

<u>CSE 3124</u>: <u>Computer Networks</u> [2 1 0 3]

Module 1: Computer Networks & the Internet

our goal:

- get "feel" and terminology
- more depth, detail
- *later* in course approach:
 - use Internet as example

overview:

- what's the Internet?
- what's a protocol?
- Network edge; Hosts, Access net, Physical media
- Network core: Packet/Circuit switching, Internet structure
- Performance: Loss, Delay, Throughput
- Protocol layers, Service Models
- History

Introduction 1-2

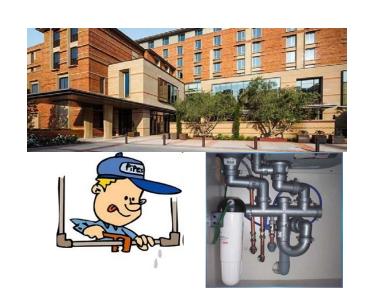
Computer Network

A **computer network** is a group of interconnected devices (such as computers, servers, and routers) that communicate and share resources using wired or wireless connections.

These networks allow data exchange, collaboration, and access to shared resources like printers, files, and the internet.

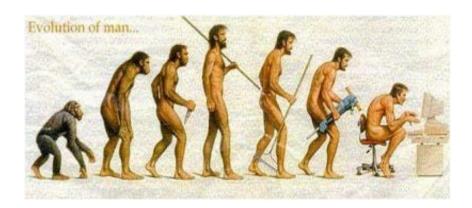
Networking = "Plumbing"

- □ Networking is the "plumbing" of computing
- □ Almost all areas of computing are network-based.
 - Distributed computing
 - Big Data
 - Cloud Computing
 - Internet of Things
 - Smart Cities
- □ Networking is the backbone of computing.



Internet Age

- Distributed Computing
- Cloud Computing
- ☐ Mobile Computing ? Smart Phones
- Streaming Video ? YouTube
- Social Networking ? FaceBook
- □ Big Data
- Machine Learning ? Artificial Intelligence
- Online Shopping 🛭 Amazon, eBay
- Most fields today Education, Health, Environment are advancing simply because of advances in networking.



Current Hot Topics in Networking



- 1. Internet of Things (IoT)
- 2. Cybersecurity
- 3. Cloud Computing
- 4. Software Defined Networking
- Wireless Networking
- 6. Streaming Media

Trend: Smart Everything



Smart Watch



Smart TV



Smart Car



Smart Health



Smart Home



Smart Kegs



Smart Space



Smart Industries



Smart Cities

Trend: Smart to Intelligent



Trend: Security & Cyber Warfare

- Security of computers, companies, smart grid, and nations
- □ Nation States are penetrating other nations' computers 5th domain of warfare (after land, sea, air, space)
- □ In 2010, the US set up US Cyber Command
- UK, China, Russia, Israel, and North Korea have similar centers
- Many cyber wars: North Korea vs. USA, Israel vs. Syria, South Korea vs.

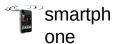
Old New

What's the Internet: "nuts and bolts" view









 millions of connected
 computing devices:

- running *network*

– running *network* – *apps*

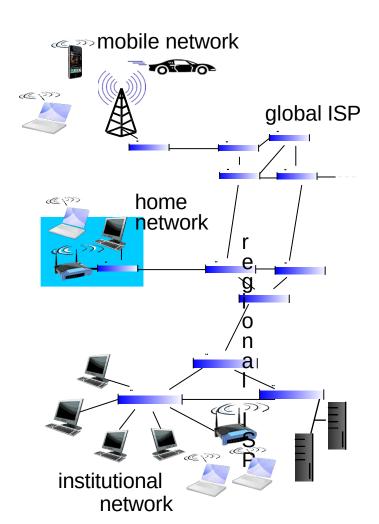


wireless^{to} communication links

- fiber, copper, radio, satellite
- transmission rate:
 bandwidth



- Packet switches: forward packets (chunks of data)
 - routers and switches



Introduction 1-10

"Fun" internet appliances



IP picture frame http://www.ceiva.com/



Web-enabled toaster + weather forecaster



Tweet-a-watt: monitor energy use



Internet refrigerator



Slingbox: watch, control cable TV remotely



Internet phones

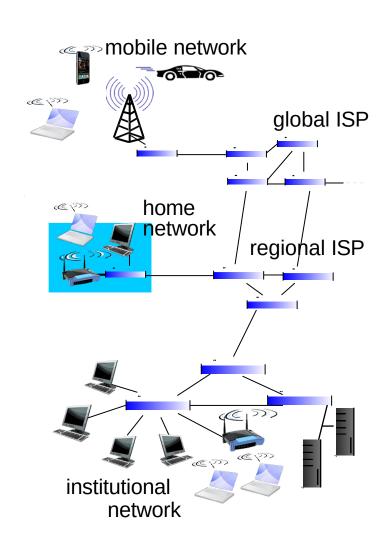
Introduction 1-11

What's the Internet: "nuts and bolts" view

- Internet: "network of networks"
 - Interconnected ISPs
- protocols control sending, receiving of msgs
 - e.g., ТСР, IP, НТТР, Skype, 802.11

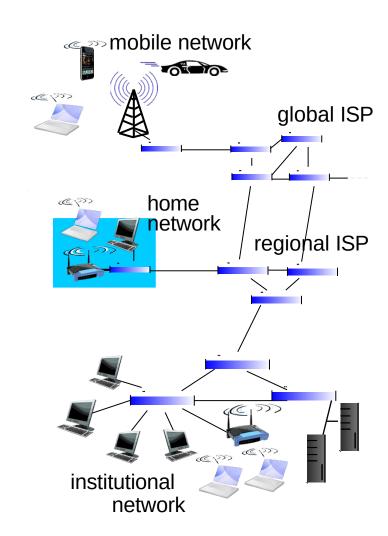
Internet standards

- RFC: Request for comments
 - IETF: Internet Engineering Task
 Force



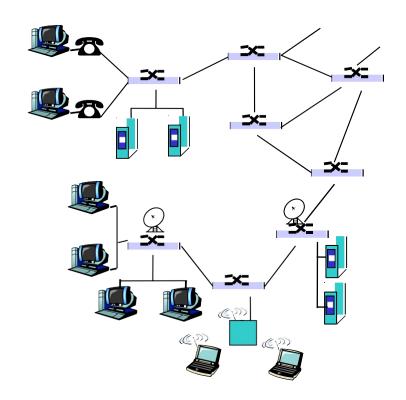
What's the Internet: a service view

- Infrastructure that provides services to applications:
 - Web, VoIP, email, games, e-commerce, social nets, ...
- provides programming interface to apps



A closer look at network structure:

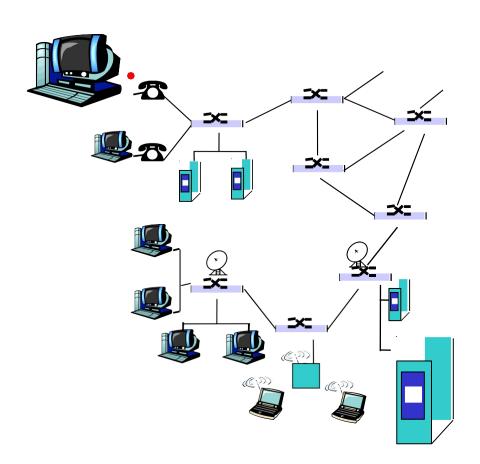
- network edge: applications and hosts
- θ network core:
 - u routers
 - network of networks
- θ access networks, physical media: communication links



The network edge:

end systems (hosts):

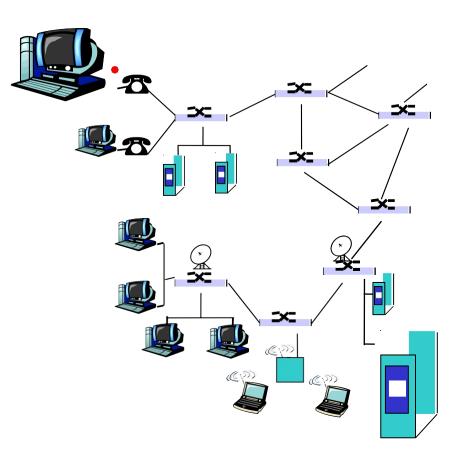
- μ run application programs
- μ e.g. Web, email
- μ at "edge of network"



The network edge:

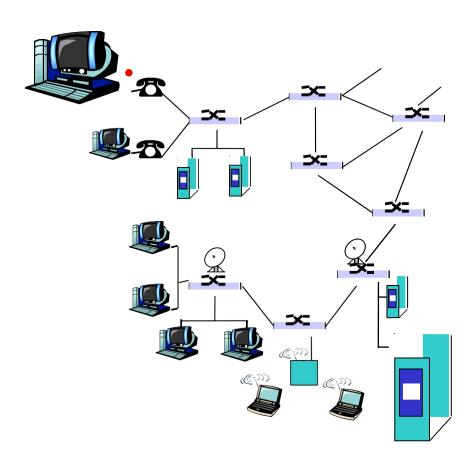
θ client/server model

- u client host requests, receives service from always-on server
- µ e.g. Web browser/server; email client/server
- θ why such a popular model?



The network edge:

- ⁹ peer-peer model:
 - minimal (or no) use of dedicated servers
 - μ e.g. Gnutella, KaZaA
 - u SETI@home?



What's a protocol?

<u>human protocols:</u>

- θ "what's the time?"
- θ "I have a question"
- θ introductions

<u>network protocols:</u>

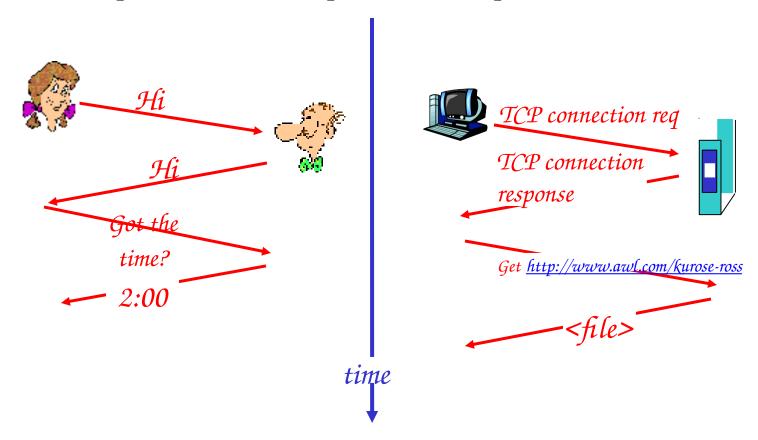
- machines rather than humans
- all communication activity in Internet governed by protocols

protocols define format, order of msgs sent and received among network entities, and actions taken on msg

transmission, receipt

What's a protocol?

a human protocol and a computer network protocol:



Q: Why are protocols so important?

Internet Services Models

⁶ Connection-oriented service

⁶ Connectionless service

- ⁰ Applications
 - 4 FTP, Internet Phone, Web, Internet radio, email

Connection-oriented service

- 6 Goal: data transfer between end systems
- handshaking: setup (prepare for) data transfer ahead of time
- ⁰ TCP Transmission Control Protocol
 - u Internet's connection-oriented service
 - u reliable, in-order byte-stream data transfer
 - loss: acknowledgements and retransmissions
 - μ flow control:
 - sender won't overwhelm receiver
 - u congestion control:
 - senders "slow down sending rate" when network

Connectionless service

Goal: data transfer between end systems

- μ same as before!
- θ **UDP** User Datagram Protocol [RFC 768]:
 - u connectionless
 - unreliable data transfer
 - uno flow control
 - u no congestion control
- ⁰ What's it good for?

A Comparison

App's using TCP:

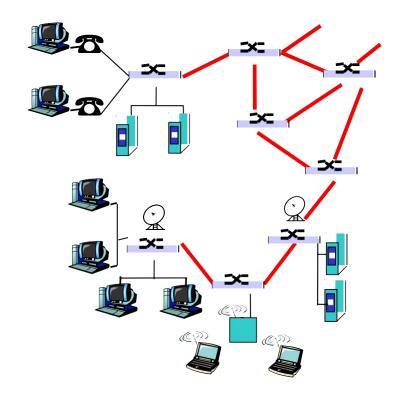
HTTP (Web), FTP (file transfer), Telnet (remote login), SMTP (email)

App's using UDP:

• streaming media, teleconferencing, DNS, Internet telephony

The Network Core

- nesh of interconnected routers
- θ <u>the</u> fundamental question: how is data transferred through net?
 - u circuit switching: dedicated circuit per call: telephone net
 - packet-switching: data sent thru net in discrete "chunks"

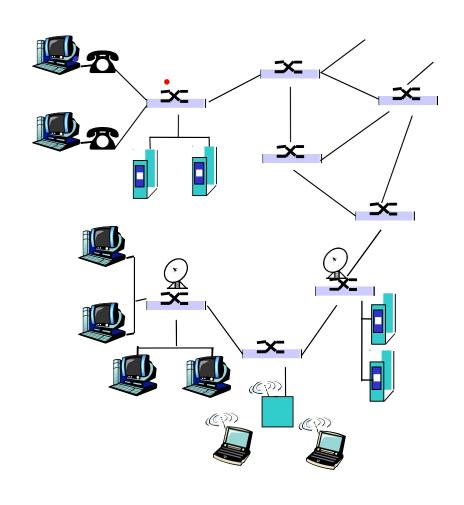


Network Core: Circuit

Switching

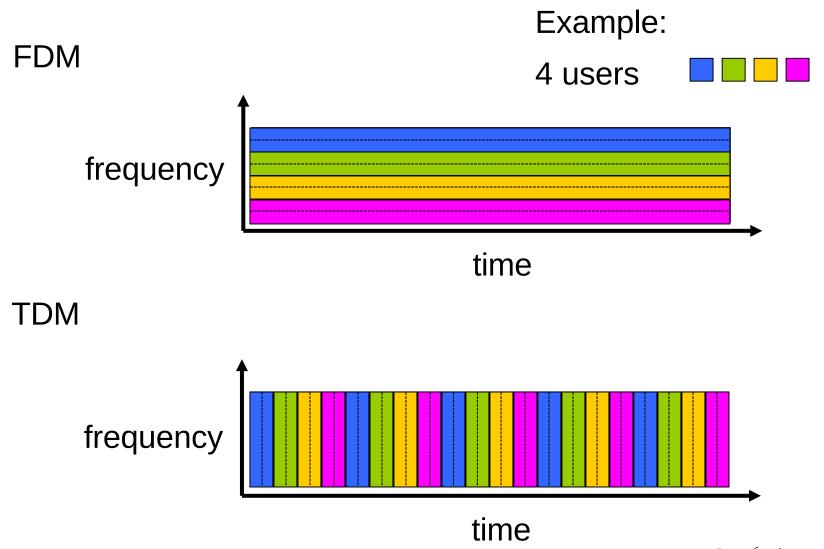
End-end resources reserved for "call"

- θ link bandwidth, switch capacity
- θ dedicated resources: no sharing
- circuit-like (guaranteed)performance
- θ call setup required



θ must divide link bw into

Circuit Switching: FDM and TDM



Circuit Switching: FDM and TDM

TDM (Time Division Multiplexing) and **FDM (Frequency Division Multiplexing)** both are multiplexing techniques where TDM is used in both analogue and digital signals.

Where as FDM used only in analogue signals.

Network Core: Packet Switching

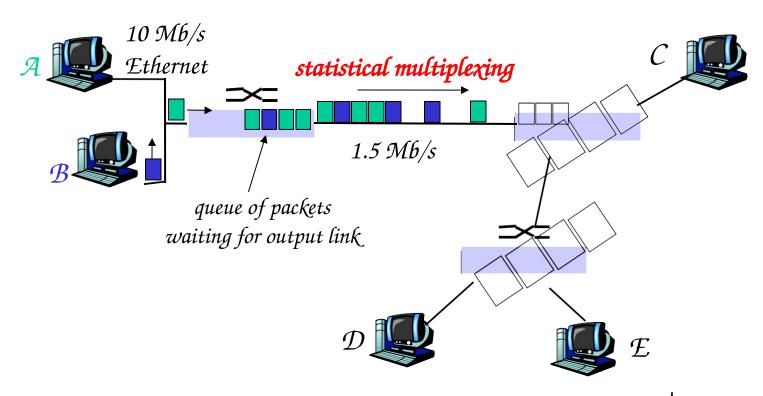
each end-end data stream divided into packets

- θ user A, B packets share network resources
- each packet uses full link bandwidth
- h resources used as needed

resource contention:

- o aggregate resource demand can exceed amount available
 - what happens if bandwidth is not available?
- 6 congestion: packets queue, wait for link use
- b store and forward: packets move one hop at a time

Packet Switching: Statistical Multiplexing



Sequence of A & B packets does not have fixed pattern statistical multiplexing.

Packet Switching: Statistical Multiplexing

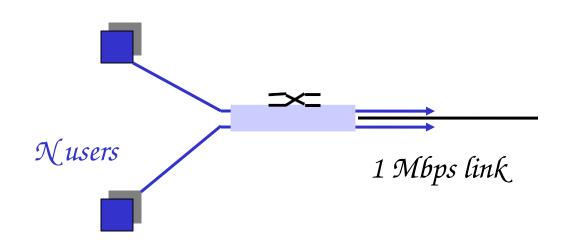
In statistical multiplexing, packets from sources A and B are sent based on demand.

The sequence of packets (A, B, A, B, etc.) does not follow a fixed pattern—this is what makes it statistical.

Packet switching versus circuit switching

Packet switching allows more users to use network!

- θ 1 Mb/s link
- θ each user:
 - μ 100 kb/s when "active"
 - μ active 10% of time
- θ circuit-switching:
 - μ 10 users
- θ packet switching:
 - with 35 users, probability > 10 active less than .0004

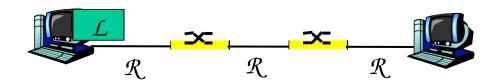


Packet switching versus circuit switching

Is packet switching a "slam dunk winner?"

- θ Great for bursty data
 - μ resource sharing
 - u simpler, no call setup
- ¹⁰ Excessive congestion: packet delay and loss
 - µ protocols needed for reliable data transfer, congestion control
- θ Circuit Switching = Guaranteed behavior
 - µ good for which apps?

Packet-switching: store-and-forward



- Takes L/R seconds to transmit (push out) packet of L bits on to link or R bps
- Entire packet must arrive at router before it can be transmitted on next link: store and forward
- θ delay = 3L/R

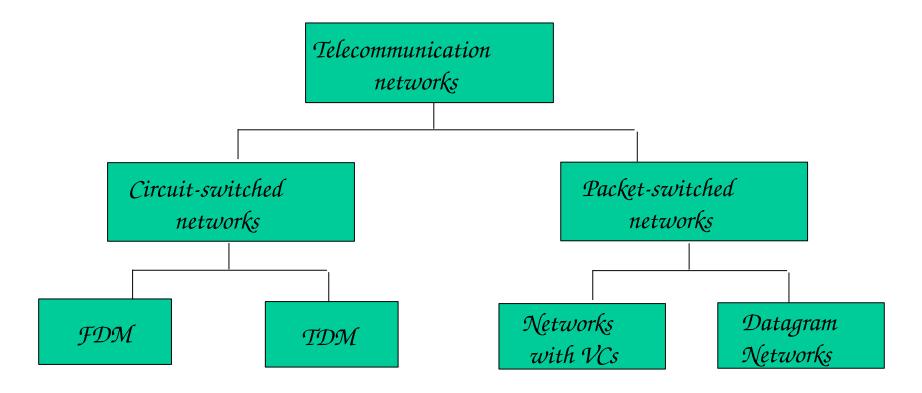
Example:

- θ $\mathcal{L} = 7.5 \text{ Mbits}$
- θ R = 1.5 Mbps
- θ delay = 15 sec

Packet-switched networks: forwarding

- 6 Goal: move packets through routers from source to destination
 - we'll study several path selection (i.e. routing) algorithms (chapter 4)
- θ datagram network:
 - u destination address in packet determines next hop
 - μ routes may change during session
 - u analogy: driving, asking directions
- θ virtual circuit network:
 - u each packet carries tag (virtual circuit ID), tag determines next hop
 - fixed path determined at call setup time, remains fixed thru call
 - u routers maintain per-call state

Network Taxonomy



- *Datagram network is <u>not</u> either connection-oriented or connectionless.
- Internet provides both connection-oriented (TCP) and connectionless services (UDP) to apps.

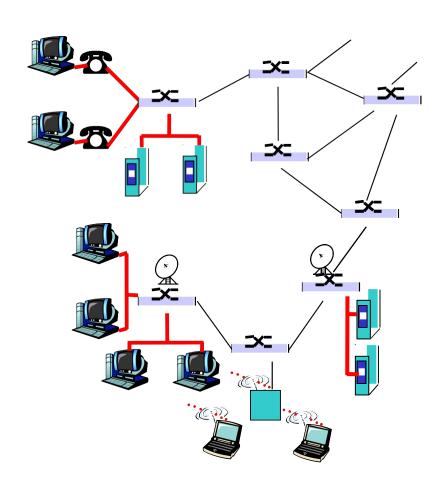
Access networks and physical media

Q: How to connect end systems to edge router?

- 6 residential access nets
- institutional access networks (school, company)
- ⁰ mobile access networks

Keep in mind:

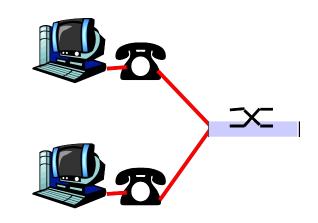
- bandwidth (bits per second) of access network?
- θ shared or dedicated?



Residential access: point to point access

θ Dialup via modem

- up to 56Kbps direct access to router (often less)
- " Can't surf and phone at same time: can't be "always on"



6 <u>ADSL:</u> asymmetric digital subscriber line

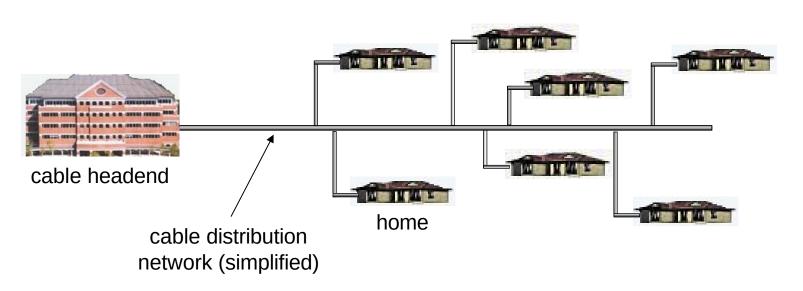
- up to 1 Mbps upstream (today typically < 256 kbps)
- up to 8 Mbps downstream (today typically < 1 Mbps)
- u FDM: 50 kHz 1 MHz for

Residential access: cable modems

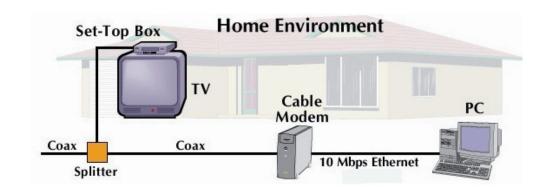
- θ HFC: hybrid fiber coax
 - upstream asymmetric: up to 30Mbps downstream, 2 Mbps
- network of cable and fiber attaches homes to ISP router
 - µ homes share access to router
- 6 deployment: available via cable TV companies

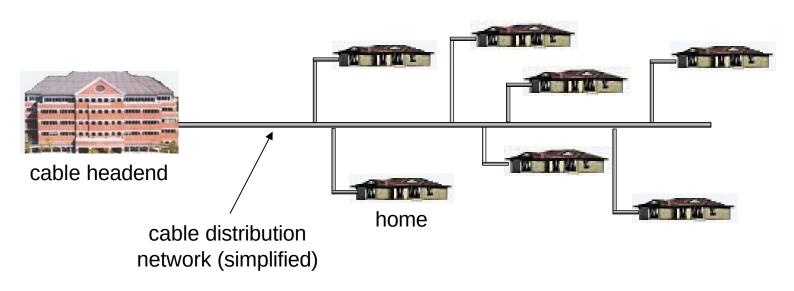
Cable Network Architecture: Overview

Typically 500 to 5,000 homes



Cable Network Architecture: Overview



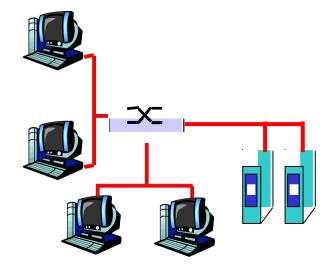


Company access: local area networks

6 company/univ local area network (LAN) connects end system to edge router

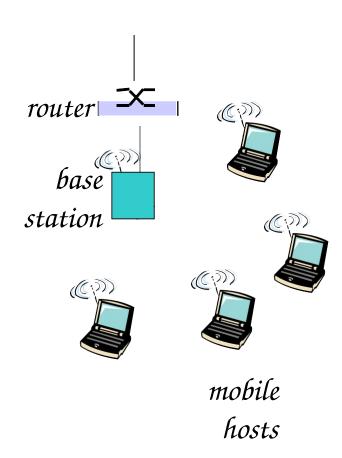
θ Ethernet:

- µ shared or dedicated link connects end system and router
- µ 10 Mbps, 100Mbps,
 Gigabit Ethernet
- θ LANs: chapter 5



Wireless access networks

- θ shared wireless access network connects end system to router
 - u via base station aka "access point"
- θ wireless LANs:
 - µ 802.11b (WiFi): 11 Мbps
- θ wider-area wireless access
 - μ provided by telco operator
 - μ 3G ~ 384 kbps
 - Will it happen??
 - и WAP/GPRS in Europe



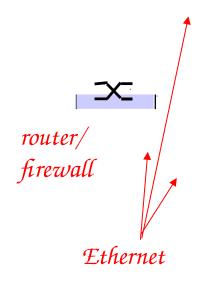
Home networks

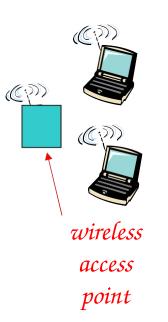
Typical home network components:

- ADSL or cable modem
- θ router/firewall/<mark>NAT</mark>
- 0 Ethernet
- θ wireless access point

to/from cable headend

cable modem





wireless laptops

Physical Media

- θ Bit: propagates between transmitter/rcvr pairs
- physical link: what lies between transmitter & receiver
- θ guided media:
 - signals propagate in solid media: copper, fiber, coax
- θ unguided media:
 - u signals propagate freely, e.g., radio

Twisted Pair (TP)

- two insulated copper wires
 - Category 3: traditional phone wires, 10 Mbps Ethernet
 - u Category 5: 100Mbps
 Ethernet



Physical Media: coax, fiber

Coaxial cable:

- two concentric copper conductors
- θ bidirectional
- baseband:
 - u single channel on cable
 - μ legacy Ethernet
- θ broadband: channel on cable HFC
 μ



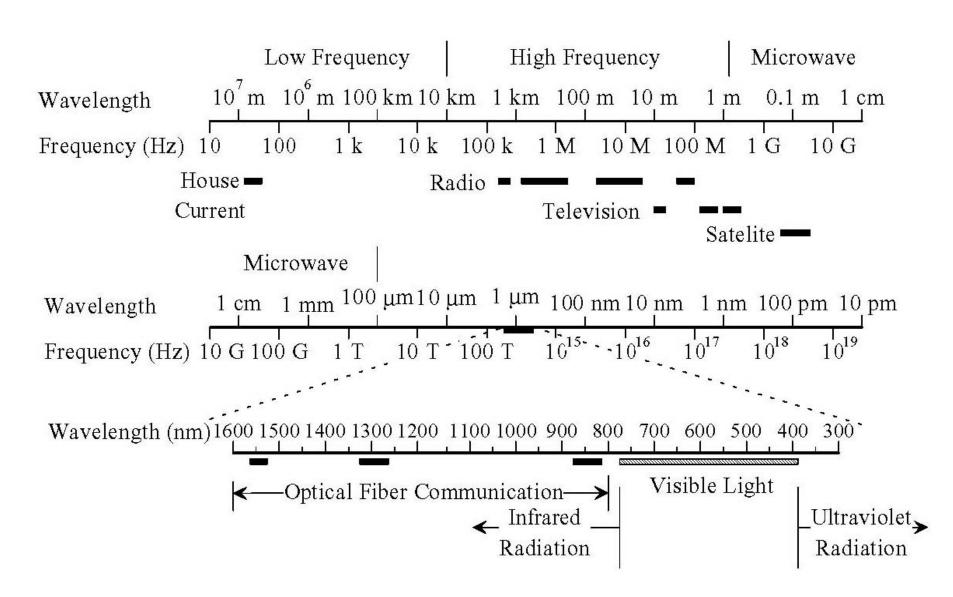
Fiber optic cable:

- θ glass fiber carrying light pulses, each pulse a bit
- θ high-speed operation:
 - high-speed point-to-point transmission (e.g., 5 Gps)
- far apart ; immune to

Physical Media: coax, fiber

Feature	Fiber Optic Cable	Copper Coaxial Cable
Speed / Bandwidth	Much higher bandwidth (10 Gbps and beyond)	Lower bandwidth; suitable for limited-speed usage
Distance	Transmits over longer distances without signal loss	Signal degrades faster, needs frequent amplifiers
Interference	Immune to electromagnetic interference	Susceptible to EMI and radio frequency interference
Signal Quality	Very low loss and high-quality signal	More signal attenuation over distance
Security	Difficult to tap without detection (more secure)	Easier to tap into (less secure)
Size and Weight	Lighter and thinner	Thicker and heavier
Cost (initial)	Higher installation cost	Cheaper initially

Electromagnetic Spectrum



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Physical media: radio

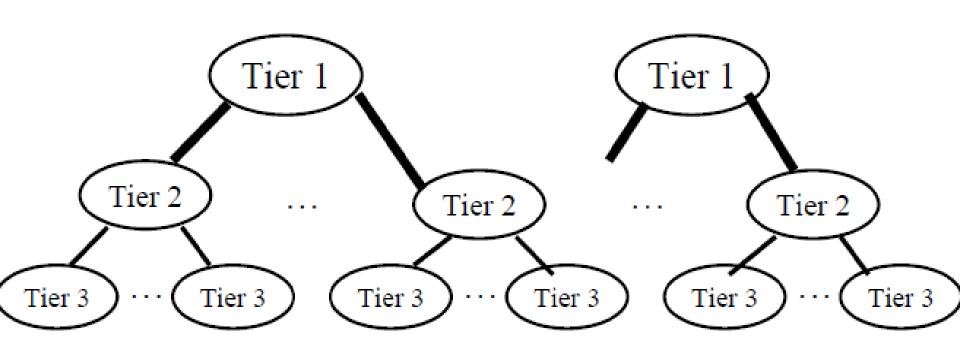
- signal carried in electromagnetic spectrum
- $^{
 m heta}$ no physical "wire"
- θ bidirectional
- θ propagation environment effects:
 - μ reflection
 - u obstruction by objects
 - μ interference

Radio link types:

- θ terrestrial microwave
 - μ e.g. up to 45 Mbps channels
 - θ LAN (e.g., Wifi)
 - µ 2Мbps, 11Мbps
 - θ wide-area (e.g., cellular)
 - μ e.g. 3G: hundreds of kbps
- θ satellite
 - up to 50Mbps channel (or multiple smaller channels)
 - μ 270 msec end-end delay Introduction

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Types of ISPs

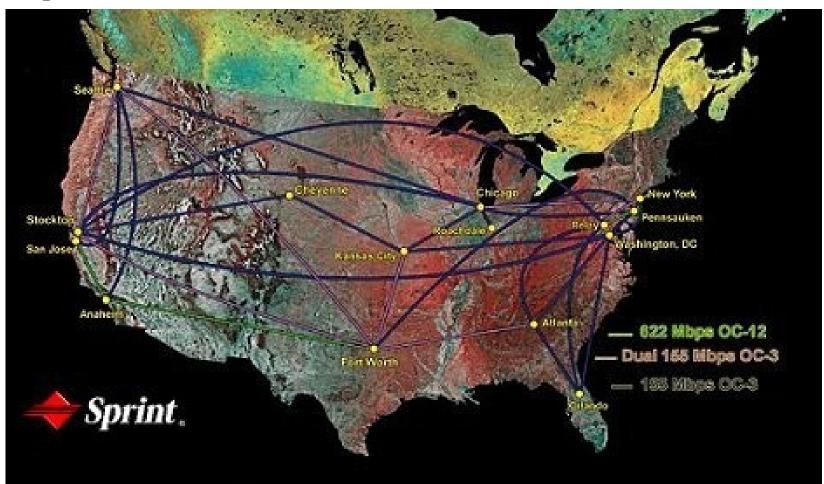


- ☐ Tier 1: Global or National, e.g., AT&T
- Tier 2: Regional
- Tier 3: Local

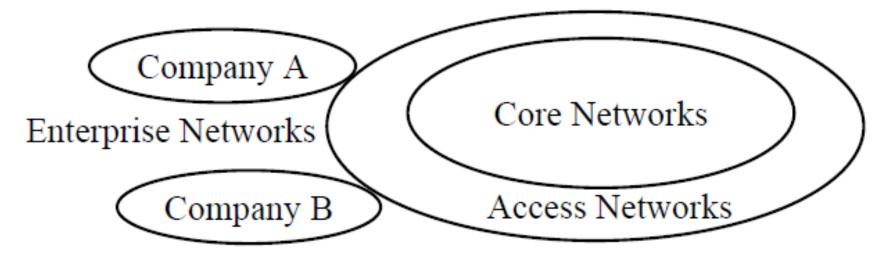
IOI

ISP: e.g., Sprint

Sprint US backbone network



Structure of the Internet

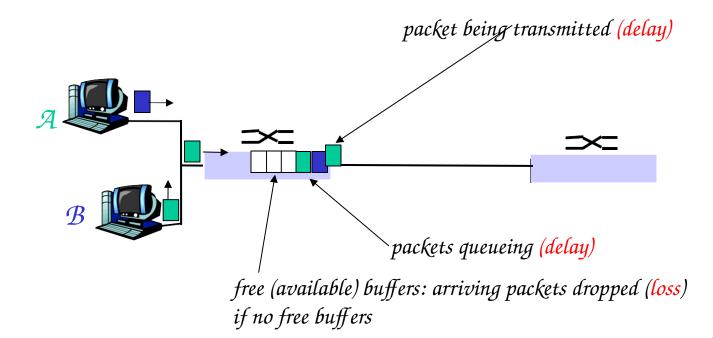


- □ Enterprise/Home Networks: Stub Networks. Privately owned ⇒ Not owned by ISP e.g., WUSTL network: Ethernet and WiFi
- Access Network: Enterprise/Users to ISP (in the city)
 WiFi, 3G/4G, DSL
- Core Network: ISP's network (between city): Optical Fiber

How do loss and delay occur?

packets queue in router buffers

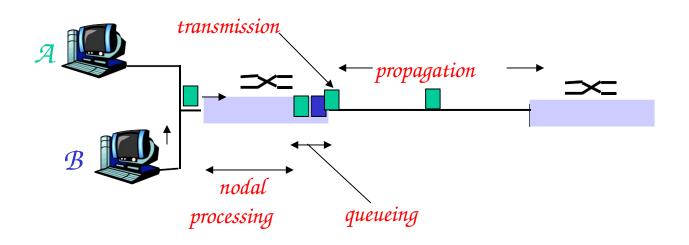
- packet arrival rate to link exceeds output link capacity
- nackets queue, wait for turn



Four sources of packet delay

- θ 1. nodal processing:
 - u check bit errors
 - u determine output link

- ⁹ 2. queueing
 - time waiting at output link for transmission
 - u depends on congestion level of router



Delay in packet-switched networks

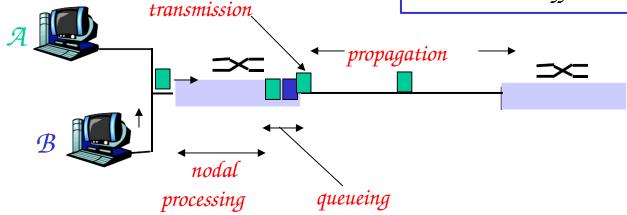
3. Transmission delay:

- θ R=link bandwidth (bps)
- θ L=packet length (bits)
- time to send bits into link = L/R.

4. Propagation delay:

- θ d = length of physical link
- $s = propagation speed in medium (~2<math>\chi$ 10 8 m/sec)
- θ propagation delay = d/s

Note: s and R are very different quantities!



Nodal delay

$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

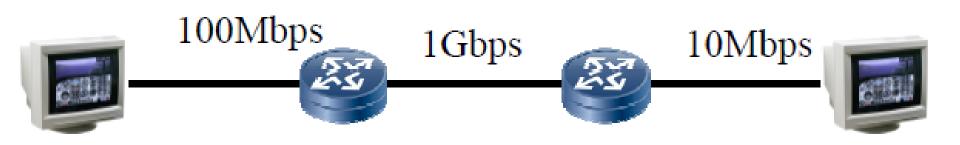
- θ $d_{proc} = processing delay$
 - μ typically a few microsecs or less
- θ $d_{queue} = queuing delay$
 - μ depends on congestion
- θ $d_{trans} = transmission delay$
 - $\mu = L/R$, significant for low-speed links
- θ $d_{prop} = propagation delay$
 - u a few microsecs to hundreds of msecs

Packet loss

- of queue (aka buffer) preceding link in buffer has finite capacity
- when packet arrives to full queue, packet is dropped (aka lost)
- lost packet may be retransmitted by previous node, by source end system, or not retransmitted at all

Throughput

- Measured in Bits/Sec
- Capacity: Nominal Throughput
- Throughput: Realistic
- Bottleneck determines the end-to-end throughput



Net end-to-end capacity = 10 Mbps

Actual throughput will be less due to sharing and overhead.

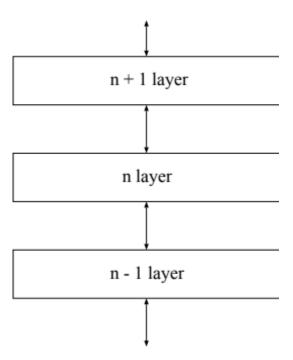
Network System Modularity

Like software modularity, but:

- Implementation is distributed across many machines (routers and hosts)
- Must decide:
 - How to break system into modules
 - Layering
 - Where modules are implemented
 - End-to-End Principle
 - Where state is stored
 - Fate-sharing

Layering Concept

- A restricted form of abstraction: system functions are divided into layers, one built upon another
- Often called a stack; but not a data structure!



Layers and Communications

- Interaction only between adjacent layers
- *layer n* uses services provided by *layer n*-1
- *layer n* provides service to *layer n+1*
- Bottom layer is physical media
- Top layer is application

Entities and Peers

Entity – a thing (an independent existence)
Entities interact with the layers above and below
Entities communicate with peer entities

 same level but different place (eg different person, different box, different host)

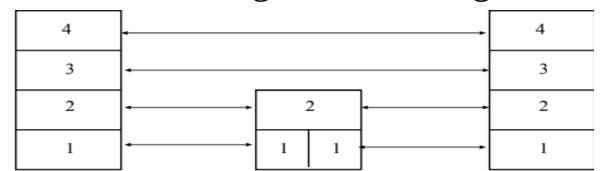
Communications between peers is supported by entities at the lower layers



Entities and Peers

Entities usually do something useful

- Encryption Error correction ReliableDelivery
- Nothing at all is also reasonable
 Not all communications is end-to-end
 Examples for things in the middle
 - IP Router Mobile Phone Cell Tower
 - Person translating French to English



Protocol "Layers"

Networks are complex!

- θ many "pieces":
 - μ hosts
 - u routers
 - u links of various media
 - μ applications
 - μ protocols
 - µ hardware, software

Organization of air travel

ticket (purchase)

baggage (check) gates

baggage (claim)

(load)

runway takeoff

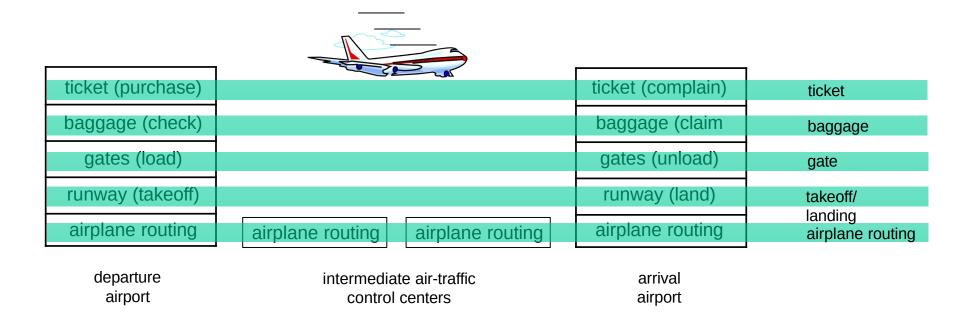
runway takeoff

airplane routing

airplane routing

θ a series of steps

Layering of airline functionality



Layers: each layer implements a service

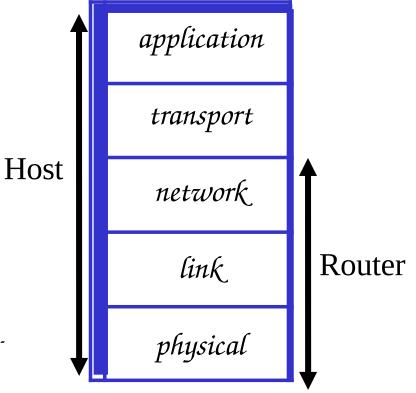
- u via its own internal-layer actions
- u relying on services provided by layer below

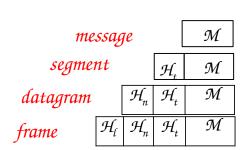
Why layering?

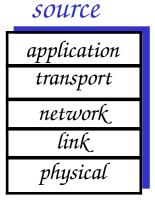
Layering is a design approach that divides the complex process of communication into simpler, well-defined steps (layers), where each layer performs a specific task.

Internet protocol stack

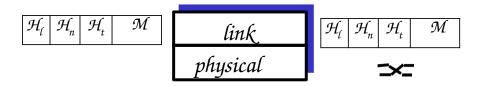
- application: supporting network applications
 - μ ftp, smtp, http
- transport: host-host data transfer
 - u tcp, udp
- network: routing of datagrams from source to destination
 - μ ip, routing protocols
- link: data transfer between neighboring network elements
 - μ ppp, ethernet
- physical: bits "on the wire", modulation scheme, linecoding format, electrical & physical specifications, etc.
- Routers in the network operate only up to the Network Layer



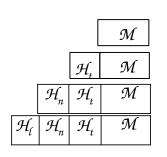


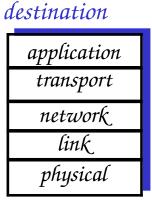


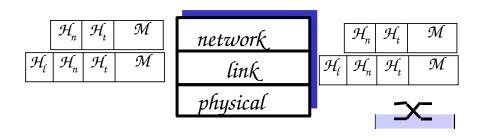
Encapsulation



switch







router

1.9 Internet History

1961-1972: Early packet-switching principles

- 1961: Kleinrock queueing theory shows effectiveness of packetswitching
 - 1964: Baran packet-switching
- θ in military nets
 - 1967: ARPAnet conceived by
- Advanced Research ProjectsAgency
 - 1969: first ARPAnet node
- $^{oldsymbol{ heta}}$ operational

- A 1972:
 - µ ARPAnet demonstrated publicly
 - NCP (Network Control Protocol) first host-host protocol
 - μ first e-mail program
 - µ ARPAnet has 15 nodes

Internet History

1972-1980: Internetworking, new and proprietary nets

- 1970: ALOHAnet satellite network in Hawaii
- θ 1973: Metcalfe's PhD thesis proposes Ethernet
- 1974: Cerf and Kahn architecture for interconnecting networks
- late 70's: proprietary architectures: DECnet, SNA, XNA
 - late 70's: switching fixed length
- packets (ATM precursor)1979: ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:

- minimalism, autonomy no internal changes required to interconnect networks
- u best effort service model
- u stateless routers
- μ decentralized control <mark>define</mark>

today's Internet architecture

<u>Internet History</u>

1980-1990: new protocols, a proliferation of networks

- θ 1983: deployment of TCP/IP
- θ 1982: smtp e-mail protocol defined
- θ 1983: DNS defined for nameto-IP-address translation
- θ 1985: ftp protocol defined
- θ 1988: TCP congestion control

- new national networks:
 Csnet, BITnet, NSFnet,
 Minitel
- 100,000 hosts

 connected to

 confederation of

 networks

Internet History

1990's: commercialization, the WWW

- θ Early 1990's: ARPAnet decommissioned
- 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- θ early 1990s: WWW
 - µ hypertext [Bush 1945, Nelson 1960's]
 - μ HTML, http: Berners-Lee
 - μ 1994: Mosaic, later Netscape
 - u late 1990's: commercialization of the WWW

Late 1990's & 2000's:

- θ est. 50 million computers on Internet
- θ est. 100 million+ users
- backbone links running at 1Gbps

You now hopefully have:

- context, overview, "feel" of networking
- p more depth, detail *later* in course