

# DYNAMIC PROGRAMMING

CSE [AIML], 6<sup>TH</sup> SEMESTER

DESIGN & ANALYSIS OF ALGORITHMS

# DYNAMIC PROGRAMMING

We follow sequence of 4 steps:

1. Characterize the structure of an optimal solution.
2. Recursively define the value of an optimal solution
3. Compute the value of an optimal solution
4. Construct an optimal solution from computed information.

# MATRIX CHAIN MULTIPLICATION

• We state the Matrix Chain Multiplication as follows:

Given a chain  $\{A_1, A_2, \dots, A_n\}$  of matrices, where for  $i = 1, 2, \dots, n$ ,  $A_i$  has dimension  $p_{i-1} \times p_i$ , fully parenthesized the product  $A_1, A_2, \dots, A_n$  in a way that minimized the number of scalar multiplication.

# MATRIX CHAIN MULTIPLICATION

1. We can build an optimal solution to an instance of the problem by splitting the problem into two sub problems.
2. Finding optimal solutions to sub problem instances.
3. Combining these optimal sub problems solutions.

# MATRIX CHAIN MULTIPLICATION

A recursive solution (Formula):

$$m[i, j] = \begin{cases} 0, & \text{if } i = j \\ \min\{m[i, k] + m[k + 1, j] + p_{i-1}p_kp_j\} & \text{if } i < j \\ & i \leq k < j \end{cases}$$

# MATRIX CHAIN MULTIPLICATION EXAMPLE

A1    A2    A3    A4

3x2        2x4        4x2        2x5

$$m[1,2] = \min(k=1)\{m[1,1] + m[2,2] + 3 \times 2 \times 4\} = 24$$

$$m[2,3] = \min(k=2)\{m[2,2] + m[3,3] + 2 \times 4 \times 2\} = 16$$

$$m[3,4] = \min(k=3)\{m[3,3] + m[4,4] + 4 \times 2 \times 5\} = 40$$

# MATRIX CHAIN MULTIPLICATION

## EXAMPLE CONTINUES...

|     |     |     |     |
|-----|-----|-----|-----|
| A1  | A2  | A3  | A4  |
| 3x2 | 2x4 | 4x2 | 2x5 |

$$\mathbf{m[1,3]} = \min(k = 1, k = 2 \{m[1,1] + m[2,3] + 3 \times 2 \times 2, \\ m[1,2] + m[3,3] + 3 \times 4 \times 2\}) = \min\{28, 48\} = 28$$

$$\mathbf{m[2,4]} = \min(k = 2, k = 3 \{m[2,2] + m[3,4] + 2 \times 4 \times 5, \\ m[2,3] + m[4,4] + 2 \times 2 \times 5\}) = \min\{80, 36\} = 36$$

$$\mathbf{m[1,4]} = \min(k = 1, k = 2, k = 3 \{m[1,1] + m[2,4] + \\ 3 \times 2 \times 5, m[1,2] + m[3,4] + 3 \times 4 \times 5, m[1,3] + m[4,4] + 3 \times 2 \times 5\}) \\ = \min\{86, 124, 58\} = 58$$

# MATRIX CHAIN MULTIPLICATION

## EXAMPLE CONTINUES...

| m |    | k  |    |
|---|----|----|----|
| 0 | 24 | 28 | 58 |
|   | 0  | 16 | 36 |
|   |    | 0  | 40 |
|   |    |    | 0  |

|  | 1 | 1 | 3 |
|--|---|---|---|
|  |   | 2 | 3 |
|  |   |   | 3 |
|  |   |   |   |

The minimum number of scalar multiplication to multiply the 4 matrices is  $m[1,4] = 58$



# TIME & SPACE COMPLEXITY

The loops are nested three times and each loop indexed at most  $n-1$  values. This yields a running time of  $O(n^3)$ .

The algorithm requires  $\theta(n^2)$  space to store  $m$  and  $k$  tables.