

#6.

B03901023

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電網導 HW1

a)  $d_{prop} = m/s$  (sec)

b)  $d_{trans} = 2/R$  (sec)

c)  $d_{end} = m/s + 2/R$  (sec)

#6.9.14.18.24

d) 此 bit 剛離開 host A

25.31

e) 第 3 bit 仍在 link 之中, 尚未抵達 host B.

f) 第 3 bit 已經抵達 host B.

g)  $m = ?$

$$m = \left(\frac{2}{R}\right) \cdot S = \frac{120}{56 \times 10^3} \cdot 2.5 \times 10^8 = 536 \text{ km}$$

#9 Link. 1 Mbps rate: 100 kbps

 $p = 0.1$ ~~1 Mbps~~  $\rightarrow$  1 Gbps

a)  $N = ?$

$$\frac{1 \text{ G}}{100 \text{ k}} \cdot \frac{10^9}{100 \times 10^3} = 10^4 = 10000$$

b) packet switching: binomial distribution

$$\Rightarrow \sum_{n=N+1}^M \binom{M}{n} p^n (1-p)^{M-n}$$

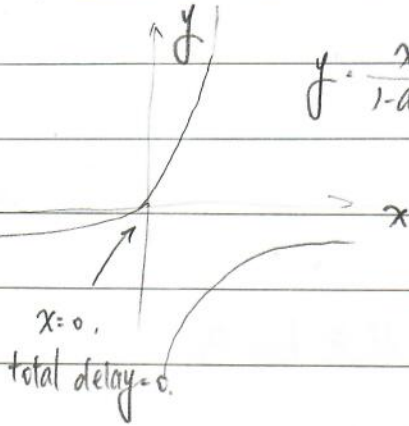
 $N+1$  是因為 "more than  $N$  users"#14 a) transmission delay:  $2/R$ 

total delay

$$= \frac{IL}{R(1-I)} + \frac{2}{R} = \frac{2/R}{1-I}$$

b) let  $x = 1/R$ , total delay =  $\frac{x}{1-ax}$ .

$$f = \frac{x}{1-ax}, \text{ a given}$$



#18

#24  $40T = 40 \times 10^{12} \times 8 \text{ bits.}$

$$\frac{40 \times 10^{12} \times 8}{100 \times 10^6} = 3200000 \text{ sec} = \frac{3200000}{60 \times 24} \approx 37 \text{ days.}$$

cf FedEx overnight delivery  $\rightarrow 24 \text{ hrs.}$  : 成本  $\leq 100 \$$   
 以内就抵達 所以用 FedEx 较好.

#25.  $d_{A \rightarrow B} = 20000 \text{ km} = 2 \times 10^7 \text{ m}$

$R = 2 \text{ Mbps}$       speed =  $2.5 \times 10^8 \text{ m/s}$

a)  $\frac{d_{A \rightarrow B}}{\text{speed}} = \frac{2 \times 10^7}{2.5 \times 10^8} = 0.08 = d_{\text{prop}}$

$R \times d_{\text{prop}} = 2 \times 10^6 \times 0.08 = 16 \times 10^4 = 160000 \text{ bits}$

b) file size = 800000 bits

max = 160000 bits

c) bandwidth-delay product of a link  
= maximum # of bits in the link

d) width of a bit =  $2 \times 10^7 \text{ m} / 160000 = 125 \text{ m} \dots$  / bit

football field length = 109.73m, - 1 bit 的 width 较短

e)  $S/R$

#3)

a)  $t = \frac{8 \times 10^6}{2 \times 10^6} = 4 \text{ sec}$ , source host  $\xrightarrow{t}$  1st packet switch

$4 \text{ sec} \times 3 \text{ hops} = 12 \text{ sec}$ , total time from source host to destination host

b)  $\frac{1 \times 10^4}{2 \times 10^6} = 5 \text{ msec} = 0.005 \text{ sec}$ , 1st packet 從 src 到 1st packet switch.

$2 \times 5 \text{ m} = 0.01 \text{ sec}$ , 2nd packet 在 1st packet switch 被 received

c)  $0.005 \text{ sec} \times 3 \text{ hops} = 15 \text{ msec} = 0.015$ , 1st packet 抵達 destination.

$0.015 + 799 \times 0.005 = 4.01 \text{ sec}$ , 800th packet

$\therefore$  delay in using message segmentation is significant less. (1/3rd)