

Binary Search on Minimax

Minimize the maximum

OR

Maximize the minimum

Farmer John has built a new long barn, with N ($2 \leq N \leq 100,000$) stalls. The stalls are located along a straight line at positions x_1, \dots, x_N ($0 \leq x_i \leq 1,000,000,000$). ¹⁰⁷

His C ($2 \leq C \leq N$) cows ¹⁰⁵ don't like this barn layout and become aggressive towards each other once put into a stall. To prevent the cows from hurting each other, FJ wants to assign the cows to the stalls, such that the minimum distance between any two of them is as large as possible. What is the largest minimum distance?

Input

t – the number of test cases, then t test cases follows.

* Line 1: Two space-separated integers: N and C

* Lines 2.. $N+1$: Line $i+1$ contains an integer stall location, x_i

Output

For each test case output one integer: the largest minimum distance.

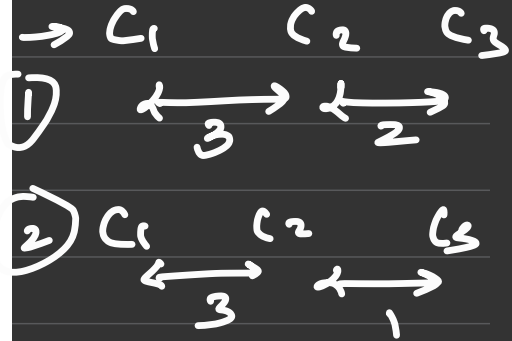
Example

Input:

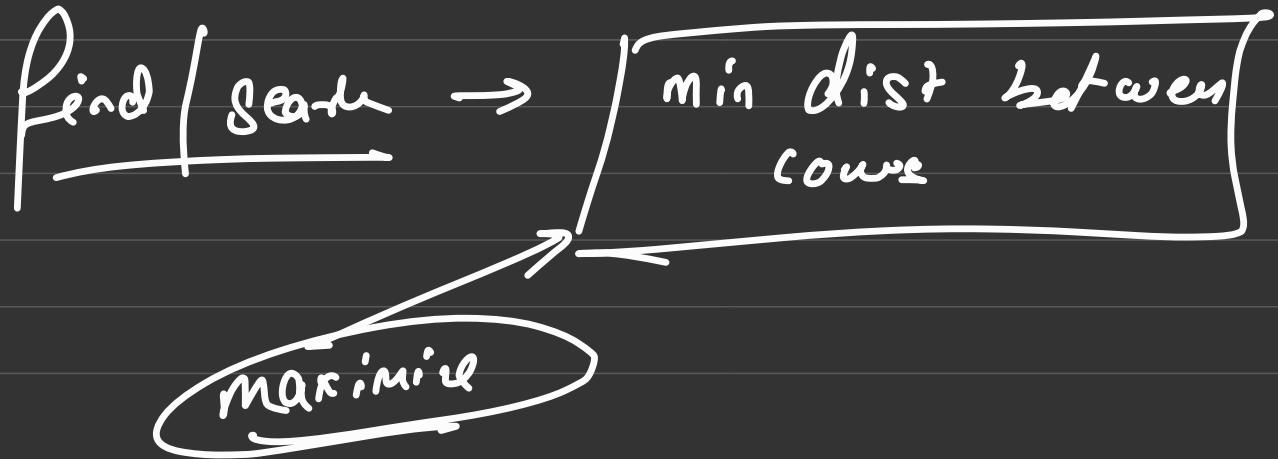
```
1
5 3
1
2
8
4
9
```

Output:

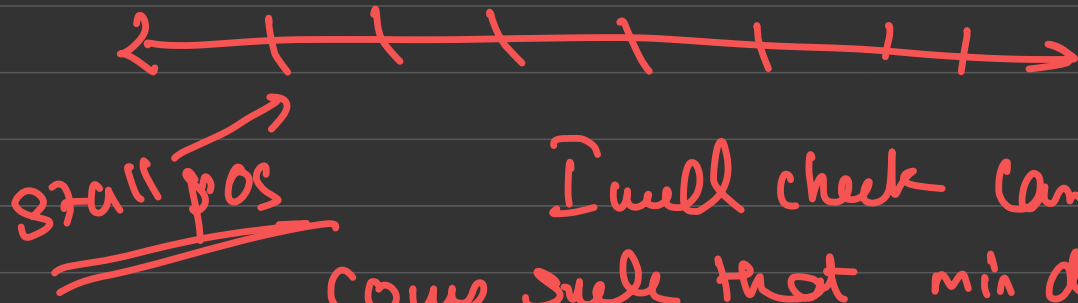
```
3
```



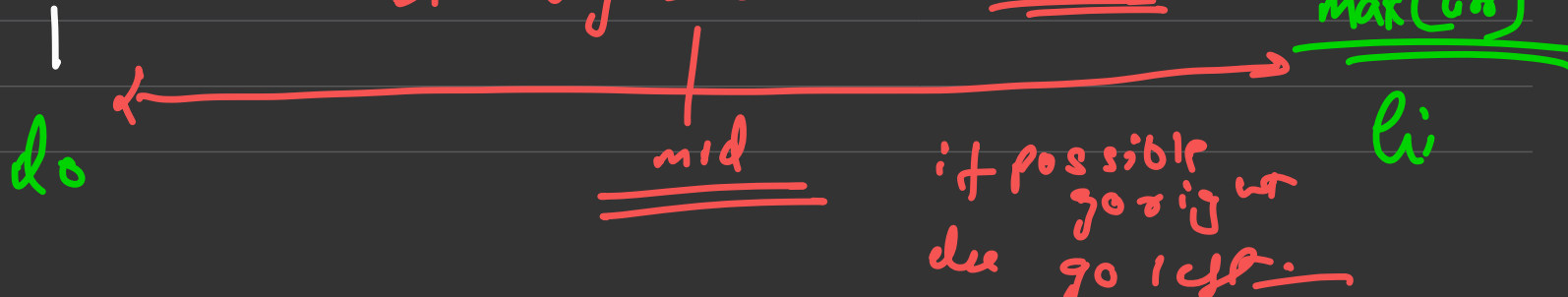
We have to calculate minimum dist btw any 2 rows. Then maximize this min dist so that every row is as far as possible.

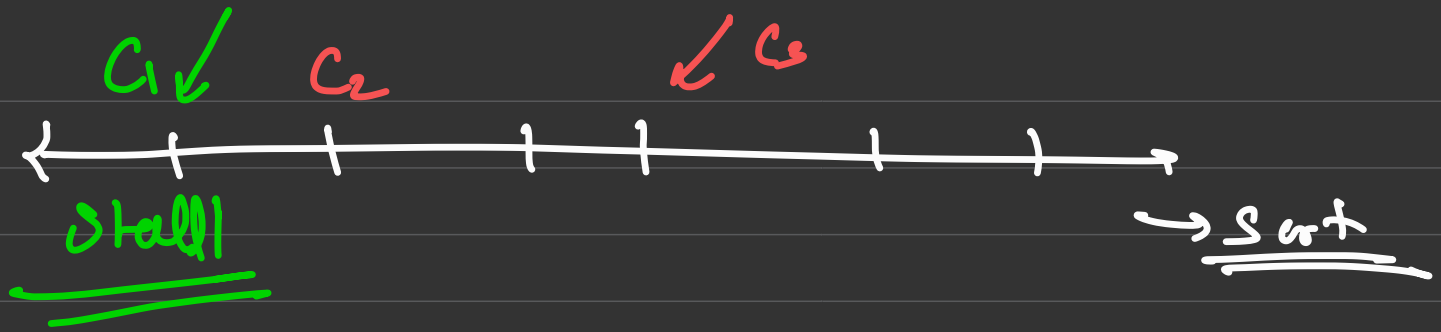


What can be in the worst case min dist
btw 2 cows.



I will check can I arrange the
cows such that min dist
btw any 2 cows is mid.





Q.1 You're given no. of pages in n books
and we have m students. Books are
arranged in sorted order based on no. of pages.

Divide the book among m students such
that maximum pages a student has to
read is minimized.

12, 34, 67, 90 \rightarrow 113

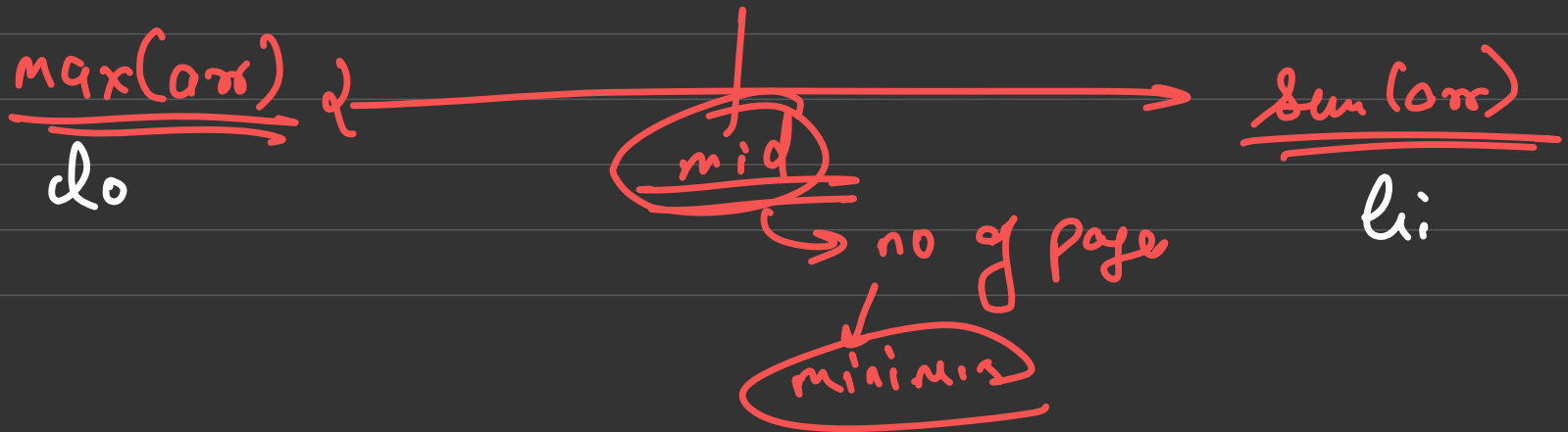
$m=2$

90

minimize the maximum no. of pages read.

Search space

will denote \rightarrow maximum no of pages for a student



mid denotes what is the maximum
no. of pages any student can read

$$O(\underbrace{n} \times \log(\underbrace{\text{sum} - \text{max}}))$$

Qⁿ You've a list of numbers, You need to divide them into K segments such that max sum on a segment is minimum possible.

[1, 3, 2, 4, 10, 8, 4, 2, 5, 3]

ans 12

$K=4$

Q₃

minimize

the

max candies

max candies \rightarrow search space

0

lo

mid

↑

$\max(a_i * b_i)$
hi

if you decide that max candies you'll

give to anyone won't be more than mid.

possible

→ M balloons

max candies any one will get
is → mid

sum

→ actual balloons yes
for any i^{th} day

q_i = balloon reqd — mid / b_i^0
balloon compensatio

sum + max(n, 0)

$$\sum x \leq M$$

→ lower

$b_i \rightarrow 1$ ballon

1 candy $\rightarrow \frac{1}{b_i}$ ballon

mid candy $\rightarrow \frac{\text{mid}}{b_i}$ ballon