

Variants of Turing Machine

1) Multitape Turing Machine (MTTM)

MTTM has multiple tapes where each tape is accessed by separate head. Each head can move independently of the other head.

A multitape TM has a finite set Q of states, an initial state q_0 , a subset F of Q called the final states, a set P of tape symbols a, b (blank)

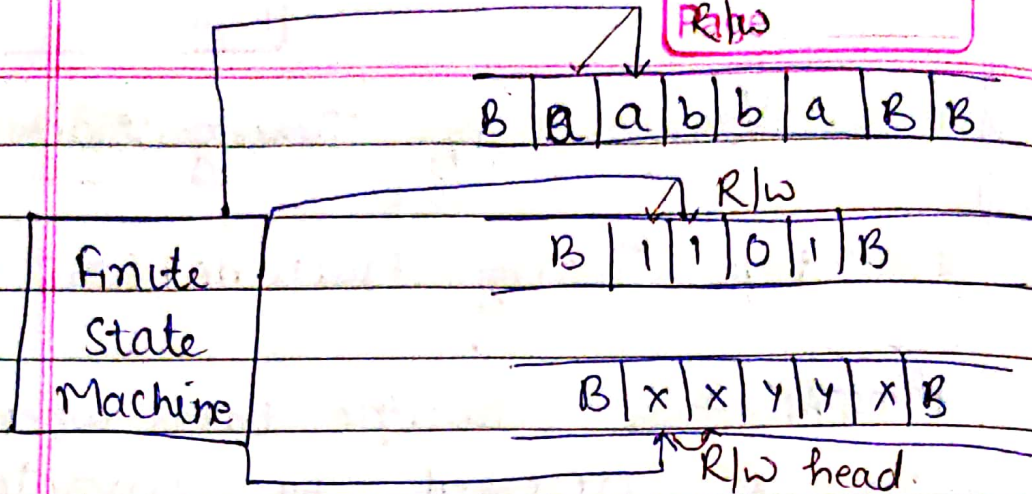
There are k tapes, each divided into cells. The first tape holds the input string w . Initially, all the other tapes hold the blank symbol.

Initially the head of the first tape (input) is at the left end of the input w . All the other heads can be placed at any cell initially.

δ (transition)

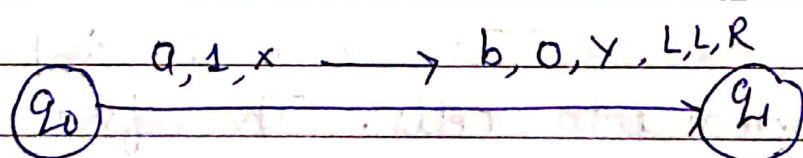
Multitape: $Q \times \Gamma^k = Q \times \Gamma^k \times \{L/R\}^k$

Singletape: $Q \times \Gamma = Q \times \Gamma \times \{R/L\}$



Several tapes - Each has its own tape head

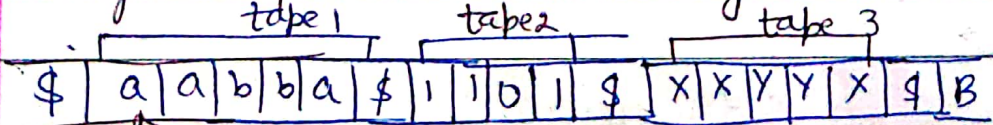
eg. of Multitape Turing machine with $R = 3$ (tapes)



$$S(q_0, a, 1, x) \rightarrow (q_1, b, 0, y, L, L, R)$$

$$Power(MTMM) = Power(STMM)$$

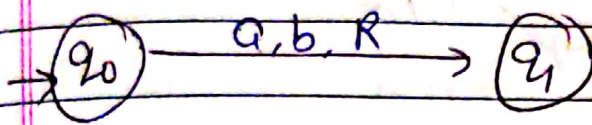
(Single tape TM simulating MTMM)



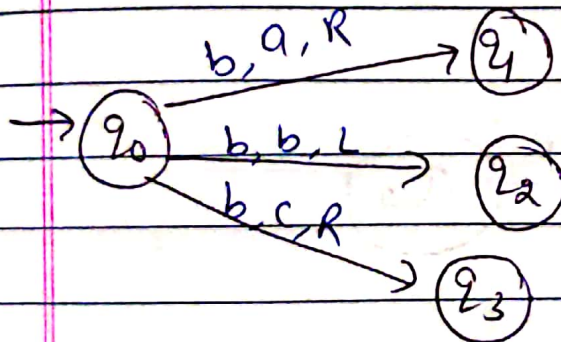
FSM

Non-Deterministic Turing Machine

In non-deterministic TM, for every state and symbol there are group of actions that T.M can have, so transitions are not deterministic.



Deterministic TM



Non-Deterministic TM

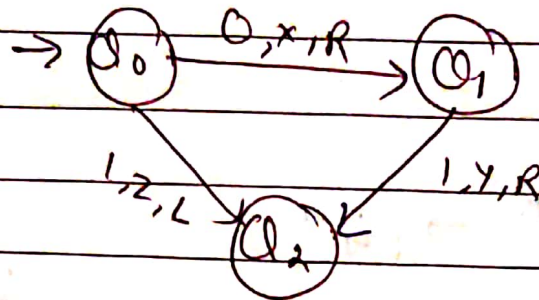
$$S(Q \times \Gamma \longrightarrow P(Q \times \Gamma \times (R/L)))$$

! Powerset of

Configuration: A way to represent the various states of TM at particular moment during compilation.

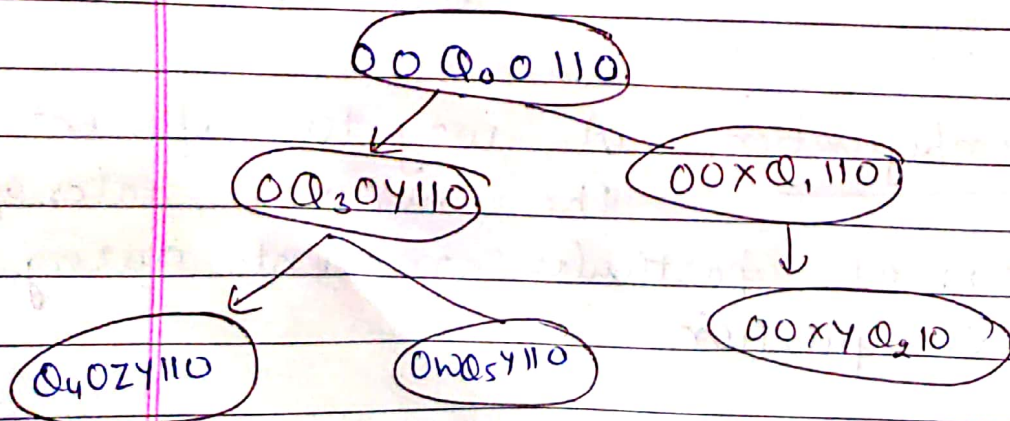
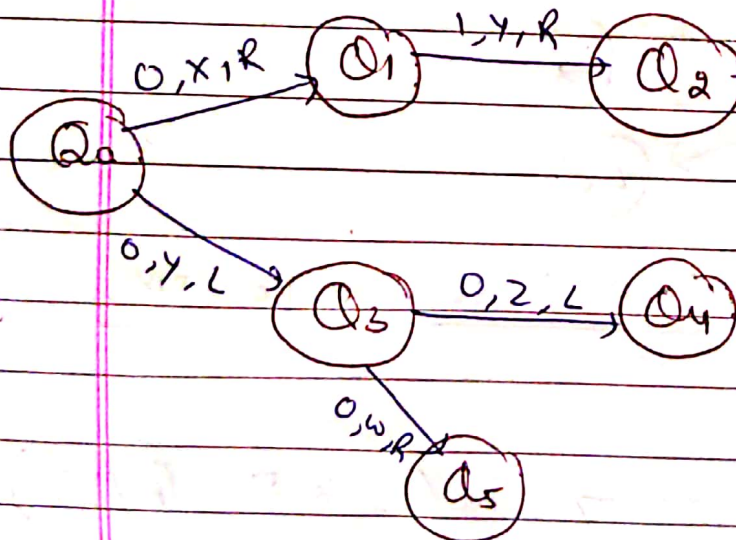
- i) Current state
- ii) Current position of head
- (iii) the entire tape content

010 Q_1 100



Deterministic

Non-Deterministic



Outcome of N.D.TM

Accept: If any branch of the computation accept then NDTM accept.

Reject: If all branches of the computation Halt and Reject (i.e. no branch in the graph accept) then NDTM Rejects.

Multitape Turing Machine is used for constructing NDTM.

