

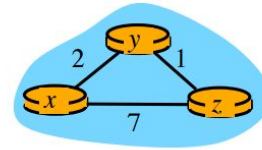
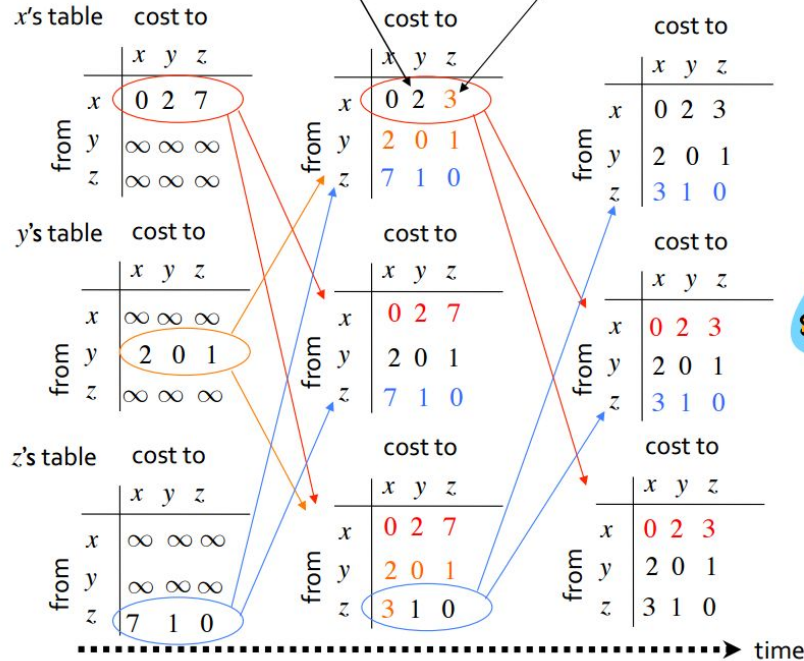
# EECS 489 Discussion 7

# UDP Demo

# Distance Vector Example

$$D[x,y] = \min\{c(x,y) + D[y,y], c(x,z) + D[z,y]\} \\ = \min\{2+0, 7+1\} = 2$$

$$D[x,z] = \min\{c(x,y) + D[y,z], c(x,z) + D[z,z]\} \\ = \min\{2+1, 7+0\} = 3$$



# Q1

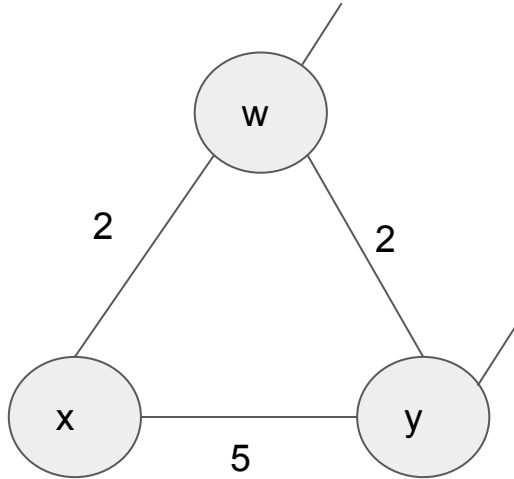
Consider the count-to-infinity problem in the distance vector routing. Will the count-to-infinity problem occur if we decrease the cost of a link? Why? How about if we connect two nodes which do not have a link?

# Q1

Consider the count-to-infinity problem in the distance vector routing. Will the count-to-infinity problem occur if we decrease the cost of a link? Why? How about if we connect two nodes which do not have a link?

No. Decreasing the cost of a link won't cause a loop

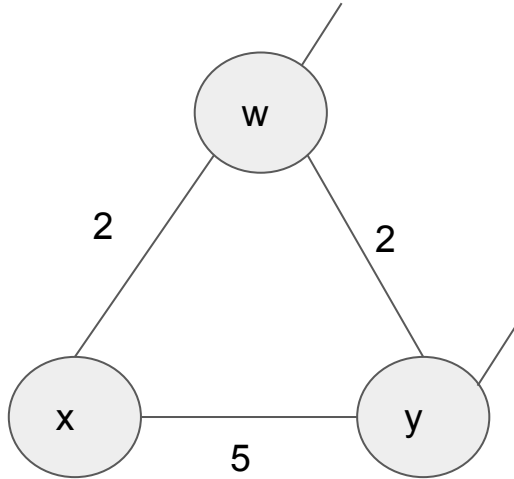
## Q2



Consider the network fragment shown below. Node  $w$  has a minimum-cost path to destination  $u$  (not shown) of 5, and node  $y$  has a minimum-cost path to  $u$  of 6. The complete paths from  $w$  and  $y$  to  $u$  are not shown. All link costs in the network have strictly positive values

- Give  $x$ 's distance vector for destinations  $w$ ,  $y$ , and  $u$
- Give a link-cost change for either  $c(x, w)$  or  $c(x, y)$  such that  $x$  will inform its neighbors of a new minimum-cost path to  $u$  as a result of executing the distance-vector algorithm
- Give a link-cost change for either  $c(x, w)$  or  $c(x, y)$  such that  $x$  will not inform its neighbors of a new minimum-cost path to  $u$  as a result of executing the distance- vector algorithm

## Q2



Consider the network fragment shown below. Node w has a minimum-cost path to destination u (not shown) of 5, and node y has a minimum-cost path to u of 6. The complete paths from w and y to u are not shown. All link costs in the network have strictly positive values

- Give x's distance vector for destinations w, y, and u  
 $D_x(w) = 2, D_x(y) = 4, D_x(u) = 7$
- Give a link-cost change for either  $c(x, w)$  or  $c(x, y)$  such that x will inform its neighbors of a new minimum-cost path to u as a result of executing the distance-vector algorithm  
Change  $c(x, y) < 1$
- Give a link-cost change for either  $c(x, w)$  or  $c(x, y)$  such that x will not inform its neighbors of a new minimum-cost path to u as a result of executing the distance-vector algorithm  
Make  $c(x, y)$  anything greater than 1