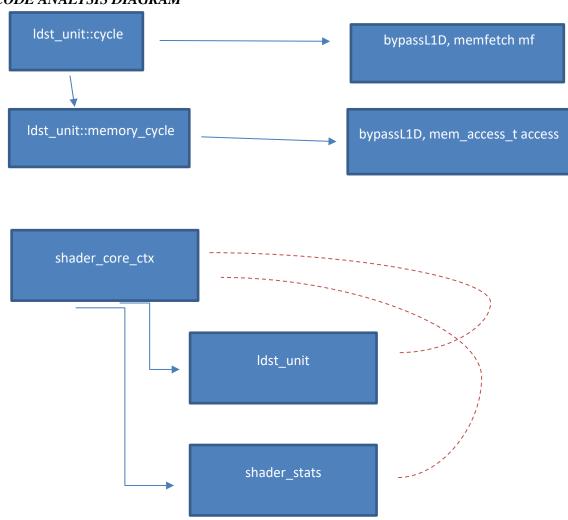
ECE786: PROGRAMMING ASSIGNMENT THREE - A

TASK: Load Bypass Mechanism

CODE ANALYSIS DIAGRAM



CODE ANALYSIS DISCUSSION

- The task requires to enforce L1 Data Cache bypass for loads with specific addresses. The load store part of gpgpu-sim is handled by the ldst_unit class that manages global, local, shared and texture memory accesses.
- The load instructions are issued for global and local memory. Investigating the methods of the ldst_unit class reveals two methods namely, memory_cycle() and cycle() that manage accesses to the L1 Data Cache within the simulator.
- The above diagram depicts the specific variables and class objects (bypassL1D, access, mf) that have been utilized to bypass the L1 Data Cache for the address range given in the assignment. This handles the first part of the assignment.
- It is also required to print the L1 bypasses for each kernel. Within a gpu each kernel is allocated to one SM. The SM is modelled by the shader_core_ctx class in the simulator.
- This SM class also contains its own instances of ldst_unit (m_ldst_uni)t and shader_stats (m_shader_stats). These classes also have a pointer to the SM object for which they are instantiated as depicted in the above diagram,

- Based on the prior information, the counters for each kernel can be modelled in the following two ways:
 - First, a counter can be declared within the shader_core_ctx class which will be
 updated in the memory_cycle and cycle methods of ldst_unit. This requires building
 public methods to access the counter as the shader_core_ctx pointer is protected
 member in ldst_unit. In a similar way, shader_stats class can access this counter in its
 print method to print the bypasses for each kernel.
 - Second, a cumulative counter and a static counter declared in the shader_stats class.
 The cumulative counter is updated in the ldst_unit methods. The print method of
 shader_stats will access the cumulative counter and subtract it with static counter to
 get the bypasses for each kernel. The static counter stores the prior cumulative
 counter value.
- Because of the simplicity of the code, the second method has been utilized to implement the count mechanism for the bypasses.

CODE SOLUTION SNIPPETS

```
// ********************************
// make and increment the counter for L1 Data Cache Load bypasses - Aditya
// ******************

if(inst.is_load())
{
    new_addr_type addr = access.get_addr();
    if((addr >= 0xc0000000) && ( addr <= 0xc000fffff)){
        bypassL1D = true;
        this->m_stats->count_cum_load_bypasses++;
    }
}
```

RESULT

Cumulative Load Bypasses	Load Bypasses for each kernel
26	26
163	137
963	800
5884	4921
34108	28224
181692	147584
609201	427509
671956	62755
674391	2435
674406	15

Note: Sample count for one of the kernels