

Optimal solution = {A1, A3, A6, A8}

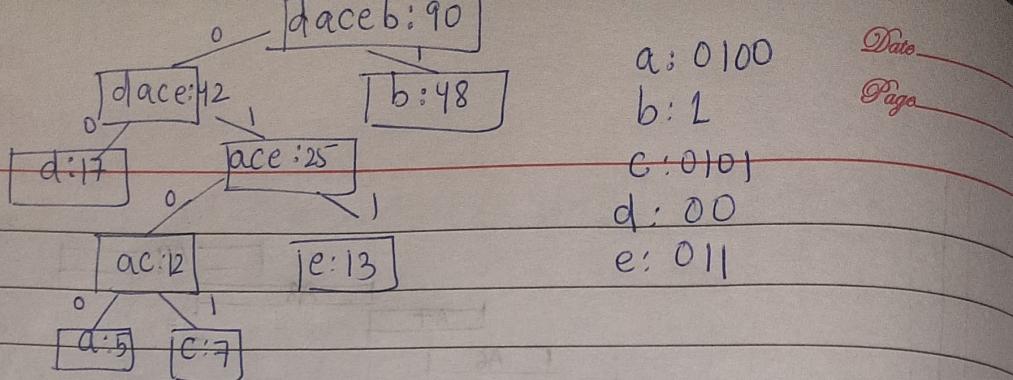
greedy Activity Selection (s, f)

1. $n \leftarrow \text{length}[S]$
2. $A \leftarrow \{a_1\}$
3. $i \leftarrow 1$
4. for $m \leftarrow 2$ to n
5. do if $s_m > f_i$
6. then $A \leftarrow A \cup \{a_m\}$
7. $i \leftarrow m$
8. return A

Hoffmann Coding (greedy).

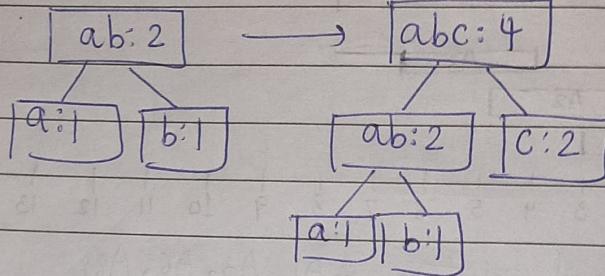
Used in compression.

Code	Fix	a: 5
a:	5	
b:	48	c: 7
c:	7	e: 13
d:	17	d: 17
e:	13	b: 48

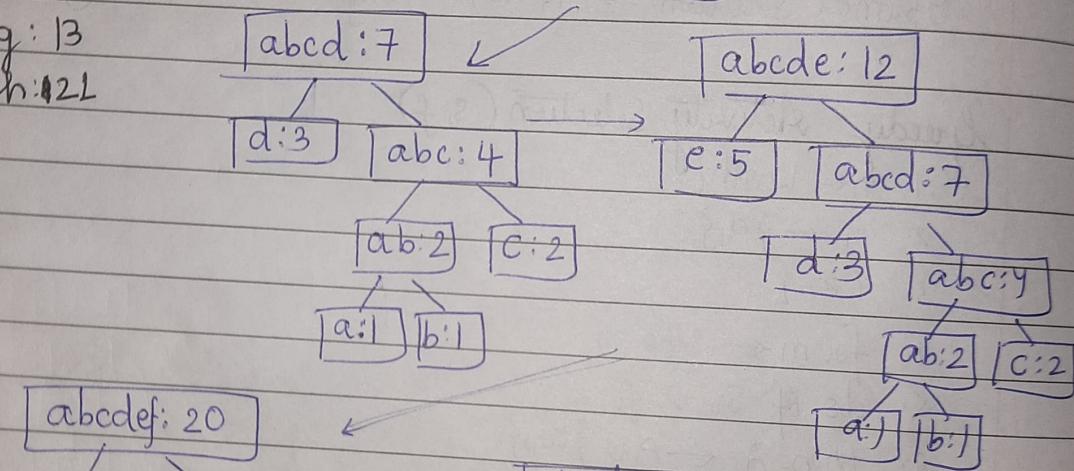


a: 0100
b: 1
c: 0101
d: 00
e: 011

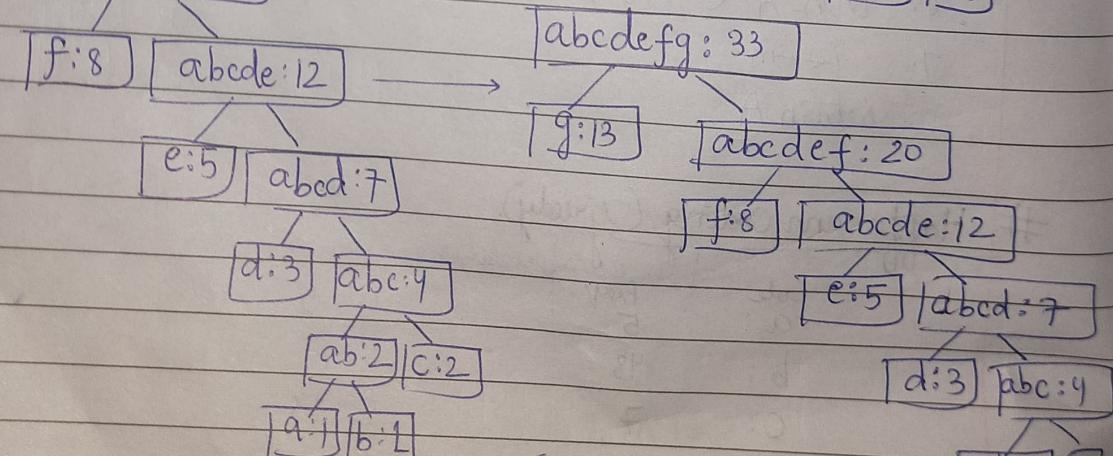
- ✓ a: 1
- ✓ b: 1
- c: 2
- ✓ d: 3
- ✓ e: 5
- ✓ f: 8



g: 13
h: 12



abcdef: 20



abcdefgh: 54

hi: 12 abcdefg: 33

g: 13 abcdef: 20

f: 8 abcde: 12

e: 5 abcd: 7

d: 3 abc: 4
a: 0 ab: 2 b: 1 c: 2

a: 111100

b: 111101

c: 111111

d: 11110

e: 1110

f: 110

g: 10

Matrix Chain Multiplication.

$$M_1 = 10 * 20$$

$$M_2 = 20 * 50$$

$$M_3 = 50 * 1$$

$$M_4 = 1 * 100$$

$$(I) \quad ((M_1, M_2) \quad (M_3, M_4))$$

$$\text{II } M_1(M_2(M_3M_4))$$

$$\text{III } (M_1 \cdot (M_2 \cdot M_3)), M_4$$

$$\text{IV} \quad ((M_1, M_2), M_3), M_4$$

A matrix chain multiplication is used in compiler design for code optimisation.

$$\begin{aligned} \text{Total Cost} &= 5000 + 100000 + 20000 \\ &= 125000. \end{aligned}$$

$(M_1 \cdot (M_2 \cdot M_3)) \cdot M_4$

Multiplication steps:

 $20 * 50 * 50 * 1 \rightarrow 20 * 50 * 1 = 2000$

 $10 * 1 \rightarrow 10 * 20 * 1 = 200$

 $10 * 200 \rightarrow 10 * 100$

Complexity

$$\begin{aligned} \text{Total Cost} &= 1000 + 200 + 1000 \\ &= 2200 \end{aligned}$$

$$\leftarrow (A_1 \cdot A_2) \cdot (A_2 \cdot A_3) \cdot (A_3 \cdot A_4)$$

$$M[1,4] = \frac{P_0}{A_1 \cdot A_2} \cdot \frac{P_2}{(A_3 \cdot A_4)} \quad (1)$$

$$= M[1,2] + M[3,4] + P_0 P_2 P_4$$

$$= 414.$$

$$A_1 \cdot \frac{P_1}{A_2 \cdot (A_3 \cdot A_4)} \quad (2)$$

$$M[1,1] + M[2,4] + M[3,4] + P_0 P_1 P_4$$
~~$$0 + 84 = 0 + 104 + 5 * 4 * 7$$~~

$$= 104 + 140$$

$$= 244.$$

$$(A_1 (A_2 \cdot A_3)) A_4 \quad (3)$$

$$M[1,3] + M[4,4] + P_0 P_3 P_4$$

$$= 158$$

$$(A_1 \cdot A_2) \cdot A_3 A_4 \quad (4)$$

$$= 158$$

$$A_1 [A_2 \rightarrow A_4]$$

$$(A_1 \rightarrow A_3) A_4$$

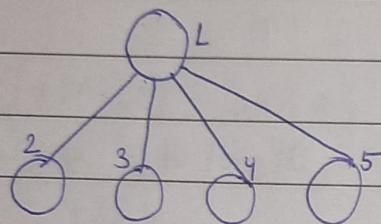
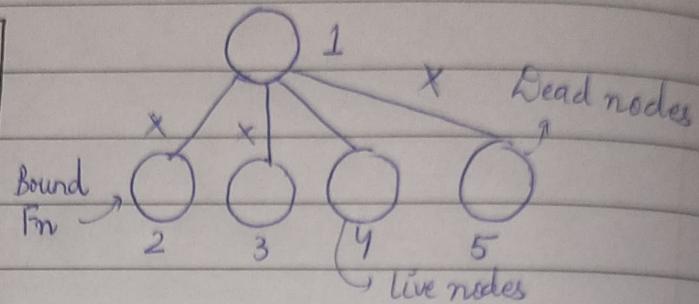
Travelling Sales Person Problem (Branch & Bound)

Cost Matrix

	1	2	3	4	5
10	∞	20	30	10	11
2	15	∞	16	4	2
3	3	5	∞	2	4
4	19	6	18	∞	3
5	16	4	7	16	∞

Row wise

State Space Tree



Step 1: Compute Reduced Cost Matrix

∞	20	30	0	1
13	∞	14	2	0
1	3	∞	0	2
16	3	15	∞	0
12	0	3	12	∞

↓ Column wise
Column - 1, 0, 3, 0, 0

∞	10	17	0	1
12	∞	11	2	0
0	3	∞	0	2
15	3	12	∞	0
11	0	0	12	∞

Reduced Cost Matrix
C[1:2] ① Make first row & second column ∞

② Make single entry
 $C_{21} = \infty$

$$10 + 2 + 2 + 3 + 4 + 1 + 3 = 25$$

③ Check each row & each column should contain atleast 1 zero

∞	∞	∞	∞	∞
∞	∞	11	2	0
0	∞	∞	0	2
15	∞	12	∞	0
11	∞	0	12	∞

$$\begin{aligned}
 C[1:2] &= \text{Reduced Cost} + C_{12} + \text{zero entry} \\
 &= 25 + 10 + 0 \\
 &= 35
 \end{aligned}$$

$$C[1:3] = \left[\begin{array}{ccccc} \infty & 10 & \cancel{\infty} & \infty & 10 \\ 12 & \infty & \infty & 2 & 0 \\ \infty & 3 & \infty & 0 & 2 \\ \cancel{15} & 4 & 3 & \infty & 0 \\ 11 & 0 & \infty & 12 & \infty \end{array} \right]$$

$$\begin{aligned}
 C[1:3] &= 25 + 17 + 0 \\
 &= 42 + 11 = 53.
 \end{aligned}$$

↓

12

$$C[1:4] = \left[\begin{array}{ccccc} \infty & \infty & \infty & \infty & 5 \\ 12 & \infty & 11 & \infty & 0 \\ 0 & 3 & \infty & \infty & 2 \\ \infty & 3 & 12 & \infty & 0 \\ 11 & 0 & 0 & \infty & \infty \end{array} \right]$$

$$C[1:4] = 25 + 0$$

$$= 25$$

$$= 25$$

$C[1:5]$

$$= \left[\begin{array}{ccccc} \infty & \infty & \infty & \infty & \infty \\ 12 & \infty & 11 & 2 & \infty \\ 0 & 3 & \infty & 0 & \infty \\ 15 & 3 & 12 & \infty & \infty \\ 0 & 0 & 0 & 12 & \infty \end{array} \right]$$

$C[1:5] + \text{Reduced Cost}$

$$= 1 + 25 + 2 + 3$$

$$= 31$$

$C[1:4:2]$

first row 4th row $\Rightarrow \infty$
 fourth column 2nd Column

$$C_{41}, C_{24}, C_{21} \Rightarrow \infty$$

$C[1:4:2] = \text{Reduced Cost of } C[1:4] + C_{42} + \text{any zero}$