Fundamentals of Data Structure

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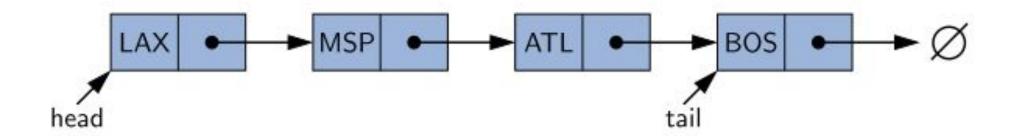
VJTI, Mumbai-19

Slides are prepared from

- 1. Data Structures and Algorithms in Java, 6th edition, by M. T. Goodrich, R. Tamassia, and M. H. Goldwasser, Wiley, 2014
- 2.Data Structures and Algorithms in Java, by Robert Lafore, Second Edition, Sams Publishing

Singly Linked Lists

- A linked list's representation relies on the collaboration of many objects
- Minimally, the linked list instance must keep a reference to the first node of the list, known as the head
- Without an explicit reference to the head, there would be no way to locate that node
- The last node of the list is known as the tail
- The tail of a list can be found by traversing the linked list— starting at the head and moving from one node to another by following each node's next reference
- We can identify the tail as the node having null as its next reference



Circularly Linked Lists

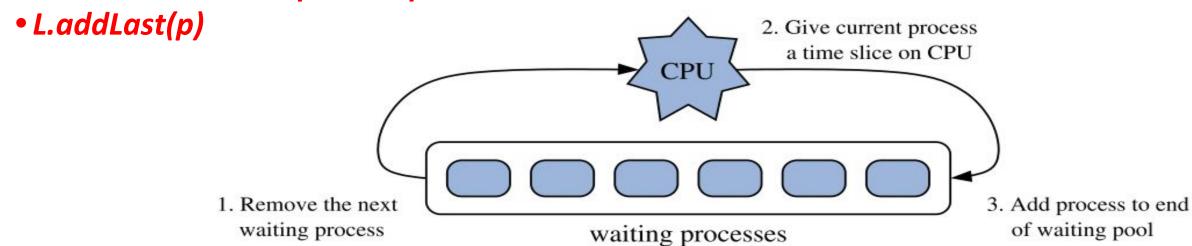
- Linked lists are traditionally viewed as storing a sequence of items in a linear order, from first to last
- Many applications in which data can be more naturally viewed as having a cyclic order
 - For example, many multiplayer games are turn-based, with player A taking a turn, then player B, then player C, and so on
 - City buses often run on a continuous loop, making stops in a scheduled order
 - Round-Robin Scheduling

Round-Robin Scheduling

- •An operating system is managing many processes that are currently active on a computer
- Most operating systems allow processes to effectively share use of the CPUs time
- •A process is given a short turn to execute, known as a *time slice*, but it is interrupted when the slice ends, even if its job is not yet complete
- •Each active process is given its own time slice, taking turns in a cyclic order
- •New processes can be added to the system, and processes that complete their work can be removed

Round-Robin Scheduler Implementation

- We can implement a round robin scheduler using a queue **Q** by repeatedly performing the following steps:
 - *e* = **Q**.dequeue()
 - Service element e
 - Q.enqueue(e)
- A round-robin scheduler could be implemented with a traditional linked list, by repeatedly performing the following steps on linked list
 - process p = L.removeFirst()
 - Give a time slice to process p

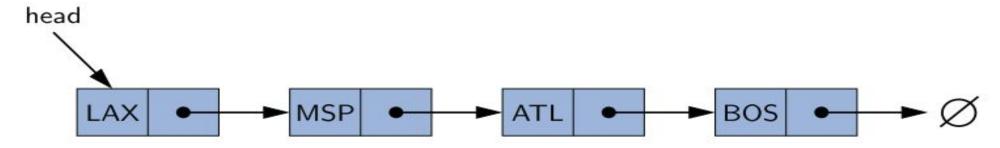


Round-Robin Scheduler Implementation Challenge

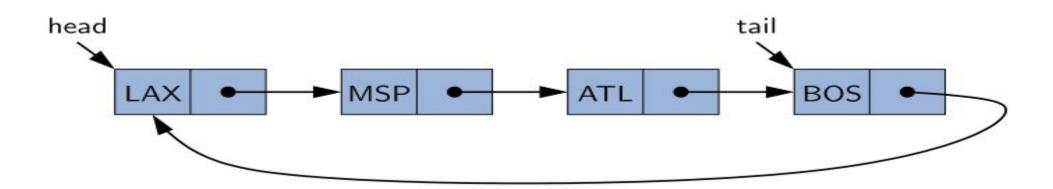
- Unfortunately, there are drawbacks to the both approaches discussed
- It is unnecessarily inefficient to repeatedly throw away a node from one end and create a new node for the same element at other end
- How to provide a more efficient data structure for representing a cyclic order?

Designing and Implementing a Circularly Linked List

Extend a singularly linked list to a Circularly Linked List



• Essentially a singularly linked list in which the next reference of the tail node is set to refer back to the head of the list (rather than null)



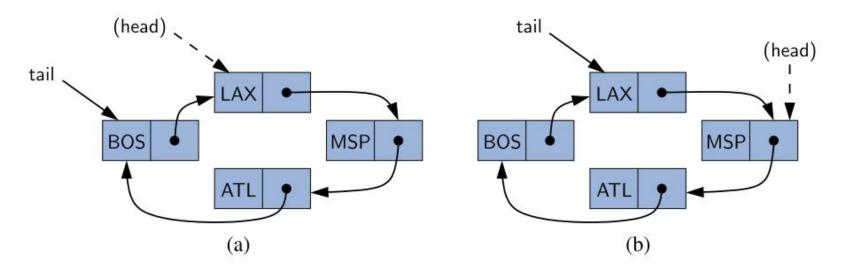
Designing and Implementing a Circularly Linked List

- We use new discussed model to design and implement a new <u>CircularlyLinkedList</u> class, which supports all of the public behaviors of SinglyLinkedList class and one additional method
 - rotate(): Moves the first element to the end of the list
- We need to following update methods of SinglyLinkedList class
 - addFirst()
 - addLast()
 - removeFirst()
- Now, round-robin scheduling can be efficiently implemented by repeatedly performing the following steps on a circularly linked list C:
 - Give a time slice to process C.first()
 - C.rotate()
- In this implementation, we remove the head reference, and get the head as tail.getNext()

```
public class CircularlyLinkedList<E> { /* node class definition*/
                                          // we store tail (but not head)
     private Node<E> tail = null;
     private int size = 0; // number of nodes in the list
     public CircularlyLinkedList() { } // constructs an initially empty list
     public int size() { return size; }
     public boolean isEmpty() { return size == 0; }
     public E first() {
          if (isEmpty()) return null;
          return tail.getNext().getElement();
                                               // the head is *after* the tail
     public E last() {
                                     // returns (but does not remove) the last element
          if (isEmpty()) return null;
          return tail.getElement();
```

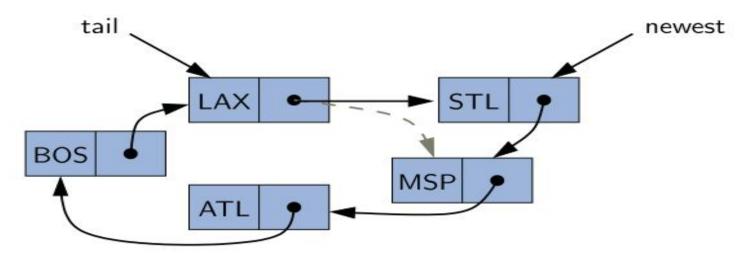
```
public void rotate() { // rotate the first element to the back of the list
  if (tail != null) // if empty, do nothing
  tail = tail.getNext(); // the old head becomes the new tail
}
```

• We do not move any nodes or elements, we simply advance the tail reference to point to the node that follows it (the implicit head of the list)



```
public void addLast(E e) { // adds element e to the
end of the list
   addFirst(e); // insert new element at front of list
   tail = tail.getNext(); // now new element becomes
the tail
}
```

- We can add a new element at the front of the list by creating a new node and linking it just after the tail
 of the list
- To implement the addLast method, we can rely on the use of a call to addFirst and then immediately advance the tail reference so that the newest node becomes the last



```
public E removeFirst() { // removes and returns the first element
   if (isEmpty()) return null; // nothing to remove
   Node<E> head = tail.getNext();
   if (head == tail)
      tail = null; // must be the only node left
   else
      tail.setNext(head.getNext()); // removes "head" from the list
   size--;
   return head.getElement();
} //end of class definition
```

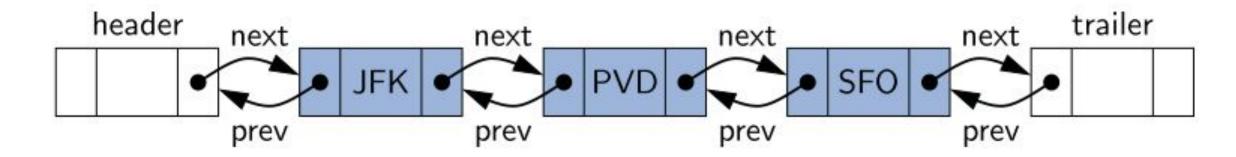
Removing the first node from a circularly linked list can be accomplished by simply updating the next field
of the tail node to bypass the implicit head

Doubly Linked Lists

- In a singly linked list, each node maintains a reference to the node that is immediately after it
- A potential problem with singly linked lists is that it's difficult to traverse backward along the list
- There are limitations that stem from the asymmetry of a singly linked list
 - We cannot efficiently delete an arbitrary node from an interior position of the list if only given a reference to that node (because we cannot determine the node that immediately precedes the node to be deleted)
- Symmetry in a linked list is each node keeps an explicit reference to the node before it and a reference to the node after it. Such a structure is known as a doubly linked list
- The doubly linked list allows you to traverse backward as well as forward through the list
- These lists allow a greater variety of O(1)-time update operations, including insertions and deletions at arbitrary positions within the list

Header and Trailer Sentinels

- In order to avoid some special cases when operating near the boundaries of a doubly linked list, it helps to add special nodes at both ends of the list:
 - a header node at the beginning of the list, and
 - a trailer node at the end of the list
- These "dummy" nodes are known as sentinels (or guards), and they do not store elements of the primary sequence



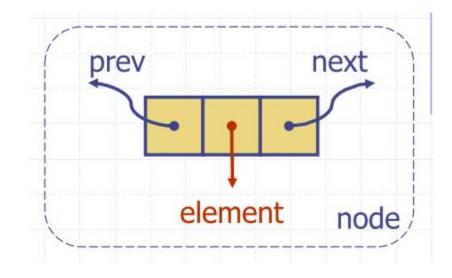
Doubly Linked List Interface

A DoublyLinkedList class supports the following public methods:

```
size(): Returns the number of elements in the list
isEmpty(): Returns true if the list is empty, and false otherwise
first(): Returns (but does not remove) the first element in the list
last(): Returns (but does not remove) the last element in the list
addFirst(e): Adds a new element to the front of the list
addLast(e): Adds a new element to the end of the list
removeFirst(): Removes and returns the first element of the list
removeLast(): Removes and returns the last element of the list
addBetween(e, predecessor, successor): Adds a new element between two nodes
predecessor and successor
```

Implementing a Doubly Linked List Class (Node)

```
private static class Node<E> {
     private E element; // reference to the element stored at this node
                             // reference to the previous node in the list
     private Node<E> prev;
     private Node<E> next; // reference to the subsequent node in the list
     public Node(E e, Node<E> p, Node<E> n) {
          element = e;
          prev = p; next = n;
     public E getElement() { return element; }
     public Node<E> getPrev() { return prev; }
     public Node<E> getNext() { return next; }
     public void setPrev(Node<E> p) { prev = p; }
     public void setNext(Node<E> n) { next = n; }
```



Nodes store:

- element
- link to the previous node
- link to the next node

Implementing a Doubly Linked List Class

trailer

next

prev

```
public class DoublyLinkedList<E> { /* Node Class*/
     private Node<E> header;
                                                                                       header
     private Node<E> trailer;
     private int size = 0;
     public DoublyLinkedList() { \\Constructs a new empty list
          header = new Node<>(null, null, null); // create header
           trailer = new Node<>(null, header, null);
                                                    // trailer is preceded by header
          header.setNext(trailer);
                                           // header is followed by trailer
     public int size() { return size; }
                                           //Returns the number of elements in the linked list
     public boolean isEmpty() { return size == 0; } //Tests whether the linked list is empty
                       //Returns (but does not remove) the first element of the list
     public E first() {
          if (isEmpty()) return null;
          return header.getNext().getElement();
     public E last() {
                       //Returns (but does not remove) the last element of the list
          if (isEmpty()) return null;
          return trailer.getPrev().getElement();
```

Implementing a Doubly Linked List Class

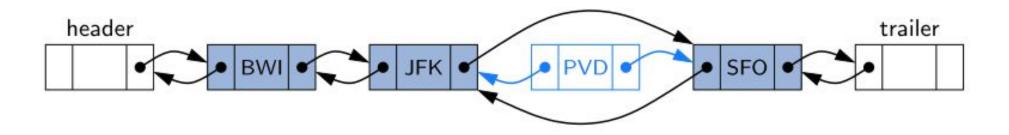
```
private void addBetween(E e, Node<E> predecessor, Node<E> successor) {
      //Adds element e to the linked list in between the given nodes
     Node<E> newest = new Node<>(e, predecessor, successor); // create and link a new node
     predecessor.setNext(newest);
     successor.setPrev(newest);
                                                                                    Create a new node,
                                                                                    Set predecessor NEXT to the new element,
     size++;
                                                                                    Set succesor PREV to the new element,
                                                                                    increase size.
public void addFirst(E e) { //Adds element e to the front of the list
     addBetween(e, header, header.getNext());
public void addLast(E e) { //Adds element e to the end of the list
     addBetween(e, trailer.getPrev(), trailer);
             header
                                                                                                    trailer
```

Implementing a Doubly Linked List Class

private E remove(Node<E> node) { \Removes the given node from the list and returns its element

```
Node<E> predecessor = node.getPrev();
     Node<E> successor = node.getNext();
     predecessor.setNext(successor);
     successor.setPrev(predecessor);
     size--;
     return node.getElement(); }
public E removeFirst() { //Removes and returns the first element of the list
     if (isEmpty()) return null;
     return remove(header.getNext());}
public E removeLast() { //Removes and returns the last element of the list
     if (isEmpty()) return null;
     return remove(trailer.getPrev()); }
} //----- end of DoublyLinkedList class -----
```

- 1. Find Predecessor and successor of node,
- 2. Set predecessor NEXT to the successor,
- 3. Set succesor PREV to the predecessor,
- 4. decrease size.



Exercise

- Implement Stack and Queue ADT using linked list storage data structure
- Stack
 - push(data): theList.addFirst(data)
 - pop(): data = theList.removeFirst()
- Queue
 - enqueue(data):theList.addLast(data)
 - dequeue():theList.removeFirst()

```
public class LinkedStack<E> implements Stack<E> {
 private SinglyLinkedList<E> list = new SinglyLinkedList<>();  // an empty list
 public LinkedStack() { } // new stack relies on the initially empty list
 public int size() { return list.size(); }
 public boolean isEmpty() { return list.isEmpty(); }
 public void push(E element) { list.addFirst(element); }
 public E top() { return list.first(); }
 public E pop() { return list.removeFirst(); }
public class LinkedQueue<E> implements Queue<E> {
  private SinglyLinkedList<E> list = new SinglyLinkedList<>();  // an empty list
  public LinkedQueue() { } // new queue relies on the initially empty list
  public int size() { return list.size(); }
  public boolean isEmpty() { return list.isEmpty(); }
  public void enqueue(E element) { list.addLast(element); }
  public E first() { return list.first(); }
  public E dequeue() { return list.removeFirst(); }
```