

CSE 12 – Basic Data Structures and Object-Oriented Design

Lecture 10

Paul Cao and Greg Miranda, Winter 2021

This lecture is being recorded

Announcements

- Quiz 10 due Monday @ 8am
- Survey 4 due Friday @ 11:59pm
- PA3 due tonight @ 11:59pm
- Exam 1 on Friday (no class)
 - Released @ 8am on Friday
 - Closes @ 10am on Saturday
 - More details on Piazza

1/29
1/30

details

Topics

- Questions on Lecture 10?
- Big O

$10^6 \in O(n)$ A yes
B no

Questions on Lecture 10?

\exists no, C

$f(n) \in O(g(n))$

cgen

When n is large,

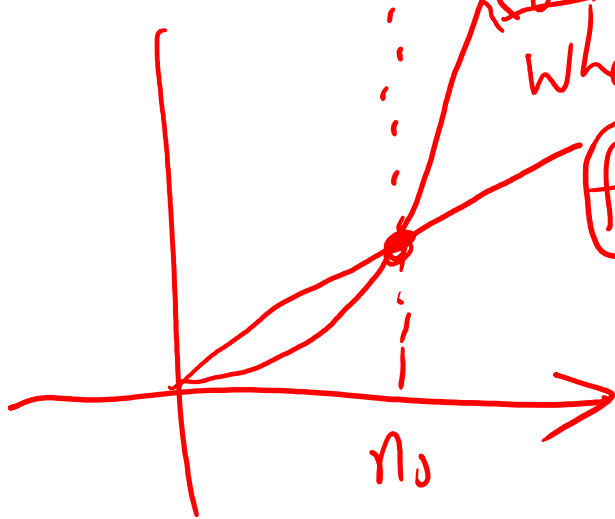
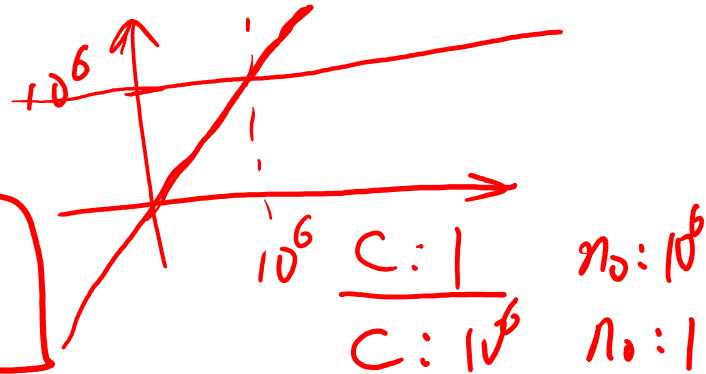
$f(n) \leq c g(n)$

where c is a const.

$O(n)$ \leftarrow set

$\frac{n}{f(n)} \in \frac{O(n)}{g(n)}$

A yes
B no



$O \iff \leq$

Let $f(n) = 100$

is : \in

- Which of the following is NOT a correct bound?

$f(n)$ is $O(1)$



A yes

is no

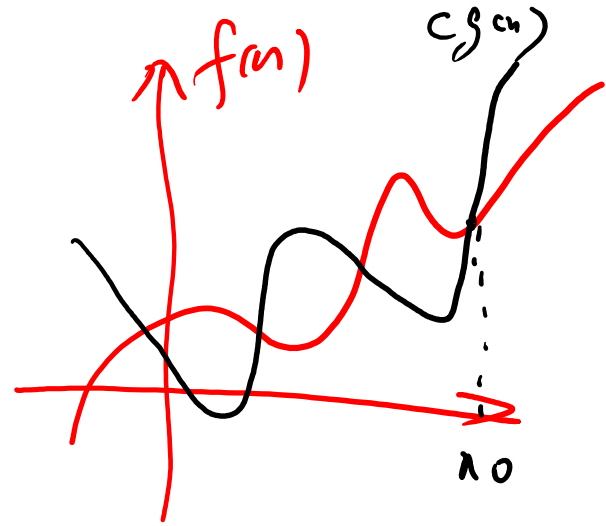
A. $f(n)$ is $O(2^n)$

B. $f(n)$ is $O(n^2)$

C. $f(n)$ is $O(n)$

D. $f(n)$ is $O(n^{100})$

E. None of these



$$O(\underline{6_2^n}) \text{ same as } O(\underline{6_3^n})$$

$$O(1)$$

For each function in the list below, it is related to the function below it by O , and the reverse is **not** true. That is, n is $O(n^2)$ but n^2 is **not** $O(n)$.

- $f(n) = 1/(n^2)$ ~~←~~
- $f(n) = 1/n$ ~~←~~
- $f(n) = 1$ ←
- $f(n) = \log(n)$ ← binary search
- $f(n) = \text{sqrt}(n)$
- $f(n) = n$
- $f(n) = n^2$
- $f(n) = n^3$
- $f(n) = n^4$
- ... and so on for constant polynomials ...
- $f(n) = 2^n$ ← very bad
- $f(n) = n!$ ← factorial really bad
- $f(n) = n^n$

polynomial

very bad
factorial really bad

$$O(2^n) \times O(3^n)$$

base 4
exp matter

$$O\left(\frac{6^n}{\sqrt{n}}\right)$$

$$O(n)$$

$$O(n \log n)$$

$$O(n^2)$$

$$O(\underline{2^n})$$

base log
doesn't matter

$$\log_2 n = \frac{\ln n}{\ln 2}$$

$$\log_3 n = \frac{\ln n}{\ln 3}$$

$$O(n!)$$

Let $f(n) = \cancel{3n^3} + \underline{2n} + \underline{7}$

dominant term
remove constant coeff

• Which of the following is a correct bound?

- A. $f(n)$ is $O(\log(n))$
- B. $f(n)$ is $O(n^2)$
- C. $f(n)$ is $O(n)$
- ☒ D. $f(n)$ is $O(n^3)$
- E. None of these

```
void printAllElementOfArray(int[] arr) {
    for (int i = 0; i < arr.length; i++) {
        printf("%d\n", arr[i]);
    }
}
```

• Which of the following is a correct bound?

A. $f(n)$ is $O(\log(n))$

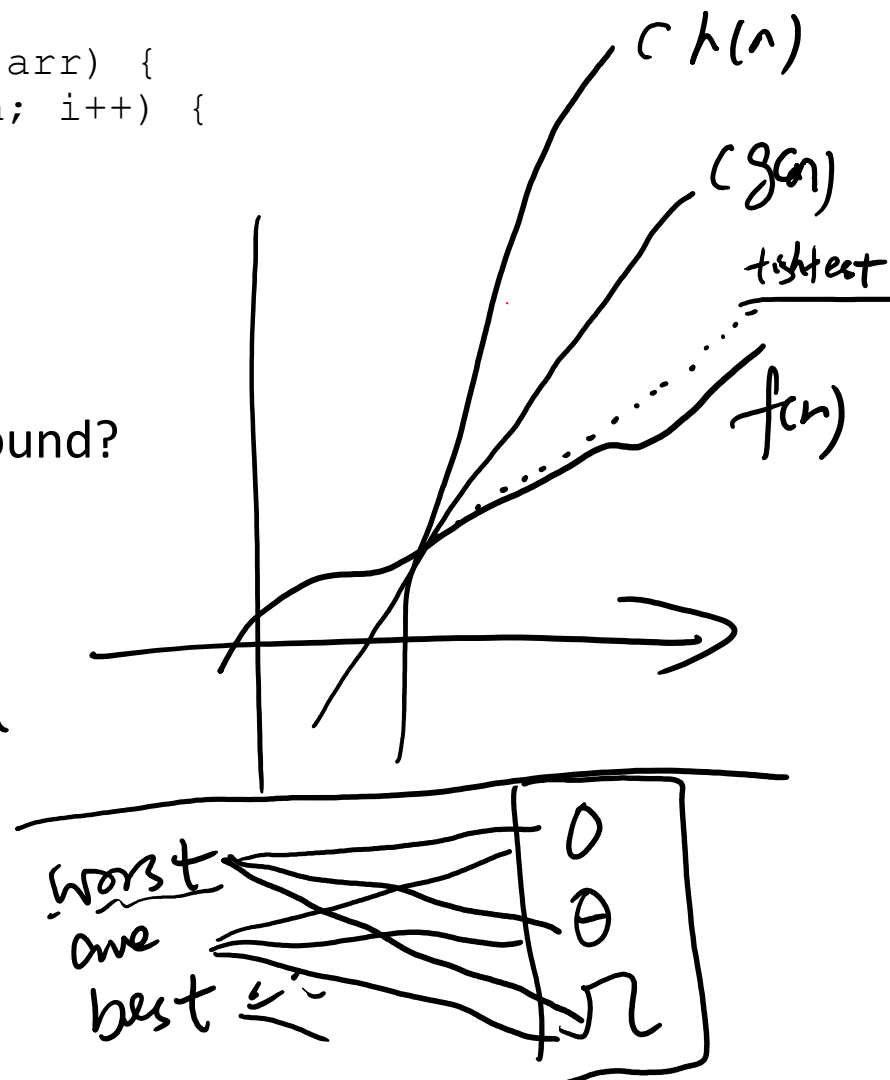
B. $f(n)$ is $O(n^2)$

C. $f(n)$ is $O(n)$

D. $f(n)$ is $O(n^3)$

E. More than one of these

← tight bound




```

void printAllPossibleOrderedPairs(int arr[]) {
    for (int i = 0; i < arr.length; i++) {
        for (int j = 0; j < arr.length; j++) {
            printf("%d = %d\n", arr[i], arr[j]);
        }
    }
}

```

$$\sum_{n} \sum_{n} 1 = n^2$$

• Which of the following is a correct bound?

A. $f(n)$ is $O(\log(n))$

B. $f(n)$ is $O(n^2)$

C. $f(n)$ is $O(n)$

D. $f(n)$ is $O(n^3)$

E. More than one of these

highest $O(n^2)$

```
int fibonacci(int num) {
    if (num <= 1) return num;
    return fibonacci(num - 2) + fibonacci(num - 1);
}
```

• Which of the following is a correct bound?

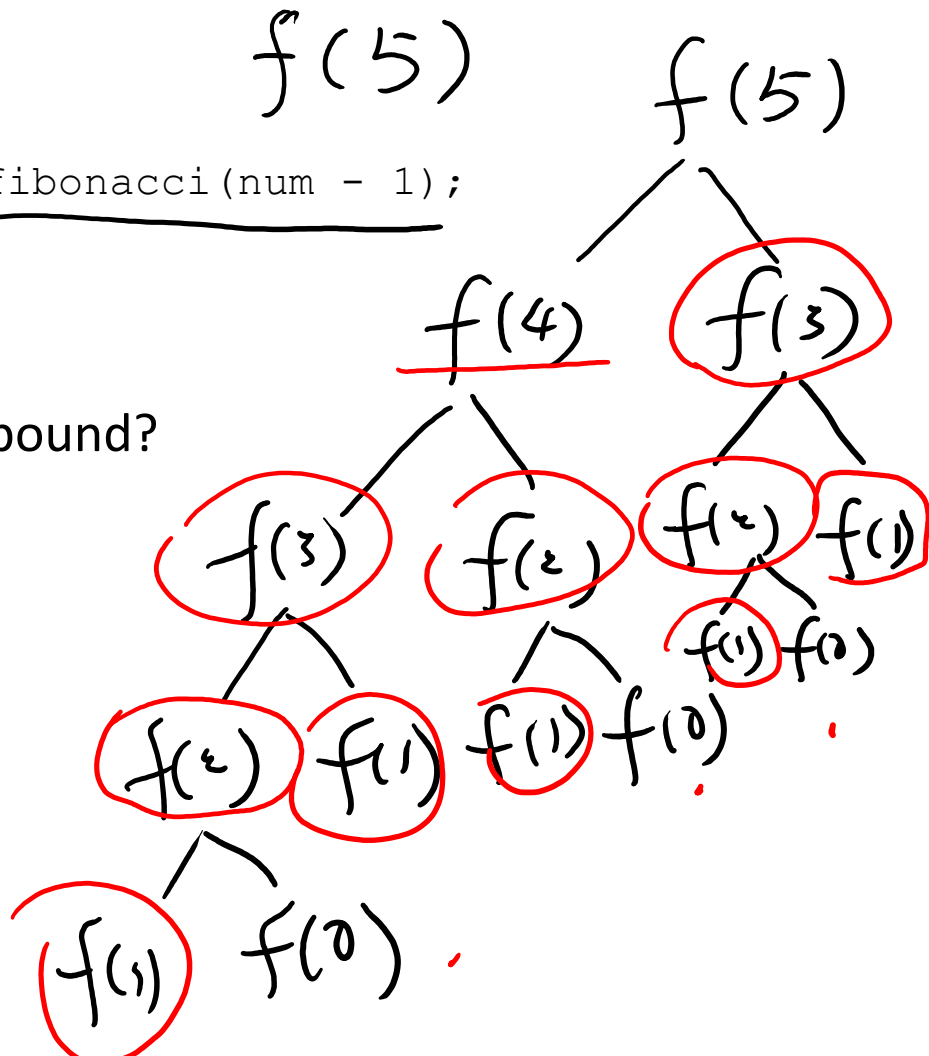
A. $f(n)$ is $O(2^n)$

B. $f(n)$ is $O(n^2)$

C. $f(n)$ is $O(n)$

D. $f(n)$ is $O(n^3)$

E. More than one of these



$$2^n = 2 \cdot 2^{n-1}$$

$$\boxed{f(n) = f(n-1) + f(n-2)}$$

$$T(n-2) \leq T(n-1)$$

$$T(n) = T(n-1) + \underline{T(n-2)} + C$$

$$T(n) \leq \boxed{2^{n-1}}$$

$$\leq 2 \underline{T(n-1)} + C$$

$$T(n) \in \Theta\left(\left(\frac{1+\sqrt{5}}{2}\right)^n\right)$$

$$\leq 2 \left[\underline{2T(n-2) + C} \right] + C$$

$$= 2 \times \underline{2T(n-2)} + 4C + C$$

$$\leq \underline{2^{n-1}} \cdot \underline{T(n-(n-1))} + \dots$$

$$\underline{O(2^{n-1})}$$