LB** = Lower limit / Lower Bound of subscript, if not specified assume 0 (zero)

Example:

Base address of an array B[1300.....1900] as 1020 and size of each element is 2 bytes in the memory. Find the address of B[1700]

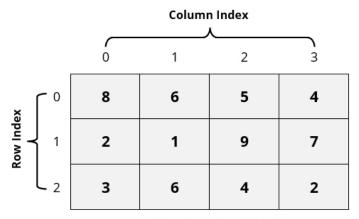
#### **Solution:**

#### Why does Array Indexing start with 0?

- Here **a** itself is a pointer which contains the memory location of the first element of the array
- first element can be accessed with a[0]
  - o internally decoded by the compiler as \*(a + 0).
- second element can be accessed by a[1] or \*(a + 1).
- As a contains the address of the first element = Base Address
- and index describes the offset from the first element, i.e. the distance from the first element.
  - If array indexing starts at 1 instead of 0
  - the first element can be accessed by a[1]
    - which is internally decoded as \*(a + 1 1).
- Thus we have to perform one extra operation i.e. subtraction by 1.
- This extra operation will greatly decrease the performance when the program is big.
- Thus to avoid this extra operation and improve the performance, array indexing starts at 0 and not at 1.

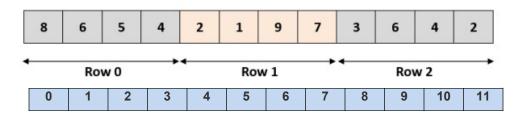
## Address Calculation in Double (Two) Dimensional Array:

- While storing the elements of a 2-D array in memory, these are allocated contiguous memory locations.
- Therefore, a 2-D array must be linearized so as to enable their storage.
- There are two alternatives to achieve linearization:
  - o Row-Major
  - o Column-Major.



Two-Dimensional Array

#### Row-Major (Row Wise Arrangement)



• Address of an element A[ I ][ J ] = = B + W \* [ N \* ( I - Lr ) + ( J - Lc ) ] for thr array declared as A[M][N]

#### Where

 $\mathbf{B} = \text{Base address}$ 

I = Row subscript of element whose address is to be found

J = Column subscript of element whose address is to be found

**W** = Storage Size of one element stored in the array (in byte)

Lr = Lower limit of row/start row index of matrix, if not given assume 0 (zero)

Lc = Lower limit of column/start column index of matrix, if not given assume 0 (zero)

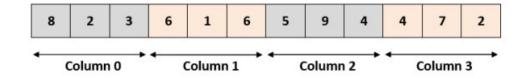
M = Number of row of the given matrix

N = Number of column of the given matrix

Example: for the given A[4][4] with A[0..3][0..3]

## **Column Oriented storage:**

#### Column-Major (Column Wise Arrangement)



Address of A[I][J] = B + W \* [M\*(J-Lc) + (I-Lr)]

Where

 $\mathbf{B} = \text{Base address}$ 

I = Row subscript of element whose address is to be found

J = Column subscript of element whose address is to be found

**W** = Storage Size of one element stored in the array (in byte)

Lr = Lower limit of row/start row index of matrix, if not given assume 0 (zero)

Lc = Lower limit of column/start column index of matrix, if not given assume 0 (zero)

M =Number of row of the given matrix

N = Number of column of the given matrix

#### • Important Note:

- Usually number of rows and columns of a matrix are given (like A[20][30] or A[40][60]) but if it is given as A[Lr----Ur, Lc----Uc]. In this case number of rows and columns are calculated using the following methods:
  - Number of rows (M) will be calculated as = (Ur Lr) + 1Number of columns (N) will be calculated as = (Uc - Lc) + 1
  - And rest of the process will remain same as per requirement (Row Major Wise or Column Major Wise).

## **Examples:**

**Q 1**. An array X [-15......10, 15......40] requires **one byte of storage**. If beginning location is 1500 determine the location of X [15][20].

**Q 1**. An array X [-15......10, 15......40] requires **one byte of storage**. If beginning location is 1500 determine the location of X [15][20].

#### **Solution:**

As you see here the number of rows and columns are not given in the question. So they are calculated as:

Number or rows say 
$$\mathbf{M} = (\mathbf{Ur} - \mathbf{Lr}) + \mathbf{1} = [10 - (-15)] + 1 = 26$$
  
Number or columns say  $\mathbf{N} = (\mathbf{Uc} - \mathbf{Lc}) + \mathbf{1} = [40 - 15)] + 1 = 26$ 

## (i) Row Major Wise Calculation of above equation

The given values are: B = 1500, W = 1 byte, I = 15, J = 20, Lr = -15, Lc = 15, N = 26

Address of A [ I ][ J ] = B + W \* [ N \* ( I - Lr ) + ( J - Lc ) ]  
= 
$$1500 + 1* [26* (15 - (-15))) + (20 - 15)]$$
  
=  $1500 + 1* [26* 30 + 5]$   
=  $1500 + 1* [780 + 5]$   
=  $1500 + 785$   
=  $2285$  [Ans]

**Q 1**. An array X [-15......10, 15......40] requires **one byte of storage**. If beginning location is 1500 determine the location of X [15][20].

#### **Solution:**

As you see here the number of rows and columns are not given in the question. So they are calculated as:

Number or rows say 
$$\mathbf{M} = (\mathbf{Ur} - \mathbf{Lr}) + \mathbf{1} = [10 - (-15)] + 1 = 26$$
  
Number or columns say  $\mathbf{N} = (\mathbf{Uc} - \mathbf{Lc}) + \mathbf{1} = [40 - 15)] + 1 = 26$ 

(i) Row Major Wise Calculation of above equation

The given values are: B = 1500, W = 1 byte, I = 15, J = 20, Lr = -15, Lc = 15, N = 26

Address of A [ I ][ J ] = B + W \* [ N \* ( I – Lr ) + ( J – Lc ) ]

= 
$$1500 + 1* [26* (15 – (-15))) + (20 – 15)]$$
=  $1500 + 1* [26* 30 + 5]$ 
=  $1500 + 1* [780 + 5]$ 
=  $1500 + 785$ 
=  $2285$  [Ans]

(ii) Column Major Wise Calculation of above equation

The given values are: B = 1500, W = 1 byte, I = 15, J = 20, Lr = -15, Lc = 15, M = 26  
Address of A [ I ][ J ] = B + W \* [ (I - Lr) + M \* (J - Lc) ]  
= 
$$1500 + 1 * [(15 - (-15)) + 26 * (20 - 15)]$$
  
=  $1500 + 1 * [30 + 26 * 5]$   
=  $1500 + 1 * [160]$   
=  $1660$  [Ans]

## Total memory allocated to an Array = Number of elements \* size of one element

## **Single Dimension:**

Total memory allocated for an Integer Array of N elements

- = Number of elements \* size of one element
- = N \* 4 bytes
- = 10 \* 4 bytes = 40 Bytes, where N = 10
- = 500 \* 4 bytes = 2000 Bytes, where N = 500

Total memory allocated for an character Array of N elements

- = Number of elements \* size of one element
- = N \* 1 Byte
- = 10 \* 1 Byte = 10 Bytes, where N = 10
- = 500 \* 1 Byte = 500 Bytes, where N=500

#### **Two Dimensions:**

#### Total memory allocated for 2D Array

- = Number of elements \* size of one element
- = Number of Rows \* Number of Columns \* Size of one element

## Total memory allocated for an Integer Array of size M X N

- = Number of elements \* size of one element
- =M Rows\* N Columns \* 4 Bytes
- = 10\*10 \* 4 bytes = 400 Bytes, where M =N = 10
- = 500\*5\*4 bytes= **10000 Bytes**, where M=500 and N= 5

#### Total memory allocated for a character Array of N elements

- = Number of elements \* size of one element
- = M Rows\* N Columns \* 1 Byte
- = 10\*10 \* 1 Byte = 100 Bytes, where N = 10
- = 500\*5\*1 Byte = **2500 Bytes**, where M=500 and N= 5

## Example 2:

Each element of an array arr[15][20] requires 'W' bytes of storage. If the address of arr[6][8] is 4440 and the base address at arr[1][1] is 4000, find the width 'W' of each cell in the array arr[][] when the array is stored as Column Major Wise.

### **Solution: Example 2:**

Each element of an array arr[15][20] requires 'W' bytes of storage. If the address of arr[6][8] is 4440 and the base address at arr[1][1] is 4000, find the width 'W' of each cell in the array arr[][] when the array is stored as Column Major Wise.

#### Given:

 $\mathbf{B} = \text{Base address} = 4000$ 

I = Row subscript of element whose address is to be found = 6

J = Column subscript of element whose address is to be found = 8

W = Storage Size of one element stored in the array (in byte) = NOT Given

Lr = Lower limit of row/start row index of matrix = 1

Lc = Lower limit of column/start column index of matrix = 1

M(or R) = Number of row of the given matrix = 15

N (or C) = Number of column of the given matrix = 20

## **Solution: Example 2:**

Each element of an array arr[15][20] requires 'W' bytes of storage. If the address of arr[6][8] is 4440 and the base address at arr[1][1] is 4000, find the width 'W' of each cell in the array arr[][] when the array is stored as Column Major Wise.

#### Given:

 $\mathbf{B} = \text{Base address} = 4000$ 

I = Row subscript of element whose address is to be found = 6

J = Column subscript of element whose address is to be found = 8

W = Storage Size of one element stored in the array (in byte) = NOT Given

Lr = Lower limit of row/start row index of matrix = 1

Lc = Lower limit of column/start column index of matrix = 1

M(or R) = Number of row of the given matrix = 15

N (or C) = Number of column of the given matrix = 20

Address of [I, J]th element in column-major

$$= B + W[R(J - L_c) + (I - L_r)]$$

$$\Rightarrow 4440 = 4000 + W[15(8-1) + (6-1)]$$

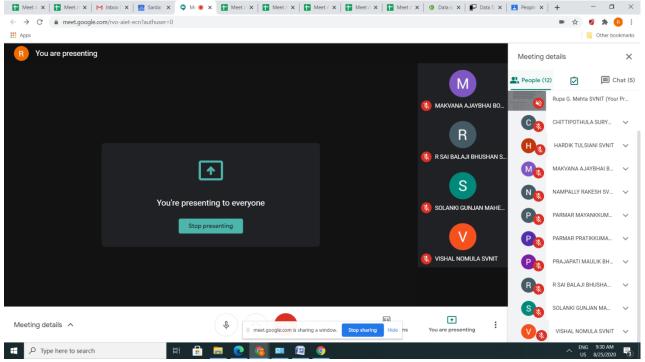
$$\Rightarrow 4440 = 4000 + W[15(7) + 5]$$

$$\Rightarrow 4440 = 4000 + W[105 + 5]$$

$$\Rightarrow 4440 = 4000 + W[110]$$

$$\Rightarrow$$
 W[110] = 440

$$\Rightarrow$$
 W = 4.



#### Example 3:

A matrix ARR[-4...6, 3...8] is stored in the memory with each element requiring 4 bytes of storage. If the base address is 1430, find the address of ARR[3][6] when the matrix is stored in Row Major Wise.

#### Given:

 $\mathbf{B} = \text{Base address} = 1430$ 

I = Row subscript of element whose address is to be found = 3

J = Column subscript of element whose address is to be found = 6

W = Storage Size of one element stored in the array (in byte) = 4

Lr = Lower limit of row/start row index of matrix = -4

Lc = Lower limit of column/start column index of matrix = 3

M(or R) = Number of row of the given matrix = 6 - (-4) + 1 = 11

N (or C) = Number of column of the given matrix = 8 - 3 + 1 = 6

## Example 3:

A matrix ARR[-4...6, 3...8] is stored in the memory with each element requiring 4 bytes of storage. If the base address is 1430, find the address of ARR[3][6] when the matrix is stored in Row Major Wise.

#### Given:

- $\mathbf{B} = \text{Base address} = 1430$
- I = Row subscript of element whose address is to be found = 3
- J = Column subscript of element whose address is to be found = 6
- W = Storage Size of one element stored in the array (in byte) = 4
- Lr = Lower limit of row/start row index of matrix = -4
- Lc = Lower limit of column/start column index of matrix = 3
- M(or R) = Number of row of the given matrix = 6 (-4) + 1 = 11
- N (or C) = Number of column of the given matrix = 8 3 + 1 = 6

## Address of [I, J]th element in row-major

- $= B + W[C(I L_r) + (J L_c)]$
- $\Rightarrow$  Address of ARR[3][6] = 1430 + 4[6(3 (-4)) + (6 3)]
- $\Rightarrow$  Address of ARR[3][6] = 1430 + 4[6(3 + 4) + 3]
- $\Rightarrow$  Address of ARR[3][6] = 1430 + 4[6(7) + 3]
- $\Rightarrow$  Address of ARR[3][6] = 1430 + 4[42 + 3]
- $\Rightarrow$  Address of ARR[3][6] = 1430 + 4[45]
- $\Rightarrow$  Address of ARR[3][6] = 1430 + 180
- $\Rightarrow$  Address of ARR[3][6] = 1610.