Team Contest Reference

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ChaosKITs

Datenstrukturen 1

Union-Find 1.1

```
1 vector<int> parent, rank2; //manche compiler verbieten Variable mit Namen rank
3 int findSet(int n) { //Pfadkompression
    if (parent[n] != n) parent[n] = findSet(parent[n]);
4
5
     return parent[n];
6 }
7
8 void linkSets(int a, int b) { //union by rank
9
    if (rank2[a] < rank2[b]) parent[a] = b;</pre>
10
     else if (rank2[b] < rank2[a]) parent[b] = a;</pre>
11
     else {
12
       parent[a] = b;
13
       rank2[b]++;
14
    }
15 }
16
17 void unionSets(int a, int b) {
18
    if (findSet(a) != findSet(b)) linkSets(findSet(a), findSet(b));
19 }
   1.2
         Segmentbaum
1 int a[MAX_N], m[4 * MAX_N];
3 int query(int x, int y, int k = 0, int X = 0, int Y = MAX_N - 1) {
    if (x <= X && Y <= y) return m[k];</pre>
    if (y < X || Y < x) return -1000000000; //ein "neutrales" Element
5
     int M = (X + Y) / 2;
7
     return max(query(x, y, 2 * k + 1, X, M), query(x, y, 2 * k + 2, M + 1, Y));
8 }
9
10 void update(int i, int v, int k = 0, int X = 0, int Y = MAX_N - 1) {
11
     if (i < X || Y < i) return;</pre>
12
     if (X == Y) {
      m[k] = v;
13
14
       a[i] = v;
15
       return;
16
    }
    int M = (X + Y) / 2;
17
18
     update(i, v, 2 * k + 1, X, M);
19
     update(i, v, 2 * k + 2, M + 1, Y);
20
    m[k] = max(m[2 * k + 1], m[2 * k + 2]);
21 }
22
23 void init(int k = 0, int X = 0, int Y = MAX_N - 1) {
24
    if (X == Y) {
25
       m[k] = a[X];
26
       return;
27
    }
28
    int M = (X + Y) / 2;
29
    init(2 * k + 1, X, M);
    init(2 * k + 2, M + 1, Y);
31
     m[k] = max(m[2 * k + 1], m[2 * k + 2]);
32 }
```

$\mathbf{2}$ Graphen

Kürzeste Wege 2.1

2.1.1 Bellmann-Ford-Algorithmus

```
1 \ //n = number of vertices, edges is vector of edges
2 dist.assign(n, INF); dist[0] = 0;
3 parent.assign(n, -1);
4 \text{ for (i = 0; i < n - 1; i++) }
```

```
5
     for (j = 0; j < (int)edges.size(); j++) {</pre>
       if (dist[edges[j].from] + edges[j].cost < dist[edges[j].to]) {</pre>
         dist[edges[j].to] = dist[edges[j].from] + edges[j].cost;
7
         parent[edges[j].to] = edges[j].from;
9
10
     }
11 }
12 //now dist and parent are correct shortest paths
13 //next lines check for negative cycles
14 for (j = 0; j < (int)edges.size(); j++) {
15
    if (dist[edges[j].from] + edges[j].cost < dist[edges[j].to]) {</pre>
16
      //NEGATIVE CYCLE found
17
18 }
```

2.2 Strongly Connected Components (Tarjans-Algorithmus)

```
1 int counter, sccCounter, n; //n == number of vertices
2 vector < bool > visited, inStack;
3 vector< vector<int> > adjlist;
4 vector <int> d, low, sccs;
5 stack<int> s;
7 void visit(int v) {
    visited[v] = true;
8
9
     d[v] = counter;
     low[v] = counter;
10
11
     counter++;
12
     inStack[v] = true;
13
     s.push(v);
14
15
     for (int i = 0; i < (int)adjlist[v].size(); i++) {</pre>
16
      int u = adjlist[v][i];
       if (!visited[u]) {
17
         visit(u);
18
19
         low[v] = min(low[v], low[u]);
       } else if (inStack[u]) {
20
21
         low[v] = min(low[v], low[u]);
22
23
24
25
     if (d[v] == low[v]) {
26
      int u;
27
       do {
28
         u = s.top();
29
         s.pop();
30
         inStack[u] = false;
31
         sccs[u] = sccCounter;
32
       } while(u != v);
33
       sccCounter++;
34
     }
35 }
36
37 \text{ void scc()}  {
38
     //read adjlist
39
40
     visited.clear(); visited.assign(n, false);
41
     d.clear(); d.resize(n);
42
     low.clear(); low.resize(n);
43
     inStack.clear(); inStack.assign(n, false);
     sccs.clear(); sccs.resize(n);
44
45
46
     counter = 0;
47
     sccCounter = 0;
48
     for (i = 0; i < n; i++) {</pre>
49
      if (!visited[i]) {
50
         visit(i);
51
52
53
     //sccs has the component for each vertex
```

2.3 Max-Flow (Edmonds-Karp-Algorithmus)

```
1 int s, t, f; //source, target, single flow
2 int res[MAX_V][MAX_V]; //adj-matrix
3 vector < vector <int> > adjList;
4 int p[MAX_V]; //bfs spanning tree
6 void augment(int v, int minEdge) {
7
     if (v == s) { f = minEdge; return; }
8
     else if (p[v] != -1) {
9
       augment(p[v], min(minEdge, res[p[v]][v]));
10
       res[p[v]][v] -= f; res[v][p[v]] += f;
11 }}
12
  int maxFlow() { //first inititalize res, adjList, s and t
13
     int mf = 0;
14
     while (true)
15
16
       f = 0;
       bitset < MAX_V > vis; vis[s] = true;
17
18
       queue < int > q; q.push(s);
       memset(p, -1, sizeof(p));
19
20
       while (!q.empty()) { //BFS
21
         int u = q.front(); q.pop();
22
         if (u == t) break;
23
         for (int j = 0; j < (int)adjList[u].size(); j++) {</pre>
           int v = adjList[u][j];
24
25
           if (res[u][v] > 0 && !vis[v]) {
26
             vis[v] = true; q.push(v); p[v] = u;
27
28
29
       augment(t, INF); //add found path to max flow
30
       if (f == 0) break;
31
       mf += f;
     }
32
33
     return mf;
34 }
```

3 Geometrie

3.1 Closest Pair

```
1 double squaredDist(point a, point b) {
2
     return (a.first-b.first) * (a.first-b.first) + (a.second-b.second) * (a.second-b.second);
3 }
4
5 bool compY(point a, point b) {
     if (a.second == b.second) return a.first < b.first;</pre>
7
     return a.second < b.second;</pre>
8 }
9
10 double shortestDist(vector < point > & points) {
     //check that points.size() > 1 and that ALL POINTS ARE DIFFERENT
11
12
     set < point, bool(*)(point, point) > status(compY);
13
     sort(points.begin(), points.end());
     double opt = 1e30, sqrtOpt = 1e15;
14
     auto left = points.begin(), right = points.begin();
15
16
     status.insert(*right); right++;
17
18
     while (right != points.end()) {
19
       if (fabs(left->first - right->first) >= sqrtOpt) {
20
         status.erase(*(left++));
21
       } else {
         auto lower = status.lower_bound(point(-1e20, right->second - sqrt0pt));
22
23
         auto upper = status.upper_bound(point(-1e20, right->second + sqrtOpt));
24
         while (lower != upper) {
           double cand = squaredDist(*right, *lower);
25
26
           if (cand < opt) {</pre>
27
             opt = cand;
28
             sqrtOpt = sqrt(opt);
29
```

30

```
++lower;
31
         }
         status.insert(*(right++));
32
33
34
     }
35
    return sqrtOpt;
36 }
         Formeln - std::complex
1 //komplexe Zahlen als Darstellung fuer Punkte
2 typedef pt complex <double >;
3 //Winkel zwischen Punkt und x-Achse in [0, 2 * PI), Winkel zwischen a und b
4 double angle = arg (a), angle_a_b = arg (a - b);
5 //Punkt rotiert um Winkel theta
6 pt a_rotated = a * exp (pt (0, theta));
7 //Mittelpunkt des Dreiecks abc
8 \text{ pt centroid} = (a + b + c) / 3;
9 //Skalarprodukt
10 double dot(pt a, pt b) {
11
    return real(conj(a) * b);
12 }
13 //Kreuzprodukt, 0, falls kollinear
14 double cross(pt a, pt b) {
15
    return imag(conj(a) * b);
16 }
17 //wenn Eckpunkte bekannt
18 double areaOfTriangle(pt a, pt b, pt c) {
19
    return abs(cross(b - a, c - a)) / 2.0;
20 }
21 //wenn Seitenlaengen bekannt
22 double areaOfTriangle(double a, double b, double c) {
23 double s = (a + b + c) / 2;
24
    return sqrt(s * (s-a) * (s-b) * (s-c));
25 }
26 // Sind die Dreiecke a1, b1, c1, and a2, b2, c2 aehnlich?
27\ //\ {\it Erste}\ {\it Zeile} testet Aehnlichkeit mit gleicher Orientierung,
28 // zweite Zeile testst Aehnlichkeit mit unterschiedlicher Orientierung
29 bool similar (pt a1, pt b1, pt c1, pt a2, pt b2, pt c2) {
30
    return (
31
       (b2 - a2) * (c1 - a1) == (b1 - a1) * (c2 - a2) | |
32
       (b2 - a2) * (conj (c1) - conj (a1)) == (conj (b1) - conj (a1)) * (c2 - a2)
33
    ):
34 }
35 //Linksknick von a->b nach a->c
36 double ccw(pt a, pt b, pt c) {
37
    return cross(b - a, c - a); //<0 => falls Rechtsknick, 0 => kollinear, >0 => Linksknick
38 }
39 //Streckenschnitt, Strecken a-b und c-d
40 bool lineSegmentIntersection(pt a, pt b, pt c, pt d) {
41
    if (ccw(a, b, c) == 0 \&\& ccw(a, b, d) == 0) { //kollinear}
42
       double dist = abs(a - b);
       return (abs(a - c) <= dist && abs(b - c) <= dist) || (abs(a - d) <= dist && abs(b - d) <= dist);
43
44
     }
    return ccw(a, b, c) * ccw(a, b, d) <= 0 && ccw(c, d, a) * ccw(c, d, b) <= 0;
45
47 //Entfernung von p zu a-b
48 double distToLine(pt a, pt b, pt p) {
49 return abs(cross(p - a, b - a)) / abs(b - a);
50 }
51 //liegt p auf a-b
52 bool pointOnLine(pt a, pt b, pt p) {
53
  return abs(distToLine(a, b, p)) < EPSILON;</pre>
54 }
55 //testet, ob d in der gleichen Ebene liegt wie a, b, und c
56 bool isCoplanar(pt a, pt b, pt c, pt d) {
57
    return (b - a) * (c - a) * (d - a) == 0;
58 }
```

4 Mathe

4.1 ggT, kgV, erweiterter euklidischer Algorithmus

```
1 11 gcd(11 a, 11 b) {
   return b == 0 ? a : gcd (b, a % b);
3 }
4
5 11 1cm(11 a, 11 b) {
    return a * (b / gcd(a, b)); //Klammern gegen Overflow
6
1 \ // \texttt{Accepted in Aufgabe mit Forderung:} \ |X| + |Y| \ \texttt{minimal (primaer) und X <= Y (sekundaer)}
2 //hab aber keinen Beweis dafuer :)
3 11 x, y, d; //a * x + b * y = d = ggT(a,b)
4 void extendedEuclid(ll a, ll b) {
     if (!b) {
       x = 1; y = 0; d = a; return;
6
7
8
     extendedEuclid(b, a % b);
     11 x1 = y; 11 y1 = x - (a / b) * y;
10
     x = x1; y = y1;
11 }
```

4.2 Binomialkoeffizienten

```
1 11 calc_binom(11 N, 11 K) {
2     11 r = 1, d;
3     if (K > N) return 0;
4     for (d = 1; d <= K; d++) {
5         r *= N--;
6         r /= d;
7     }
8     return r;
9 }</pre>
```

5 Sonstiges

5.1 2-SAT

- 1. Bedingungen in 2-CNF formulieren.
- 2. Implikationsgraph bauen, $(a \lor b)$ wird zu $\neg a \Rightarrow b$ und $\neg b \Rightarrow a$.
- 3. Finde die starken Zusammenhangskomponenten.
- 4. Genau dann lösbar, wenn keine Variable mit ihrer Negation in einer SCC liegt.