

Part IA - Numbers and Sets

Theorems

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Introduction to number systems and logic

Overview of the natural numbers, integers, real numbers, rational and irrational numbers, algebraic and transcendental numbers. Brief discussion of complex numbers; statement of the Fundamental Theorem of Algebra.

Ideas of axiomatic systems and proof within mathematics; the need for proof; the role of counter-examples in mathematics. Elementary logic; implication and negation; examples of negation of compound statements. Proof by contradiction. [2]

Sets, relations and functions

Union, intersection and equality of sets. Indicator (characteristic) functions; their use in establishing set identities. Functions; injections, surjections and bijections. Relations, and equivalence relations. Counting the combinations or permutations of a set. The Inclusion-Exclusion Principle. [4]

The integers

The natural numbers: mathematical induction and the well-ordering principle. Examples, including the Binomial Theorem. [2]

Elementary number theory

Prime numbers: existence and uniqueness of prime factorisation into primes; highest common factors and least common multiples. Euclid's proof of the infinity of primes. Euclid's algorithm. Solution in integers of $ax + by = c$.

Modular arithmetic (congruences). Units modulo n . Chinese Remainder Theorem. Wilson's Theorem; the Fermat-Euler Theorem. Public key cryptography and the RSA algorithm. [8]

The real numbers

Least upper bounds; simple examples. Least upper bound axiom. Sequences and series; convergence of bounded monotonic sequences. Irrationality of $\sqrt{2}$ and e . Decimal expansions. Construction of a transcendental number. [4]

Countability and uncountability

Definitions of finite, infinite, countable and uncountable sets. A countable union of countable sets is countable. Uncountability of \mathbb{R} . Non-existence of a bijection from a set to its power set. Indirect proof of existence of transcendental numbers. [4]

Contents

1	Sets, functions and relations	4
1.1	Sets	4
1.2	Functions	4
1.3	Relations	4
2	Division	5
2.1	Euclid's Algorithm	5
2.2	Primes	5
3	Counting and Integers	6
3.1	Combinations	6
3.2	Well-ordering and induction	7
4	Modular arithmetic	8
4.1	Multiple moduli	8
4.2	Prime moduli	8
4.3	Public-key (Asymmetric) cryptography	9
4.3.1	RSA encryption	9
5	Real numbers	10
5.1	Construction of natural numbers	10
5.2	Construction of integers	10
5.3	Construction of rationals	10
5.4	Construction of real numbers	10
5.5	Sequences	10
5.6	Series	11
5.6.1	Decimal expansions	11
5.7	Irrational numbers	11
5.8	Euler's number	11
5.9	Algebraic numbers	11
6	Countability	12

1 Sets, functions and relations

1.1 Sets

Theorem. $(A = B) \Leftrightarrow (A \subseteq B \text{ and } B \subseteq A)$

Proposition.

$$- (A \cap B) \cap C = A \cap (B \cap C)$$

$$- (A \cup B) \cup C = A \cup (B \cup C)$$

$$- A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$$

1.2 Functions

Theorem. The left inverse of A exists iff f is injective.

Theorem. The right inverse of A exists iff f is surjective.

1.3 Relations

Theorem. If \sim is an equivalence relation on A , the equivalence classes of \sim form a partition of A .

2 Division

2.1 Euclid's Algorithm

Theorem (Division Algorithm). Given $a, b \in \mathbb{Z}$, there are unique $q, r \in \mathbb{Z}$ with $a = qb + r$ and $0 \leq r < b$.

Proposition. If $c|a$ and $c|b$, $c|(ua + vb)$ for all $u, v \in \mathbb{Z}$.

Theorem. Let $a, b \in \mathbb{N}$. Then (a, b) exists.

Corollary. (from the proof) Let $d = (a, b)$, then d is the smallest positive linear combination of a and b .

Corollary (Bézout's identity). Let $a, b \in \mathbb{N}$ and $c \in \mathbb{Z}$. Then there exists $u, v \in \mathbb{Z}$ with $c = ua + vb$ iff $(a, b)|c$.

Proposition (Euclid's Algorithm). If we continuously break down a and b by the following procedure:

$$\begin{aligned} a &= q_1b + r_1 \\ b &= q_2r_1 + r_2 \\ r_1 &= q_3r_2 + r_3 \\ &\vdots \\ r_{n-2} &= q_nr_{n-1} \end{aligned}$$

then the highest common factor is r_{n-1} .

2.2 Primes

Theorem. Every number can be written as a product of primes.

Theorem. There are infinitely many primes.

Theorem. If $a|bc$ and $(a, b) = 1$, then $a|c$.

Corollary. If p is a prime and $p|ab$, then $p|a$ or $p|b$. (True for all p, a, b)

Corollary. If p is a prime and $p|n_1n_2 \cdots n_i$, then $p|n_i$ for some i .

Theorem (Fundamental Theorem of Arithmetic). Every natural number is expressible as a product of primes in exactly one way. In particular, if $p_1p_2 \cdots p_k = q_1q_2 \cdots q_l$, where p_i, q_i are primes but not necessarily distinct, then $k = l$. q_1, \dots, q_l are p_1, \dots, p_k in some order. Since we already showed at least one way above, we only need to show uniqueness.

Corollary. If $a = p_1^{i_1}p_2^{i_2} \cdots p_r^{i_r}$ and $b = p_1^{j_1}p_2^{j_2} \cdots p_r^{j_r}$, where p_i are distinct primes (exponents can be zero). Then $(a, b) = \prod p_k^{\min\{i_k, j_k\}}$. However, this is not an efficient way to calculate (a, b) . Likewise, $\text{lcm}(a, b) = \prod p_k^{\max\{i_k, j_k\}}$. We have $\text{hcf}(a, b) \times \text{lcm}(a, b) = ab$.

3 Counting and Integers

Theorem (Pigeonhole Principle). Given $(m-1)n+1$ pigeons and n pigeonholes, some pigeonhole has at least m pigeons.

Proposition.

- (i) $i_A = i_B \Leftrightarrow A = B$
- (ii) $i_{A \cap B} = i_A i_B$
- (iii) $i_{\bar{A}} = 1 - i_A$
- (iv) $i_{A \cup B} = 1 - i_{\overline{A \cup B}} = 1 - i_{\bar{A} \cap \bar{B}} = 1 - i_{\bar{A}} i_{\bar{B}} = 1 - (1 - i_A)(1 - i_B) = i_A + i_B - i_{A \cap B}$.
- (v) $i_{A \setminus B} = i_{A \cap \bar{B}} = i_A i_{\bar{B}} = i_A(1 - i_B) = i_A - i_{A \cap B}$

Proposition. $|A \cup B| = |A| + |B| - |A \cap B|$

Theorem (Inclusion-Exclusion Principle). Let A_i be subsets of a finite set X , for $1 \leq i \leq n$. Then

$$|\bar{A}_1 \cap \dots \cap \bar{A}_n| = |X| - \sum_i |A_i| + \sum_{i < j} |A_i \cap A_j| - \dots + (-1)^n |A_1 \cap \dots \cap A_n|.$$

Equivalently,

$$|A_1 \cup \dots \cup A_n| = \sum_i |A_i| - \sum_{i < j} |A_i \cap A_j| + \dots + (-1)^{n-1} |A_1 \cap \dots \cap A_n|.$$

Note: The two forms are equivalent since $|A_1 \cup \dots \cup A_n| = |X| - |\bar{A}_1 \cap \dots \cap \bar{A}_n|$.

3.1 Combinations

Proposition. By definition,

$$\binom{n}{0} + \binom{n}{1} + \dots + \binom{n}{n} = 2^n$$

Theorem (Binomial theorem). For $n \in \mathbb{N}$ with $a, b \in \mathbb{R}$, we have

$$(a+b)^n = \binom{n}{0} a^n b^0 + \binom{n}{1} a^{n-1} b^1 + \dots + \binom{n}{r} a^{n-r} b^r + \dots + \binom{n}{n} a^0 b^n$$

Proposition.

- (i) $\binom{n}{r} = \binom{n}{n-r}$. This is because choosing r things to keep is the same as choosing $n-r$ things to throw away.
- (ii) $\binom{n}{r-1} + \binom{n}{r} = \binom{n+1}{r}$ (Pascal's identity) The RHS counts the number of ways to choose a team of r players from $n+1$ available players, one of whom is Pietersen. If Pietersen is chosen, there are $\binom{n}{r-1}$ ways to choose the remaining players. Otherwise, there are $\binom{n}{r}$ ways. The total number of ways is $\binom{n}{r-1} + \binom{n}{r}$.

Now given that $\binom{n}{0} = \binom{n}{n} = 1$, since there is only one way to choose nothing or everything, we can construct *Pascal's triangle*:

$$\begin{array}{ccccccc}
& & & & 1 & & \\
& & & & 1 & & 1 \\
& & & 1 & 2 & & 1 \\
& & 1 & 3 & 3 & & 1 \\
1 & & 4 & 6 & 4 & & 1
\end{array}$$

where each number is the sum of the two numbers above it, and the r th item of the n th row is $\binom{n}{r}$ (first row is row 0).

- (iii) $\binom{n}{k}\binom{k}{r} = \binom{n}{r}\binom{n-r}{k-r}$. LHS counts pairs of sets (Y, Z) with $|Y| = k$ and $|Z| = r$ with $Z \subseteq Y$. We first choose Y then choose $Z \subseteq Y$. The RHS chooses Z first and then choose the remaining $Y \setminus Z$ from $\{1, 2, \dots, n\} \setminus Z$.
- (iv) $\binom{a}{r}\binom{b}{0} + \binom{a}{r-1}\binom{b}{1} + \dots + \binom{a}{r-k}\binom{b}{k} + \dots + \binom{a}{0}\binom{b}{r} = \binom{a+b}{r}$
(Vandermonde's convolution) Suppose we have a men and b women, and we need to choose a committee of r people. The right hand side is the total number of choices. The left hand side breaks the choices up according to the number of men vs women.

Proposition. $\binom{n}{r} = \frac{n!}{(n-r)!r!}$.

3.2 Well-ordering and induction

Theorem (Weak Principle of Induction). Let $P(n)$ be a statement about the natural number n . Suppose that

- (i) $P(1)$ is true
- (ii) $\forall n(P(n) \Rightarrow P(n+1))$

Then $P(n)$ is true for all $n \geq 1$.

Theorem. Inclusion-exclusion principle

Theorem (Strong principle of induction). Let $P(n)$ be a statement about $n \in \mathbb{N}$. Suppose that

- (i) $P(1)$ is true
- (ii) $\forall n \in \mathbb{N}$, if $P(k)$ is true $\forall k < n$ then $P(n)$ is true.

Then $P(n)$ is true for all $n \in \mathbb{N}$.

Theorem. The strong principle of induction is equivalent to the weak principle of induction.

Theorem (Well-ordering principle). \mathbb{N} is well-ordered under the usual order, i.e. every non-empty subset of \mathbb{N} has a minimal element.

Theorem. The well-ordering principle is equivalent to the strong principle of induction.

4 Modular arithmetic

Proposition. If $a \equiv b \pmod{m}$, and $d|m$, then $a \equiv b \pmod{d}$.

Proposition. If $a \equiv b \pmod{m}$ and $u \equiv v \pmod{m}$, then $a + u \equiv b + v \pmod{m}$ and $au \equiv bv \pmod{m}$.

Theorem. There are infinitely many primes that are $\equiv -1 \pmod{4}$.

Theorem. u is a unit modulo m if and only if $(u, m) = 1$.

Corollary. If $(a, m) = 1$, then the congruence $ax \equiv b \pmod{m}$ has a unique (modulo m) solution.

Proposition. There is a solution to $ax \equiv b \pmod{m}$ if and only if $(a, m)|b$.

If $d = (a, m)|b$, then the solution is the unique solution to $\frac{a}{d}x \equiv \frac{b}{d} \pmod{\frac{m}{d}}$

4.1 Multiple moduli

Theorem (Chinese remainder theorem). Let $(m, n) = 1$ and $a, b \in \mathbb{Z}$. Then there is a unique solution (modulo mn) to the simultaneous congruences

$$\begin{cases} x \equiv a \pmod{m} \\ x \equiv b \pmod{n} \end{cases},$$

i.e. $\exists x$ satisfying both and every other solution is $\equiv x \pmod{mn}$

Proposition. Given any $(m, n) = 1$, c is a unit \pmod{mn} iff c is a unit both mod m and mod n .

Proposition.

- (i) $\phi(m)\phi(n) = \phi(mn)$ if $(m, n) = 1$, i.e. ϕ is multiplicative.
- (ii) If p is a prime, $\phi(p) = p - 1$
- (iii) If p is a prime, $\phi(p^k) = p^k - p^{k-1} = p^k(1 - 1/p)$
- (iv) $\phi(m) = m \prod_{p|m} (1 - 1/p)$.

4.2 Prime moduli

Theorem (Wilson's theorem). $(p - 1)! \equiv -1 \pmod{p}$ if p is a prime.

Theorem (Fermat's little theorem). Let p be a prime. Then $a^p \equiv a \pmod{p}$ for all $a \in \mathbb{Z}$. Equivalently, $a^{p-1} \equiv 1 \pmod{p}$ if $a \not\equiv 0 \pmod{p}$.

Theorem (Fermat-Euler Theorem). Let a, m be coprime. Then

$$a^{\phi(m)} \equiv 1 \pmod{m}.$$

Proposition. If p is an odd prime, then -1 is a quadratic residue if and only if $p \equiv 1 \pmod{4}$.

Proposition. (Unproven) A prime p is the sum of two squares if and only if $p \equiv 1 \pmod{4}$.

Proposition. There are infinitely many primes $\equiv 1 \pmod{4}$.

Proposition. Let $p = 4k + 3$ be a prime. Then if a is a quadratic residue, i.e. $a \equiv z^2 \pmod{p}$ for some z , then $z = \pm a^{k+1}$.

4.3 Public-key (Asymmetric) cryptography

4.3.1 RSA encryption

Theorem (RSA Encryption). We want people to be able to send a message to Bob without Eve eavesdropping. So the message must be encrypted. We want an algorithm that allows anyone to encrypt, but only Bob to decrypt (e.g. many parties sending passwords with the bank).

Let us first agree to write messages as sequences of numbers, e.g. in ASCII or UTF-8.

This is often done with RSA encryption (Rivest, Shamier, Adleman). Bob thinks of two large primes p, q . Let $n = pq$ and pick e coprime to $\phi(n) = (p-1)(q-1)$. Then work out d with $de \equiv 1 \pmod{\phi(n)}$ (i.e. $de = k\phi(n) + 1$). Bob then publishes the pair (n, e) .

For Alice to encrypt a message, Alice splits the message into numbers $M < n$. Alice sends $M^e \pmod{n}$ to Bob.

Bob then computes $(M^e)^d = M^{k\phi(n)+1} \equiv M \pmod{n}$ by Fermat-Euler theorem.

How can Eve find M ? We can, of course, factorize n and be in the same position as Bob. However, it is currently assumed that this is hard. Is there any other way? Currently we do not know iff RSA can be broken without factorizing (c.f. RSA problem).

5 Real numbers

5.1 Construction of natural numbers

5.2 Construction of integers

5.3 Construction of rationals

Proposition. \mathbb{Q} is a totally ordered-field.

Proposition. \mathbb{Q} is densely ordered, i.e. $\forall p, q \in \mathbb{Q}$ with $p < q$, then $\exists r \in \mathbb{Q}(p < r < q)$, e.g. $r = \frac{p+q}{2}$.

Proposition. There is no rational $q \in \mathbb{Q}$ with $q^2 = 2$.

5.4 Construction of real numbers

Axiom (Least upper bound axiom). Every non-empty set of the real numbers that has an upper bound has a least upper bound.

Corollary. Every non-empty set of the real numbers bounded below has an infimum

Theorem. (Axiom of Archimedes) Given $r \in \mathbb{R}$, there exists $n \in \mathbb{N}$ with $n > r$.

Proposition. $\inf\{\frac{1}{n} : n \in \mathbb{N}\} = 0$. Certainly 0 is a lower bound for S . If $t > 0$, there exists $n \in \mathbb{N}$ such that $n \geq 1/t$. So $t \geq 1/n \in S$. So t is not a lower bound for S .

Theorem. \mathbb{Q} is dense in \mathbb{R} , i.e. given $r, s \in \mathbb{R}$, with $r < s$, $\exists q \in \mathbb{Q}$ with $r < q < s$.

Theorem. There exists $x \in \mathbb{R}$ with $x^2 = 2$.

5.5 Sequences

Theorem. Every bounded monotonic sequence converges.

Theorem. Every sequence has a monotonic subsequence.

Theorem.

- (i) If $a_n \rightarrow a$ and $a_n \rightarrow b$, then $a = b$ (i.e. limits are unique)
- (ii) If $a_n \rightarrow a$ and $b_n = a_n$ for all but finitely many n , then $b_n \rightarrow a$.
- (iii) If $a_n = a$ for all n , then $a_n = a$.
- (iv) If $a_n \rightarrow a$ and $b_n \rightarrow b$, then $a_n + b_n \rightarrow a + b$
- (v) If $a_n \rightarrow a$ and $b_n \rightarrow b$, then $a_n b_n \rightarrow ab$
- (vi) If $a_n \rightarrow a \neq 0$, and $\forall n(a_n \neq 0)$. Then $1/a_n \rightarrow 1/a$.
- (vii) If $a_n \rightarrow a$ and $b_n \rightarrow a$, and $\forall n(a_n \leq c_n \leq b_n)$, then $c_n \rightarrow a$. (Sandwich theorem)

5.6 Series

5.6.1 Decimal expansions

5.7 Irrational numbers

Proposition. A number is periodic iff it is rational.

5.8 Euler's number

Proposition. e is irrational.

5.9 Algebraic numbers

Proposition. All rational numbers are algebraic.

Theorem. (Liouville 1851; Non-examinable) L is transcendental, where

$$L = \sum_{n=1}^{\infty} \frac{1}{10^{n!}} = 0.11000100 \dots$$

with 1s in the factorial positions.

Theorem. (Hermite 1873) e is transcendental.

Theorem. (Lindemann 1882) π is transcendental.

6 Countability

Lemma. If $f : [n] \rightarrow [n]$ is injective, then f is surjective.

Corollary. If A is a set and $f : A \rightarrow [n]$ and $g : A \rightarrow [m]$ are both bijections, then $m = n$.

Lemma. Let $S \subseteq \mathbb{N}$. Then either S is finite or there is a bijection $g : \mathbb{N} \rightarrow S$.

Theorem. The following are equivalent:

- (i) A is countable
- (ii) There is an injection from $A \rightarrow \mathbb{N}$
- (iii) $A = \emptyset$ or there is a surjection from $\mathbb{N} \rightarrow A$

Proposition. The integers \mathbb{Z} are countable.

Proposition. $\mathbb{N} \times \mathbb{N}$ is countable.

Proposition. If $A \rightarrow B$ is injective and B is countable, then A is countable (since we can inject $B \rightarrow \mathbb{N}$).

Proposition. \mathbb{Z}^k is countable for all $k \in \mathbb{N}$

Theorem. A countable union of countable sets is countable.

Proposition. \mathbb{Q} is countable.

Theorem. The set of algebraic numbers is countable.

Theorem. The set of real numbers \mathbb{R} is uncountable.

Corollary. There are uncountable many transcendental numbers.

Theorem. Let A be a set. Then there is no surjection from $A \rightarrow \mathcal{P}A$.

Theorem (Cantor-Schröder-Bernstein theorem). Suppose there are injections $A \rightarrow B$ and $B \rightarrow A$. Then there's a bijection $A \leftrightarrow B$.