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## Assignment 4

### 1. Distributed Deadlock Detection Simulation

→ Fragments :-

- $S_1 : P_1 \rightarrow P_2, P_3 \rightarrow P_4$
- $S_2 : P_2 \rightarrow P_5, P_5 \rightarrow P_6$
- $S_3 : P_6 \rightarrow P_1$

(a) Global wait for graph (combined)  
 $P_1 \rightarrow P_2 \rightarrow P_5 \rightarrow P_6 \rightarrow P_1$  (cycle)  
Also,  $P_3 \rightarrow P_4$  (separate)

(b) deadlock?

→ Yes, Processes involved in deadlock cycle :  $P_1, P_2, P_5, P_6$

### c suggested distributed algorithms

→ use the Chandy → Mishra → Hass edge-chasing (probe) Algorithm for distributed deadlock detection. Each site sends probes along wait-for edges to detect cycles without centralized graph assembly.

### 2. Distributed file system performances.

Given: local = 5ms, remote = 75ms  
prob(remote) = 0.3

(a) Expected access time:-

$$\begin{aligned} E &= 0.7 \times 5\text{ms} + 0.3 \times 75\text{ms} \\ &= 3.5 + 22.5 \\ &= 26\text{ms} \end{aligned}$$



(b) coaching strategy:-  
client side read cache with LRU + TTL  
based validation

→ Justification:- Frequently-read remote files will be served locally reducing remote access (0.3 fraction), LRU evicts less used items. TTL keeps staleness bounded. Improves average latency while keeping consistency manageable

3 check pointing mix to meet  $RPO = 1s$

→ Given

full = 200ms, incremental = 80ms,  $RPO = 1s$

(a) proposed mix (over 10s):-

- Take one full checkpoint every 10s (at  $t=0$  in period)
- Take incremental checkpoints every 1s (at  $t = 1, 2, \dots, 9$ )

$$\begin{aligned} \text{Total overhead (per 10s)} &: 1 \times 200\text{ms} \\ &+ 9 \times 80\text{ms} = 200 + 720 \\ &= 920\text{ms} \end{aligned}$$

(b) reasoning:-

- with incremental 1s, the maximum work lost on failure  $\leq 1s \Rightarrow$  meets RPO.
- full once per 10s bound recovery time
- This mix minimizes full checkpoint cost while keeping incremental frequency high enough to meet RPO



## 9 case study - Global e-commerce platform -

### (a) distributed scheduling challenges for flash sales :-

- massive, sudden spike in request
  - geographic distribution & latency - data locality matters for latency and inventory correctness
  - heterogeneous nodes
  - stateful server
- suggest algorithms for load balancing
- hybrid approach

### (b) Fault-tolerance strategy (RTO & RPO) :-

- Active :- active multi-region deployment : service run both currently in multiple regions so failure is seamless ( $RTO \approx \text{near-zero at service level}$ )
  - Data strategy
- critical transactional data :- use synchronous replication within the region to guarantee consistency and low RPO, cross region replication can be asynchronous but with frequent replication to keep RPO small.
- catalog / less critical data :- use multiple region eventual consistency with frequent async replication and cache



- operational measures :- chaos testing, backups, runbook, and automata scaling to minimize RPO, use snapshot + incremental backup to bend RPO

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