CSE 12 – Basic Data Structures and Object-Oriented Design
Lecture 13

Greg Miranda and Paul Cao, Winter 2021

This lecture is being recorded

Announcements

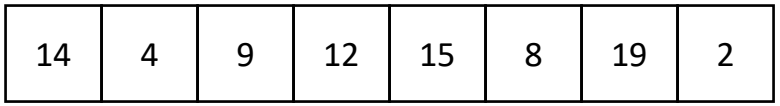
- Quiz 13 due Monday @ 8am
- Survey 5 due tonight @ 11:59pm
- PA5 due Wednesday @ 11:59pm

Topics

- Quick Partition/Sort
- Questions on Lecture 13?

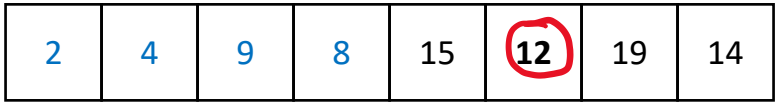
Quicksort: Another magical (recursive) algorithm

https://www.youtube.com/watch?v=ywWBy6J5gz8



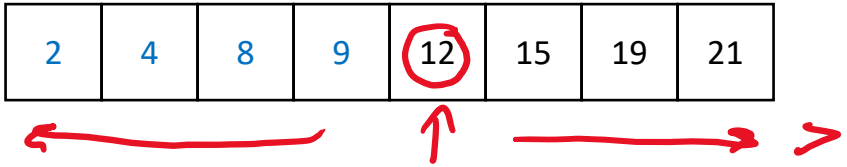
Select a **pivot** element:

"Partition" the elements in the array (smaller or equal to pivot, larger or equal to pivot)



Magically sort the smaller elements and the larger elements (Quicksort)

pivot index: 4
pivot value = 12



```
Quicksort(numbers, lowIndex, highIndex) {
   if (lowIndex >= highIndex) {
      return
   }

   lowEndIndex = Partition(numbers, lowIndex, highIndex)
   Quicksort(numbers, lowIndex, lowEndIndex)
   Quicksort(numbers, lowEndIndex + 1, highIndex)
}
```

There are many ways to partition!

```
Partition(numbers, lowIndex, highIndex) {
   // Pick middle element as pivot
   midpoint = lowIndex + (highIndex - lowIndex) / 2
   pivot = numbers[midpoint]

   done = false
   while (!done) {
      // Increment lowIndex while numbers[lowIndex] < pivot
      while (numbers[lowIndex] < pivot) {
         lowIndex += 1
      }

      // Decrement highIndex while pivot < numbers[highIndex]
      while (pivot < numbers[highIndex]) {
         highIndex -= 1
      }

      // If zero or one elements remain, then all numbers are
      // partitioned. Return highIndex.
      if (lowIndex >= highIndex) {
         done = true
      }
      else {
         // Swap numbers[lowIndex] and numbers[highIndex]
         temp = numbers[lowIndex]
         numbers[lowIndex] = numbers[highIndex]
         numbers[highIndex] = temp

         // Update lowIndex and highIndex
         lowIndex += 1
         highIndex -= 1
      }
   }

   return highIndex
}
```

Quick sort

sort {12, 4, 9, 3, 15, 8, 19, 2}

lowIndex = 0
highIndex = 7
midpoint = 3
pivot = 3

```
Partition(numbers, lowIndex, highIndex) {
   // Pick middle element as pivot
   midpoint = lowIndex + (highIndex - lowIndex) / 2
   pivot = numbers[midpoint]

   done = false
   while (!done) {
      // Increment lowIndex while numbers[lowIndex] < pivot
      while (numbers[lowIndex] < pivot) {
         lowIndex += 1
      }

      // Decrement highIndex while pivot < numbers[highIndex]
      while (pivot < numbers[highIndex]) {
         highIndex -= 1
      }

      // If zero or one elements remain, then all numbers are
      // partitioned. Return highIndex.
      if (lowIndex >= highIndex) {
         done = true
      }
      else {
         // Swap numbers[lowIndex] and numbers[highIndex]
         temp = numbers[lowIndex]
         numbers[lowIndex] = numbers[highIndex]
         numbers[highIndex] = temp

         // Update lowIndex and highIndex
         lowIndex += 1
         highIndex -= 1
      }
   }

   return highIndex
}
```

Quick Sort Details

1. We always pick the middle location as pivot
2. The data we sort is {2, 3, 1, 5, 4, 6, 7}

After the first split, what is the order of elements in the list that was <= pivot?

A. 1 2 3 4

B. 2 3 1 4

C. 4 3 2 1

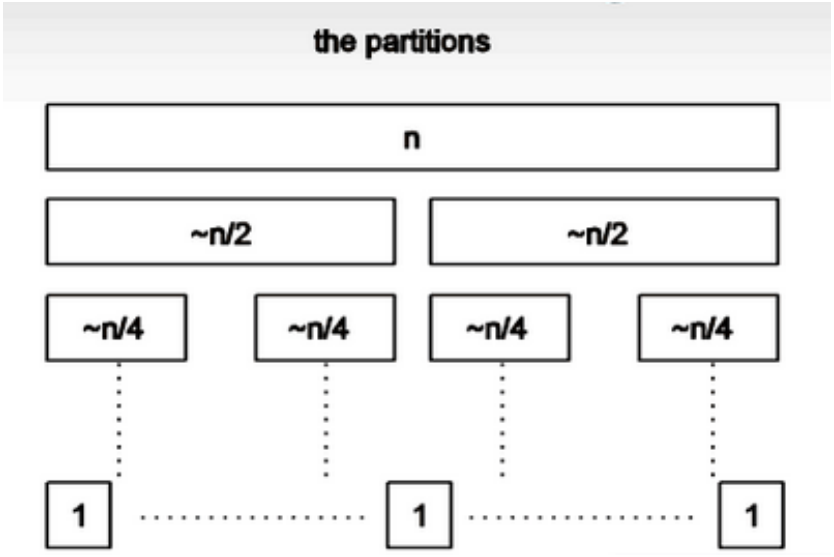
D. 3 4 1 2

E. None of the above

Quick Sort: Using a "good" pivot

$\frac{7}{2} \rightarrow \frac{3}{2} \rightarrow \frac{1}{2} \rightarrow 0$

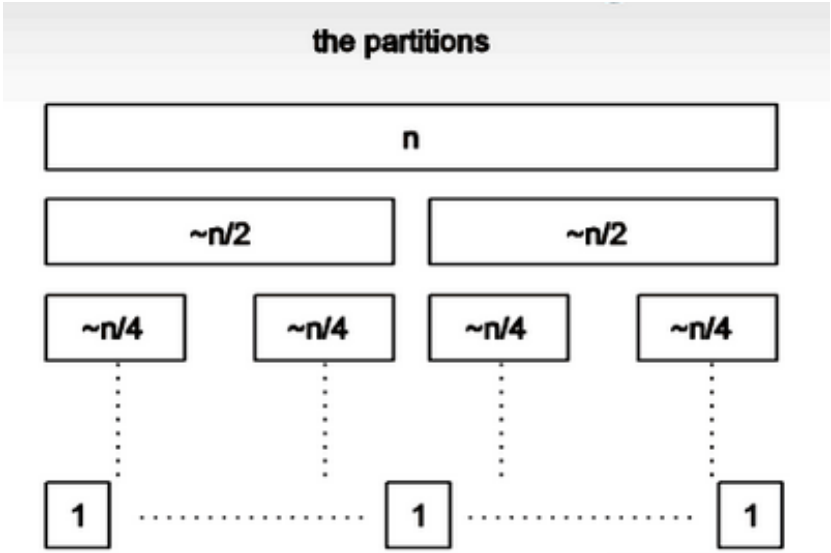
3 levels



How many levels will there be if you choose a pivot that divides the list in half?

A. 1
B. log(N)
C. N
D. N*log(N)
E. N^2

Quick Sort: Using a "good" pivot



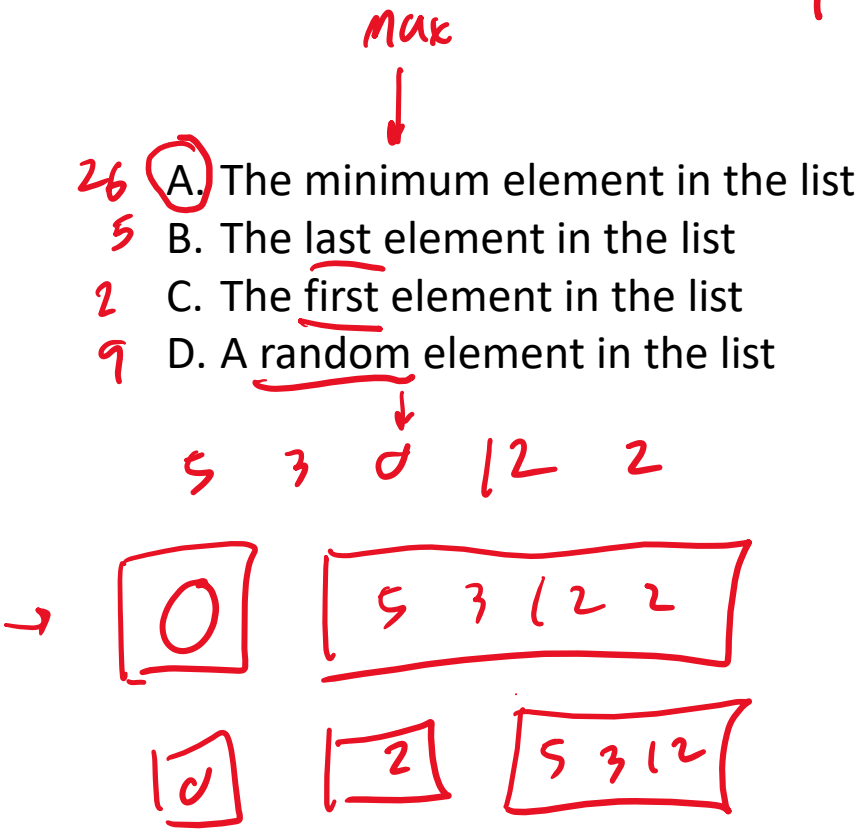
$\log_2(N)$

If the time to partition on each level takes N comparisons, how long does Quicksort take with a good partition?

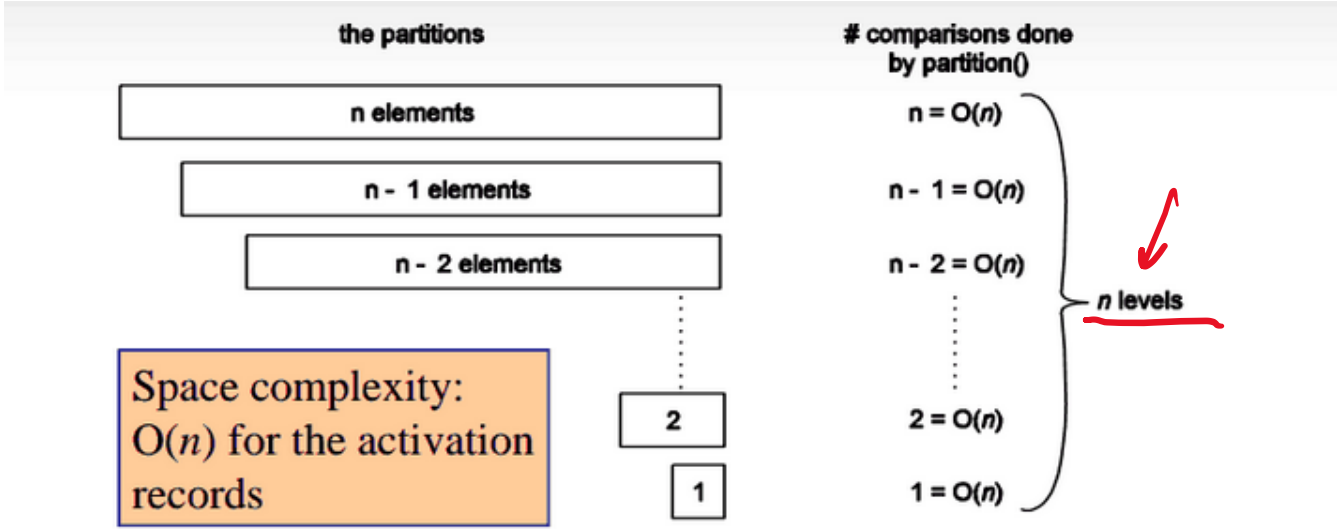
A. O(1)
B. O(log(N))
C. O(N)
D. O(N*log(N))
E. $O(N^2)$

$N * \log_2(N)$

Which of these choices would be the *worst* choice for the pivot?



Quick sort with a bad pivot



If the pivot always produces one empty partition and one with $n - 1$ elements, there will be n levels, each of which requires $O(n)$ comparisons: $O(n^2)$ time complexity

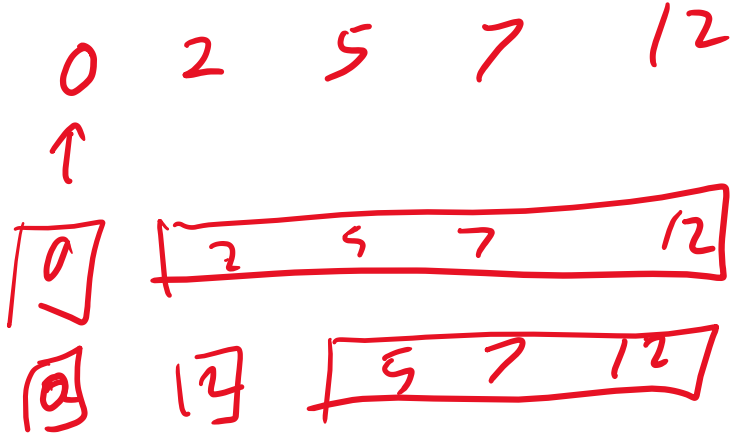
Which of these choices is a better choice for the pivot?

A. The first element in the list

B. A random element in the list

C. They are about the same

except when sorted



QuickSort – Draw the picture of sort()

```
public class Sort {
  public static void swap(String[] array, int i1, int i2) {
    String temp = array[i1];
    array[i1] = array[i2];
    array[i2] = temp;
  }
  public static int partition(String[] array, int low, int high) {
    int pivotStartIndex = high - 1;
    String pivot = array[pivotStartIndex];
    int smallerBefore = low, largerAfter = high - 2;

    while (smallerBefore <= largerAfter) {
      if (array[smallerBefore].compareTo(pivot) < 0) {
        smallerBefore += 1;
      }
      else {
        swap(array, smallerBefore, largerAfter);
        largerAfter -= 1;
      }
    }

    swap(array, smallerBefore, pivotStartIndex);
    return smallerBefore;
  }
```

```
  public static void qsort(String[] array, int low, int high) {
    if (high - low <= 1) { return; }
    int splitAt = partition(array, low, high);
    qsort(array, low, splitAt);
    qsort(array, splitAt + 1, high);
  }

  public static void sort(String[] array) {
    qsort(array, 0, array.length);
  }

main() {
  String[] str = {"f", "b", "a", "e", "d", "c" };
  int[] result = Sort.sort(str);
  System.out.println(Arrays.deepToString(result));
}
```