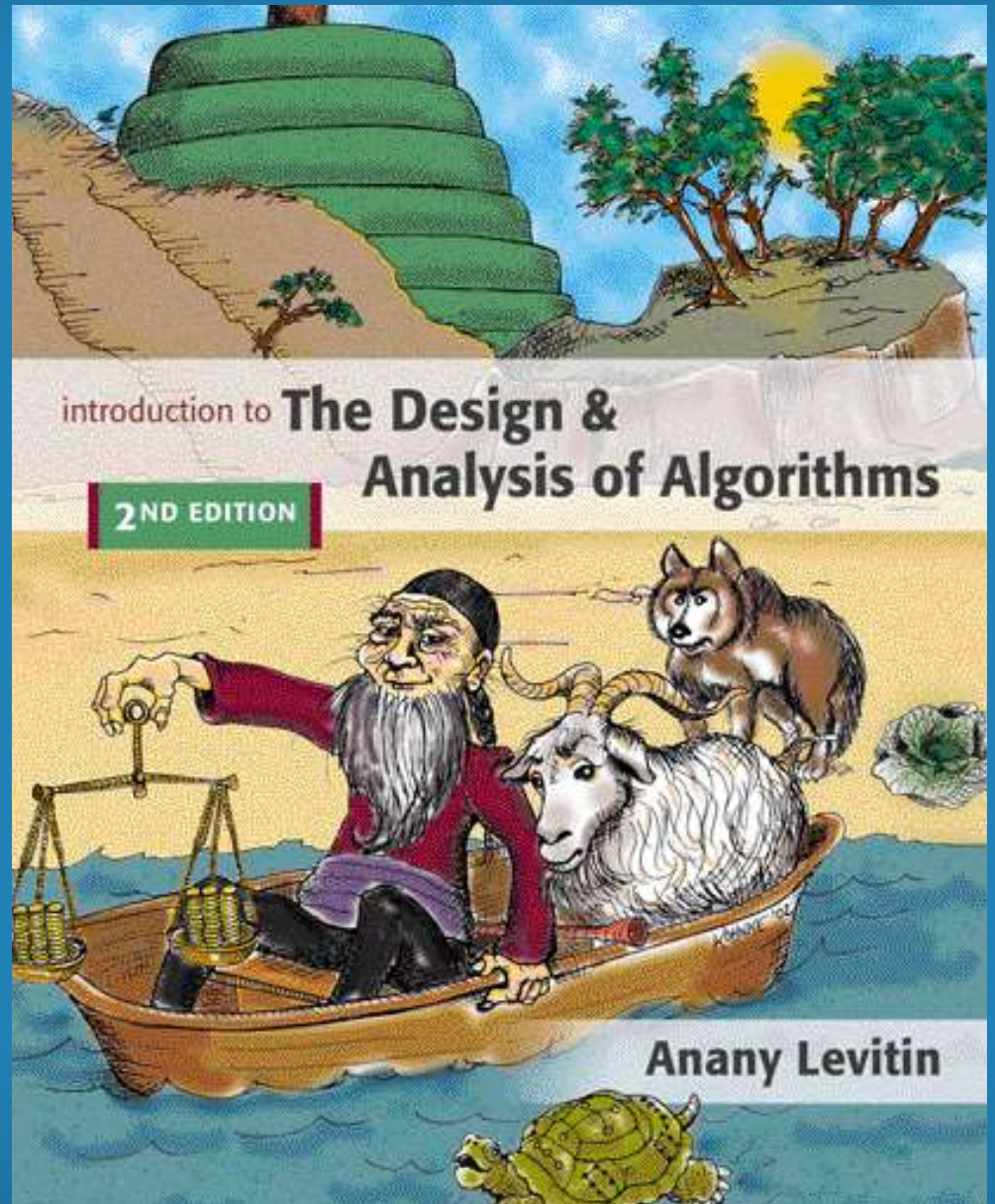


# Chapter 11

## Limitations of Algorithm Power

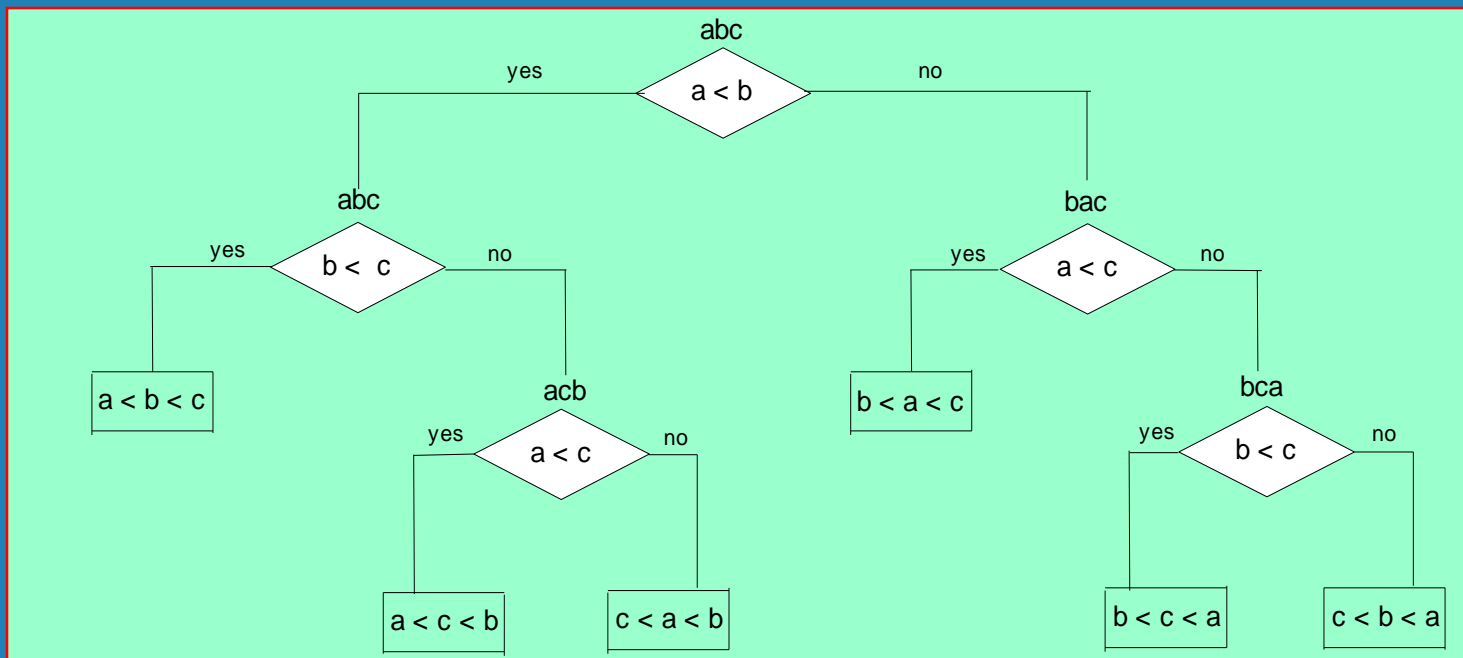


# Decision Trees

**Decision tree** — a convenient model of algorithms involving comparisons in which:

- internal nodes represent comparisons
- leaves represent outcomes (or input cases)

## Decision tree for 3-element insertion sort



# Decision Trees and Sorting Algorithms

- Any comparison-based sorting algorithm can be represented by a decision tree (for each fixed  $n$ )
  - Number of leaves (outcomes)  $\geq n!$
  - Height of binary tree with  $n!$  leaves  $\geq \lceil \log_2 n! \rceil$
  - Minimum number of comparisons in the worst case  $\geq \lceil \log_2 n! \rceil$  for any comparison-based sorting algorithm, since the longest path represents the worst case and its length is the height
  - $\lceil \log_2 n! \rceil \approx n \log_2 n$  (by Sterling approximation)
  - This lower bound is tight (mergesort or heapsort)
- Ex. Prove that 5 (or 7) comparisons are necessary and sufficient for sorting 4 keys (or 5 keys, respectively).

# Class $P$



$P$ : the class of decision problems that are solvable in  $O(p(n))$  time, where  $p(n)$  is a polynomial of problem's input size  $n$

## Examples:

- ⌚ searching
- ⌚ element uniqueness
- ⌚ graph connectivity
- ⌚ graph acyclicity





# Class *NP*



*NP* (*nondeterministic polynomial*): class of decision problems whose proposed solutions can be verified in polynomial time = solvable by a *nondeterministic polynomial algorithm*

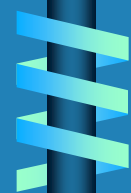
A *nondeterministic polynomial algorithm* is an abstract two-stage procedure that:

- ⌚ generates a **solution** of the problem (on some input) by **guessing**
- ⌚ checks whether this solution is correct in polynomial time

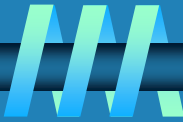
By definition, it solves the problem if it's capable of generating and verifying a solution on one of its tries

Why this definition?

- ⌚ led to development of the rich theory called “computational complexity”



# Backtracking



## ❧ n- Queens Problem

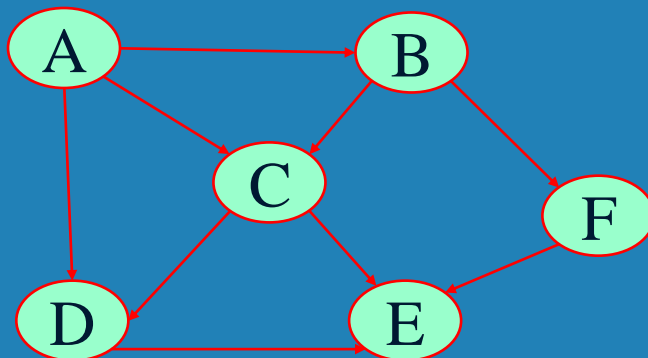
❧  $n=1 \rightarrow$  trivial solution

❧  $n=2$  &  $n=3 \rightarrow$  no solution

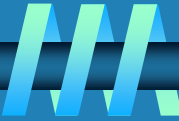
## ❧ Hamiltonian Circuit Problem

❧ Starts and ends with same vertex

❧ Visits exactly once



# Continued.....

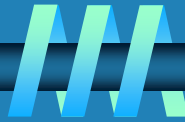


## ⌚ Subset-Sum Problem

⌚ Find a subset of a given set  $S = \{1, 2, 5, 6, 8\}$  with sum  $d=9$



# Branch-and-Bound



## Assignment Problem –(lower bound)

	Job1	Job2	Job3	Job4
Person a	9	2	7	8
Person b	6	4	3	7
Person c	5	8	1	8
Person d	7	6	9	4

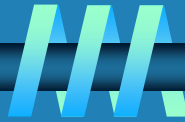
## Knapsack Problem -(Upper bound)

Item	Weight	Value	Value/Weight
1	4	40	10
2	7	42	6
3	5	25	5
4	3	12	4

**W=10**



# Continued.....



⌚  $ub = v + (W-w)(v_i+1/w_i+1)$

⌚ **Travelling Salesman Problem (lower bound)**

