

Rigidity Theory

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1. Preliminaries

1.1. Basic definitions

Def 1.1) Framework

A **framework** in \mathbb{R}^d is a pair (G, p) , where G is a finite simple graph and $p : V(G) \rightarrow \mathbb{R}^d$ is a (not necessarily injective) map. We also say that p is a d -dimensional **realization** of G . Sometimes, we denote $f_p = (G, p)$.

Def 1.2) Joint, Bar

For $v_i \in V(G)$, we say that $p_i := p(v_i)$ is the **joint** of (G, p) corresponding to v_i .
Also, for $v_i v_j \in E(G)$, we say that $\{p_i, p_j\}$ is the **bar** of (G, p) corresponding to $v_i v_j$.
The **length of the bar** $\{p_i, p_j\}$ is defined as $L(\{p_i, p_j\}) := \|p(v_i) - p(v_j)\|_2$.

Def 1.3) Motion, rigid and flexible framework

A (**finite**) **motion** of (G, p) is an indexed family $\{P_i\}_{i \in [n]}$ of functions $P_i : [0, 1] \rightarrow \mathbb{R}^d$ s.t.

1. (Initial Condition) $P_i(0) = p(v_i)$, $\forall i \in [n]$.
2. (Differentiability) $P_i(t)$ is differentiable on $[0, 1]$, $\forall i \in [n]$.
3. (Distance-preserving) $\|P_i(t) - P_j(t)\| = \|p_i - p_j\|$ for all $t \in [0, 1]$ and $v_i v_j \in E(G)$.

A motion $\{P_i\}$ is called **rigid** if $\|P_i(t) - P_j(t)\| = \|p_i - p_j\|$ for all $t \in [0, 1]$ and $(i, j) \in [n]^2$.

Also, $\{P_i\}$ is called **flexible** if $\exists (i, j) \in [n]^2 \setminus E(G)$ s.t. $\|P_i(t) - P_j(t)\| \neq \|p_i - p_j\|$ for all $t \in (0, 1]$.

With these definitions, (G, p) is called rigid if there is no flexible motion. Otherwise, (G, p) is called flexible.

Q. Why "for all $t \in (0, 1]$ " for a framework to be flexible?

1.2. Equivalence & Congruence

So far, we defined a framework with a tuple of a simple graph and a realization. However, for a formation system(engineering aspect), the set of edges $E(G)$ could be defined by a set of constraints. Here, the constraint means a task or a goal that agents work together to achieve. (e.g. communication network, distance-preserving, ...)

So, we can alternatively describe a formation by $f_p = (\mathbf{V}, \mathbf{C}, p)$ where $\mathbf{V} = [n]$ and \mathbf{C} be a set of constraints. Obviously, $|\mathbf{C}| = e(G)$.

In this aspect, we can introduce more definitions.

Def 1.4) Equivalence / Congruence of frameworks

Two frameworks $f_p = (G, p) = (\mathbf{V}, \mathbf{C}, p)$ and $f_q = (\mathbf{V}, \mathbf{C}, q) = (G, q)$ with different realizations p, q , are called...

- **Equivalent** if $\text{mag}[c_i(p(x_i))] = \text{mag}[c_i(q(x_i))]$ for all $c_i \in \mathbf{C}$ and denoted by $f_p \sim f_q$ where $x_i \subset \mathbf{V}$ is the set of nodes involved in the i^{th} constraint c_i . In this context, $p(x_i)$ and $q(x_i)$ represent the realized points(coordinates) of these nodes in each framework, and $\text{mag}[\cdot]$ is the exact value of the argument.
- **Congruent** if $\text{mag}[c_i(p(x_i))] = \text{mag}(c_i(q(x_i)))$ for all $c_i \in \mathbf{C}_K$ and denoted by $f_p \equiv f_q$ where \mathbf{C}_K is the complete constraint set containing constraints for every possible pair of vertices in the graph G . This corresponds to the set of all edges in the complete graph K_n . In this case, x_i represents the set of nodes involved in each constraint c_i within the expanded set \mathbf{C}_K .

For example, if the constraint is given as a distance between each vertices, then $c_i \in \mathbf{C}$ is the distance constraint on the i^{th} edge and x_i is the set of endpoints of the i^{th} edge.

So, if $uv \in E(G)$ is the i^{th} edge, then $x_i = \{u, v\}$ and $c_i(p(x_i)) = \|p(u) - p(v)\|$. In this sense, we can describe the equivalence and congruence for distance constraint only.

Two frameworks f_p and f_q are called...

- **Distance equivalent** if $\|p_i - p_j\| = \|q_i - q_j\|$ for all $v_i v_j \in E(G)$
- **Distance congruent** if $\|p_i - p_j\| = \|q_i - q_j\|$ for all $(i, j) \in [n]^2$.

Def 1.5) Constraint function

For a framework $f_p = (G, p) = (\mathbf{V}, \mathbf{C}, p)$, the constraint function $h_G : \mathbb{R}^{dn} \rightarrow \mathbb{R}^m$ ($m = e(G)$) is defined by

$$h_G(p) := \begin{bmatrix} \text{mag}[c_1(p(x_1))] \\ \vdots \\ \text{mag}[c_m(p(x_m))] \end{bmatrix}$$

for square of distance constraint,

$$h_G(p) = \begin{bmatrix} \text{mag}[c_1(p(x_1))] \\ \vdots \\ \text{mag}[c_m(p(x_m))] \end{bmatrix} = \begin{bmatrix} \|p_{j_1} - p_{i_1}\|^2 \\ \vdots \\ \|p_{j_m} - p_{i_m}\|^2 \end{bmatrix}$$

2. Rigidity and infinitesimal rigidity

2.1. Basic rigidity concepts

Q. What is the "Rigid framework" ?

Def 2.1) Rigidity (General definition)

A framework $f_p = (G, p)$ is called **rigid** in a given Euclidean space \mathbb{R}^d if there exists a nbd. U_p of p s.t. every framework $f_q = (G, q)$ with $q \in U_p$ satisfy $f_p \sim f_q \Rightarrow f_p \equiv f_q$.

Algebraically, $f_p = (G, p)$ is rigid in \mathbb{R}^d if there exists a nbd. U_p of p s.t.

$$h_G^{-1}(h_G(p)) \cap U_p = h_K^{-1}(h_K(p)) \cap U_p$$

Here,

- LHS = the set of frameworks in U_p satisfying the constraint C .
- RHS = the set of frameworks in U_p satisfying the constraint C_K .

Obviously, there is only one congruent framework satisfying C_K . Also, since $C \subset C_K$, $LHS \supset RHS$. So, this expression means that, if $LHS \neq RHS$, then there is at least one framework that is distinct (non-congruent) from that of in RHS and hence f_p is not rigid.

Def 2.2) Distance rigidity

A framework f_p is called **distance rigid** in \mathbb{R}^d if there exists $\varepsilon > 0$ s.t. for any framework f_q that is locally perturbed from f_p (i.e. $q_i \in B_\varepsilon(p_i), \forall i$), f_p and f_q are distance equivalent only if f_p and f_q are distance congruent.

This means that f_p is distance rigid if every framework satisfying same edge-distance constraints to f_p near p is congruent to f_p (Locally unique embedding satisfying edge-distance constraint).

Q. Are there two definition of distance rigidity? \rightarrow No.

Thm 2.3) Equivalent of definitions of rigidity

Two definitions of distance rigidity (in the view of motion and set-theoretic definition) of framework is equivalent.

Pf)

Def 2.4) Global distance rigidity

A framework (G, p) is called **globally distance rigid** in \mathbb{R}^d if "every" framework (G, q) equivalent to (G, p) is also congruent to (G, p) in \mathbb{R}^d .

Henceforth, unless otherwise specified, the term 'rigidity' will refer exclusively to '**distance rigidity**'.

2.2. Infinitesimal rigidity

Motivation: It is in general very difficult to determine whether a given framework is rigid or not since it requires solving a system of quadratic equations. It is therefore common to **linearize** the problem by differentiating the equations.

Let $\{P_i\}$ be a finite motion of a 2-dimensional framework (G, p) and let $\{p_i, p_j\}$ be a bar of length l . Then,

$$\|P_i(t) - P_j(t)\|^2 = \left[P_i^{(1)}(t) - P_j^{(1)}(t)\right]^2 + \left[P_i^{(2)}(t) - P_j^{(2)}(t)\right]^2 = l^2 \quad \forall t \in [0, 1]$$

where $P_i^{(k)}(t)$ is the k^{th} coordinate of p_i at t .

By differentiating both sides w.r.t. t ,

$$\left(P_1^{(1)}(t) - P_j^{(1)}(t), P_1^{(2)}(t) - P_j^{(2)}(t)\right) \cdot \left(\left[P_1^{(1)}(t) - P_j^{(1)}(t)\right]', \left[P_1^{(2)}(t) - P_j^{(2)}(t)\right]'\right) = 0$$

Evaluating at $t = 0$, we obtain

$$\left(p_i^{(1)} - p_j^{(1)}, p_i^{(2)} - p_j^{(2)}\right) \cdot \left(u_i^{(1)} - u_j^{(1)}, u_i^{(2)} - u_j^{(2)}\right) = 0$$

or in short $(p_i - p_j) \cdot (u_i - u_j)$. Here, u_i, u_j are initial velocity vectors of p_i and p_j , respectively. So, this implies that the difference between initial velocity vectors is orthogonal to the line through p_i and p_j .

These vectors preserve the length of $\{p_i, p_j\}$ at first order.

Def 2.5) Constricted framework

A d -dimensional framework (G, p) is called **constricted** if all points in $\mathcal{P} = \{p_1, \dots, p_n\}$ lie in $k \geq d - 1$ dimensional affine subspace. Otherwise, (G, p) is called **normal**.

Thus, a framework is constricted if it is geometrically degenerated.

Def 2.6) Infinitesimal motion / Infinitesimal rigidity

An **infinitesimal motion** of a framework (G, p) is a function $u : V(G) \rightarrow \mathbb{R}^d$ s.t.

$$(p_i - p_j) \cdot (u_i - u_j) = 0 \quad \forall v_i v_j \in E(G)$$

That is, for a distance constrained graph G , a function consists of dn instantaneous velocity vectors of the nodes that do not change the equivalence conditions between nodes.

Here, an infinitesimal motion is called **rigid** if

$$(p_i - p_j) \cdot (u_i - u_j) = 0 \quad \forall (i, j) \in [n]^2$$

Also, an infinitesimal motion is called **flex** if

$$(p_i - p_j) \cdot (u_i - u_j) \neq 0 \quad \exists v_i v_j \notin E(G)$$

In this sense, (G, p) is called **infinitesimally rigid** if every infinitesimal motion of (G, p) is an infinitesimal rigid motion. Otherwise, (G, p) is said to be **infinitesimally flexible**.

Note: In some context, infinitesimal rigid motion also can be called by a "Trivial motion" or a "Special Euclidean motion".

- **Trivial motions** is an infinitesimal motions include only the translational and rotational motions. (infinitesimal rigid motion)
- **Special Euclidean motion** is defined by the motions of nodes that keep the congruence conditions among the nodes. It contains the translations and rotations, and hence Special Euclidean motion is equivalent to the **direct isometries**.

In short,

Equivalence of definitions

- Infinitesimal rigid motion = Trivial motion
- Special Euclidean motion = Direct isometry

However, **in the infinitesimal viewpoint, the four of them are all equivalent.**

Remark: All constricted frameworks infinitesimally flexible.

- Now, we can come up with a very natural question.

Q. Is there any relationship between the rigidity and the infinitesimal rigidity?

Thm 2.7) Infinitesimal rigidity is stronger than rigidity

If a framework f_p is infinitesimally rigid, then f_p is rigid.

Proof)

2.2.1. Rigidity Matrix

Def 2.8) Edge function / Rigidity Matrix

Let (G, p) be a framework with graph G and realization p . Then, the **edge function** $g_G : \mathbb{R}^{dn} \rightarrow \mathbb{R}^m$ where $m = e(G)$ is defined by

$$g_G(p) = \frac{1}{2} \left[\|p_i - p_j\|^2 \right]_{(i,j) \in E(G)}$$

Then, the **rigidity matrix** is defined by

$$\mathbf{R}(G, p) = \frac{\partial g_G(p)}{\partial p} \in \mathbb{R}^{m \times dn}$$

Here, each row and column corresponds to each edges and nodes of the framework, respectively.

Explicitly, the rigidity matrix is expressed as

$$\mathbf{R}(G, p) = \begin{pmatrix} & & & & \vdots & & & & \\ 0 & \dots & 0 & p_i - p_j & 0 & \dots & 0 & p_j - p_i & 0 & \dots & 0 \\ & & & & \vdots & & & & & & \end{pmatrix}$$

Let $e_k = v_i v_j \in E(G)$ is the k^{th} edge with endpoints v_i and v_j . Then, it occupies the $d(i-1) + 1^{\text{th}}$ to $d i^{\text{th}}$ columns and $d(j-1) + 1^{\text{th}}$ to $d j^{\text{th}}$ columns of the k^{th} row.

Here, $(p_i^{(1)} - p_j^{(1)}), \dots, (p_i^{(d)} - p_j^{(d)})$ are in the $d(i-1) + 1^{\text{th}}$ to $d i^{\text{th}}$ columns of the k^{th} row.

Similarly, $(p_j^{(1)} - p_i^{(1)}), \dots, (p_j^{(d)} - p_i^{(d)})$ are in the $d(j-1) + 1^{\text{th}}$ to $d j^{\text{th}}$ columns of the k^{th} row.

Example (Tetrahedron)

For a formation (G, p) of regular tetrahedron in \mathbb{R}^3 ,

$V(G) = \{v_1, v_2, v_3, v_4\}$ and $E(G) = \{v_1v_2, v_1v_3, v_1v_4, v_2v_3, v_2v_4, v_3v_4\}$ and $p = [p_1^\top p_2^\top p_3^\top p_4^\top]^\top$ where $p_i = [x_i \ y_i \ z_i]$.

Note that $|V(G)| = n(G) = 4$, $|E(G)| = e(G) = 6$, $d = 3$.

Then, by definition, the rigidity matrix be

$$\mathbf{R}(p) = \begin{bmatrix} (p_1 - p_2)^\top & (p_2 - p_1)^\top & 0^\top & 0^\top \\ (p_1 - p_3)^\top & 0^\top & (p_3 - p_1)^\top & 0^\top \\ (p_1 - p_4)^\top & 0^\top & 0^\top & (p_4 - p_1)^\top \\ 0^\top & (p_2 - p_3)^\top & (p_3 - p_2)^\top & 0^\top \\ 0^\top & (p_2 - p_4)^\top & 0^\top & (p_4 - p_2)^\top \\ 0^\top & 0^\top & (p_3 - p_4)^\top & (p_4 - p_3)^\top \end{bmatrix} \in \mathbb{R}^{6 \times (3 \times 4)}$$

Let $p_1 = (1, 1, 1)$, $p_2 = (1, -1, -1)$, $p_3 = (-1, 1, -1)$, $p_4 = (-1, -1, 1)$ be coordinates of each vertex. Then, the rigidity matrix be

$$\mathbf{R}(p) = \begin{bmatrix} 0 & 2 & 2 & 0 & -2 & -2 & 0 & 0 & 0 & 0 & 0 & 0 \\ 2 & 0 & 2 & 0 & 0 & 0 & -2 & 0 & -2 & 0 & 0 & 0 \\ 2 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & -2 & -2 & 0 \\ 0 & 0 & 0 & 2 & -2 & 0 & -2 & 2 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 2 & 0 & -2 & 0 & 0 & 0 & -2 & 0 & 2 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 2 & -2 & 0 & -2 & 2 \end{bmatrix}$$

Note that infinitesimal motion is a function of velocity vectors of each nodes. So, there are d velocity components for each nodes and hence one infinitesimal motion can be identified as a dn -dimensional column vector \mathbf{u} .

- Recall: $u : V(G) \rightarrow \mathbb{R}^d$ is an infinitesimal motion if $(p_i - p_j) \cdot (u_i - u_j) = 0$ for all $v_i v_j \in E(G)$.

From here, we can find some useful property.

Prop 2.9) Infinitesimal motion with rigidity matrix

The set of infinitesimal motion of framework (G, p) is the kernel of $\mathbf{R}(G, p)$.

Proof

Proof) Let $\mathcal{M}(G, p)$ be the set of infinitesimal motion of (G, p) and $\mathbf{u} \in \mathcal{M}(G, p)$.

Let $e_k = v_i v_j$. Then, the k^{th} row of $\mathbf{R}(G, p)\mathbf{u}$ be

$$[\mathbf{R}(G, p)\mathbf{u}]_k = (p_i - p_j) \cdot u_i + (p_j - p_i) \cdot u_j = (p_i - p_j) \cdot (u_i - u_j)$$

Here, since \mathbf{u} is a column vector representation of an infinitesimal motion, $\text{RHS} = 0$ and hence $\mathbf{u} \in \text{Ker}(\mathbf{R}(G, p))$.

Thus, $\mathcal{M}(G, p) \subset \text{Ker}(\mathbf{R}(G, p))$.

Conversely, given $\mathbf{x} \in \text{Ker}(\mathbf{R}(G, p))$, $[\mathbf{R}(G, p)\mathbf{x}]_k = 0$ for each k and hence $(p_i - p_j) \cdot (x_i - x_j) = 0$.

Thus, \mathbf{x} be an infinitesimal motion of (G, p) and hence $\text{Ker}(\mathbf{R}(G, p)) \subset \mathcal{M}(G, p)$.

$\therefore \mathcal{M}(G, p) = \text{Ker}(\mathbf{R}(G, p))$. □

From this fact, by letting $\mathcal{R}(G, p)$ be the set of all infinitesimal rigid motion,

$$\mathcal{R}(G, p) \subset \mathcal{M}(G, p)$$

by definition.

Also, (G, p) is infinitesimally rigid if and only if $\mathcal{R}(G, p) = \mathcal{M}(G, p)$.

By the previous proposition, the following fact directly follows.

Lem 2.10) Dimension and infinitesimal rigidity

Let (G, p) be a normal framework in \mathbb{R}^d . Then,

$$\dim \mathcal{R}(G, p) \leq \dim \mathcal{M}(G, p) = \text{nullity}(\mathbf{R}(G, p))$$

Equivalently, by rank-nullity theorem,

$$\text{rank}(\mathbf{R}(G, p)) \leq dn - \dim \mathcal{R}(G, p)$$

The equality holds if and only if (G, p) is infinitesimally rigid.

Here, while we know d, n but we don't know about $\dim \mathcal{R}(G, p)$ and hence we need to find it.

To find it, we assume **2-dimensional general framework** for the following arguments. They can be directly generalized to an arbitrary dimension.

Def 2.11) General Position

$\mathcal{P} = \{v_1, \dots, v_n\} \subseteq \mathbb{R}^d$ is in **general position** if for $1 \leq m \leq d$, no $m + 1$ points in \mathcal{P} lie in an $(m - 1)$ -dimensional affine subspace of \mathbb{R}^d . A framework whose joints are in general position is called a **general framework**.

The following definition explains about two types of infinitesimal motions.

Def 2.12) Infinitesimal translation / Infinitesimal rotation (\mathbb{R}^2 ver.)

An **infinitesimal translation of the plane** is of the form

$$\tau_{(u,v)} : \mathbb{R}^2 \rightarrow \mathbb{R}^2 ; \quad \tau_{(u,v)}(x, y) = (x + u, y + v)$$

This implies that every point (x, y) in \mathbb{R}^2 moves with same velocity (u, v) .

An **infinitesimal rotation of the plane** is of the form

$$a\rho_{(x_0, y_0)} : \mathbb{R}^2 \rightarrow \mathbb{R}^2 ; \quad a\rho_{(x_0, y_0)}(x, y) = a(y_0 - y, x - x_0)$$

Let r be a radius vector from a center (x_0, y_0) to (x, y) , i.e. $r = (x - x_0, y - y_0)$. Then, the rotation vector v of (x, y) must satisfy $r \perp v$ and hence $r \cdot v = 0$. So, v must be of the form $a(-(y - y_0), x - x_0) = a(y_0 - y, x - x_0)$ where $a \in \mathbb{R}$.

Lem 2.13) Dimension of infinitesimal translations and rotations

The set of all infinitesimal translations and rotations of \mathbb{R}^2 forms a 3-dimensional vector space.

Proof)

Now, recall that for infinitesimal viewpoint, infinitesimal rigid motion is equal to direct isometry and hence **infinitesimal rigid motion consists of infinitesimal translation and rotation only.**

The following lemma describes about this fact.

Lem 2.14) Classifications fo infinitesimal rigid motions

Let (G, p) be a 2-dimensional general framework with $V(G) = \{v_1, \dots, v_n\}$, $n \geq 2$. Then $\mathcal{R}(G, p)$ is the 3-dimensional space of infinitesimal translations and rotations restricted to the points $\mathcal{P} = \{v_1, \dots, v_n\}$ of (G, p) .

Proof

The proof consists of two steps.

For a restriction map $(\cdot)|_{\mathcal{P}} : I'_d(\mathbb{R}^2) \rightarrow \mathcal{R}(G, p)$; $f \mapsto f|_{\mathcal{P}}$, (where $I'_d(\mathbb{R}^2)$ is a set of infinitesimal direct isometries, i.e. infinitesimal translations and rotations)

- **Surjectivity:** Given $u \in \mathcal{R}(G, p)$, $\exists \bar{u} \in I'_d(\mathbb{R}^2)$ s.t. $\bar{u}|_{\mathcal{P}} = u$.
- **Injectivity:** Given $\kappa, \lambda \in I'_d(\mathbb{R}^2)$ and corresponding restriction $\kappa|_{\mathcal{P}}, \lambda|_{\mathcal{P}} \in \mathcal{R}(G, p)$, $\kappa|_{\mathcal{P}} = \lambda|_{\mathcal{P}} \Rightarrow \kappa = \lambda$.

Surjectivity

Let $u = (u_1, \dots, u_n) \in \mathcal{R}(G, p)$ with $u_i = (u_i^{(1)}, u_i^{(2)})$ for all $i = 1, \dots, n$.

Clearly, the restriction of τ_{u_1} to \mathcal{P} is an infinitesimal rigid motion of (G, p) .

Hence $\hat{u} = u - \tau_{u_1}|_{\mathcal{P}} \in \mathcal{R}(G, p)$.

Here, note that $\hat{u} = (u_1 - u_1, u_2 - u_1, \dots, u_n - u_1)$ and hence $\hat{u}_1 = (0, 0)$. (fixing p_1)

Next, from the definition of infinitesimally rigid motion $(p_i - p_j) \cdot (u_i - u_j) = 0$ for all $(i, j) \in [n]^2$, we have

$$(p_2 - p_1) \cdot (\hat{u}_2 - \hat{u}_1) = (p_2 - p_1) \cdot \hat{u}_2 = 0$$

Since $(p_2 - p_1) \cdot \hat{u}_2 = (p_2^{(1)} - p_1^{(1)}, p_2^{(2)} - p_1^{(2)}) \cdot \hat{u}_2 = 0$,

\hat{u}_2 must be of the form $\hat{u}_2 = a (p_2^{(1)} - p_1^{(1)}, -(p_2^{(2)} - p_1^{(2)}))$.

So, $\hat{u}_2^{(1)} = a (p_2^{(1)} - p_1^{(1)})$ and $\hat{u}_2^{(2)} = a (p_1^{(2)} - p_2^{(2)})$ and hence

$$a = \frac{\hat{u}_2^{(1)}}{p_2^{(1)} - p_1^{(1)}} = \frac{\hat{u}_2^{(2)}}{p_1^{(2)} - p_2^{(2)}} \quad \text{where } p_1^{(1)} \neq p_2^{(1)} \text{ or } p_1^{(2)} \neq p_2^{(2)}$$

Now, define another motion $\bar{u} = \hat{u} - a\rho_{p_1}|_{\mathcal{P}} = u - \tau_{u_1}|_{\mathcal{P}} - a\rho_{p_1}|_{\mathcal{P}}$.

Here, since p_1 is fixed by \hat{u} and the rotation $a\rho_{p_1}|_{\mathcal{P}}$ is a rotation through p_1 , \bar{u} fixes p_1 as well.

Also, by choosing proper rotation, we can make \bar{u} fixes p_2 . Then, since \bar{u} fixes both p_1 and p_2 , $\bar{u}_1 = \bar{u}_2 = (0, 0)$.

Now, claim: **If p_1 and p_2 be fixed, then p_i are fixed for all $i > 2$.**

From the definition of infinitesimal rigid motion, we have

$$(p_i - p_1) \cdot (\bar{u}_i - \bar{u}_1) = (p_i - p_1) \cdot \bar{u}_i = 0 \quad \forall i > 2 \quad \dots (*)$$

$$(p_i - p_2) \cdot (\bar{u}_i - \bar{u}_2) = (p_i - p_2) \cdot \bar{u}_i = 0 \quad \forall i > 2 \quad \dots (**)$$

Since we assumed the general position, p_i is not on the line through p_1 and p_2 and hence the vectors $p_i - p_1$ and $p_i - p_2$ are not parallel.

However, since \bar{u}_i is perpendicular to both of them by (*) and (**).

Note that in the two-dimensional plane, the only vector that is perpendicular to two vectors that are not parallel to each other is the zero vector $(0, 0)$ and hence $\bar{u}_i = (0, 0)$ for all $i > 2$.

Thus, $\bar{u} = u - \tau_{u_1}|_{\mathcal{P}} - a\rho_{p_1}|_{\mathcal{P}} = \mathbf{0}$ and hence $u = \tau_{u_1}|_{\mathcal{P}} + a\rho_{p_1}|_{\mathcal{P}} = (\tau_{u_1} + a\rho_{p_1})|_{\mathcal{P}}$.

This implies that any infinitesimal rigid motion u of framework (G, p) is the restriction of combination of infinitesimal translation and rotation of \mathbb{R}^2 to \mathcal{P} .

Injectivity

Suppose that κ and λ are two infinitesimal rigid motions of \mathbb{R}^2 s.t.

$$\kappa|_{\mathcal{P}} = \lambda|_{\mathcal{P}} = u$$

Let $x \in \mathbb{R}^2$ that is not on the line through p_1 and p_2 .

Consider $\mu = \kappa - \lambda$. Then, obviously, μ is an infinitesimal rigid motion of \mathbb{R}^2 .

Then, by the definition of infinitesimal rigid motion,

$$(\mu(x) - \mu(p_1)) \cdot (x - p_1) = 0$$

$$(\mu(x) - \mu(p_2)) \cdot (x - p_2) = 0$$

Here, since $\kappa|_{\mathcal{P}} = \lambda|_{\mathcal{P}}$, for any joint p_i , $\kappa(p_i) = \lambda(p_i)$ and hence $\mu(p_i) = (0, 0)$ for all i .

So,

$$\mu(x) \cdot (x - p_1) = 0$$

$$\mu(x) \cdot (x - p_2) = 0$$

By assumption, x is not on the line through p_1 and p_2 and hence two vectors $x - p_1$ and $x - p_2$ are not parallel. Again, since the only vector that is perpendicular to two vectors that are not parallel to each other is the zero vector $(0, 0)$, we can find that $\mu(x) = (0, 0)$.

So far, we've shown that $\mu(x) = (0, 0)$ for $x \in \mathbb{R}^2 \setminus L(p_1, p_2)$ where $L(p_1, p_2)$ is the straight line through p_1 and p_2 .

Finally, let $y \in L(p_1, p_2)$. Then, obviously, y does not lie on a line through x and p_1 or x and p_2 (or both).

WLOG, let y is not on a line through x and p_1 . By the definition of infinitesimal rigid motion,

$$(\mu(y) - \mu(p_1)) \cdot (y - p_1) = 0$$

$$(\mu(y) - \mu(x)) \cdot (y - x) = 0$$

and since $\mu(p_1) = \mu(x) = 0$, $\mu(y)$ is perpendicular to both $y - p_1$ and $y - x$ that are not in parallel. Thus, $\mu(y) = (0, 0)$.

Therefore, we proved that $\mu(x) = 0$ for all $x \in \mathbb{R}^2$ and hence $\kappa = \lambda$.

Hence, proved. □

Using previous lemma, we can obtain an important result.

Before observing it, introduce a definition first.

Def 2.15) Degree of freedom

For a framework (G, p) , a **degree of freedom** of (G, p) is nullity $R(G, p)$.

Also, $\text{nullity}(\mathbf{R}(G, p)) - \dim \mathcal{R}(G, p)$ is called an **internal degrees of freedom** of (G, p) .

Then, we can observe a degree of freedom by the following theorem.

Thm 2.16) Lower bound of DOF

Let (G, p) be a general framework in \mathbb{R}^d with $n(G) \geq d$. Then,

$$\text{nullity}(\mathbf{R}(G, p)) \geq d + \binom{d}{2} = \binom{d+1}{2}$$

or equivalently,

$$\text{rank}(R(G, p)) \leq dn - \binom{d+1}{2}$$

Here, equality holds if and only if (G, p) is infinitesimally rigid.

Proof

Proof) Recall that $\dim \mathcal{R}(G, p) \leq \dim \mathcal{M}(G, p) = \text{nullity}(\mathbf{R}(G, p))$ where $\mathcal{R}(G, p)$ is the set of infinitesimally rigid motions of (G, p) and $\mathcal{M}(G, p)$ is the set of infinitesimal motions of (G, p) by [lemma 2.10](#).

By [lemma 2.14](#), a set of infinitesimal rigid motion consists of infinitesimal translations and rotations and hence $\dim \mathcal{R}(G, p)$ is the number of basis, i.e. the number of independent translations and rotations.

Note that there are d translations $\tau_{e_1}, \dots, \tau_{e_d}$ where e_i is a standard basis vector and there are $\binom{d}{2}$ rotations R_{ij} where for all $i, j \in [d]$, $R_{ij} : \mathcal{P} \rightarrow \mathbb{R}^d$; $R_{ij}(p_k) = p_k^{(i)} e_j - p_k^{(j)} e_i$ for all k .

So, the total number of basis is $d + \binom{d}{2} = \binom{d+1}{2}$ and hence

$$\binom{d+1}{2} \leq \text{nullity}(\mathbf{R}(G, p))$$

Hence, proved. □

Cor 2.17) Maximum number of edges of flexible framework

Let (G, p) be a general framework in \mathbb{R}^d with $n(G) \geq d$. If $e(G) < dn - \binom{d+1}{2}$, then (G, p) is infinitesimally flexible.

This corollary provides the lower bound of number of edges for a framework to be rigid. **The following definition explains about a framework that has exactly this minimum threshold of edges required for rigidity.**

Def 2.18) Minimally rigid framework

A rigid framework (G, p) is called **minimally rigid** if the removal of any edges makes it flexible. Similarly, an infinitesimally rigid framework (G, p) is called **minimally infinitesimally rigid** if the removal of any edges makes it not infinitesimally rigid.

Def 2.19) Independent / Isostatic framework

A framework (G, p) is **independent** if the rows of its rigidity matrix $\mathbf{R}(G, p)$ are linearly independent. (G, p) is **isostatic** if it is infinitesimally rigid and independent.

Thm 2.20) Equivalent conditions of isostatic framework

For a d -dimensional general framework (G, p) with $n(G) \geq d$, TFAE:

1. (G, p) is isostatic.
2. (G, p) is infinitesimally rigid and $e(G) = dn - \binom{d+1}{2}$.
3. (G, p) is independent and $e(G) = dn - \binom{d+1}{2}$.
4. (G, p) is minimal infinitesimally rigid.

Def 2.21) Regular points

Given a framework (G, p) , a realization p is called **regular** if it satisfies

$$\text{rank}(\mathbf{R}(G, p)) = \max_{p \in \mathbb{R}^{dn}} \text{rank}(\mathbf{R}(G, p))$$

While this definition might be confused with [General position](#), they represent fundamentally different perspectives on a framework:

General position is a purely geometric property of the point coordinates, whereas **Regularity** is an algebraic property concerning the rank of the rigidity matrix.

To break this down more precisely:

- **General position** focuses on the **geometry of the joints**; it requires that no $m + 1$ points lie in an $(m - 1)$ -dimensional affine subspace, a **condition that depends only on the point locations and is independent of the graph's edges.**
- **Regular realization** focuses on the **algebraic rank of the structure**; a realization p is regular if its rigidity matrix R_p attains the maximum possible rank k for that specific graph, meaning it **inherently**

depends on the graph's connectivity.

Therefore, a "general framework" is defined by its well-distributed geometry, while a "regular realization" is defined by its ability to reach the full structural potential (maximum rank) allowed by its specific graph topology.

3. Combinatorial aspects of rigidity

Motivation

Up to this point, our analysis has **focused on the rigidity of frameworks based on their specific geometric configurations**—that is, the exact coordinates of vertices in \mathbb{R}^d .

Now, a fundamental question arises: to what extent can these **properties be characterized solely by the underlying graph's connectivity**, independent of its specific realization?

So, we'll investigate the possibility of **treating rigidity as an intrinsic property** of the underlying graph, leading us to the study of generic frameworks where infinitesimal rigidity becomes purely combinatorial.

Introduce some definitions first.

Def 3.1) Indeterminate rigidity matrix

Let K_n be the complete graph on n vertices. For each $i = 1, 2, \dots, n$, we introduce a d -tuple of variables $p'_i = ((p_i^{(1)})', \dots, (p_i^{(d)})')$ (d unknown variables).

Then, let

$$\mathbf{R}(n, d) = \begin{pmatrix} 0 & \dots & 0 & p'_i - p'_j & 0 & \dots & 0 & p'_j - p'_i & 0 & \dots & 0 \\ \vdots & & & & & & & & & & \\ \vdots & & & & & & & & & & \end{pmatrix}$$

be the matrix that is obtained from the rigidity matrix $\mathbf{R}(K_n, p)$ by replacing each $p_i^{(j)} \in \mathbb{R}$ with the unknown variable $(p_i^{(j)})'$. We call $\mathbf{R}(n, d)$ the d -dimensional **indeterminate rigidity matrix of K_n** .

Def 3.2) Generic framework

Let $V = \{v_1, \dots, v_n\}$ and $p : V \rightarrow \mathbb{R}^d$ be a map.

We say that p is **generic realization** if the determinant of any submatrix of $\mathbf{R}(K_n, p)$ is zero only if the determinant of the corresponding submatrix of $\mathbf{R}(n, d)$ is (identically) zero.

A framework (G, p) is said to be **generic** if p is a generic map.

Lem 3.3) Rank of generic frameworks

Let $p : V(G) \rightarrow \mathbb{R}^d$ be a generic map. Then, the rank of the rigidity matrix $\mathbf{R}(G, p)$ is equal to the maximum possible rank attainable by any realization of G in \mathbb{R}^d . That is,

$$\text{rank}(\mathbf{R}(G, p)) = \max_{q \in \mathbb{R}^{dn}} \text{rank}(\mathbf{R}(G, q))$$

Proof

Let $k = \max_{q \in \mathbb{R}^{dn}} \text{rank}(R(G, q))$ be the maximum rank.

By the definition of matrix rank, there exists at least one realization $q_0 \in \mathbb{R}^{dn}$ s.t. $\mathbf{R}(G, q_0)$ contains a $k \times k$ submatrix S with a non-zero determinant, i.e. $\det S(q_0) \neq 0$.

The determinant $\det S$ is a polynomial in the coordinates of the joints. Since there exists a point q_0 where $\det S(q_0) \neq 0$, this polynomial is not identically zero.

IAW [definition 3.2](#), a map p is generic if the determinant of any submatrix is zero only if the corresponding determinant polynomial is identically zero.

Since our $k \times k$ determinant polynomial is not identically zero, it must be non-zero for any generic map p . Thus, $\text{rank}(\mathbf{R}(G, p)) \geq k$.

As k is defined as the maximum possible rank, we can conclude that $\text{rank}(\mathbf{R}(G, p)) = k$. □

Then, there is a result directly follows.

Cor) Genericity implies Regularity

Every generic realization p of a graph G in \mathbb{R}^d is a regular point.

Example (Genericity of triangle frameworks)

We want to determine whether frameworks $(G, p), (G, q)$ in \mathbb{R}^2 is generic where

$$V(G) = \{v_1, v_2, v_3\}$$

with distinct realizations p, q where

$$p_1 = (0, 0), p_2 = (1, 0), p_3 = (2, 0) \quad \text{and} \quad q_1 = (0, 0), q_2 = (1, 0), q_3 = (0, 1)$$

Before algebraically determining them, we can guess.

1. p seems NOT generic because three points are colinear.
2. q seems generic because three points are not colinear.

Let's check.

First, construct an indeterminate rigidity matrix with $n = 3$.

$$\mathbf{R}(3, 2) = \left(\begin{array}{cc|cc|cc} x_1 - x_2 & y_1 - y_2 & x_2 - x_1 & y_2 - y_1 & 0 & 0 \\ x_1 - x_3 & y_1 - y_3 & 0 & 0 & x_3 - x_1 & y_3 - y_1 \\ 0 & 0 & x_2 - x_3 & y_2 - y_3 & x_3 - x_2 & y_3 - y_2 \end{array} \right)$$

and construct $\mathbf{R}(K_3, p)$.

$$\mathbf{R}(K_3, p) = \left(\begin{array}{cc|cc|cc} 0 - 1 & 0 - 0 & 1 - 0 & 0 - 0 & 0 & 0 \\ 0 - 2 & 0 - 0 & 0 & 0 & 2 - 0 & 0 - 0 \\ 0 & 0 & 1 - 2 & 0 - 0 & 2 - 1 & 0 - 0 \end{array} \right) = \left(\begin{array}{cc|cc|cc} -1 & 0 & 1 & 0 & 0 & 0 \\ -2 & 0 & 0 & 0 & 2 & 0 \\ 0 & 0 & -1 & 0 & 1 & 0 \end{array} \right)$$

Taking submatrix $\mathbf{R}_{(145)}$ of $\mathbf{R}(K_3, p)$ by taking column 1, 4 and 5. Then,

$$\det \mathbf{R}_{(145)} = \det \begin{pmatrix} -1 & 0 & 0 \\ -2 & 0 & 2 \\ 0 & 0 & 1 \end{pmatrix} = 0$$

However, a determinant of corresponding submatrix $\mathbf{R}(3, 2)_{145}$ is

$$\begin{aligned} \det \mathbf{R}(3, 2)_{(145)} &= \det \begin{pmatrix} x_1 - x_2 & y_2 - y_1 & 0 \\ x_1 - x_3 & 0 & x_3 - x_1 \\ 0 & y_2 - y_3 & x_3 - x_2 \end{pmatrix} \\ &= (x_1 - x_3)[x_1(y_2 - y_3) + x_2(y_3 - y_1) + x_3(y_1 - y_2)] \end{aligned}$$

This is not necessarily zero (may not be zero by choosing coordinates properly) and hence, p is **NOT a generic realization**. (\rightarrow We correctly guessed)

Now, again, to determine whether q is generic or not, construct $\mathbf{R}(K_3, q)$.

$$\mathbf{R}(K_3, q) = \left(\begin{array}{cc|cc|cc} 0-1 & 0-0 & 1-0 & 0-0 & 0 & 0 \\ 0-0 & 0-1 & 0 & 0 & 0-0 & 1-0 \\ 0 & 0 & 1-0 & 0-1 & 0-1 & 1-0 \end{array} \right) = \left(\begin{array}{cc|cc|cc} -1 & 0 & 1 & 0 & 0 & 0 \\ 0 & -1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & -1 & -1 & 1 \end{array} \right)$$

Taking submatrix $\mathbf{R}_{(145)}$ of $\mathbf{R}(K_3, q)$ by taking column 1, 4 and 5. Then,

$$\det \mathbf{R}_{(145)} = \det \begin{pmatrix} -1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & -1 & -1 \end{pmatrix} = 0$$


However, a determinant of corresponding submatrix $\mathbf{R}(3, 2)_{145}$ is

$$\begin{aligned} \det \mathbf{R}(3, 2)_{(145)} &= \det \begin{pmatrix} x_1 - x_2 & y_2 - y_1 & 0 \\ x_1 - x_3 & 0 & x_3 - x_1 \\ 0 & y_2 - y_3 & x_3 - x_2 \end{pmatrix} \\ &= (x_1 - x_3)[x_1(y_2 - y_3) + x_2(y_3 - y_1) + x_3(y_1 - y_2)] \end{aligned}$$

This is not necessarily zero (may not be zero by choosing coordinates properly) and hence, q is **NOT a generic realization**. (\rightarrow Our guess was Wrong!)

As anticipated, the realization p with collinear points proved to be non-generic. However, realization q - which we initially presumed to be generic because it is a general framework - was also found to be non-generic upon verification.

This leads us to an important question: **What is the fundamental distinction between being 'general' and 'generic'?**

 **Remark) Generic is stronger than general**

Generic frameworks are general. But the converse doesn't hold.

There are some fundamental facts regarding this definition of generic.

Thm 3.3) Set theoretic properties of non-generic realizations

The set of all non-generic realizations $p : V \rightarrow \mathbb{R}^d$ where $V = \{v_1, \dots, v_n\}$ is a **closed set of measure zero**.

Thm 3.4) Combinatorial Rigidity Theorem

For a fixed dimension d and a graph G with $n(G) \geq d$, TFAE:

1. (G, p) is infinitesimally rigid (or independent, or isostatic) for at least one injective map
 $p : V(G) \rightarrow \mathbb{R}^d$.
2. (G, p) is infinitesimally rigid (or independent, or isostatic) for every d -dimensional generic realization p .

Intuitively, success in a single normal configuration implies success in almost every possible configuration.

The three structural properties (**infinitesimal rigidity**, **independence**, and **isostaticity**) only depend on the **underlying graph G** ; specifically its **combinatorial structure**.

Proof

(1) \Rightarrow (2): Suppose that (G, p) is infinitesimally rigid for an injective $p : V(G) \rightarrow \mathbb{R}^d$ (not necessarily generic).

By [theorem 2.16](#), for infinitesimally rigid framework,

$$\text{rank}(\mathbf{R}(G, p)) = dn - \binom{d+1}{2}$$

where $n(G) = n$.

Let \hat{p} be a d -dimensional generic framework for a graph G . By [lemma 3.3](#), the rigidity matrix of \hat{p} has the maximum rank and hence $\text{rank}(\mathbf{R}(G, \hat{p})) \geq \text{rank}(\mathbf{R}(G, p))$.

Again, by [theorem 2.16](#), a rank of rigidity matrix of d -dimensional framework cannot exceed $dn - \binom{d+1}{2}$ and hence $\text{rank}(\mathbf{R}(G, \hat{p})) \leq dn - \binom{d+1}{2}$.

Combining the equation and inequalities, we obtain

$$dn - \binom{d+1}{2} = \text{rank}(\mathbf{R}(G, p)) \leq \text{rank}(\mathbf{R}(G, \hat{p})) \leq dn - \binom{d+1}{2}$$

Thus, $\text{rank}(\mathbf{R}(G, \hat{p})) = dn - \binom{d+1}{2}$ and hence we can conclude that (G, \hat{p}) is infinitesimally rigid by [theorem 2.16](#).

(2) \Rightarrow (1): Suppose that (G, p) is infinitesimally rigid for every generic realization p .

Existence of p : Since the set of generic realizations is open dense set in \mathbb{R}^{dn} , there exists a map $p : V(G) \rightarrow \mathbb{R}^d$ s.t. (G, p) is infinitesimally rigid.

Injectivity: Also, since generic framework is general, the vertices never overlap and hence p must be injective.

Thus, for any generic realization p , it is an injective realization that satisfies (G, p) is infinitesimally rigid.

This proof can be applied directly for independence or isostaticity instead of infinitesimal rigidity. \square

Thm 3.5) Equivalence of rigidity

For generic frameworks, infinitesimal rigidity is equivalent to rigidity.

Def 3.6) Generically rigid/independent/isostatic

A graph G is called **generically rigid/independent/isostatic in dimension d** or generically d -rigid/independent/isostatic if d -dimensional generic realizations of G are infinitesimally rigid/independent/isostatic.

Here, by [theorem 3.5](#), A graph G is generically rigid if a generic realizations of G is rigid (\because infinitesimally rigid \equiv rigid for generic).

Note: Generically d -isostatic graphs can be constructed using inductive construction techniques. The following definition 3.7 and theorem 3.8 describe about it.

Introduce a definition first.

Def 3.7) Vertex addition

Let G be a graph, $U \subseteq V(G)$ with $|U| = d$ and $v \notin V(G)$. Then the graph \hat{G} with $V(\hat{G}) = V(G) \cup \{v\}$ and $E(\hat{G}) = E(G) \cup \{\{v, u\} : u \in U\}$ is called **a vertex d -addition** (by v) of G .

Intuitively, a vertex d -addition is the process of expanding a graph by introducing **one new vertex** and connecting it to exactly d **distinct existing vertices** via d new edges.

This operation serves as a fundamental building block for inductive constructions because it preserves the structural properties of the graph, as stated in the following theorem.

Thm 3.8) Vertex Addition Theorem

A vertex d -addition of a generically d -isostatic graph is generically d -isostatic. Conversely, deleting a vertex of degree d from a generically d -isostatic graph results in a generically d -isostatic graph.

Proof) By [definition 2.19](#), to show that a vertex d -addition of a generically d -isostatic graph is generically d -isostatic, we need to show that it is infinitesimally rigid and independent.

First, show the infinitesimal rigidity.

Now, show the independence.

Def 3.9) Edge split

Let G be a graph, and let $U \subseteq V(G)$ with $|U| = d + 1$ and $\{u_1, u_2\} \in E(G)$ for some $u_1, u_2 \in U$. Further, let $v \notin V(G)$. Then, the graph \hat{G} with $V(\hat{G}) = V(G) \cup \{v\}$ and $E(\hat{G}) = (E(G) \setminus \{u_1 u_2\}) \cup \{vu : u \in U\}$ is called an **edge d -split (on $u_1, u_2; v$)** of G .