

CSE 12 – Basic Data Structures and Object-Oriented Design

Lecture 9

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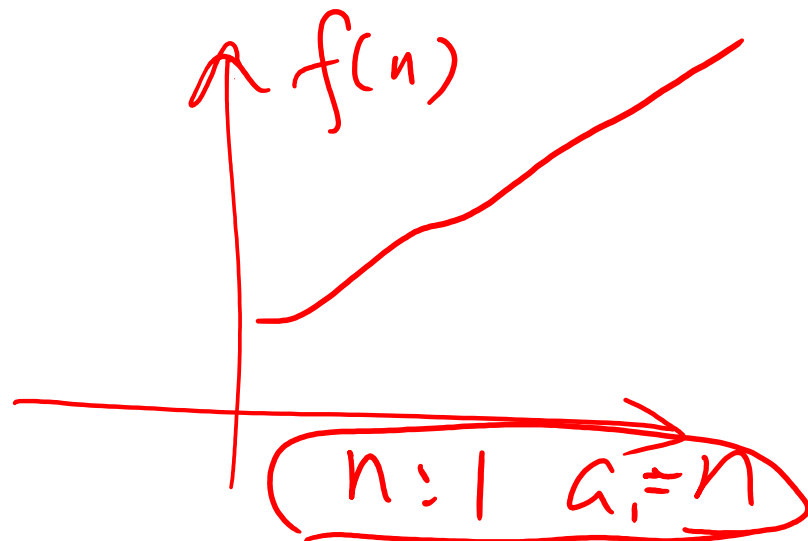
Announcements

- Quiz 9 due Wednesday @ 8am
- PA3 due Wednesday @ 11:59pm
- Survey 4 due Friday @ 11:59pm
- Exam 1 released Friday @ 8am, due Saturday before ~~8am~~ 12

Topics

- Questions on Lecture 9?
- Counting Steps

Questions on Lecture 9?



$$\left\{ \begin{array}{l} a_n = a_{n-1} + 7 \\ a_{n+1} = a_n + 7 \end{array} \right.$$

$n=0, a_0=0$

Analyzing the worst case

bool find(arr, ele)
for index = 0 -- n-1
if (arr[index] == ele)
return true

```
boolean find( String[] theList, String toFind ) {  
    for ( int i = 0; i < theList.length; i++ ) {  
        if ( theList[i].equals( toFind ) ) return false  
        return true;  
    }  
    return false;  
}
```

How many instructions do you have to execute to find out if the element is in the list in the worst case, if n represents the length of the list?

Analyzing the worst case

$i = 0$
①

$i < \text{length}$
 $n+1$

$i++$
 $n+1$

Code	# of instr.
<code>boolean find(String[] theList, String toFind) {</code>	N/A
<code> for (int i = 0; <u>$i < \text{theList.length}$</u>; i++) {</code>	$i = 0$ ① $i < \text{length}$ $n+1$ $i++$ n
<code> <u>\Rightarrow if (theList[i].equals(toFind))</u> 1</code>	① $\times n$
<code> return true;</code>	0
<code> }</code>	
<code> return false; 4</code>	1
<code>}</code>	

How many instructions do you have to execute to find out if the element is in the list in the worst case, if n represents the length of the list?

$$1 + (n+1) + n$$

n

~~$$2n + 3$$~~

$i: 0-1 \quad 1-2 \quad \dots \quad n-1 \dots n$
linear algorithm

Analyzing the worst case

n is large
fundamental difference

```
boolean find( String[] theList, String toFind ) {  
    for ( int i = 0; i < theList.length; i++ ) {  
        if ( theList[i].equals(toFind) )  
            return true;  
    }  
    return false;  
}
```

$$\underline{3n + 3}$$

```
boolean mysteryFind( String[] theList, String toFind ) {  
    int count = 0; ①  
    for ( int i = 0; i < theList.length; i++ ) {  
        count = count + 1;  
        if ( theList[i].equals(toFind) )  
            return true;  
    }  
    return false;  
}
```

$$\underline{4n + 4}$$

Which method is faster?

- A. find
- B. mysteryFind
- ☒ C. They are about the same

Analyzing the worst case : n large

```
boolean find( String[] theList, String toFind ) {  
    for ( int i = 0; i < theList.length; i++ ) {  
        if ( theList[i].equals(toFind) )  
            return true;  
    }  
    return false;  
}
```

n

```
boolean fastFind( String[] theList, String toFind ) {  
    return false;  
}
```

for (i=0 ; i < 1000 ; i++) {
 }
return false

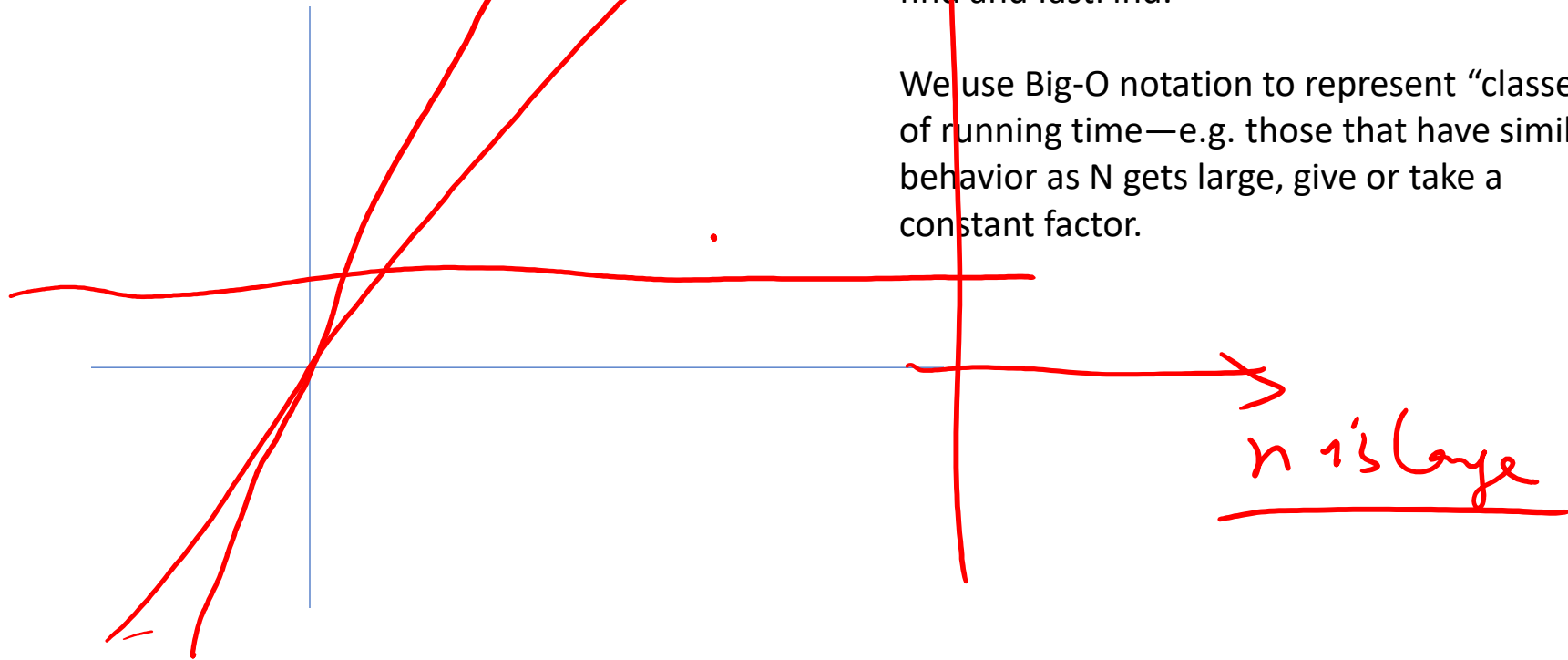
10^6

constant

Which method is faster?

- A. find
- B. fastFind
- C. They are about the same

$3n$ vs $2n$ vs 1



(Play)lists typically have 1000s (or more!) elements, so we only care about very large n .

find and mysteryFind are “more similar” than find and fastFind.

We use Big-O notation to represent “classes” of running time—e.g. those that have similar behavior as N gets large, give or take a constant factor.

Steps for calculating the Big O (Theta, Omega) bound on code or algorithms

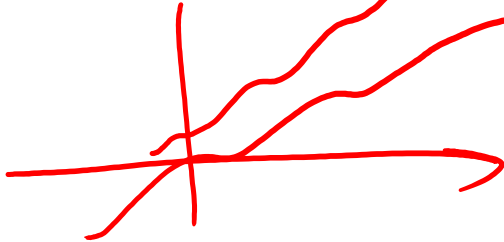
$$f(n) = \cancel{1}n^2 + 7n - 6$$

1. Count the number of instructions in your code (or algorithm) as precisely as possible as a function of n , which represents the size of your input (e.g. the length of the array). This is your $f(n)$.

- Make sure you know if you are counting best case, worst case or average case – could be any of these!

2. Simplify your $f(n)$ to find a simple $g(n)$ such that $f(n) = O(g(n))$ (or $\Omega(g(n))$ or $\Theta(g(n))$)

$$f(n) \in O(n^2)$$
$$\in O(n^3)$$



Counting Steps – ArrayList Insert – ignore EC

n : size of ArrayList

```
public void insert(int index, String s) {  
    expandCapacity(); //ignore  
    for (int i = size - 1; i >= index ; i--) {  
        this.elements[i+1] = this.elements[i];  
    }  
    2 this.elements[index] = s;  
    this.size += 1;  
}
```

Linear
 n

A: about
 n

B: n^2

C: n^3

about

Counting Steps – ArrayList Expand Capacity

```
private void expandCapacity() {  
    | int currentCapacity = this.elements.length;  
    | if(this.size < currentCapacity) { return; }  
    | String[] expanded = new String[currentCapacity * 2];  
    | for(int i = 0; i < this.size; i += 1) {  
n |     expanded[i] = this.elements[i];  
    | }  
    | this.elements = expanded;  
    | }  
}
```

A: n

B: n^2

C: $n \lg n$

D: $2dk$

series of actions: worst part dominates about

Counting Steps – ArrayList Insert w/ EC

```
public void insert(int index, String s) {  
     $n \rightarrow$  expandCapacity();  
    for (int i = size - 1; i >= index ; i--) {  
         $n$  this.elements[i+1] = this.elements[i];  
    }  
    this.elements[index] = s;  
    this.size += 1;  
}
```

n

n

n

n

A: n

B: n^2

C: n^3

D: $n \lg n$

E: $2dk$

Challenge

```
void tricky(int n) {  
    int operations = 0;  
    while(n > 0) {  
        for(int i = 0; i < n; i++) {  
            cout << "Operations: " << operations++ << endl;  
        }  
        n /= 2;  
    }  
}
```

What is the *step counting result* for the algorithm above?

A: n

B: $n \lg n$

C: n^2

D: n^3

E: $2nk$