

CSE 12 Week 7 Discussion

2-16-21

Focus: Maps

Reminders

- PA7 is a **closed** assignment - no collaborating!
 - Due Tuesday, March 2nd 11:59 PM
- PA4 Resubmission due Friday, February 19th 11:59 PM
- PA5 Resubmission due Friday, February 26th 11:59 PM

Maps

Maps

Maps are an Abstract Data Type (ADT)

Assign a **key** to each **value** we are trying to keep track of.

Key 1 ---> Some value a

Key 2 ---> Some value b

Key 3 ---> Some value c

etc...

Map<K,V> Interface

- Implemented in Java by AbstractMap, HashMap, TreeMap etc.
- Index for an entry is determined by a hash function that calculates an index using the key value (useful for quick lookup and insert)
- Contains methods such as get(Object key), put(K key, V value), size(), replace(K key, V value) etc.
- Keys need to be unique
- Existing data structures can be used to implement this - ArrayList!

Inserting Into Map

Key	Value
"a"	"apple"
"o"	"orange"

Insert some entries into map:

```
put("a", "apple");
```

```
put("o", "orange");
```

Note - this is a simplified view of map entries.
May not be in this exact order

Inserting Into Map

Key	Value
"a"	"apple"
"o"	"orange"

Can we insert the following as a new entry?

```
put("a", "apricot")
```

- A. Yes
- B. No
- C. Will only work this time since values are different at key = 1

Note - this is a simplified view of map entries.
May not be in this exact order

Answer - B

Key	Value
"a"	"apricot"
"o"	"orange"

Maps can only hold unique key values. If you try inserting value to an existing key, the old value will be overwritten.

Note - this is a simplified view of map entries.
May not be in this exact order

Inserting Into Map

Key	Value
"a"	"apple"
"o"	"orange"

Can we insert the following as a new entry?

```
put("c", "orange")
```

- A. Yes
- B. No
- C. I don't know

Note - this is a simplified view of map entries.
May not be in this exact order

Answer - A

Key	Value
"a"	"apple"
"o"	"orange"
"c"	"orange"

Values across key-value pairs in maps do not need to be unique.

Note - this is a simplified view of map entries.
May not be in this exact order

HashMaps - Separate Chaining

Example Hash Functions

```
int hash1(String s) {  
    return s.length();  
}
```

```
int hash2(String s) {  
    int hash = 0;  
    for(int i = 0; i < s.length(); i += 1) {  
        hash += Character.codePointAt(s, i);  
    }  
    return hash;  
}
```

```
public int hash3(String s) {  
    int h = 0;  
    for (int i = 0; i < s.length(); i++) {  
        h = 31 * h + Character.codePointAt(s, i);  
    }  
    return h;  
}
```

What is a bucket?

A: The collection of key values at a specific index

B: The sum of the keys in a given hash map

C: The collection of elements at a specific index

D: The collection of keys at a specific index

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B: The sum of the keys in a given hash map

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D: The collection of keys at a specific index

Given the example below, what does the HashMap look like after line: `set("red", 70)`?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	
1	
2	
3	

A `HashMap<Key, Value>` using Separate Chaining has:

- size: an int
- buckets: an array of lists of Entries
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

```
void set(key, value):
    hashed = hash(key)
    index = hashed % this.buckets.length
    if this.buckets[index] contains an Entry with key:
        update that Entry to contain value
    else:
        increment size
        bucket = buckets[index]
        add {key: value} to end of bucket
```

```
Value get(key):
    hashed = hash(key)
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```
set("red", 70)
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```

0	
1	
2	
3	— {red: 70}

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```


Given the example below, what does the HashMap look like after line: `set("blue", 90)`?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	
1	
2	
3	— {red: 70}

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```

Given the example below, what does the HashMap look like after line: `set("blue", 90)`?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	— {blue: 90}
1	
2	
3	— {red: 70}

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    else:
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```

Given the example below, what does the HashMap look like after line: `set("pink", 100)`?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	— {blue: 90}
1	
2	
3	— {red: 70}

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- size: an int
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- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

```
void set(key, value):
    hashed = hash(key)
    index = hashed % this.buckets.length
    if this.buckets[index] contains an Entry with key:
        update that Entry to contain value
    else:
        increment size
        bucket = buckets[index]
        add {key: value} to end of bucket
```

```
Value get(key):
    hashed = hash(key)
    index = hashed % this.buckets.length
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        return the value of that entry
    else:
        return null/report an error
```

Given the example below, what does the HashMap look like after line: `set("pink", 100)`?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	— {blue: 90} — {pink: 100}
1	
2	
3	— {red: 70}

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void set(key, value):
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        add {key: value} to end of bucket
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```

Example:

Start buckets array with size 4

Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

Assuming the above example has been

executed, **how many elements are in bucket 0?**

A: 0

B: 1

C: 2

D: 3

E: more than 3

A HashMap<Key, Value> using Separate Chaining has:

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executed, **how many elements are in bucket 0?**

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B: 1

C: 2

D: 3

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```
set("red", 70)
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set("pink", 100)
set("orange", 40)
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```

Assuming the above example has been

executed, **how many elements are in bucket 1?**

A: 0

B: 1

C: 2

D: 3

E: more than 3

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B: 1

C: 2

D: 3

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Example:

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```

Assuming the above example has been

executed, **how many elements are in bucket 2?**

A: 0

B: 1

C: 2

D: 3

E: more than 3

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C: 2

D: 3

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```
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set("orange", 40)
set("purplish", 30)
```

Assuming the above example has been
executed, **how many entries are checked for
get("purplish")?**

A: 0

B: 1

C: 2

D: 3

E: more than 3

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Please complete the below HashMap, assuming the entire example code has executed.

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```

0	— {blue: 90} — {pink: 100}
1	
2	
3	— {red: 70}

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set("purplish", 30)
```

0	— {blue: 90} — {pink: 100} — {purplish: 30}
1	
2	— {orange: 40}
3	— {red: 70}

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```

What is the load factor?

A: $\# \text{ elements} * \# \text{ buckets} / 2$

B: $\# \text{ buckets} * \# \text{ elements}$

C: $\# \text{ buckets} / \# \text{ elements}$

D: $\# \text{ elements} / \# \text{ buckets}$

What is the load factor?

A: $\# \text{ elements} * \# \text{ buckets} / 2$

B: $\# \text{ buckets} * \# \text{ elements}$

C: $\# \text{ buckets} / \# \text{ elements}$

D: $\# \text{ elements} / \# \text{ buckets}$

What is the load factor of the HashMap below?

0	— {blue: 90} — {pink: 100} — {purplish: 30}
1	
2	— {orange: 40}
3	— {red: 70}

What is the load factor of the HashMap below?

0	— {blue: 90} — {pink: 100} — {purplish: 30}
1	
2	— {orange: 40}
3	— {red: 70}

Load Factor: 5/4

What is the load factor after the line `set("red", 70)` is executed?

Example:

Start buckets array with size 4

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0	
1	
2	
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A `HashMap<Key, Value>` using Separate Chaining has:

- size: an int
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An Entry is a single {key: value} pair.

void set(key, value):

if LoadFactor > 0.5: expandCapacity()

... as before ...

void expandCapacity():

newBuckets = new List[this.buckets.length * 2];

oldBuckets = this.buckets

this.buckets = newBuckets

this.size = 0

for each list of entries in oldBuckets:

for each {k: v} in the list:

this.set(k, v)

What is the load factor after the line `set("red", 70)` is executed?

Example:

Start buckets array with size 4
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set("red", 70)
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```

0	
1	
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3	

Load Factor: 1/4

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Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	
1	
2	
3	— {red: 70}

Load Factor: 1/2

A `HashMap<Key, Value>` using Separate Chaining has:

- size: an int
- buckets: an array of lists of Entries
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

if LoadFactor > 0.5: expandCapacity()

... as before ...

void expandCapacity():

newBuckets = new List[this.buckets.length * 2];

oldBuckets = this.buckets

this.buckets = newBuckets

this.size = 0

for each list of entries in oldBuckets:

for each {k: v} in the list:

this.set(k, v)

What is the load factor after the line `set("pink", 100)` is executed?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	— {blue: 90}
1	
2	
3	— {red: 70}

A `HashMap<Key, Value>` using Separate Chaining has:

- size: an int
- buckets: an array of lists of Entries
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

if LoadFactor > 0.5: expandCapacity()

 ... as before ...

void expandCapacity():

 newBuckets = new List[this.buckets.length * 2];

 oldBuckets = this.buckets

 this.buckets = newBuckets

 this.size = 0

for each list of entries in oldBuckets:

for each {k: v} in the list:

 this.set(k, v)

What is the load factor after the line `set("pink", 100)` is executed?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	— {blue: 90}
1	
2	
3	— {red: 70}

Load Factor: 3/4

A `HashMap<Key, Value>` using Separate Chaining has:

- size: an int
- buckets: an array of lists of Entries
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

if LoadFactor > 0.5: expandCapacity()

... as before ...

void expandCapacity():

newBuckets = new List[this.buckets.length * 2];

oldBuckets = this.buckets

this.buckets = newBuckets

this.size = 0

for each list of entries in oldBuckets:

for each {k: v} in the list:

this.set(k, v)

What is different when the line `set("orange", 40)` is executed?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	— {blue: 90} — {pink: 100}
1	
2	
3	— {red: 70}

A `HashMap<Key, Value>` using Separate Chaining has:

- size: an int
- buckets: an array of lists of Entries
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

if LoadFactor > 0.5: expandCapacity()

... as before ...

void expandCapacity():

newBuckets = new List[this.buckets.length * 2];

oldBuckets = this.buckets

this.buckets = newBuckets

this.size = 0

for each list of entries in oldBuckets:

for each {k: v} in the list:

this.set(k, v)

What is different when the line `set("orange", 40)` is executed?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	— {blue: 90} — {pink: 100}
1	
2	
3	— {red: 70}

**expandCapacity()
is called!**

A `HashMap<Key, Value>` using Separate Chaining has:

- size: an int
- buckets: an array of lists of Entries
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

if LoadFactor > 0.5: expandCapacity()

... as before ...

void expandCapacity():

newBuckets = new List[this.buckets.length * 2];

oldBuckets = this.buckets

this.buckets = newBuckets

this.size = 0

for each list of entries in oldBuckets:

for each {k: v} in the list:

this.set(k, v)

What does the HashMap look like after `expandCapacity` is called in `set("orange", 40)`?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	
1	
2	
3	

4	
5	
6	
7	

A `HashMap<Key, Value>` using Separate Chaining has:

- `size`: an `int`
- `buckets`: an array of lists of `Entries`
- `hash`: a hash function for the `Key` type

An `Entry` is a single `{key: value}` pair.

```
void set(key, value):
```

```
    if LoadFactor > 0.5: expandCapacity()
```

```
    ... as before ...
```

```
void expandCapacity():
```

```
    newBuckets = new List[this.buckets.length * 2];
```

```
    oldBuckets = this.buckets
```

```
    this.buckets = newBuckets
```

```
    this.size = 0
```

```
    for each list of entries in oldBuckets:
```

```
        for each {k: v} in the list:
```

```
            this.set(k, v)
```

What does the HashMap look like after `expandCapacity` is called in `set("orange", 40)`?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	
1	
2	
3	— {red: 70}

4	— {blue: 90} - {pink: 100}
5	
6	
7	

A `HashMap<Key, Value>` using Separate Chaining has:

- size: an int
- buckets: an array of lists of Entries
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

if LoadFactor > 0.5: expandCapacity()

... as before ...

void expandCapacity():

newBuckets = new List[this.buckets.length * 2];

oldBuckets = this.buckets

this.buckets = newBuckets

this.size = 0

for each list of entries in oldBuckets:

for each {k: v} in the list:

this.set(k, v)

What does the HashMap look like after set("orange", 40) is called?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	
1	
2	
3	— {red: 70}

4	— {blue: 90} - {pink: 100}
5	
6	
7	

A HashMap<Key, Value> using Separate Chaining has:

- size: an int
- buckets: an array of lists of Entries
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

if LoadFactor > 0.5: expandCapacity()

... as before ...

void expandCapacity():

newBuckets = new List[this.buckets.length * 2];

oldBuckets = this.buckets

this.buckets = newBuckets

this.size = 0

for each list of entries in oldBuckets:

for each {k: v} in the list:

this.set(k, v)

What does the HashMap look like after set("orange", 40) is called?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	
1	
2	
3	{red: 70}

4	{blue: 90} - {pink: 100}
5	
6	{orange: 40}
7	

A HashMap<Key, Value> using Separate Chaining has:

- size: an int
- buckets: an array of lists of Entries
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

if LoadFactor > 0.5: expandCapacity()

... as before ...

void expandCapacity():

newBuckets = new List[this.buckets.length * 2];

oldBuckets = this.buckets

this.buckets = newBuckets

this.size = 0

for each list of entries in oldBuckets:

for each {k: v} in the list:

this.set(k, v)

What does the HashMap look like after set("purplish", 40) is called?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```

0	
1	
2	
3	{red: 70}

4	{blue: 90} - {pink: 100}
5	
6	{orange: 40}
7	

A HashMap<Key, Value> using Separate Chaining has:

- size: an int
- buckets: an array of lists of Entries
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

if LoadFactor > 0.5: expandCapacity()

... as before ...

void expandCapacity():

newBuckets = new List[this.buckets.length * 2];

oldBuckets = this.buckets

this.buckets = newBuckets

this.size = 0

for each list of entries in oldBuckets:

for each {k: v} in the list:

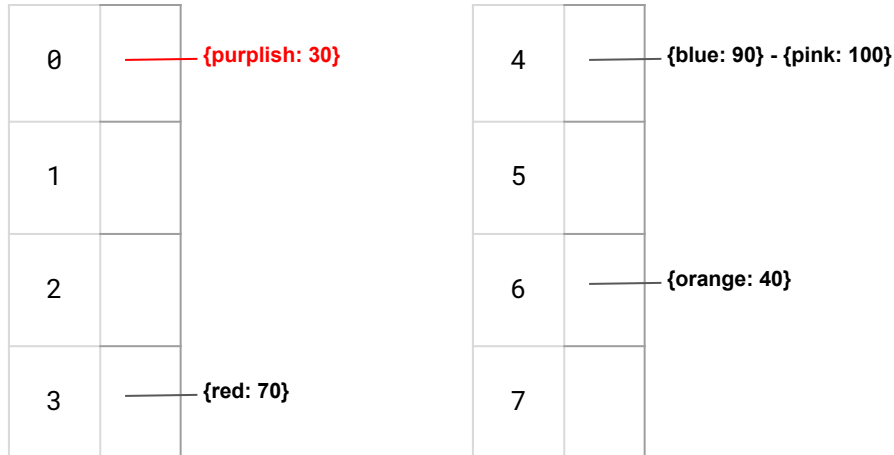
this.set(k, v)

What does the HashMap look like after set("purplish", 30) is called?

Example:

Start buckets array with size 4
Use string length as the hash function (**In general this is a BAD hash function**)

```
set("red", 70)
set("blue", 90)
set("pink", 100)
set("orange", 40)
set("purplish", 30)
```



A HashMap<Key, Value> using Separate Chaining has:

- size: an int
- buckets: an array of lists of Entries
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

if LoadFactor > 0.5: expandCapacity()

... as before ...

void expandCapacity():

newBuckets = new List[this.buckets.length * 2];

oldBuckets = this.buckets

this.buckets = newBuckets

this.size = 0

for each list of entries in oldBuckets:

for each {k: v} in the list:

this.set(k, v)

HashMaps - Linear Probing

What does the HashMap below look like after the example code has executed?

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2  
set("f", 90)  
set("f", 100)
```

0	
1	
2	
3	

A `HashMap<Key, Value>` using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

```
void set(key, value):  
    if loadFactor > 0.67: expandCapacity()  
    hashed = hash(key)  
    index = hashed % array length  
    while this.buckets[index] != null:  
        b = this.buckets[index]  
        if b.key.equals(key):  
            b.value = value  
            return  
        index += 1  
  
    // key not in table, add it at first index containing null  
    this.buckets[index] = {key: value}
```

```
Value get(key):  
    hashed = hash(key)  
    index = hashed % this.buckets.length  
    while this.buckets[index] != null:  
        b = this.buckets[index]  
        if b.key.equals(key): return b.value  
        index += 1  
  
    // haven't found the key  
    return null/throw exception
```

```
void expandCapacity():  
    newEntries = new Entry[this.buckets.length * 2];  
    oldEntries = this.buckets  
    this.buckets = newEntries  
    this.size = 0  
    for each entry {k:v} in oldEntries:  
        this.set(k, v)
```

What does the HashMap below look like after the example code has executed?

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2  
set("f", 90)  
set("f", 100)
```

0	
1	
2	— {b: 70}
3	— {f: 100}

A HashMap<Key, Value> using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

```
void set(key, value):  
    if loadFactor > 0.67: expandCapacity()  
    hashed = hash(key)  
    index = hashed % array length  
    while this.buckets[index] != null:  
        b = this.buckets[index]  
        if b.key.equals(key):  
            b.value = value  
            return  
        index += 1  
  
    // key not in table, add it at first index containing null  
    this.buckets[index] = {key: value}
```

```
Value get(key):  
    hashed = hash(key)  
    index = hashed % this.buckets.length  
    while this.buckets[index] != null:  
        b = this.buckets[index]  
        if b.key.equals(key): return b.value  
        index += 1  
  
    // haven't found the key  
    return null/throw exception
```

```
void expandCapacity():  
    newEntries = new Entry[this.buckets.length * 2];  
    oldEntries = this.buckets  
    this.buckets = newEntries  
    this.size = 0  
    for each entry {k:v} in oldEntries:  
        this.set(k, v)
```

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2
set("f", 90) # note 'f' = 102
set("f", 100)
```

Assuming the above example has been
executed, **how many elements are in bucket 1?**

A: 0

B: 1

C: 2

D: 3

E: more than 3

A HashMap<Key, Value> using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

```
if loadFactor > 0.67: expandCapacity()
hashed = hash(key)
index = hashed % array length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key):
        b.value = value
        return
    index += 1
```

```
// key not in table, add it at first index containing null
this.buckets[index] = {key: value}
```

Value get(key):

```
hashed = hash(key)
index = hashed % this.buckets.length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key): return b.value
    index += 1
```

```
// haven't found the key
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];
oldEntries = this.buckets
this.buckets = newEntries
this.size = 0
for each entry {k:v} in oldEntries:
    this.set(k, v)
```

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2
set("f", 90)
set("f", 100)
```

Assuming the above example has been
executed, **how many elements are in bucket 1?**

A: 0

B: 1

C: 2

D: 3

E: more than 3

A HashMap<Key, Value> using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

```
if loadFactor > 0.67: expandCapacity()
hashed = hash(key)
index = hashed % array length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key):
        b.value = value
        return
    index += 1
```

```
// key not in table, add it at first index containing null
this.buckets[index] = {key: value}
```

Value get(key):

```
hashed = hash(key)
index = hashed % this.buckets.length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key): return b.value
    index += 1
```

```
// haven't found the key
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];
oldEntries = this.buckets
this.buckets = newEntries
this.size = 0
for each entry {k:v} in oldEntries:
    this.set(k, v)
```

Example:
Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2  
set("f", 90)  
set("f", 100)
```

Assuming the above example has been
executed, **how many elements are in bucket 2?**

- A: 0
- B: 1
- C: 2
- D: 3
- E: more than 3

A HashMap<Key, Value> using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

```
if loadFactor > 0.67: expandCapacity()  
hashed = hash(key)  
index = hashed % array length  
while this.buckets[index] != null:  
    b = this.buckets[index]  
    if b.key.equals(key):  
        b.value = value  
        return  
index += 1
```

```
// key not in table, add it at first index containing null  
this.buckets[index] = {key: value}
```

Value get(key):

```
hashed = hash(key)  
index = hashed % this.buckets.length  
while this.buckets[index] != null:  
    b = this.buckets[index]  
    if b.key.equals(key): return b.value  
index += 1
```

```
// haven't found the key  
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];  
oldEntries = this.buckets  
this.buckets = newEntries  
this.size = 0  
for each entry {k:v} in oldEntries:  
    this.set(k, v)
```

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2
set("f", 90)
set("f", 100)
```

Assuming the above example has been
executed, **how many elements are in bucket 2?**

A: 0

B: 1

C: 2

D: 3

E: more than 3

A HashMap<Key, Value> using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

```
if loadFactor > 0.67: expandCapacity()
hashed = hash(key)
index = hashed % array length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key):
        b.value = value
        return
    index += 1
```

```
// key not in table, add it at first index containing null
this.buckets[index] = {key: value}
```

Value get(key):

```
hashed = hash(key)
index = hashed % this.buckets.length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key): return b.value
    index += 1
```

```
// haven't found the key
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];
oldEntries = this.buckets
this.buckets = newEntries
this.size = 0
for each entry {k:v} in oldEntries:
    this.set(k, v)
```

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2
set("f", 90)
set("f", 100)
```

Assuming the above example has been
executed, **how many elements are in bucket 3?**

A: 0

B: 1

C: 2

D: 3

E: more than 3

A HashMap<Key, Value> using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

```
if loadFactor > 0.67: expandCapacity()
hashed = hash(key)
index = hashed % array length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key):
        b.value = value
        return
    index += 1
```

```
// key not in table, add it at first index containing null
this.buckets[index] = {key: value}
```

Value get(key):

```
hashed = hash(key)
index = hashed % this.buckets.length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key): return b.value
    index += 1
```

```
// haven't found the key
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];
oldEntries = this.buckets
this.buckets = newEntries
this.size = 0
for each entry {k:v} in oldEntries:
    this.set(k, v)
```


Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2
set("f", 90)
set("f", 100)
```

Assuming the above example has been
executed, **how many elements are in bucket 3?**

A: 0

B: 1

C: 2

D: 3

E: more than 3

A HashMap<Key, Value> using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

```
if loadFactor > 0.67: expandCapacity()
hashed = hash(key)
index = hashed % array length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key):
        b.value = value
        return
    index += 1
```

```
// key not in table, add it at first index containing null
this.buckets[index] = {key: value}
```

Value get(key):

```
hashed = hash(key)
index = hashed % this.buckets.length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key): return b.value
    index += 1
```

```
// haven't found the key
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];
oldEntries = this.buckets
this.buckets = newEntries
this.size = 0
for each entry {k:v} in oldEntries:
    this.set(k, v)
```

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2
set("f", 90)
set("f", 100)
```

Assuming the above example has been executed,
how many entries are checked when doing
`set("f", 100)`?

A: 0

B: 1

C: 2

D: 3

E: more than 3

A `HashMap<Key, Value>` using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

```
if loadFactor > 0.67: expandCapacity()
hashed = hash(key)
index = hashed % array length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key):
        b.value = value
        return
    index += 1
```

```
// key not in table, add it at first index containing null
this.buckets[index] = {key: value}
```

Value get(key):

```
hashed = hash(key)
index = hashed % this.buckets.length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key): return b.value
    index += 1
```

```
// haven't found the key
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];
oldEntries = this.buckets
this.buckets = newEntries
this.size = 0
for each entry {k:v} in oldEntries:
    this.set(k, v)
```

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2  
set("f", 90)  
set("f", 100)
```

Assuming the above example has been executed,
how many entries are checked when doing
`set("f", 100)`?

A: 0

B: 1

C: 2

D: 3

E: more than 3

A HashMap<Key, Value> using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

```
if loadFactor > 0.67: expandCapacity()  
hashed = hash(key)  
index = hashed % array length  
while this.buckets[index] != null:  
    b = this.buckets[index]  
    if b.key.equals(key):  
        b.value = value  
        return  
    index += 1
```

```
// key not in table, add it at first index containing null  
this.buckets[index] = {key: value}
```

Value get(key):

```
hashed = hash(key)  
index = hashed % this.buckets.length  
while this.buckets[index] != null:  
    b = this.buckets[index]  
    if b.key.equals(key): return b.value  
    index += 1
```

```
// haven't found the key  
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];  
oldEntries = this.buckets  
this.buckets = newEntries  
this.size = 0  
for each entry {k:v} in oldEntries:  
    this.set(k, v)
```

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2
set("f", 90)
set("f", 100)
```

Assuming the above example has been executed,
what will the result of `get("f")` be after this
sequence?

A: 70

B: 90

C: 100

D: null

E: error

A `HashMap<Key, Value>` using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

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while this.buckets[index] != null:
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        b.value = value
        return
    index += 1
```

```
// key not in table, add it at first index containing null
this.buckets[index] = {key: value}
```

Value get(key):

```
hashed = hash(key)
index = hashed % this.buckets.length
while this.buckets[index] != null:
    b = this.buckets[index]
    if b.key.equals(key): return b.value
    index += 1
```

```
// haven't found the key
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];
oldEntries = this.buckets
this.buckets = newEntries
this.size = 0
for each entry {k:v} in oldEntries:
    this.set(k, v)
```

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2  
set("f", 90)  
set("f", 100)
```

Assuming the above example has been executed,
what will the result of `get("f", 100)` be after
this sequence?

A: 70

B: 90

C: 100

D: null

E: error

A `HashMap<Key, Value>` using Linear Probing has:

- size: an int
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// key not in table, add it at first index containing null  
this.buckets[index] = {key: value}
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Value get(key):

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hashed = hash(key)  
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while this.buckets[index] != null:  
    b = this.buckets[index]  
    if b.key.equals(key): return b.value  
    index += 1
```

```
// haven't found the key  
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];  
oldEntries = this.buckets  
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this.size = 0  
for each entry {k:v} in oldEntries:  
    this.set(k, v)
```

What does the HashMap below look like after the example code has executed?

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2  
set("f", 90)  
set("f", 100)
```

0	
1	
2	
3	

A `HashMap<Key, Value>` using Linear Probing has:

- size: an int
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- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

```
void set(key, value):  
    if loadFactor > 0.67: expandCapacity()  
    hashed = hash(key)  
    index = hashed % array length  
    while this.buckets[index] != null:  
        b = this.buckets[index]  
        if b.key.equals(key):  
            b.value = value  
            return  
        index += 1  
  
    // key not in table, add it at first index containing null  
    this.buckets[index] = {key: value}
```

```
Value get(key):  
    hashed = hash(key)  
    index = hashed % this.buckets.length  
    while this.buckets[index] != null:  
        b = this.buckets[index]  
        if b.key.equals(key): return b.value  
        index += 1  
  
    // haven't found the key  
    return null/throw exception
```

```
void expandCapacity():  
    newEntries = new Entry[this.buckets.length * 2];  
    oldEntries = this.buckets  
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    this.size = 0  
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```

What does the HashMap below look like after the example code has executed?

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2  
set("f", 90)  
set("f", 100)
```

0	
1	
2	— {b: 70}
3	— {f: 100}

A `HashMap<Key, Value>` using Linear Probing has:

- size: an int
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An Entry is a single {key: value} pair.

```
void set(key, value):  
    if loadFactor > 0.67: expandCapacity()  
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    index = hashed % array length  
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        b = this.buckets[index]  
        if b.key.equals(key):  
            b.value = value  
            return  
        index += 1  
  
    // key not in table, add it at first index containing null  
    this.buckets[index] = {key: value}
```

```
Value get(key):  
    hashed = hash(key)  
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    // haven't found the key  
    return null/throw exception
```

```
void expandCapacity():  
    newEntries = new Entry[this.buckets.length * 2];  
    oldEntries = this.buckets  
    this.buckets = newEntries  
    this.size = 0  
    for each entry {k:v} in oldEntries:  
        this.set(k, v)
```

Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2  
set("f", 90)  
set("f", 100)
```

Assuming the above example has been executed,
and an additional line is added below: `set("c", 40)`,
Which bucket is "c" stored in?

A: 0

B: 1

C: 2

D: 3

E: it causes an error

A `HashMap<Key, Value>` using Linear Probing has:

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- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

void set(key, value):

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if loadFactor > 0.67: expandCapacity()  
hashed = hash(key)  
index = hashed % array length  
while this.buckets[index] != null:  
    b = this.buckets[index]  
    if b.key.equals(key):  
        b.value = value  
        return  
    index += 1
```

```
// key not in table, add it at first index containing null  
this.buckets[index] = {key: value}
```

Value get(key):

```
hashed = hash(key)  
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while this.buckets[index] != null:  
    b = this.buckets[index]  
    if b.key.equals(key): return b.value  
    index += 1
```

```
// haven't found the key  
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];  
oldEntries = this.buckets  
this.buckets = newEntries  
this.size = 0  
for each entry {k:v} in oldEntries:  
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```


Example:

Start buckets array with size 4, containing null
ASCII code as hash function ("a" = 97)

```
set("b", 70) # note 98 % 4 is 2
set("f", 90)
set("f", 100)
```

Assuming the above example has been executed,
and an additional line is added below: `set("c", 40)`,
Which bucket is "c" stored in?

A: 0

B: 1

C: 2

D: 3

E: it causes an error (`ArrayIndexOutOfBoundsException`)

A `HashMap<Key, Value>` using Linear Probing has:

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An Entry is a single {key: value} pair.

void set(key, value):

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    b = this.buckets[index]
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        b.value = value
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    index += 1
```

```
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this.buckets[index] = {key: value}
```

Value get(key):

```
hashed = hash(key)
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while this.buckets[index] != null:
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    if b.key.equals(key): return b.value
    index += 1
```

```
// haven't found the key
return null/throw exception
```

void expandCapacity():

```
newEntries = new Entry[this.buckets.length * 2];
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this.buckets = newEntries
this.size = 0
for each entry {k:v} in oldEntries:
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How can we fix the ArrayOutOfBounds issue?

A `HashMap<Key, Value>` using Linear Probing has:

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An Entry is a single {key: value} pair.

```
void set(key, value):
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    if loadFactor > 0.67: expandCapacity()
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```

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this.buckets[index] = {key: value}
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Value get(key):
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    hashed = hash(key)
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        if b.key.equals(key): return b.value
        index += 1
```

```
// haven't found the key
return null/throw exception
```

```
void expandCapacity():
```

```
    newEntries = new Entry[this.buckets.length * 2];
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A `HashMap<Key, Value>` using Linear Probing has:

- size: an int
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An Entry is a single {key: value} pair.

When you get to the end of the array just fall off the end, wrap around to the beginning, and starting searching again at 0.

We would no longer have `ArrayIndexOutOfBoundsException` issue!

Loadfactor - never update size!!! Where should we increment size?

Asume load factor is `size/currentlength` (helper method)

```
void set(key, value):
    if loadFactor > 0.67: expandCapacity()
    hashed = hash(key)
    index = hashed % array length
    while this.buckets[index] != null:
        b = this.buckets[index]
        if b.key.equals(key):
            b.value = value
            return
        index += 1
    index = index % buckets.length
    // key not in table, add it at first index containing null
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```

```
Value get(key):
    hashed = hash(key)
    index = hashed % this.buckets.length
    while this.buckets[index] != null:
        b = this.buckets[index]
        if b.key.equals(key): return b.value
        index += 1
    index = index % buckets.length
    // haven't found the key
    return null/throw exception
```

```
void expandCapacity():
    newEntries = new Entry[this.buckets.length * 2];
    oldEntries = this.buckets
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    this.size = 0
    for each entry {k:v} in oldEntries:
        this.set(k, v)
```

Are there any other issues that need to be fixed?

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    while this.buckets[index] != null:
        b = this.buckets[index]
        if b.key.equals(key):
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```
void expandCapacity():
    newEntries = new Entry[this.buckets.length * 2];
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Are there any other issues that need to be fixed?

A `HashMap<Key, Value>` using Linear Probing has:

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An Entry is a single {key: value} pair.

YES! size is never being updated!

(Let's assume `loadFactor` is actually a helper method that returns the current size divided by the current length.)

```
void set(key, value):
    if loadFactor() > 0.67: expandCapacity()
    hashed = hash(key)
    index = hashed % array length
    while this.buckets[index] != null:
        b = this.buckets[index]
        if b.key.equals(key):
            b.value = value
            return
        index += 1
    index = index % buckets.length
    // key not in table, add it at first index containing null
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    size += 1
```

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Value get(key):
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        b = this.buckets[index]
        if b.key.equals(key): return b.value
        index += 1
    index = index % buckets.length
    // haven't found the key
    return null/throw exception
```

```
void expandCapacity():
    newEntries = new Entry[this.buckets.length * 2];
    oldEntries = this.buckets
    this.buckets = newEntries
    this.size = 0
    for each entry {k:v} in oldEntries:
        this.set(k, v)
```

What happens if we set loadFactor to be 1 instead of 0.67?

A `HashMap<Key, Value>` using Linear Probing has:

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        if b.key.equals(key):
            b.value = value
            return
        index += 1
    index = index % buckets.length
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    size += 1
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```
Value get(key):
    hashed = hash(key)
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        b = this.buckets[index]
        if b.key.equals(key): return b.value
        index += 1
    index = index % buckets.length
    // haven't found the key
    return null/throw exception
```

```
void expandCapacity():
    newEntries = new Entry[this.buckets.length * 2];
    oldEntries = this.buckets
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A `HashMap<Key, Value>` using Linear Probing has:

- size: an int
- buckets: an array of Entries (not of lists of Entries!)
- hash: a hash function for the Key type

An Entry is a single {key: value} pair.

INFINITE LOOP!

There would be an infinite loop once the array is full. If the array is full of entries the method will search until it finds a bucket equal to null and there is no null to find.

```
void set(key, value):
```

```
    if loadFactor() > 0.67: expandCapacity()
    hashed = hash(key)
    index = hashed % array length
    while this.buckets[index] != null:
        b = this.buckets[index]
        if b.key.equals(key):
            b.value = value
            return
        index += 1
    index = index % buckets.length
    // key not in table, add it at first index containing null
    this.buckets[index] = {key: value}
    size += 1
```

```
Value get(key):
```

```
    hashed = hash(key)
    index = hashed % this.buckets.length
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        b = this.buckets[index]
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void expandCapacity():
```

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    newEntries = new Entry[this.buckets.length * 2];
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    this.size = 0
    for each entry {k:v} in oldEntries:
        this.set(k, v)
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