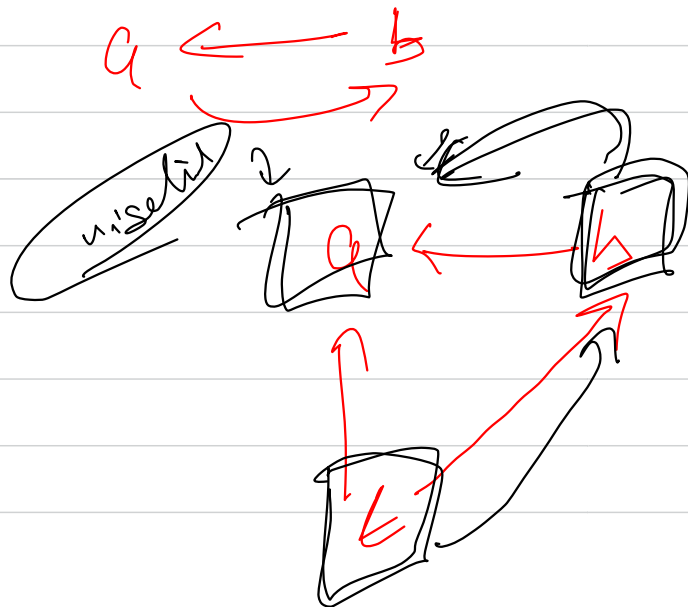
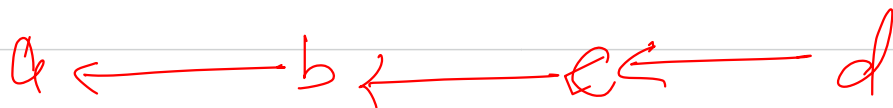


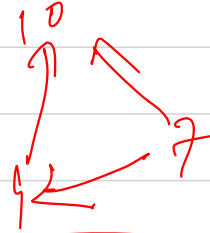
Parent



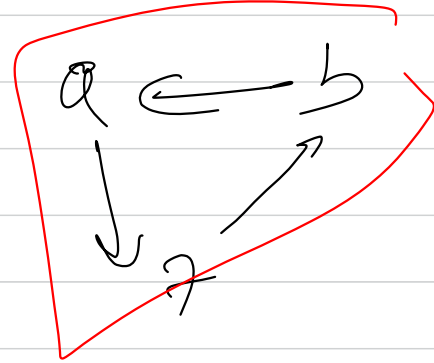
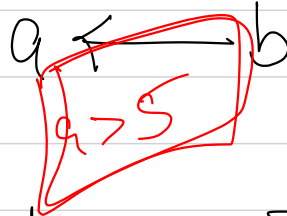
integer

$$q < 4$$

$$q > 5$$



$$a >$$



$q = 2$

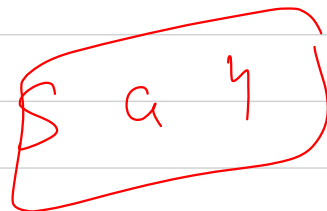
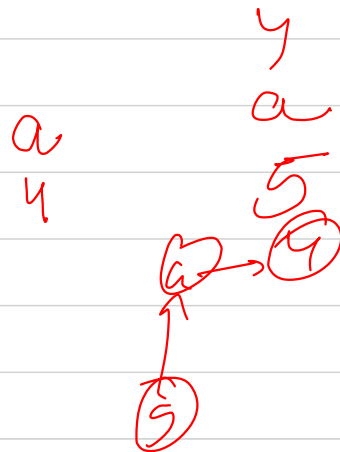
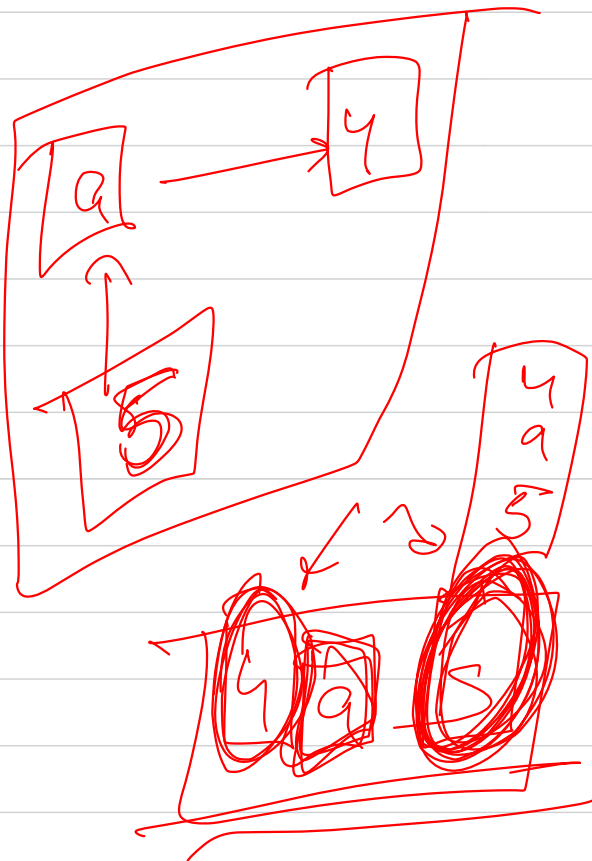
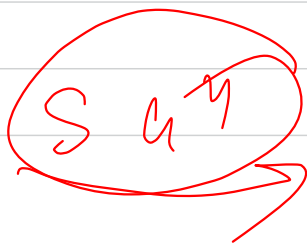
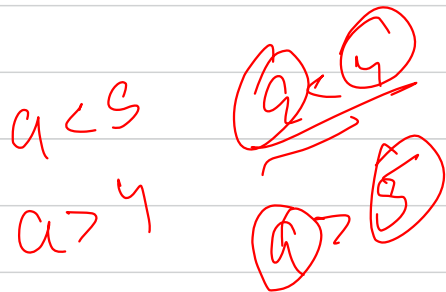
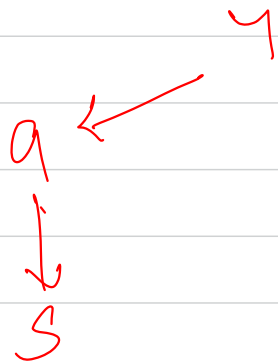
$$q > 7$$

$$10 > 7$$

$$10 > 9$$



$a > 7$   
 $b > 7$   
 $a > b$



$$a > s$$

$$a < y$$

$$a \rightarrow s$$

$$\uparrow$$

$$y \leftarrow$$

$$a \rightarrow s$$

$$\downarrow$$

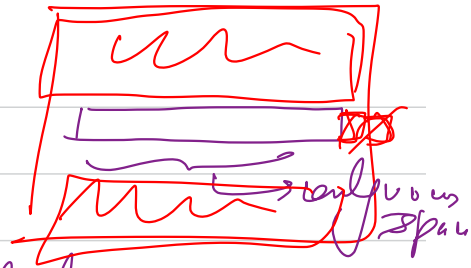
$$y \checkmark$$

$$a > s$$

$$a < y$$

$$s > y$$

# Linked Lists



Why do we even need a new DS called as LL.

Arrays

we need to give size of arrays  
arrays always have contiguous  
memory alloc

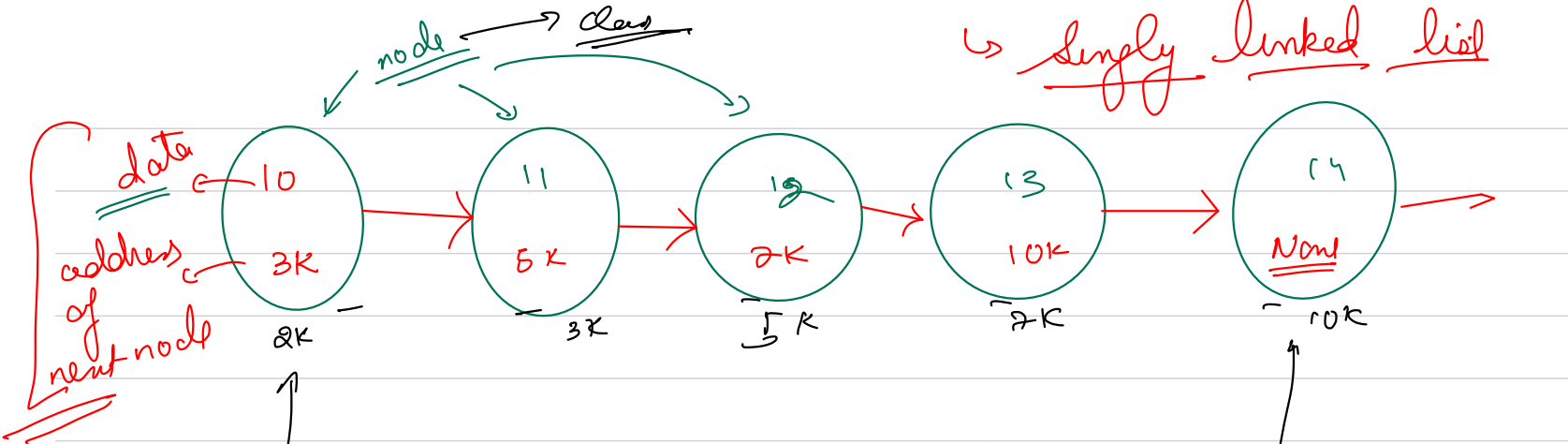
list/vectors/stack

they also consume contiguous  
memory spots

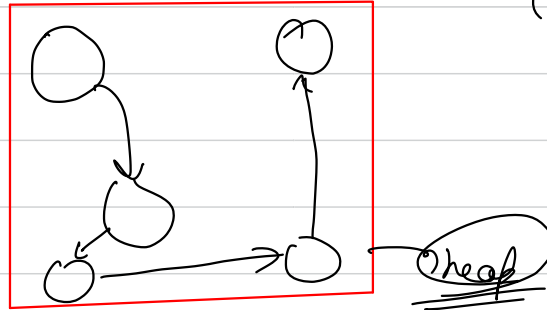
Arrays doesn't consume the available memory efficiently  
→ To optimize this we have got linked list

What is a LL??

→ LL are linear data structures that store one data element in an entity called as a Node and multiple such nodes are connected to form a chained list. The nodes are not reqd to be created in continuous memory blocks, so LL never consumes a continuous memory block.



The first node of LL is called as "Head"



→ memory view

the last node of LL is called as "Tail"

The button of previous and next page in browser are implemented using LL.

→ Music player software

→ cache mechanism

→ Hash Map

⋮





class Node :

data # value of node

next # address of next node

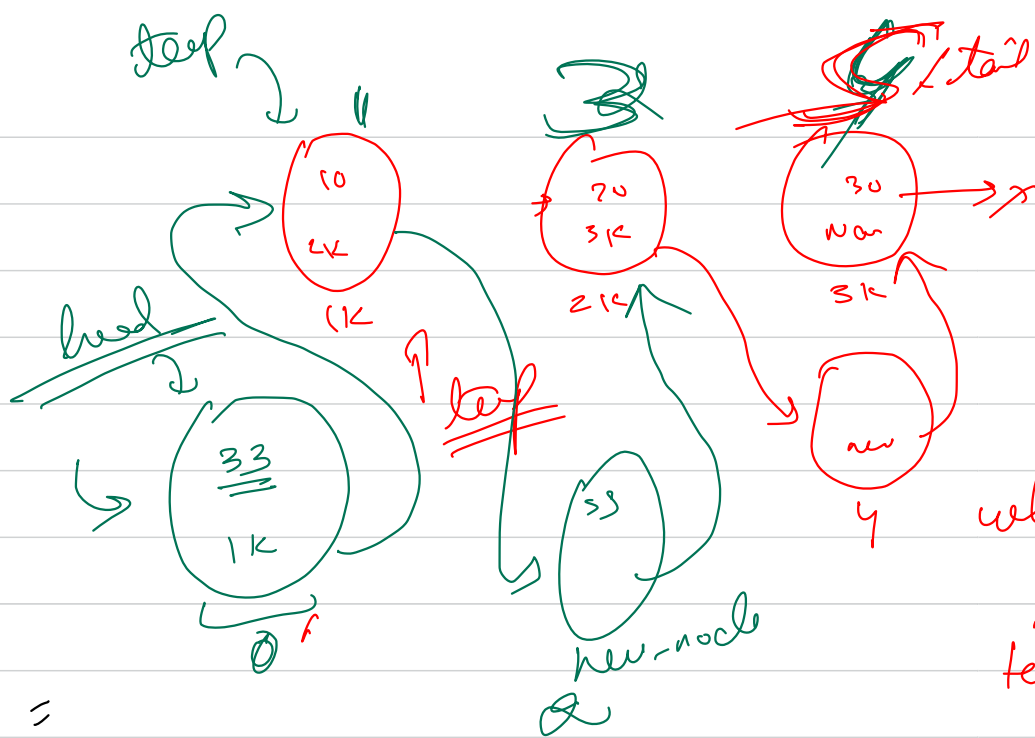
Adding data in LL

→ at the head →  $O(1)$

→ at the tail →  $O(n)$

→ in middle somewhere →  $O(n)$

dest → at at  $O^{th}$  node →  $O(n)$   
add at last node →  $O(1)$   
middle →  $O(n)$



while (temp != Null)

temp = temp->next

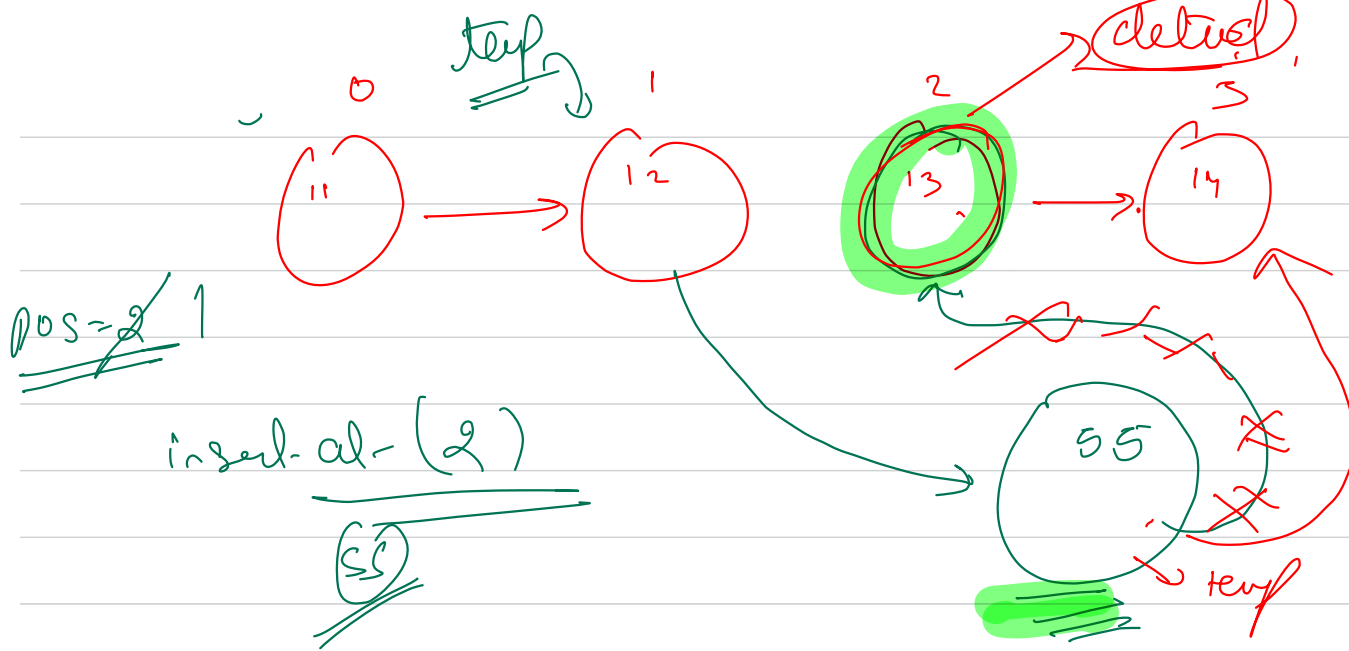
temp = temp->next

new-node->next = temp->next  
temp->next = new-node

→ fetch the node previous to the idn where we want to insert



→ mark the next of new node = to the node present at the idn currently (Temp.next)

→ update temp.next = new\_node.



temp.next = temp.next.next  
temp.next = null

delete from head  
delete free tail  
delete in between

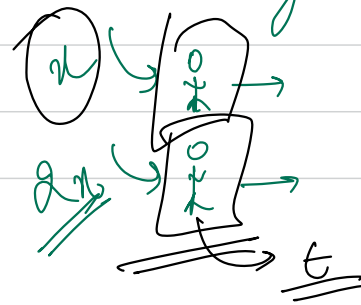


Q2 find mid node of a ll without using length function & by iterativ only once on the ll.

$$TC \rightarrow \underline{\underline{O(n)}}$$

$$SC \rightarrow \underline{\underline{O(1)}}$$

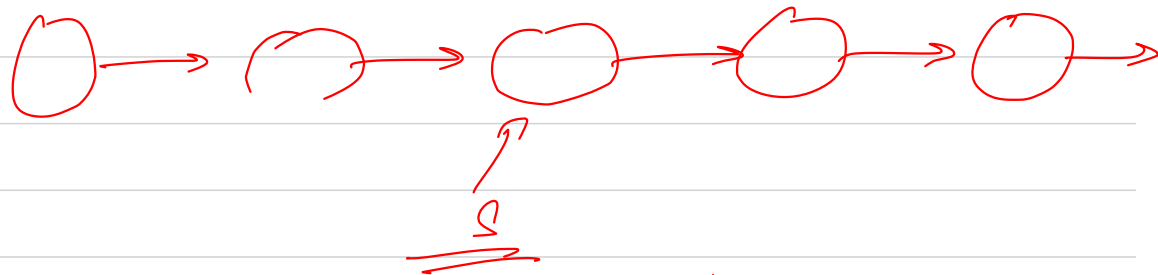
→ length of ll → L



$d = \underline{\underline{S \times t}} \rightarrow \cancel{x} \times \frac{L}{\cancel{2x}} \rightarrow \underline{\underline{\frac{L}{2}}}$

$\textcircled{t} = \frac{L}{2x}$

Q



Q You're given 2 linked lists, of diff size.

Add mem & return a new ll.

single pass



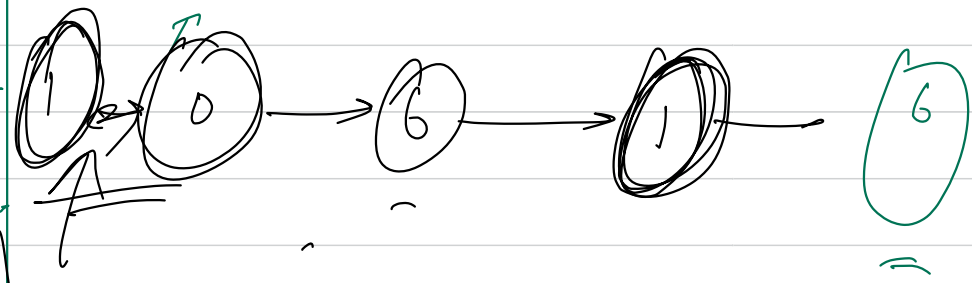
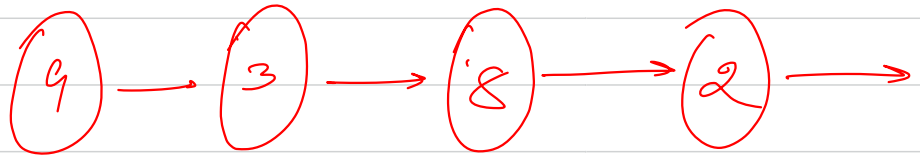
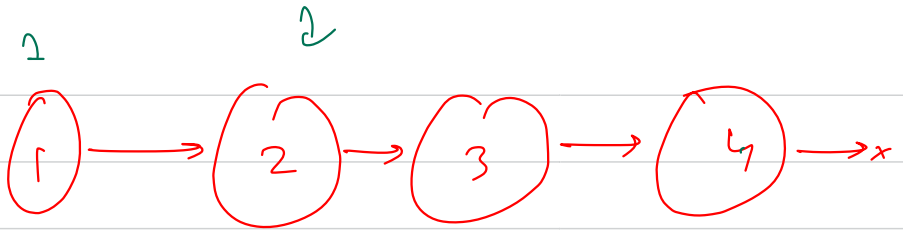
$$\begin{array}{r}
 \text{C} \\
 1 \ 2 \ 3 \ 4 \\
 2 \ 8 \ 1 \ 3 \\
 \hline
 9 \ 0 \ 4 \ 7 \\
 \hline
 \hline
 \end{array}$$





# Recursive

↓  
get the  
prev  
state  
from which  
we are  
calling



1, 4, 0

carry

