

# Rigidity Theory

Created: Jan. 10. 2026

Last Update: Feb. 17. 2026

Tag: #Distance\_Geometry/Rigidity

## 1. Preliminaries

### 1.1. Basic definitions

#### Def 1.1) Framework

A **framework** in  $\mathbb{R}^d$  is a pair  $(G, p)$ , where  $G$  is a finite simple graph and  $p : V(G) \rightarrow \mathbb{R}^d$  is a (not necessarily injective) map. We also say that  $p$  is a  $d$ -dimensional **realization** of  $G$ . Sometimes, we denote  $f_p = (G, p)$ .

#### Def 1.2) Joint, Bar

For  $v_i \in V(G)$ , we say that  $p_i := p(v_i)$  is the **joint** of  $(G, p)$  corresponding to  $v_i$ .

Also, for  $v_i v_j \in E(G)$ , we say that  $\{p_i, p_j\}$  is the **bar** of  $(G, p)$  corresponding to  $v_i v_j$ .

The **length of the bar**  $\{p_i, p_j\}$  is defined as  $L(\{p_i, p_j\}) := \|p(v_i) - p(v_j)\|_2$ .

#### Def 1.3) Motion, rigid and flexible framework

A (**finite**) **motion** of  $(G, p)$  is an indexed family  $\{P_i\}_{i \in [n]}$  of functions  $P_i : [0, 1] \rightarrow \mathbb{R}^d$  s.t.

1. (Initial Condition)  $P_i(0) = p(v_i), \forall i \in [n]$ .
2. (Differentiability)  $P_i(t)$  is differentiable on  $[0, 1], \forall i \in [n]$ .
3. (Distance-preserving)  $\|P_i(t) - P_j(t)\| = \|p_i - p_j\|$  for all  $t \in [0, 1]$  and  $v_i v_j \in E(G)$ .

A motion  $\{P_i\}$  is called **rigid** if  $\|P_i(t) - P_j(t)\| = \|p_i - p_j\|$  for all  $t \in [0, 1]$  and  $(i, j) \in [n]^2$ .

Also,  $\{P_i\}$  is called **flexible** if  $\exists (i, j) \in [n]^2 \setminus E(G)$  s.t.  $\|P_i(t) - P_j(t)\| \neq \|p_i - p_j\|$  for all  $t \in (0, 1]$ .

With these definitions,  $(G, p)$  is called rigid if there is no flexible motion. Otherwise,  $(G, p)$  is called flexible.

Q. Why "for all  $t \in (0, 1]$ " for a framework to be flexible?

### 1.2. Equivalence & Congruence

So far, we defined a framework with a tuple of a simple graph and a realization. However, for a formation system(engineering aspect), the set of edges  $E(G)$  could be defined by a set of constraints. Here, the constraint means a task or a goal that agents work together to achieve. (e.g. communication network, distance-preserving, ...)

So, we can alternatively describe a formation by  $f_p = (\mathbf{V}, \mathbf{C}, p)$  where  $\mathbf{V} = [n]$  and  $\mathbf{C}$  be a set of constraints. Obviously,  $|\mathbf{C}| = e(G)$ .

In this aspect, we can introduce more definitions.

#### Def 1.4) Equivalence / Congruence of frameworks

Two frameworks  $f_p = (G, p) = (\mathbf{V}, \mathbf{C}, p)$  and  $f_q = (\mathbf{V}, \mathbf{C}, q) = (G, q)$  with different realizations  $p, q$ , are called...

- **Equivalent** if  $\text{mag}[c_i(p(x_i))] = \text{mag}[c_i(q(x_i))]$  for all  $c_i \in \mathbf{C}$  and denoted by  $f_p \sim f_q$  where  $x_i \subset \mathbf{V}$  is the set of nodes involved in the  $i^{\text{th}}$  constraint  $c_i$ . In this context,  $p(x_i)$  and  $q(x_i)$  represent the realized points(coordinates) of these nodes in each framework, and  $\text{mag}[\cdot]$  is the exact value of the argument.
- **Congruent** if  $\text{mag}[c_i(p(x_i))] = \text{mag}[c_i(q(x_i))]$  for all  $c_i \in \mathbf{C}_K$  and denoted by  $f_p \equiv f_q$  where  $\mathbf{C}_K$  is the complete constraint set containing constraints for every possible pair of vertices in the graph  $G$ . This corresponds to the set of all edges in the complete graph  $K_n$ . In this case,  $x_i$  represents the set of nodes involved in each constraint  $c_i$  within the expanded set  $\mathbf{C}_K$ .

For example, if the constraint is given as a distance between each vertices, then  $c_i \in \mathbf{C}$  is the distance constraint on the  $i^{\text{th}}$  edge and  $x_i$  is the set of endpoints of the  $i^{\text{th}}$  edge.

So, if  $uv \in E(G)$  is the  $i^{\text{th}}$  edge, then  $x_i = \{u, v\}$  and  $c_i(p(x_i)) = \|p(u) - p(v)\|$ . In this sense, we can describe the equivalence and congruence for distance constraint only.

Two frameworks  $f_p$  and  $f_q$  are called...

- **Distance equivalent** if  $\|p_i - p_j\| = \|q_i - q_j\|$  for all  $v_i v_j \in E(G)$
- **Distance congruent** if  $\|p_i - p_j\| = \|q_i - q_j\|$  for all  $(i, j) \in [n]^2$ .

#### Def 1.5) Constraint function

For a framework  $f_p = (G, p) = (\mathbf{V}, \mathbf{C}, p)$ , the constraint function  $h_G : \mathbb{R}^{dn} \rightarrow \mathbb{R}^m$  ( $m = e(G)$ ) is defined by

$$h_G(p) := \begin{bmatrix} \text{mag}[c_1(p(x_1))] \\ \vdots \\ \text{mag}[c_m(p(x_m))] \end{bmatrix}$$

for square of distance constraint,

$$h_G(p) = \begin{bmatrix} \text{mag}[c_1(p(x_1))] \\ \vdots \\ \text{mag}[c_m(p(x_m))] \end{bmatrix} = \begin{bmatrix} \|p_{j_1} - p_{i_1}\|^2 \\ \vdots \\ \|p_{j_m} - p_{i_m}\|^2 \end{bmatrix}$$

## 2. Rigidity and infinitesimal rigidity

### 2.1. Basic rigidity concepts

Q. What is the "Rigid framework" ?

## Def 2.1) Rigidity (General definition)

A framework  $f_p = (G, p)$  is called **rigid** in a given Euclidean space  $\mathbb{R}^d$  if there exists a nbd.  $U_p$  of  $p$  s.t. every framework  $f_q = (G, q)$  with  $q \in U_p$  satisfy  $f_p \sim f_q \Rightarrow f_p \equiv f_q$ .

Algebraically,  $f_p = (G, p)$  is rigid in  $\mathbb{R}^d$  if there exists a nbd.  $U_p$  of  $p$  s.t.

$$h_G^{-1}(h_G(p)) \cap U_p = h_K^{-1}(h_K(p)) \cap U_p$$

Here,

- LHS = the set of frameworks in  $U_p$  satisfying the constraint  $C$ .
- RHS = the set of frameworks in  $U_p$  satisfying the constraint  $C_K$ .

Obviously, there is only one congruent framework satisfying  $C_K$ . Also, since  $C \subset C_K$ , LHS  $\supset$  RHS. So, this expression means that, if LHS  $\neq$  RHS, then there is at least one framework that is distinct (non-congruent) from that of in RHS and hence  $f_p$  is not rigid.

## Def 2.2) Distance rigidity

A framework  $f_p$  is called **distance rigid** in  $\mathbb{R}^d$  if there exists  $\varepsilon > 0$  s.t. for any framework  $f_q$  that is locally perturbed from  $f_p$  (i.e.  $q_i \in B_\varepsilon(p_i)$ ,  $\forall i$ ),  $f_p$  and  $f_q$  are distance equivalent only if  $f_p$  and  $f_q$  are distance congruent.

This means that  $f_p$  is distance rigid if every framework satisfying same edge-distance constraints to  $f_p$  near  $p$  is congruent to  $f_p$  (Locally unique embedding satisfying edge-distance constraint).

Q. Are there two definition of distance rigidity?  $\rightarrow$  No.

## Thm 2.3) Equivalent of definitions of rigidity

Two definitions of distance rigidity (in the view of motion and set-theoretic definition) of framework is equivalent.

Pf)

## Def 2.4) Global distance rigidity

A framework  $(G, p)$  is called **globally distance rigid** in  $\mathbb{R}^d$  if "every" framework  $(G, q)$  equivalent to  $(G, p)$  is also congruent to  $(G, p)$  in  $\mathbb{R}^d$ .

Henceforth, unless otherwise specified, the term 'rigidity' will refer exclusively to '**distance rigidity**'.

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## 2.2. Infinitesimal rigidity

**Motivation:** It is in general very difficult to determine whether a given framework is rigid or not since it requires solving a system of quadratic equations. It is therefore common to **linearize** the problem by differentiating the equations.

Let  $\{P_i\}$  be a finite motion of a 2-dimensional framework  $(G, p)$  and let  $\{p_i, p_j\}$  be a bar of length  $l$ . Then,

$$\|P_i(t) - P_j(t)\|^2 = \left[ P_i^{(1)}(t) - P_j^{(1)}(t) \right]^2 + \left[ P_i^{(2)}(t) - P_j^{(2)}(t) \right]^2 = l^2 \quad \forall t \in [0, 1]$$

where  $P_i^{(k)}(t)$  is the  $k^{\text{th}}$  coordinate of  $p_i$  at  $t$ .

By differentiating both sides w.r.t.  $t$ ,

$$\left( P_1^{(1)}(t) - P_j^{(1)}(t), P_1^{(2)}(t) - P_j^{(2)}(t) \right) \cdot \left( \left[ P_1^{(1)}(t) - P_j^{(1)}(t) \right]', \left[ P_1^{(2)}(t) - P_j^{(2)}(t) \right]' \right) = 0$$

Evaluating at  $t = 0$ , we obtain

$$\left( p_i^{(1)} - p_j^{(1)}, p_i^{(2)} - p_j^{(2)} \right) \cdot \left( u_i^{(1)} - u_j^{(1)}, u_i^{(2)} - u_j^{(2)} \right) = 0$$

or in short  $(p_i - p_j) \cdot (u_i - u_j)$ . Here,  $u_i, u_j$  are initial velocity vectors of  $p_i$  and  $p_j$ , respectively. So, this implies that the difference between initial velocity vectors is orthogonal to the line through  $p_i$  and  $p_j$ .

These vectors preserve the length of  $\{p_i, p_j\}$  at first order.

### Def 2.5) Constricted framework

A  $d$ -dimensional framework  $(G, p)$  is called **constricted** if all points in  $\mathcal{P} = \{p_1, \dots, p_n\}$  lie in  $k \geq d - 1$  dimensional affine subspace. Otherwise,  $(G, p)$  is called **normal**.

Thus, a framework is constricted if it is geometrically degenerated.

### Def 2.6) Infinitesimal motion / Infinitesimal rigidity

An **infinitesimal motion** of a framework  $(G, p)$  is a function  $u : V(G) \rightarrow \mathbb{R}^d$  s.t.

$$(p_i - p_j) \cdot (u_i - u_j) = 0 \quad \forall v_i v_j \in E(G)$$

That is, for a distance constrained graph  $G$ , a function consists of  $dn$  instantaneous velocity vectors of the nodes that do not change the equivalence conditions between nodes.

Here, an infinitesimal motion is called **rigid** if

$$(p_i - p_j) \cdot (u_i - u_j) = 0 \quad \forall (i, j) \in [n]^2$$

Also, an infinitesimal motion is called **flex** if

$$(p_i - p_j) \cdot (u_i - u_j) \neq 0 \quad \exists v_i v_j \notin E(G)$$

In this sense,  $(G, p)$  is called **infinitesimally rigid** if every infinitesimal motion of  $(G, p)$  is an infinitesimal rigid motion. Otherwise,  $(G, p)$  is said to be **infinitesimally flexible**.

**Note:** In some context, infinitesimal rigid motion also can be called by a "Trivial motion" or a "Special Euclidean motion".

- **Trivial motions** is an infinitesimal motions include only the translational and rotational motions. (infinitesimal rigid motion)
- **Special Euclidean motion** is defined by the motions of nodes that keep the congruence conditions among the nodes. It contains the translations and rotations, and hence Special Euclidean motion is equivalent to the **direct isometries**.

In short,

## 🔗 Equivalence of definitions

- Infinitesimal rigid motion = Trivial motion
- Special Euclidean motion = Direct isometry

However, **in the infinitesimal viewpoint, the four of them are all equivalent.**

**Remark:** All constricted frameworks infinitesimally flexible.

- Now, we can come up with a very natural question.

Q. Is there any relationship between the rigidity and the infinitesimal rigidity?

### Thm 2.7) Infinitesimal rigidity is stronger than rigidity

If a framework  $f_p$  is infinitesimally rigid, then  $f_p$  is rigid.

**Proof)**

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#### 2.2.1. Rigidity Matrix

##### Def 2.8) Edge function / Rigidity Matrix

Let  $(G, p)$  be a framework with graph  $G$  and realization  $p$ . Then, the **edge function**  $g_G : \mathbb{R}^{dn} \rightarrow \mathbb{R}^m$  where  $m = e(G)$  is defined by

$$g_G(p) = \frac{1}{2} \left[ \|p_i - p_j\|^2 \right]_{(i,j) \in E(G)}$$

Then, the **rigidity matrix** is defined by

$$\mathbf{R}(G, p) = \frac{\partial g_G(p)}{\partial p} \in \mathbb{R}^{m \times dn}$$

Here, each row and column corresponds to each edges and nodes of the framework, respectively.

Explicitly, the rigidity matrix is expressed as

$$\mathbf{R}(G, p) = \begin{pmatrix} & & & & & \vdots & & & \\ 0 & \dots & 0 & p_i - p_j & 0 & \dots & 0 & p_j - p_i & 0 & \dots & 0 \\ & & & & & \vdots & & & \\ & & & & & & & & & & \end{pmatrix}$$

Let  $e_k = v_i v_j \in E(G)$  is the  $k^{\text{th}}$  edge with endpoints  $v_i$  and  $v_j$ . Then, it occupies the  $d(i-1)+1^{\text{th}}$  to  $di^{\text{th}}$  columns and  $d(j-1)+1^{\text{th}}$  to  $dj^{\text{th}}$  columns of the  $k^{\text{th}}$  row.

Here,  $(p_i^{(1)} - p_j^{(1)}), \dots, (p_i^{(d)} - p_j^{(d)})$  are in the  $d(i-1)+1^{\text{th}}$  to  $di^{\text{th}}$  columns of the  $k^{\text{th}}$  row.

Similarly,  $(p_j^{(1)} - p_i^{(1)}), \dots, (p_j^{(d)} - p_i^{(d)})$  are in the  $d(j-1)+1^{\text{th}}$  to  $dj^{\text{th}}$  columns of the  $k^{\text{th}}$  row.

### ☰ Example (Tetrahedron)

For a formation  $(G, p)$  of regular tetrahedron in  $\mathbb{R}^3$ ,

$V(G) = \{v_1, v_2, v_3, v_4\}$  and  $E(G) = \{v_1v_2, v_1v_3, v_1v_4, v_2v_3, v_2v_4, v_3v_4\}$  and

$p = [p_1^\top \ p_2^\top \ p_3^\top \ p_4^\top]^\top$  where  $p_i = [x_i \ y_i \ z_i]$ .

Note that  $|V(G)| = n(G) = 4$ ,  $|E(G)| = e(G) = 6$ ,  $d = 3$ .

Then, by definition, the rigidity matrix be

$$\mathbf{R}(p) = \begin{bmatrix} (p_1 - p_2)^\top & (p_2 - p_1)^\top & 0^\top & 0^\top \\ (p_1 - p_3)^\top & 0^\top & (p_3 - p_1)^\top & 0^\top \\ (p_1 - p_4)^\top & 0^\top & 0^\top & (p_4 - p_1)^\top \\ 0^\top & (p_2 - p_3)^\top & (p_3 - p_2)^\top & 0^\top \\ 0^\top & (p_2 - p_4)^\top & 0^\top & (p_4 - p_2)^\top \\ 0^\top & 0^\top & (p_3 - p_4)^\top & (p_4 - p_3)^\top \end{bmatrix} \in \mathbb{R}^{6 \times (3 \times 4)}$$

Let  $p_1 = (1, 1, 1)$ ,  $p_2 = (1, -1, -1)$ ,  $p_3 = (-1, 1, -1)$ ,  $p_4 = (-1, -1, 1)$  be coordinates of each vertex. Then, the rigidity matrix be

$$\mathbf{R}(p) = \left[ \begin{array}{ccc|ccc|ccc|ccc} 0 & 2 & 2 & 0 & -2 & -2 & 0 & 0 & 0 & 0 & 0 & 0 \\ 2 & 0 & 2 & 0 & 0 & 0 & -2 & 0 & -2 & 0 & 0 & 0 \\ 2 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & -2 & -2 & 0 \\ 0 & 0 & 0 & 2 & -2 & 0 & -2 & 2 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 2 & 0 & -2 & 0 & 0 & 0 & -2 & 0 & 2 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 2 & -2 & 0 & -2 & 2 \end{array} \right]$$

Note that infinitesimal motion is a function of velocity vectors of each nodes. So, there are  $d$  velocity components for each nodes and hence one infinitesimal motion can be identified as a  $dn$ -dimensional column vector  $\mathbf{u}$ .

- Recall:  $u : V(G) \rightarrow \mathbb{R}^d$  is an infinitesimal motion if  $(p_i - p_j) \cdot (u_i - u_j) = 0$  for all  $v_i v_j \in E(G)$ .

From here, we can find some useful property.

### Prop 2.9) Infinitesimal motion with rigidity matrix

The set of infinitesimal motion of framework  $(G, p)$  is the kernel of  $\mathbf{R}(G, p)$ .

#### Proof

**Proof)** Let  $\mathcal{M}(G, p)$  be the set of infinitesimal motion of  $(G, p)$  and  $\mathbf{u} \in \mathcal{M}(G, p)$ .

Let  $e_k = v_i v_j$ . Then, the  $k^{\text{th}}$  row of  $\mathbf{R}(G, p)\mathbf{u}$  be

$$[\mathbf{R}(G, p)\mathbf{u}]_k = (p_i - p_j) \cdot u_i + (p_j - p_i) \cdot u_j = (p_i - p_j) \cdot (u_i - u_j)$$

Here, since  $\mathbf{u}$  is a column vector representation of an infinitesimal motion, RHS = 0 and hence  $\mathbf{u} \in \text{Ker}(\mathbf{R}(G, p))$ .

Thus,  $\mathcal{M}(G, p) \subset \text{Ker}(\mathbf{R}(G, p))$ .

Conversely, given  $\mathbf{x} \in \text{Ker}(\mathbf{R}(G, p))$ ,  $[\mathbf{R}(G, p)\mathbf{x}]_k = 0$  for each  $k$  and hence  $(p_i - p_j) \cdot (x_i - x_j) = 0$ .

Thus,  $\mathbf{x}$  be an infinitesimal motion of  $(G, p)$  and hence  $\text{Ker}(\mathbf{R}(G, p)) \subset \mathcal{M}(G, p)$ .

$$\therefore \mathcal{M}(G, p) = \text{Ker}(\mathbf{R}(G, p)). \quad \square$$

From this fact, by letting  $\mathcal{R}(G, p)$  be the set of all infinitesimal rigid motion,

$$\mathcal{R}(G, p) \subset \mathcal{M}(G, p)$$

by definition.

Also,  $(G, p)$  is infinitesimally rigid if and only if  $\mathcal{R}(G, p) = \mathcal{M}(G, p)$ .

By the previous proposition, the following fact directly follows.

### Lem 2.10) Dimension and infinitesimal rigidity

Let  $(G, p)$  be a normal framework in  $\mathbb{R}^d$ . Then,

$$\dim \mathcal{R}(G, p) \leq \dim \mathcal{M}(G, p) = \text{nullity}(\mathbf{R}(G, p))$$

Equivalently, by rank-nullity theorem,

$$\text{rank}(\mathbf{R}(G, p)) \leq dn - \dim \mathcal{R}(G, p)$$

The equality holds if and only if  $(G, p)$  is infinitesimally rigid.

Here, while we know  $d, n$  but we don't know about  $\dim \mathcal{R}(G, p)$  and hence we need to find it.

To find it, we assume **2-dimensional general framework** for the following arguments. They can be directly generalized to an arbitrary dimension.

### Def 2.11) General Position

$\mathcal{P} = \{v_1, \dots, v_n\} \subseteq \mathbb{R}^d$  is in **general position** if for  $1 \leq m \leq d$ , no  $m + 1$  points in  $\mathcal{P}$  lie in an  $(m - 1)$ -dimensional affine subspace of  $\mathbb{R}^d$ . A framework whose joints are in general position is called a **general framework**.

The following definition explains about two types of infinitesimal motions.

### Def 2.12) Infinitesimal translation / Infinitesimal rotation ( $\mathbb{R}^2$ ver.)

An **infinitesimal translation of the plane** is of the form

$$\tau_{(u,v)} : \mathbb{R}^2 \rightarrow \mathbb{R}^2 ; \quad \tau_{(u,v)}(x, y) = (u, v)$$

This implies that every point  $(x, y)$  in  $\mathbb{R}^2$  moves with same velocity  $(u, v)$ .

An **infinitesimal rotation of the plane** is of the form

$$a\rho_{(x_0,y_0)} : \mathbb{R}^2 \rightarrow \mathbb{R}^2 ; \quad a\rho_{(x_0,y_0)}(x, y) = a(y_0 - y, x - x_0)$$

Let  $r$  be a radius vector from a center  $(x_0, y_0)$  to  $(x, y)$ , i.e.  $r = (x - x_0, y - y_0)$ . Then, the rotation vector  $v$  of  $(x, y)$  must satisfy  $r \perp v$  and hence  $r \cdot v = 0$ . So,  $v$  must be of the form

$$a(-(y - y_0), x - x_0) = a(y_0 - y, x - x_0) \text{ where } a \in \mathbb{R}.$$

### Lem 2.13) Dimension of infinitesimal translations and rotations

The set of all infinitesimal translations and rotations of  $\mathbb{R}^2$  forms a 3-dimensional vector space.

#### Proof)

Now, recall that for infinitesimal viewpoint, infinitesimal rigid motion is equal to direct isometry and hence **infinitesimal rigid motion consists of infinitesimal translation and rotation only**.

The following lemma describes about this fact.

### Lem 2.14) Classifications fo infinitesimal rigid motions

Let  $(G, p)$  be a 2-dimensional general framework with  $V(G) = \{v_1, \dots, v_n\}$ ,  $n \geq 2$ . Then  $\mathcal{R}(G, p)$  is the 3-dimensional space of infinitesimal translations and rotations restricted to the points  $\mathcal{P} = \{v_1, \dots, v_n\}$  of  $(G, p)$ .

#### Proof

The proof consists of two steps.

For a restriction map  $(\cdot)|_{\mathcal{P}} : I'_d(\mathbb{R}^2) \rightarrow \mathcal{R}(G, p)$ ;  $f \mapsto f|_{\mathcal{P}}$ , (where  $I'_d(\mathbb{R}^2)$  is a set of infinitesimal direct isometries, i.e. infinitesimal translations and rotations)

- **Surjectivity:** Given  $u \in \mathcal{R}(G, p)$ ,  $\exists \bar{u} \in I'_d(\mathbb{R}^2)$  s.t.  $\bar{u}|_{\mathcal{P}} = u$ .
- **Injectivity:** Given  $\kappa, \lambda \in I'_d(\mathbb{R}^2)$  and corresponding restriction  $\kappa|_{\mathcal{P}}, \lambda|_{\mathcal{P}} \in \mathcal{R}(G, p)$ ,  $\kappa|_{\mathcal{P}} = \lambda|_{\mathcal{P}} \Rightarrow \kappa = \lambda$ .

#### Surjectivity

Let  $u = (u_1, \dots, u_n) \in \mathcal{R}(G, p)$  with  $u_i = (u_i^{(1)}, u_i^{(2)})$  for all  $i = 1, \dots, n$ .

Clearly, the restriction of  $\tau_{u_1}$  to  $\mathcal{P}$  is an infinitesimal rigid motion of  $(G, p)$ .

Hence  $\hat{u} = u - \tau_{u_1}|_{\mathcal{P}} \in \mathcal{R}(G, p)$ .

Here, note that  $\hat{u} = (u_1 - u_1, u_2 - u_1, \dots, u_n - u_1)$  and hence  $\hat{u}_1 = (0, 0)$ . (fixing  $p_1$ )

Next, from the definition of infinitesimally rigid motion  $(p_i - p_j) \cdot (u_i - u_j) = 0$  for all  $(i, j) \in [n]^2$ , we have

$$(p_2 - p_1) \cdot (\hat{u}_2 - \hat{u}_1) = (p_2 - p_1) \cdot \hat{u}_2 = 0$$

Since  $(p_2 - p_1) \cdot \hat{u}_2 = (p_2^{(1)} - p_1^{(1)}, p_2^{(2)} - p_1^{(2)}) \cdot \hat{u}_2 = 0$ ,

$\hat{u}_2$  must be of the form  $\hat{u}_2 = a(p_2^{(1)} - p_1^{(1)}, -(p_2^{(2)} - p_1^{(2)}))$ .

So,  $\hat{u}_2^{(1)} = a(p_2^{(1)} - p_1^{(1)})$  and  $\hat{u}_2^{(2)} = a(p_1^{(2)} - p_2^{(2)})$  and hence

$$a = \frac{\hat{u}_2^{(1)}}{p_2^{(2)} - p_1^{(2)}} = \frac{\hat{u}_2^{(2)}}{p_1^{(1)} - p_2^{(1)}} \quad \text{where } p_1^{(1)} \neq p_2^{(1)} \text{ or } p_1^{(2)} \neq p_2^{(2)}$$

Now, define another motion  $\bar{u} = \hat{u} - a\rho_{p_1}|_{\mathcal{P}} = u - \tau_{u_1}|_{\mathcal{P}} - a\rho_{p_1}|_{\mathcal{P}}$ .

Here, since  $p_1$  is fixed by  $\hat{u}$  and the rotation  $a\rho_{p_1}|_{\mathcal{P}}$  is a rotation through  $p_1$ ,  $\bar{u}$  fixes  $p_1$  as well.

Also, by choosing proper rotation, we can make  $\bar{u}$  fixes  $p_2$ . Then, since  $\bar{u}$  fixes both  $p_1$  and  $p_2$ ,  $\bar{u}_1 = \bar{u}_2 = (0, 0)$ .

Now, claim: If  $p_1$  and  $p_2$  be fixed, then  $p_i$  are fixed for all  $i > 2$ .

From the definition of infinitesimal rigid motion, we have

$$(p_i - p_1) \cdot (\bar{u}_i - \bar{u}_1) = (p_i - p_1) \cdot \bar{u}_i = 0 \quad \forall i > 2 \quad \dots (*)$$

$$(p_i - p_2) \cdot (\bar{u}_i - \bar{u}_2) = (p_i - p_2) \cdot \bar{u}_i = 0 \quad \forall i > 2 \quad \dots (**)$$

Since we assumed the general position,  $p_i$  is not on the line through  $p_1$  and  $p_2$  and hence the vectors  $p_i - p_1$  and  $p_i - p_2$  are not parallel.

However, since  $\bar{u}_i$  is perpendicular to both of them by  $(*)$  and  $(**)$ .

Note that in the two-dimensional plane, the only vector that is perpendicular to two vectors that are not parallel to each other is the zero vector  $(0, 0)$  and hence  $\bar{u}_i = (0, 0)$  for all  $i > 2$ .

Thus,  $\bar{u} = u - \tau_{u_1}|_{\mathcal{P}} - a\rho_{p_1}|_{\mathcal{P}} = \mathbf{0}$  and hence  $u = \tau_{u_1}|_{\mathcal{P}} + a\rho_{p_1}|_{\mathcal{P}} = (\tau_{u_1} + a\rho_{p_1})|_{\mathcal{P}}$ .

This implies that any infinitesimal rigid motion  $u$  of framework  $(G, p)$  is the restriction of combination of infinitesimal translation and rotation of  $\mathbb{R}^2$  to  $\mathcal{P}$ .

---

## Injectivity

Suppose that  $\kappa$  and  $\lambda$  are two infinitesimal rigid motions of  $\mathbb{R}^2$  s.t.

$$\kappa|_{\mathcal{P}} = \lambda|_{\mathcal{P}} = u$$

Let  $x \in \mathbb{R}^2$  that is not on the line through  $p_1$  and  $p_2$ .

Consider  $\mu = \kappa - \lambda$ . Then, obviously,  $\mu$  is an infinitesimal rigid motion of  $\mathbb{R}^2$ .

Then, by the definition of infinitesimal rigid motion,

$$(\mu(x) - \mu(p_1)) \cdot (x - p_1) = 0$$

$$(\mu(x) - \mu(p_2)) \cdot (x - p_2) = 0$$

Here, since  $\kappa|_{\mathcal{P}} = \lambda|_{\mathcal{P}}$ , for any joint  $p_i$ ,  $\kappa(p_i) = \lambda(p_i)$  and hence  $\mu(p_i) = (0, 0)$  for all  $i$ .

So,

$$\mu(x) \cdot (x - p_1) = 0$$

$$\mu(x) \cdot (x - p_2) = 0$$

By assumption,  $x$  is not on the line through  $p_1$  and  $p_2$  and hence two vectors  $x - p_1$  and  $x - p_2$  are not parallel. Again, since the only vector that is perpendicular to two vectors that are not parallel to each other is the zero vector  $(0, 0)$ , we can find that  $\mu(x) = (0, 0)$ .

So far, we've shown that  $\mu(x) = (0, 0)$  for  $x \in \mathbb{R}^2 \setminus L(p_1, p_2)$  where  $L(p_1, p_2)$  is the straight line through  $p_1$  and  $p_2$ .

Finally, let  $y \in L(p_1, p_2)$ . Then, obviously,  $y$  does not lie on a line through  $x$  and  $p_1$  or  $x$  and  $p_2$  (or both).

WLOG, let  $y$  is not on a line through  $x$  and  $p_1$ . By the definition of infinitesimal rigid motion,

$$(\mu(y) - \mu(p_1)) \cdot (y - p_1) = 0$$

$$(\mu(y) - \mu(x)) \cdot (y - x) = 0$$

and since  $\mu(p_1) = \mu(x) = 0$ ,  $\mu(y)$  is perpendicular to both  $y - p_1$  and  $y - x$  that are not in parallel. Thus,  $\mu(y) = (0, 0)$ .

Therefore, we proved that  $\mu(x) = 0$  for all  $x \in \mathbb{R}^2$  and hence  $\kappa = \lambda$ .

Hence, proved. □

Using previous lemma, we can obtain an important result.

Before observing it, introduce a definition first.

### Def 2.15) Degree of freedom

For a framework  $(G, p)$ , a **degree of freedom** of  $(G, p)$  is nullity  $R(G, p)$ .

Also,  $\text{nullity}(\mathbf{R}(G, p)) - \dim \mathcal{R}(G, p)$  is called an **internal degrees of freedom** of  $(G, p)$ .

Then, we can observe a degree of freedom by the following theorem.

### Thm 2.16) Lower bound of DOF

Let  $(G, p)$  be a general framework in  $\mathbb{R}^d$  with  $n(G) \geq d$ . Then,

$$\text{nullity}(\mathbf{R}(G, p)) \geq d + \binom{d}{2} = \binom{d+1}{2}$$

or equivalently,

$$\text{rank}(R(G, p)) \leq dn - \binom{d+1}{2}$$

Here, equality holds if and only if  $(G, p)$  is infinitesimally rigid.

#### Proof

**Proof** Recall that  $\dim \mathcal{R}(G, p) \leq \dim \mathcal{M}(G, p) = \text{nullity}(\mathbf{R}(G, p))$  where  $\mathcal{R}(G, p)$  is the set of infinitesimally rigid motions of  $(G, p)$  and  $\mathcal{M}(G, p)$  is the set of infinitesimal motions of  $(G, p)$  by [lemma 2.10](#).

By [lemma 2.14](#), a set of infinitesimal rigid motion consists of infinitesimal translations and rotations and hence  $\dim \mathcal{R}(G, p)$  is the number of basis, i.e. the number of independent translations and rotations.

Note that there are  $d$  translations  $\tau_{e_1}, \dots, \tau_{e_d}$  where  $e_i$  is a standard basis vector and there are  $\binom{d}{2}$  rotations  $R_{ij}$  where for all  $i, j \in [d]$ ,  $R_{ij} : \mathcal{P} \rightarrow \mathbb{R}^d$ ;  $R_{ij}(p_k) = p_k^{(i)} e_j - p_k^{(j)} e_i$  for all  $k$ .

So, the total number of basis is  $d + \binom{d}{2} = \binom{d+1}{2}$  and hence

$$\binom{d+1}{2} \leq \text{nullity}(\mathbf{R}(G, p))$$

Hence, proved. □

## Cor 2.17) Maximum number of edges of flexible framework

Let  $(G, p)$  be a general framework in  $\mathbb{R}^d$  with  $n(G) \geq d$ . If  $e(G) < dn - \binom{d+1}{2}$ , then  $(G, p)$  is infinitesimally flexible.

This corollary provides the lower bound of number of edges for a framework to be rigid. **The following definition explains about a framework that has exactly this minimum threshold of edges required for rigidity.**

## Def 2.18) Minimally rigid framework

A rigid framework  $(G, p)$  is called **minimally rigid** if the removal of any edges makes it flexible. Similarly, an infinitesimally rigid framework  $(G, p)$  is called **minimally infinitesimally rigid** if the removal of any edges makes it not infinitesimally rigid.

## Def 2.19) Independent / Isostatic framework

A framework  $(G, p)$  is **independent** if the rows of its rigidity matrix  $\mathbf{R}(G, p)$  are linearly independent.  $(G, p)$  is **isostatic** if it is infinitesimally rigid and independent.

## Thm 2.20) Equivalent conditions of isostatic framework

For a  $d$ -dimensional general framework  $(G, p)$  with  $n(G) \geq d$ , TFAE:

1.  $(G, p)$  is isostatic.
2.  $(G, p)$  is infinitesimally rigid and  $e(G) = dn - \binom{d+1}{2}$ .
3.  $(G, p)$  is independent and  $e(G) = dn - \binom{d+1}{2}$ .
4.  $(G, p)$  is minimal infinitesimally rigid.

## Def 2.21) Regular points

Given a framework  $(G, p)$ , a realization  $p$  is called **regular** if it satisfies

$$\text{rank}(\mathbf{R}(G, p)) = \max_{p \in \mathbb{R}^{dn}} \text{rank}(\mathbf{R}(G, p))$$

While this definition might be confused with [General position](#), they represent fundamentally different perspectives on a framework:

**General position** is a purely geometric property of the point coordinates, whereas **Regularity** is an algebraic property concerning the rank of the rigidity matrix.

To break this down more precisely:

- **General position** focuses on the **geometry of the joints**; it requires that no  $m + 1$  points lie in an  $(m - 1)$ -dimensional affine subspace, a condition that depends only on the point locations and is independent of the graph's edges.
- **Regular realization** focuses on the **algebraic rank of the structure**; a realization  $p$  is regular if its rigidity matrix  $R_p$  attains the maximum possible rank  $k$  for that specific graph, meaning it inherently

depends on the graph's connectivity.

Therefore, a "general framework" is defined by its well-distributed geometry, while a "regular realization" is defined by its ability to reach the full structural potential (maximum rank) allowed by its specific graph topology.

### 3. Combinatorial aspects of rigidity

#### Motivation

Up to this point, our analysis has **focused on the rigidity of frameworks based on their specific geometric configurations**—that is, the exact coordinates of vertices in  $\mathbb{R}^d$ .

Now, a fundamental question arises: to what extent can these **properties be characterized solely by the underlying graph's connectivity**, independent of its specific realization?

So, we'll investigate the possibility of **treating rigidity as an intrinsic property** of the underlying graph, leading us to the study of generic frameworks where infinitesimal rigidity becomes purely combinatorial.

Introduce some definitions first.

#### Def 3.1) Indeterminate rigidity matrix

Let  $K_n$  be the complete graph on  $n$  vertices. For each  $i = 1, 2, \dots, n$ , we introduce a  $d$ -tuple of variables  $p'_i = ((p_i^{(1)})', \dots, (p_i^{(d)})')$  ( $d$  unknown variables).

Then, let

$$\mathbf{R}(n, d) = \begin{pmatrix} & & & & & \vdots & & & \\ 0 & \dots & 0 & p'_i - p'_j & 0 & \dots & 0 & p'_j - p'_i & 0 & \dots & 0 \\ & & & & & \vdots & & & \\ & & & & & & & & & & \end{pmatrix}$$

be the matrix that is obtained from the rigidity matrix  $\mathbf{R}(K_n, p)$  by replacing each  $p_i^{(j)} \in \mathbb{R}$  with the unknown variable  $(p_i^{(j)})'$ . We call  $\mathbf{R}(n, d)$  the  $d$ -dimensional **indeterminate rigidity matrix of  $K_n$** .

#### Def 3.2) Generic framework

Let  $V = \{v_1, \dots, v_n\}$  and  $p : V \rightarrow \mathbb{R}^d$  be a map.

We say that  $p$  is **generic realization** if the determinant of any submatrix of  $\mathbf{R}(K_n, p)$  is zero only if the determinant of the corresponding submatrix of  $\mathbf{R}(n, d)$  is (identically) zero.

A framework  $(G, p)$  is said to be **generic** if  $p$  is a generic map.

#### Lem 3.3) Rank of generic frameworks

Let  $p : V(G) \rightarrow \mathbb{R}^d$  be a generic map. Then, the rank of the rigidity matrix  $\mathbf{R}(G, p)$  is equal to the maximum possible rank attainable by any realization of  $G$  in  $\mathbb{R}^d$ . That is,

$$\text{rank}(\mathbf{R}(G, p)) = \max_{q \in \mathbb{R}^{dn}} \text{rank}(\mathbf{R}(G, q))$$

## Proof

Let  $k = \max_{q \in \mathbb{R}^{dn}} \text{rank}(R(G, q))$  be the maximum rank.

By the definition of matrix rank, there exists at least one realization  $q_0 \in \mathbb{R}^{dn}$  s.t.  $\mathbf{R}(G, q_0)$  contains a  $k \times k$  submatrix  $S$  with a non-zero determinant, i.e.  $\det S(q_0) \neq 0$ .

The determinant  $\det S$  is a polynomial in the coordinates of the joints. Since there exists a point  $q_0$  where  $\det S(q_0) \neq 0$ , this polynomial is not identically zero.

IAW [definition 3.2](#), a map  $p$  is generic if the determinant of any submatrix is zero only if the corresponding determinant polynomial is identically zero.

Since our  $k \times k$  determinant polynomial is not identically zero, it must be non-zero for any generic map  $p$ . Thus,  $\text{rank}(\mathbf{R}(G, p)) \geq k$ .

As  $k$  is defined as the maximum possible rank, we can conclude that  $\text{rank}(\mathbf{R}(G, p)) = k$ . □

Then, there is a result directly follows.

### Cor) Genericity implies Regularity

Every generic realization  $p$  of a graph  $G$  in  $\mathbb{R}^d$  is a regular point.

#### Example (Genericity of triangle frameworks)

We want to determine whether frameworks  $(G, p), (G, q)$  in  $\mathbb{R}^2$  is generic where

$$V(G) = \{v_1, v_2, v_3\}$$

with distinct realizations  $p, q$  where

$$p_1 = (0, 0), p_2 = (1, 0), p_3 = (2, 0) \quad \text{and} \quad q_1 = (0, 0), q_2 = (1, 0), q_3 = (0, 1)$$

Before algebraically determining them, we can guess.

1.  $p$  seems NOT generic because three points are colinear.

2.  $q$  seems generic because three points are not colinear.

Let's check.

First, construct an indeterminate rigidity matrix with  $n = 3$ .

$$\mathbf{R}(3, 2) = \left( \begin{array}{cc|cc|cc} x_1 - x_2 & y_1 - y_2 & x_2 - x_1 & y_2 - y_1 & 0 & 0 \\ x_1 - x_3 & y_1 - y_3 & 0 & 0 & x_3 - x_1 & y_3 - y_1 \\ \hline 0 & 0 & x_2 - x_3 & y_2 - y_3 & x_3 - x_2 & y_3 - y_2 \end{array} \right)$$

and construct  $\mathbf{R}(K_3, p)$ .

$$\mathbf{R}(K_3, p) = \left( \begin{array}{cc|cc|cc} 0 - 1 & 0 - 0 & 1 - 0 & 0 - 0 & 0 & 0 \\ 0 - 2 & 0 - 0 & 0 & 0 & 2 - 0 & 0 - 0 \\ \hline 0 & 0 & 1 - 2 & 0 - 0 & 2 - 1 & 0 - 0 \end{array} \right) = \left( \begin{array}{cc|cc|cc} -1 & 0 & 1 & 0 & 0 & 0 \\ -2 & 0 & 0 & 0 & 2 & 0 \\ \hline 0 & 0 & -1 & 0 & 1 & 0 \end{array} \right)$$

Taking submatrix  $\mathbf{R}_{(145)}$  of  $\mathbf{R}(K_3, p)$  by taking column 1, 4 and 5. Then,

$$\det \mathbf{R}_{(145)} = \det \begin{pmatrix} -1 & 0 & 0 \\ -2 & 0 & 2 \\ 0 & 0 & 1 \end{pmatrix} = 0$$

However, a determinant of corresponding submatrix  $\mathbf{R}(3, 2)_{145}$  is

$$\begin{aligned} \det \mathbf{R}(3, 2)_{(145)} &= \det \begin{pmatrix} x_1 - x_2 & y_2 - y_1 & 0 \\ x_1 - x_3 & 0 & x_3 - x_1 \\ 0 & y_2 - y_3 & x_3 - x_2 \end{pmatrix} \\ &= (x_1 - x_3)[x_1(y_2 - y_3) + x_2(y_3 - y_1) + x_3(y_1 - y_2)] \end{aligned}$$

This is not necessarily zero (may not be zero by choosing coordinates properly) and hence,  $p$  is **NOT a generic realization.** ( $\rightarrow$  We correctly guessed)

Now, again, to determine whether  $q$  is generic or not, construct  $\mathbf{R}(K_3, q)$ .

$$\mathbf{R}(K_3, q) = \left( \begin{array}{ccc|ccc|cc} 0-1 & 0-0 & | & 1-0 & 0-0 & | & 0 & 0 \\ 0-0 & 0-1 & | & 0 & 0 & | & 0-0 & 1-0 \\ \hline 0 & 0 & | & 1-0 & 0-1 & | & 0-1 & 1-0 \end{array} \right) = \left( \begin{array}{cc|cc|cc} -1 & 0 & | & 1 & 0 & | & 0 & 0 \\ 0 & -1 & | & 0 & 0 & | & 0 & 1 \\ \hline 0 & 0 & | & 1 & -1 & | & -1 & 1 \end{array} \right)$$

Taking submatrix  $\mathbf{R}_{(145)}$  of  $\mathbf{R}(K_3, q)$  by taking column 1, 4 and 5. Then,

$$\det \mathbf{R}_{(145)} = \det \begin{pmatrix} -1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & -1 & -1 \end{pmatrix} = 0$$

However, a determinant of corresponding submatrix  $\mathbf{R}(3, 2)_{145}$  is

$$\begin{aligned} \det \mathbf{R}(3, 2)_{(145)} &= \det \begin{pmatrix} x_1 - x_2 & y_2 - y_1 & 0 \\ x_1 - x_3 & 0 & x_3 - x_1 \\ 0 & y_2 - y_3 & x_3 - x_2 \end{pmatrix} \\ &= (x_1 - x_3)[x_1(y_2 - y_3) + x_2(y_3 - y_1) + x_3(y_1 - y_2)] \end{aligned}$$

This is not necessarily zero (may not be zero by choosing coordinates properly) and hence,  $q$  is **NOT a generic realization.** ( $\rightarrow$  Our guess was Wrong!)

As anticipated, the realization  $p$  with collinear points proved to be non-generic. However, realization  $q$  - which we initially presumed to be generic because it is a general framework -was also found to be non-generic upon verification.

This leads us to an important question: **What is the fundamental distinction between being 'general' and 'generic'?**

### Remark) Generic is stronger than general

Generic frameworks are general. But the converse doesn't hold.

There are some fundamental facts regarding this definition of generic.

### Thm 3.3) Set theoretic properties of non-generic realizations

The set of all non-generic realizations  $p : V \rightarrow \mathbb{R}^d$  where  $V = \{v_1, \dots, v_n\}$  is a **closed set of measure zero**.

### Thm 3.4) Combinatorial Rigidity Theorem

For a fixed dimension  $d$  and a graph  $G$  with  $n(G) \geq d$ , TFAE:

1.  $(G, p)$  is infinitesimally rigid (or independent, or isostatic) for at least one injective map  $p : V(G) \rightarrow \mathbb{R}^d$ .
2.  $(G, p)$  is infinitesimally rigid (or independent, or isostatic) for every  $d$ -dimensional generic realization  $p$ .

Intuitively, success in a single normal configuration implies success in almost every possible configuration.

The three structural properties (**infinitesimal rigidity**, **independence**, and **isostaticity**) only depend on the **underlying graph**  $G$ ; specifically its **combinatorial structure**.

#### Proof

(1)  $\Rightarrow$  (2): Suppose that  $(G, p)$  is infinitesimally rigid for an injective  $p : V(G) \rightarrow \mathbb{R}^d$  (not necessarily generic).

By [theorem 2.16](#), for infinitesimally rigid framework,

$$\text{rank}(\mathbf{R}(G, p)) = dn - \binom{d+1}{2}$$

where  $n(G) = n$ .

Let  $\hat{p}$  be a  $d$ -dimensional generic framework for a graph  $G$ . By [lemma 3.3](#), the rigidity matrix of  $\hat{p}$  has the maximum rank and hence  $\text{rank}(\mathbf{R}(G, \hat{p})) \geq \text{rank}(\mathbf{R}(G, p))$ .

Again, by [theorem 2.16](#), a rank of rigidity matrix of  $d$ -dimensional framework cannot exceed  $dn - \binom{d+1}{2}$  and hence  $\text{rank}(\mathbf{R}(G, \hat{p})) \leq dn - \binom{d+1}{2}$ .

Combining the equation and inequalities, we obtain

$$dn - \binom{d+1}{2} = \text{rank}(\mathbf{R}(G, p)) \leq \text{rank}(\mathbf{R}(G, \hat{p})) \leq dn - \binom{d+1}{2}$$

Thus,  $\text{rank}(\mathbf{R}(G, \hat{p})) = dn - \binom{d+1}{2}$  and hence we can conclude that  $(G, \hat{p})$  is infinitesimally rigid by [theorem 2.16](#).

(2)  $\Rightarrow$  (1): Suppose that  $(G, p)$  is infinitesimally rigid for every generic realization  $p$ .

**Existence of  $p$ :** Since the set of generic realizations is open dense set in  $\mathbb{R}^{dn}$ , there exists a map  $p : V(G) \rightarrow \mathbb{R}^d$  s.t.  $(G, p)$  is infinitesimally rigid.

**Injectivity:** Also, since generic framework is general, the vertices never overlap and hence  $p$  must be injective.

Thus, for any generic realization  $p$ , it is an injective realization that satisfies  $(G, p)$  is infinitesimally rigid.

This proof can be applied directly for independence or isostaticity instead of infinitesimal rigidity.  $\square$

### Thm 3.5) Equivalence of rigidity

For generic frameworks, infinitesimal rigidity is equivalent to rigidity.

### Def 3.6) Generically rigid/independent/isostatic

A graph  $G$  is called **generically rigid/independent/isostatic in dimension  $d$**  or generically  $d$ -rigid/independent/isostatic if  $d$ -dimensional generic realizations of  $G$  are infinitesimally rigid/independent/isostatic.

Here, by [theorem 3.5](#), A graph  $G$  is generically rigid if a generic realizations of  $G$  is rigid ( $\because$  infinitesimally rigid  $\equiv$  rigid for generic).

**Note:** Generically  $d$ -isostatic graphs can be constructed using inductive construction techniques. The following definition 3.7 and theorem 3.8 describe about it.

Introduce a definition first.

### Def 3.7) Vertex addition

Let  $G$  be a graph,  $U \subseteq V(G)$  with  $|U| = d$  and  $v \notin V(G)$ . Then the graph  $\hat{G}$  with  $V(\hat{G}) = V(G) \cup \{v\}$  and  $E(\hat{G}) = E(G) \cup \{\{v, u\} : u \in U\}$  is called **a vertex  $d$ -addition** (by  $v$ ) of  $G$ .

Intuitively, a vertex  $d$ -addition is the process of expanding a graph by introducing **one new vertex** and connecting it to exactly  $d$  **distinct existing vertices** via  $d$  new edges.

This operation serves as a fundamental building block for inductive constructions because it preserves the structural properties of the graph, as stated in the following theorem.

### Thm 3.8) Vertex Addition Theorem

A vertex  $d$ -addition of a generically  $d$ -isostatic graph is generically  $d$ -isostatic. Conversely, deleting a vertex of degree  $d$  from a generically  $d$ -isostatic graph results in a generically  $d$ -isostatic graph.

**Proof)** By [definition 2.19](#), to show that a vertex  $d$ -addition of a generically  $d$ -isostatic graph is generically  $d$ -isostatic, we need to show that it is infinitesimally rigid and independent.

First, show the infinitesimal rigidity.

Now, show the independence.

### Def 3.9) Edge split

Let  $G$  be a graph, and let  $U \subseteq V(G)$  with  $|U| = d + 1$  and  $\{u_1, u_2\} \in E(G)$  for some  $u_1, u_2 \in U$ . Further, let  $v \notin V(G)$ . Then, the graph  $\hat{G}$  with  $V(\hat{G}) = V(G) \cup \{v\}$  and  $E(\hat{G}) = (E(G) \setminus \{u_1 u_2\}) \cup \{vu : u \in U\}$  is called an **edge  $d$ -split (on  $u_1, u_2; v$ )** of  $G$ .