

```
class Trie:
    root = None
```

```
class TrieNode:
    parent = None
    children = None
    key = None
    isEndOfWord = False
```

```
def newNode(key):
    new = TrieNode()
    new.children = []
    new.key = key
    return new
```

"""EJERCICIO 1"""

```
def isLetterOnChildren(node: TrieNode, letter: str) -> [bool, TrieNode | None]:
    """Returns a tuple: bool is found and the node itself (None if not found)."""
    for n in node.children:
        if n.key == letter:
            return True, n
    return False, None
```

```
def insert(T, element):
    if T.root is None:
        T.root = newNode(None)
    node = T.root
    for letter in element:
        letterFound, new = isLetterOnChildren(node, letter)
        if not letterFound:
            new = newNode(letter)
            node.children.append(new)
            new.parent = node
            node = new
    node.isEndOfWord = True
    return
```

```
def searchWordLastLetterNode(T, element) -> [bool, TrieNode | None]:
    """Given a word, returns True if it is on the Trie and the last letter node, otw False and None"""
    if T is None or T.root is None or not element:
        raise Warning("Empty Trie, root or word")
    node = T.root
    for letter in element:
        letterFound, node = isLetterOnChildren(node, letter)
        if not letterFound:
            return False, None
    return True, node
```

```
def search(T, element):
    return searchWordLastLetterNode(T, element)[0]
```

"""EJERCICIO 2:

Para que la operación `search()` tenga un orden de complejidad  $O(m)$ , es decir, que la cardinalidad del alfabeto ( $|S|$ ) no influya en el tiempo de búsqueda, podríamos modificar la estructura Trie de forma tal que cada campo `children` se exprese mediante un TAD en la que la búsqueda de un elemento dentro de ella sea de  $O(1)$ . De esta forma, lo único que resta verificar es que todas las  $m$  letras/partes del elemento estén en

el Trie. Esto puede lograrse utilizando, por ejemplo, un TAD que en clase hemos empezado a ver y satisface esta necesidad, hablamos del Diccionario, de esta forma el acceso por key es inmediato."""

"""EJERCICIO 3"""

```
def delete(T, element):
    wordFound, lastLetterNode = searchWordLastLetterNode(T, element)
    if not wordFound:
        return False
    lastLetterNode.isEndOfWord = False

    # A word can be a leaf nor an inner node
    if lastLetterNode.children is []: # is inner, job done
        return False
    node = lastLetterNode # is leaf
    while True:
        if node is None or node.isEndOfWord: # reached the root or a new word ending
            return True
        nodeParent = node.parent
        nodeParent.children.remove(node)
        node = nodeParent
```

"""EJERCICIO 7"""

```
def autoComplete(T, prefijo):
    """Devuelve la parte restante de una raíz de 2 o + palabras que tengan a -prefijo-
    como parte de su raíz."""
    cadenaEncontrada, nodoUltimaLetra = searchWordLastLetterNode(T, prefijo)
    if not cadenaEncontrada:
        return ''
    resto = ''
    while True:
        if len(nodoUltimaLetra.children) > 1:
            return '' if not resto else resto
        nodoSigue = nodoUltimaLetra.children[0]
        resto += nodoSigue.key
        nodoUltimaLetra = nodoSigue
```

"""EJERCICIO 4"""

```
def withPrefix(T: Trie, prefix: str, size: int) -> list:
    """Returns a list with all words with -prefix- and length -size-."""

    def withPrefixR(node, size, finalList, remaining):
        if size == 0 and node.isEndOfWord:
            finalList.append(remaining)

        for child in node.children:
            withPrefixR(child, size - 1, finalList, remaining + child.key) # decrease
            remaining while adding key
        return finalList # return when loop finishes (all possible situations where
            verified)

    wordFound, lastLetterNode = searchWordLastLetterNode(T, prefix)
    if not wordFound or size <= 0 or size <= len(prefix):
        return [] # prefix not found or input errors

    remainingSize = size - len(prefix) # so to see when it reaches zero
    return withPrefixR(lastLetterNode, remainingSize, [], prefix)
```

"""EJERCICIO 5"""

```

def sameDocument(T1: Trie, T2: Trie) -> bool:
    """Tells if T1 and T2 have all same words, iow, they are the same Trie.
    El costo es de  $O(m*n*\log(n))$ , donde m es el número total de nodos y n es el número de
    hijos por cada nodo m.
    El costo  $n*\log(n)$  se debe a que cada llamada aplica dos ordenamientos de listas."""

    def nodeKey(node):
        return node.key

    def sameDocumentR(node1, node2):
        if len(node1.children) != len( # here are all the logical cases where a doc-node
            node2.children) or node1.key != node2.key or node1.isEndOfWord !=
            node2.isEndOfWord:
            return False
        keySorted1 = sorted(node1.children, key=nodeKey)
        keySorted2 = sorted(node2.children, key=nodeKey)
        for i in range(len(keySorted1)): # checks that all nodes satisfy sameDocument
            condition
            if not sameDocumentR(keySorted1[i], keySorted2[i]): # only if a False
                condition is found, return
            return False
        return True # else, continue up to here - all nodes and their children checked

    if T1 and T2:
        return sameDocumentR(T1.root, T2.root)
    return False

```

"""EJERCICIO 6"""

```

def hasReversedStrings(T):
    node = T.root
    stack = [(node, '')]
    while stack: # while there is still a node to be searched
        node, word = stack.pop() # bring back the last node
        if node.isEndOfWord and search(T, word[::-1]): # when an EOW was found, search
            its reversed form
            return True
        for child in node.children:
            stack.append((child, word + child.key)) # append a node so to analyze it
    return False

```