

Tugas Praktikum Analisis Algoritma
KOMPLEKSITAS WAKTU



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Mencari Nilai Max

```
procedure CariMaks(input  $x_1, x_2, \dots, x_n$ : integer, output maks: integer)
{  Mencari elemen terbesar dari sekumpulan elemen larik integer  $x_1, x_2, \dots, x_n$ . Elemen terbesar akan
   disimpan di dalam maks
   Input:  $x_1, x_2, \dots, x_n$ 
   Output: maks (nilai terbesar)
}
```

Deklarasi

i : integer

Algoritma

```
maks  $\leftarrow x_1$ 
 $i \leftarrow 2$ 
while  $i \leq n$  do
  if  $x_i > \text{maks}$  then
    maks  $\leftarrow x_i$ 
  endif
   $i \leftarrow i + 1$ 
endwhile
{ $i > n$ }
```

PROGRAM C++

```
#include <iostream>
using namespace std;

int main(){
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++){
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }

    int maks = x[0];
    int i = 1;
    while (i <= n) {
        if (x[i] > maks)
            maks = x[i];
        i++;
    }
```

```

}
cout << "Maksimum Number : " << maks << endl;

return 0;
}

```

Kompleksitas Waktu

```

maks ← x1      1 kali
i ← 2          1 kali
maks ← xi      n kali
i ← i + 1      n kali

```

$$T(n) = 1 + 1 + n + n = 2n + 2$$

Sequential Search

```

procedure SequentialSearch(input  $x_1, x_2, \dots, x_n$  : integer, y : integer, output idx : integer)
{   Mencari y di dalam elemen  $x_1, x_2, \dots, x_n$ . Lokasi (indeks elemen) tempat y ditemukan diisi ke dalam
    idx. Jika y tidak ditemukan, maka idx diisi dengan 0.
    Input:  $x_1, x_2, \dots, x_n$ 
    Output: idx
}

```

Deklarasi

```

i : integer
found : boolean { bernilai true jika y ditemukan atau false jika y tidak ditemukan }

```

Algoritma

```

i ← 1
found ← false
while (i ≤ n) and (not found) do
    if  $x_i = y$  then
        found ← true
    else
        i ← i + 1
    endif
endwhile

```

PROGRAM C++

```
#include <iostream>
using namespace std;

int main() {
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++){
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }

    int y;
    cout << "Masukkan yang dicari : ";
    cin >> y;

    int i = 0;
    bool found = false;
    int idx;
    while ((i < n) && (!found)){
        if (x[i] == y)
            found = true;
        else
            i++;
    }
    if (found)
        idx = i+1;
    else
        idx = 0;

    cout << "Yang dicari berada di urutan : " << idx << endl;

    return 0;
}
```

Kompleksitas Waktu

- Best Case

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow l$	1 kali

$$T_{\min}(n) = 1 + 1 + 1 + 1 = 4$$

- Average Case

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$i \leftarrow i + 1$	$\frac{1}{2} n$ kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow l$	1 kali

$$T_{\text{avg}}(n) = 1 + 1 + \frac{1}{2} n + 1 + 1 = \frac{1}{2} n + 4$$

- Worst Case

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$i \leftarrow i + 1$	n kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow l$	1 kali

$$T_{\max}(n) = 1 + 1 + n + 1 + 1 = n + 4$$

Binary Search

```
procedure BinarySearch(input  $x_1, x_2, \dots, x_n$  : integer,  $x$  : integer, output :  $idx$  : integer)
{ Mencari  $y$  di dalam elemen  $x_1, x_2, \dots, x_n$ . Lokasi (indeks elemen) tempat  $y$  ditemukan diisi ke dalam
   $idx$ . Jika  $y$  tidak ditemukan maka  $idx$  diisi dengan 0.
  Input:  $x_1, x_2, \dots, x_n$ 
  Output:  $idx$ 
}

Deklarasi
   $i, j, mid$  : integer
  found : Boolean

Algoritma
   $i \leftarrow 1$ 
   $j \leftarrow n$ 
  found  $\leftarrow$  false
  while (not found) and ( $i \leq j$ ) do
     $mid \leftarrow (i + j) \text{ div } 2$ 
    if  $x_{mid} = y$  then
      found  $\leftarrow$  true
    else
      if  $x_{mid} < y$  then {mencari di bagian kanan}
         $i \leftarrow mid + 1$ 
      else {mencari di bagian kiri}
         $j \leftarrow mid - 1$ 
      endif
    endif
  endwhile
  {found or  $i > j$ }

  If found then
     $Idx \leftarrow mid$ 
  else
     $Idx \leftarrow 0$ 
  endif
  { $i < n$  or found}

  If found then { $y$  ditemukan}
     $idx \leftarrow i$ 
  else
     $idx \leftarrow 0$  { $y$  tidak ditemukan}
  endif
```

PROGRAM C++

```
#include <iostream>
using namespace std;

int main() {
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++){
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }

    int y;
    cout << "Masukkan yang dicari : ";
    cin >> y;

    int i = 0;
    int j = n-1;
    bool found = false;
    int idx;
    int mid;
    while ((i <= j) && (!found)){
        mid = (i + j)/2;
        if (x[mid] == y)
            found = true;
        else{
            if (x[mid] < y)
                i = mid + 1;
            else
                j = mid - 1;
        }
    }

    if (found)
        idx = mid+1;
    else
        idx = 0;

    cout << "Yang dicari berada di urutan : " << idx << endl;

    return 0;
}
```

Kompleksitas Waktu

- Best Case

$i \leftarrow 1$	1 kali
$j \leftarrow n$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$\text{mid} \leftarrow (i + j) \text{ div } 2$	1 kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{Idx} \leftarrow \text{mid}$	1 kali

$$T_{\min}(n) = 1 + 1 + 1 + 1 + 1 + 1 = 6$$

- Average Case

$i \leftarrow 1$	1 kali
$j \leftarrow n$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$\text{mid} \leftarrow (i + j) \text{ div } 2$	$\frac{1}{2}n + 1$ kali
$i \leftarrow \text{mid} + 1$ or $j \leftarrow \text{mid} - 1$	$\frac{1}{2}n$ kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{Idx} \leftarrow \text{mid}$	1 kali

$$T_{\text{avg}}(n) = 1 + 1 + 1 + \frac{1}{2}n + 1 + \frac{1}{2}n + 1 + 1 = n + 6$$

- Worst Case

$i \leftarrow 1$	1 kali
$j \leftarrow n$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$\text{mid} \leftarrow (i + j) \text{ div } 2$	$n + 1$ kali
$i \leftarrow \text{mid} + 1$ or $j \leftarrow \text{mid} - 1$	n kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{Idx} \leftarrow \text{mid}$	1 kali

$$T_{\max}(n) = 1 + 1 + 1 + n + 1 + n + 1 + 1 = 2n + 6$$

Insertion Sort

```
procedure InsertionSort(input/output  $x_1, x_2, \dots, x_n$  : integer)
{
    Mengurutkan elemen-elemen  $x_1, x_2, \dots, x_n$  dengan metode insertion sort.
    Input:  $x_1, x_2, \dots, x_n$ 
    Output:  $x_1, x_2, \dots, x_n$  (sudah terurut menaik)
}
Deklarasi
    i, j, insert : integer
Algoritma
    for i  $\leftarrow$  2 to n do
        insert  $\leftarrow$   $x_i$ 
        j  $\leftarrow$  i
        while (j < i) and ( $x[j-i] >$  insert) do
             $x[j] \leftarrow x[j-1]$ 
            j  $\leftarrow$  j-1
        endwhile
         $x[j] =$  insert
    endfor
    {i < n or found}
    If found then {y ditemukan}
```

PROGRAM C++

```
#include <iostream>
using namespace std;

int main()
{
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++)
    {
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }
    cout << "Data Sebelum di Sorting : ";
    for (int i = 0; i < n; i++)
        cout << x[i] << " ";
    cout << endl;

    int insert;
```

```

int j;
for (int i = 1; i < n; i++)
{
    insert = x[i];
    j = i-1;
    while ((j >= 0) && (x[j] > insert))
    {
        x[j+1] = x[j];
        j--;
    }
    x[j+1] = insert;
}

cout << "Data setelah di Sorting : ";
for (int i = 0; i < n; i++)
    cout << x[i] << " ";

return 0;
}

```

Kompleksitas Waktu

- Best Case

For i ← 2 to n do	1 kali
insert ← x _i	n kali
j ← i	n kali
x[j] = insert	n kali

$$T_{min}(n) = 1 + n + n + n = 3n + 1$$

- Average Case

For i ← 2 to n do	1 kali
insert ← x _i	n kali
j ← i	n kali
x[j] ← x[j-1]	$n * \frac{1}{2} n$ kali
j ← j-1	$n * \frac{1}{2} n$ kali
x[j] = insert	n kali

$$T_{avg}(n) = 1 + n + n + \frac{1}{2} n^2 + \frac{1}{2} n^2 + n = n^2 + 3n + 1$$

- Worst Case

For i \leftarrow 2 to n do	1 kali
insert \leftarrow x _i	n kali
j \leftarrow i	n kali
x[j] \leftarrow x[j-1]	n * n kali
j \leftarrow j-1	n * n kali
x[j] = insert	n kali

$$T_{max}(n) = 1 + n + n + n^2 + n^2 + n = 2n^2 + 3n + 1$$

Selection Sort

```

procedure SelectionSort(input/output  $x_1, x_2, \dots, x_n$  : integer)
{  Mengurutkan elemen-elemen  $x_1, x_2, \dots, x_n$  dengan metode selection sort.
  Input:  $x_1, x_2, \dots, x_n$ 
  Output:  $x_1, x_2, \dots, x_n$  (sudah terurut menaik)
}

Deklarasi
  i, j, imaks, temp : integer

Algoritma
  for i  $\leftarrow$  n downto 2 do {pass sebanyak n-1 kali}
    imaks  $\leftarrow$  1
    for j  $\leftarrow$  2 to i do
      if  $x_j > x_{imaks}$  then
        imaks  $\leftarrow$  j
      endif
    endfor
    {pertukarkan  $x_{imaks}$  dengan  $x_i$ }
    temp  $\leftarrow$   $x_i$ 
     $x_i \leftarrow x_{imaks}$ 
     $x_{imaks} \leftarrow$  temp
  endfor

```

PROGRAM C++

```
#include <iostream>
using namespace std;

int main(){
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++){
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }
    cout << "Data Sebelum di Sorting : ";
    for (int i = 0; i < n; i++)
        cout << x[i] << " ";
    cout << endl;

    int imaks;
    int temp;
    for (int i = n-1; i >= 1; i--){
        imaks = 0;
        for (int j = 1; j <= i; j++){
            if (x[j] > x[imaks])
                imaks = j;
        }
        temp = x[i];
        x[i] = x[imaks];
        x[imaks] = temp;
    }

    cout << "Data setelah di Sorting : ";
    for (int i = 0; i < n; i++)
        cout << x[i] << " ";

    return 0;
}
```

Kompleksitas Waktu

- Best Case

for i ← n downto 2 do	1 kali
imaks ← 1	n kali
for j ← 2 to i do	n kali
imaks ← j	n * 1 kali
temp ← xi	n kali
xi ← imaks	n kali
ximaks ← temp	n kali

$$T_{min}(n) = 1 + n + n + n * 1 + n + n + n = 6n + 1$$

- Average Case

for i ← n downto 2 do	1 kali
imaks ← 1	n kali
for j ← 2 to i do	n kali
imaks ← j	n * ½ n kali
temp ← xi	n kali
xi ← imaks	n kali
ximaks ← temp	n kali

$$T_{avg}(n) = 1 + n + n + \frac{1}{2}n^2 + n + n + n = \frac{1}{2}n^2 + 5n + 1$$

- Worst Case

for i ← n downto 2 do	1 kali
imaks ← 1	n kali
for j ← 2 to i do	n kali
imaks ← j	n * n kali
temp ← xi	n kali
xi ← imaks	n kali
ximaks ← temp	n kali

$$T_{max}(n) = 1 + n + n + n^2 + n + n + n = n^2 + 5n + 1$$