CSE 12 — Basic Data Structures and Object-Oriented Design Lecture 14

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Announcements

- Quiz 14 due Wednesday @ 8am
- PA5 due Wednesday @ 11:59pm
- Survey 6 due Friday @ 11:59pm

Topics

- Sorting Wrap-up
- Questions on Lecture 14?

Questions on Lecture 14?

```
import java.util.Arrays;
public class Sort {
static void selectionSort(int[] arr) {
 for(int i = 0; i < arr.length; i += 1) {
  int minIndex = i;
  for(int j = i; j < arr.length; j += 1) {
   if(arr[minIndex] > arr[j]) { minIndex = j; }
  int temp = arr[i];
  arr[i] = arr[minIndex];
  arr[minIndex] = temp;
```

```
static void insertionSort(int[] arr) {
 for(int i = 0; i < arr.length; i += 1) {
  for(int j = i; j > 0; j = 1) {
    if(arr[i] < arr[i-1]) {
     int temp = arr[j-1];
     arr[j-1] = arr[j];
     arr[j] = temp;
    else { break; } // new! exit inner loop early
```

```
import java.util.Arrays;
public class SortFaster {
 static int[] combine(int[] p1, int[] p2) {...}
 static int[] mergeSort(int[] arr) {
  int len = arr.length
  if(len <= 1) { return arr; }
  else {
   int[] p1 = Arrays.copyOfRange(arr, 0, len / 2);
   int[] p2= Arrays.copyOfRange(arr, len / 2, len);
   int[] sortedPart1 = mergeSort(p1);
   int[] sortedPart2 = mergeSort(p2);
    int[] sorted = combine(sortedPart1, sortedPart2);
   return sorted;
```

```
static int partition(String[] array, int I, int h) {...}
static void qsort(String[] array, int low, int high) {
 if(high - low <= 1) { return; }
 int splitAt = partition(array, low, high);
 qsort(array, low, splitAt);
 qsort(array, splitAt + 1, high);
public static void sort(String[] array) {
 qsort(array, 0, array.length);
```

Quick sort

sort {12, 4, 9, 3, 15, 8, 19, 2}

```
Quicksort(Arr, low, high)
  if (low < high)
    pivotPos = partition(Arr, low, high)
    Quicksort(Arr, low, pivotPos - 1)
    Quicksort(Arr, pivotPos + 1, high)</pre>
```

```
pivot = Arr[high]
i = low - 1
for (j = low; j < high, j++) {
   if (Arr[j] < pivot)
        i++
        swap(A[i], A[j])
swap(A[i+1], A[high])
return i + 1</pre>
```

partition(Arr, low, high)

	Insertion	Selection	Merge	Quick
Best case time				
Worst case time				
Key operations	swap(a, j, j-1) (until in the right	swap(a, i, indexOfMin) (after finding minimum	I = copy(a, 0, len/2) r = copy(a, len/2, len)	p = partition(a, l, h) sort(a, l, p)
	place)	value)	s = sort(I) rs = sort(r) merge(Is, rs)	sort(a, p + 1, h)

Non-comparison based sorting

- Normally it is for integer sorting
- Count sort
 - Assume that we have n positive integers and we know that all of them are

```
Count sort (A, B, k):
  let C[0, ..., k] be a new array and initialized to be 0
  for j = 1 to A.length
    C[A[j]]++
                                                    Example {12, 4, 9, 3, 15, 8, 19, 2}
  for i = 1 to k
    C[i] = C[i] + C[i-1]
  for j = A.length down to 1
    B[C[A[\dot{j}]]] = A[\dot{j}]
    C[A[j]] = C[A[j]] - 1
```

Last note about sorting

- Not only do we care about runtime, we also care about
 - Space: do we need extra storage?
 - stable: if we have duplicates, do we maintain the same ordering?

Algorithm	Space	Stable
Bubble sort	O(1)	Yes
Selection sort	O(1)	No
Insertion sort	O(1)	Yes
Heap sort	O(1)	No
Merge sort	O(n)	Yes
Quick sort	O(logn)	No
Count sort	O(n+k)	Yes

Array Sorting Algorithms (wiki)

A loo with no	Time Complexity				
Algorithm	Best	Average	Worst		
Quicksort	$\Omega(n \log(n))$	Θ(n log(n))	O(n^2)		
Mergesort	$\Omega(n \log(n))$	Θ(n log(n))	O(n log(n))		
<u>Timsort</u>	Ω(n)	Θ(n log(n))	O(n log(n))		
<u>Heapsort</u>	Ω(n log(n))	Θ(n log(n))	O(n log(n))		
Bubble Sort	Ω(n)	Θ(n^2)	O(n^2)		
Insertion Sort	Ω(n)	Θ(n^2)	O(n^2)		
Selection Sort	Ω(n^2)	Θ(n^2)	O(n^2)		
Tree Sort	Ω(n log(n))	Θ(n log(n))	O(n^2)		
Shell Sort	Ω(n log(n))	Θ(n(log(n))^2)	O(n(log(n))^2)		
Bucket Sort	Ω(n+k)	Θ(n+k)	O(n^2)		
Radix Sort	Ω(nk)	Θ(nk)	O(nk)		
Counting Sort	Ω(n+k)	Θ(n+k)	O(n+k)		
Cubesort	Ω(n)	Θ(n log(n))	O(n log(n))		