Assignment No.-2

Title: - Dictionary using BST

Problem Statement:- Dictionary stores keywords & its meanings. Provide facility for adding new keywords, deleting keywords, updating values of any entry. Provide facility to display whole data sorted in ascending/ Descending order. Also find how many maximum comparisons may require for finding any keyword. Use Binary Search Tree for implementation.

Learning Objectives:-

1. To understand the binary search tree data structure.
2. To understand the various operations on binary search tree.
3. To implement dictionary using binary search tree and object oriented concepts.

Learning Outcomes: - After successfully completing this assignment you should be able to:

1. Define the classes required to implement dictionary using binary search tree and object oriented features in C++.
2. Analyze and trace the working of recursive functions used in implementation.
3. Understand and implement various operations on BST such as Insert, Search, Update, Delete etc.

Software Tools Used:- linux GCC compiler, Eclipse Editor on 64-bit Ubuntu 14.04 operating system.

Hardware Required:- PC configuration- Intel Core2duo 2.93 GHZ CPU, 4GB RAM, 320GB HDD, 17" Monitor, Keyboard and Mouse.

Theory:- Binary search tree is the binary tree which satisfies the following condition:

d) All elements should be distinct.
e) For any node its left subtree should contain the less values than the node and its right subtree should contain the larger values than the node.
f) All subtree should be BST.

Searching time in the BST is O(log n).

e.g.

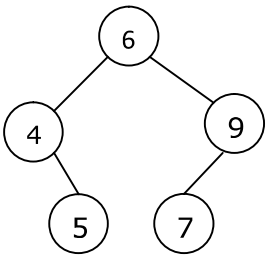

Fig.(a) Valid BST

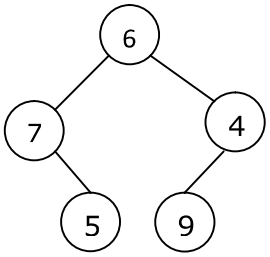

Fig.(b) Invalid BST

Representation Used for BST:
In this assignment linked representation of BST is used. For that we need to define node structure to be linked in BST structure. The BST node structure has the following fields.

*Left	Data	*Right

The node contains 3 fields as follows:

d) ***Left** holds the address of left child/subtree if present otherwise NULL.
e) **Data** field hold data of node may contain single value or record.
f) ***Right** holds the address of right child/subtree if present otherwise NULL.

Algorithm:-The various algorithms are as follows:

A. Create ()
1. Start
2. Read the value.
3. If value is already present in BST
   Do not insert.
   Otherwise
   Insert(Value);
4. Go to step 2 if you want to add anymore values.
5. Stop.

B. Insert(Val)
1. Start
2. Create the new node P initialized with value given.
3. If Root is empty
   Root=P; Go to
4. Initialize T to Root.
5. If Val< T->Val
   If left child is empty
   T->left=P; Go to 6
   Otherwise
   Descend the tree in left direction by one level.go to5
   Otherwise
   If right child is empty

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                T->right=P; Go to 6
            Otherwise
                Descend the tree in right direction by one level.goto5
    6. Stop

C. Search(Val)
    1. Start
    2. Initialize T with Root
    3. If T is NULL go to 7
    4. If T->data==Val
            return T go to 7
    5. If Val<T->data
            Descend the tree in left direction by one level go to 4
    6. If Val>T->data
            Descend the tree in right direction by one level go to 4
    7. Stop

D. Delete(*T)
    1. Start
    2. IF T is NULL go to  7
    3. If T is leaf node then
            If T is Root then
                Root=NULL  go to 7
            Otherwise
                P <- Parent(T)
                If T is left child of P
                    P->left = NULL
                Otherwise
                    P->right = NULL
            Release T go to 7
    4. If T is One degree node then
            If T is Root then
                If left child of T present
                    Descend the tree in left by one level
                Otherwise
                    Descend the tree in right by one level
                Release T go to 7
            Otherwise
                P <- Parent(T)
                If T is left child of P then
                    Connect the child of T to P as left child
                Otherwise
                    Connect the child of T to P as right child
            Release T go to 7
    5. Locate the Inorder successor(S) of T
    6. Replace the value of Y by value of S and call Delete(S)
    7. Stop
```

Program Code with Sample Output : –

Analysis of Algorithm : –

1. Create() = O(n)
2. Insert() = O(log n)
3. Search() = O(log n)
4. Inorder() = O(n)
5. Delete() = O(n)

Applications : –

1. BST can be used to implement Dictionary.
2. Used in *many* search applications where data is constantly entering/leaving, such as the map and set objects in many languages' libraries.
3. It is used to implement multilevel indexing in DATABASE.
4. It is used to implement searching Algorithm.

Conclusion : –

References : –

1. Horowitz, Sahani, Dinesh Mehata, –Fundamentals of Data Structures in C++, Galgotia Publisher, ISBN: 8175152788, 9788175152786.
2. Peter Brass, –Advanced Data Structures‖, Cambridge University Press, ISBN: 978-1-107-43982-5.

Date of Completion : –

FAQ's : –

1. What are advantages of dynamic memory allocation over static memory allocation?
2. What is necessity of friend class in oop?
3. If tree is complete binary search tree having total number of nodes are 1000 then what is the height of it?
4. What are the maximum no of comparison required to search value in BST of n nodes?
5. What are the applications of BST?