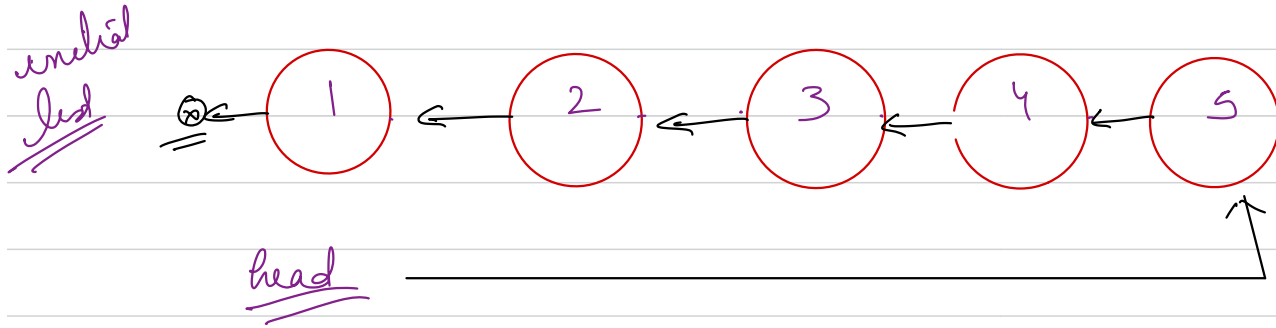
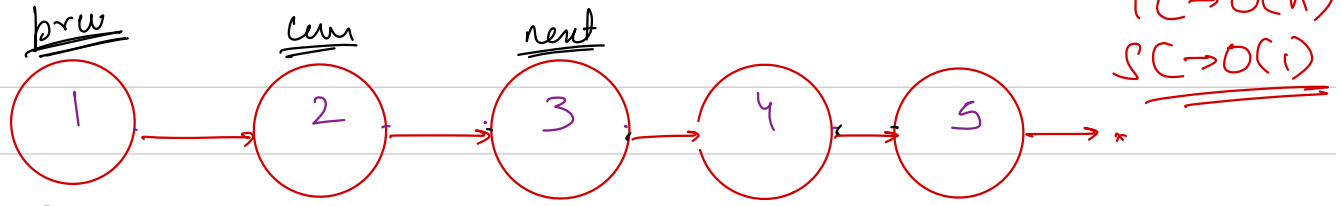


Q.1 Given a linked list of integer values in a node,
Reverse the linked list, by only manipulating the
next pointers. (Do not create a new LL).



Try to solve
both iteratively
and
recursively

Reverse Pointers Iterative



head

At last (5) should become the head.

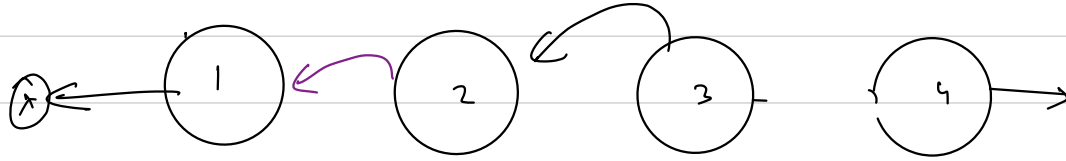
→ Before manipulating the ptr of any node, we should store the rest of the LL somewhere

→ Any node after the process should point to no prev node

prev, cur, next

cur.next = prev
prev = cur
cur = next
next = cur.next

← core logic →



prev = None

curr = 1

next = 2

initial
state

1

2

3

1st iteration

2

3

4

2nd iteration

3

4

None

insert stops
here

Reverse
Pointer
Recursive

└→ Base Case

└→ Self work

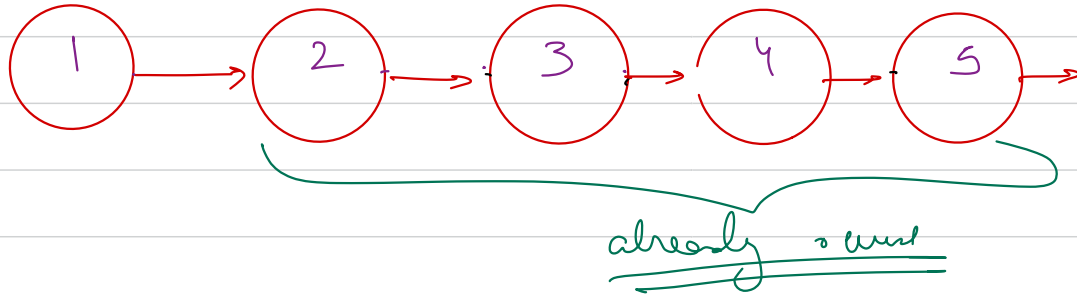
└→ Recursive Intuition

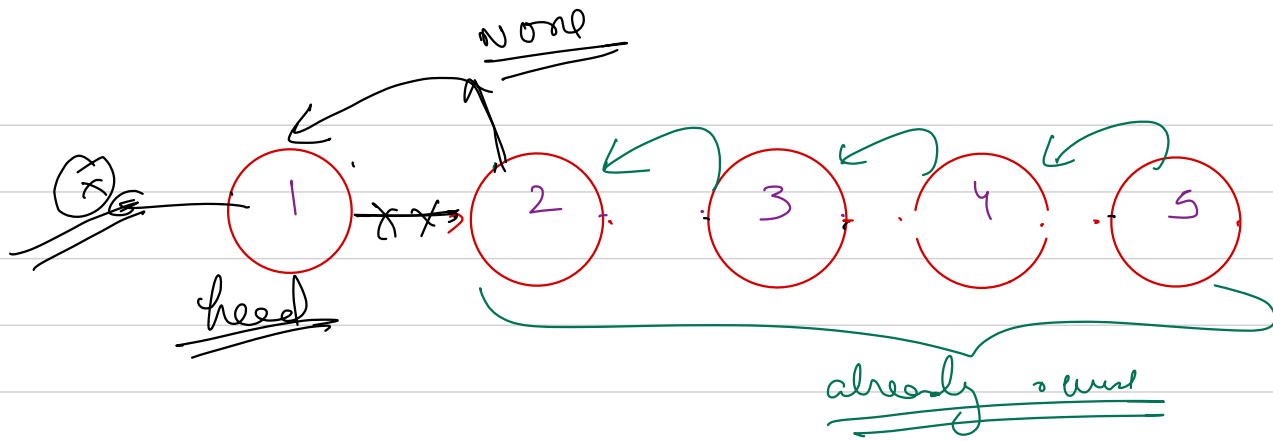
TC → $O(n)$

SC → $O(n)$

↓
due to
recursion call
stack

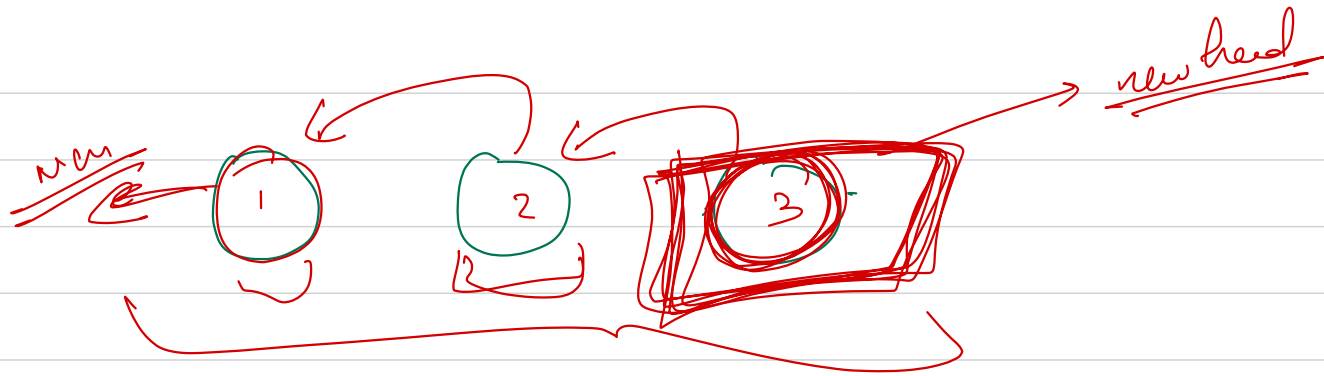
Base Case → Empty LL or LL of size 1





$f(\text{head}) = f(\text{head.next})$ → recurse until
 ↓
 reverse the ll

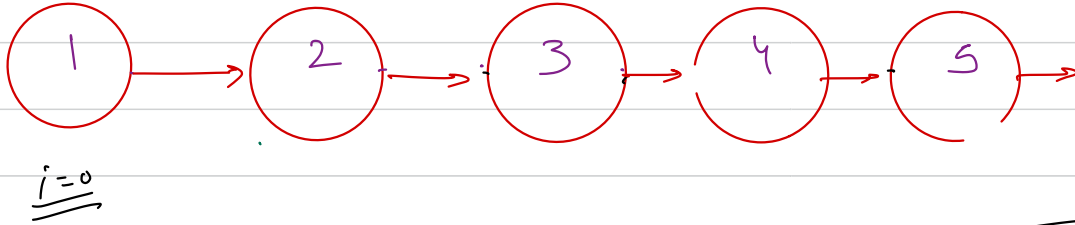
↪ self work as → $\text{head.next.next} = \text{head}$
 $\text{head.next} = \underline{\underline{\text{Null}}}$



Reverse
LL
Data
Iterative

You have to reverse a ll, without manipulating pointers & only replacing data

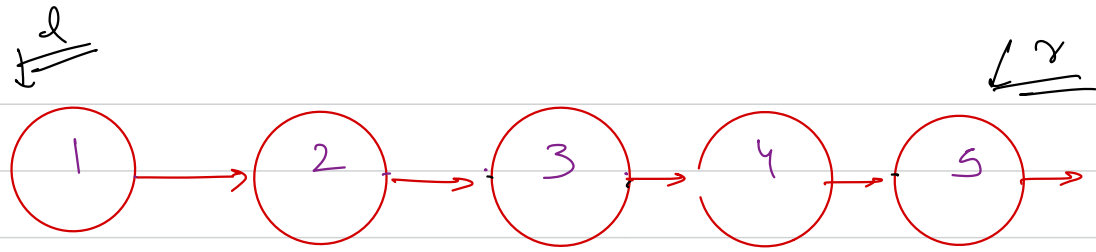
(Modify the original LL)



TC $\rightarrow \underline{\underline{O(n^2)}}$
SC $\rightarrow \underline{\underline{O(1)}}$

$\boxed{n-i-1}$

Reverse
Data
Recursively



We have a benefit with recursion (palindromic LL)
when you should stop??

odd \rightarrow $l == 0$ \leftarrow stop here

even \rightarrow $l \cdot next == 0$ stop

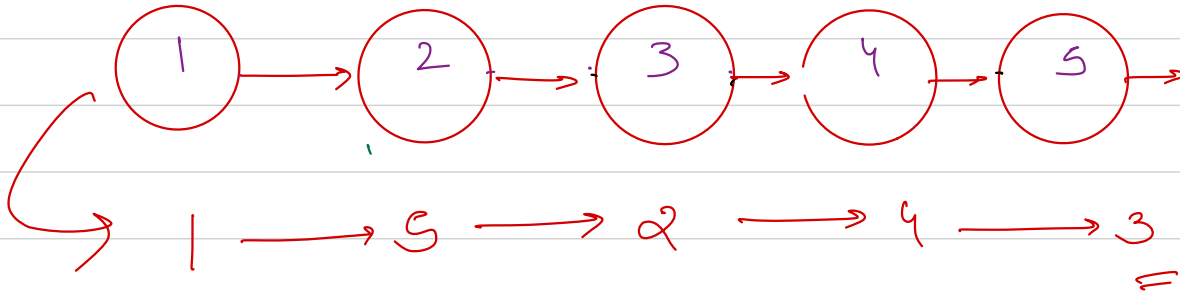
Q Given a linked list, rearrange the ll in-place

from $L_0 \rightarrow L_1 \rightarrow L_2 \rightarrow L_3 \dots \dots L_{n-1} \rightarrow L_n$ it

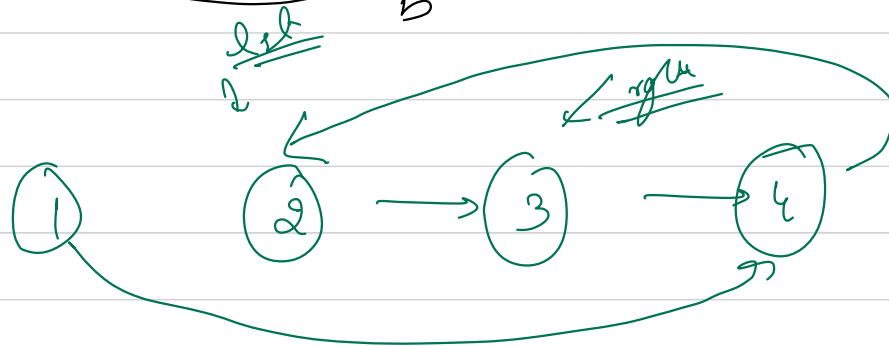
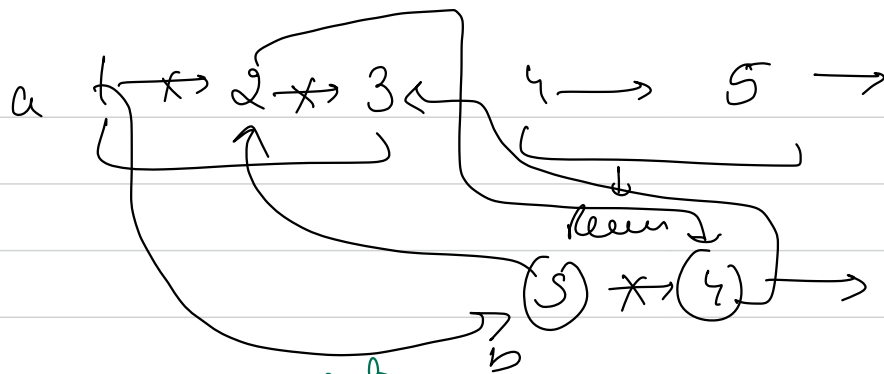
transf comes to

Hint \rightarrow Palindromic LL

$L_0 \rightarrow L_n \rightarrow L_1 \rightarrow L_{n-1} \rightarrow L_2 \rightarrow L_{n-2} \dots \dots$

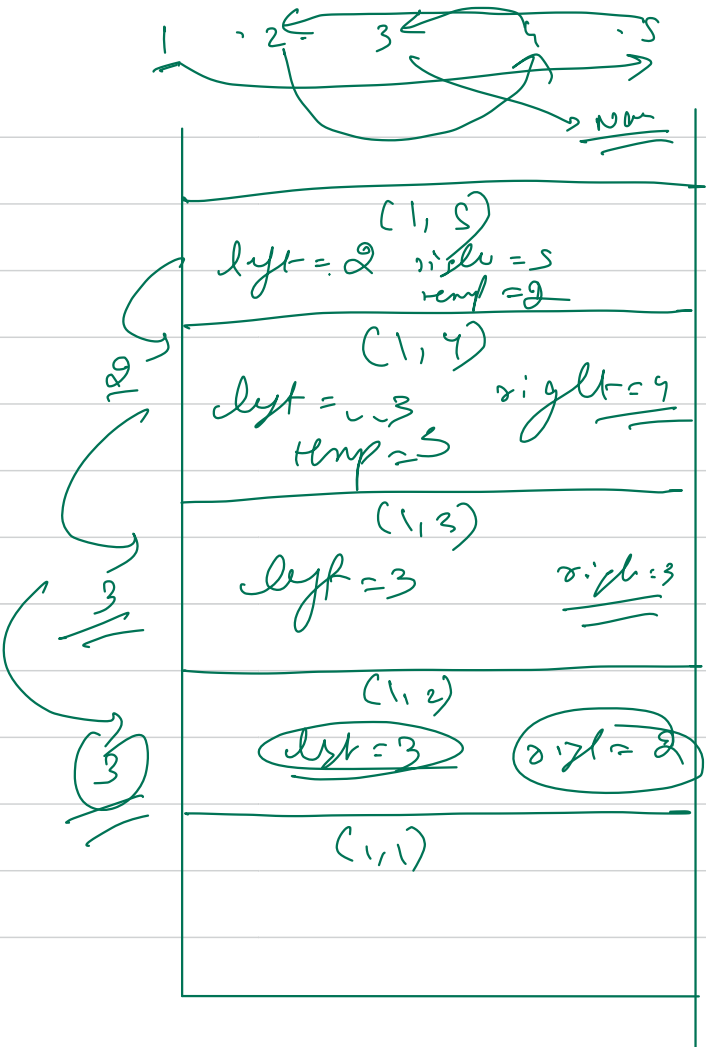


we can't
do data
swaps

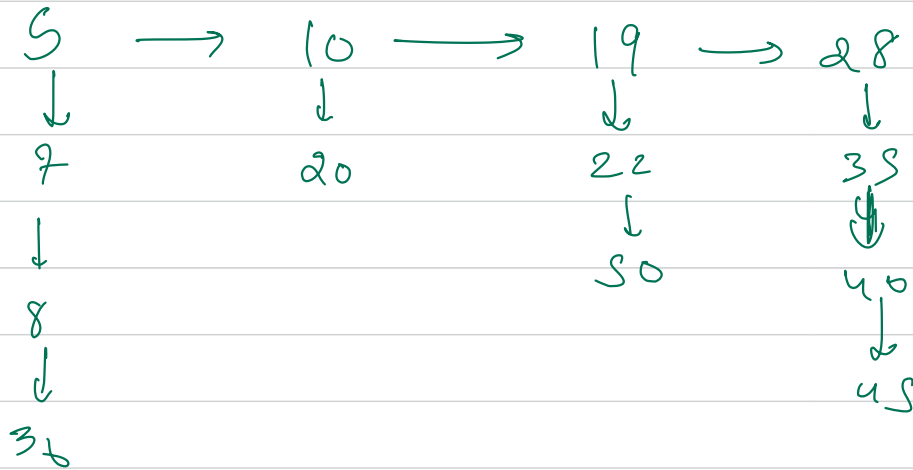


$$L_0 \rightarrow L_n \rightarrow L_1 \rightarrow L_{n-1}$$

```
199 def fold_helper(left, right):
200     if right == None:
201         return left
202
203     left = fold_helper(left, right.next)
204     if left == None:
205         return left
206
207     if left != right and left.next != right:
208         temp = left.next
209         left.next = right
210         right.next = temp
211         left = temp
212     else:
213         right.next = None
214
215     return left
```



Q2 You've a ll, with every node having a next ptr & bottom ptr. flatten the ll. all new ll is sorted



1. place

Merge 2 sorted LL

5 → 7 → 8 → 10 → 19 → 20 → 22 → 28 → 30 → 35 → 40 → 45 → 50

→ flat → head.next

5 → 10 → 19 → 20 → 22 → 28 → 35 → 40 → 25 → 50
↓
7
↓
8
↓
30

$O(n)$
 $O(n)$

$n \rightarrow$ no of nodes

(1) → head.next
(2) → head.next