

CSI 3350: PROGRAMMING LANGUAGES

Department of Computer Science &
Engineering
Oakland University

Anonymous Functions

- $F(x) = (x + 1)$

- $g(x) = (x + 1)$

Lambda Expression

• $F(x) = (x + 1)$

function name binding variable function body

function definition

`(lambda (x) (+ x 1))`

Lambda Expression

• $F(x) = (x + 1)$

function name input parameter output

function definition

`(lambda (x) (+ x 1))`

Lambda Expression

• $F(x) = (x + 1)$

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`(lambda (x) (+ x 1))`

Lambda Expression

• $F(x) = (x + 1)$

function name input parameter output

function definition

`(lambda (x) (+ x 1))`

Lambda Expression

• $F(x) = (x + 1)$

function name input parameter output

anonymous function

`(lambda (x) (+ x 1))`

A Function of Two Forms of Definition

- $F(x) = (x + 1)$

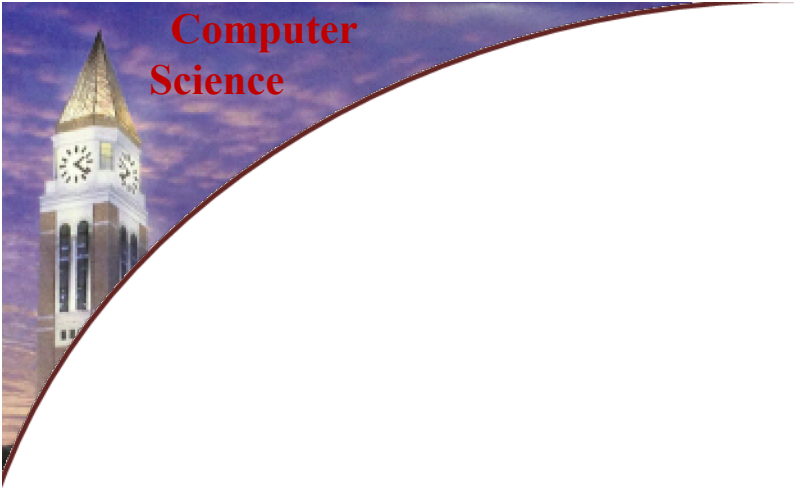
```
(define ( F x )  
  (+ x 1)  
)
```

```
(define F  
  ( lambda (x) (+ x 1) )  
)
```

- $G(x) = (x + 1)$

```
(define ( G x )  
  (+ x 1)  
)
```

```
(define G F)
```

In general, proving any two functions are the same (equivalent) or not is **undecidable** !

```
(define  
  (fact-new n )  
  (fact-tail n 1)  
)
```

```
(define  
  (fact-tail n prod)  
    (if  
      (= n 0)  
      prod  
      (fact-tail (- n 1) (* n prod) )  
    )  
)
```

tail recursion!

```
(define  
  (fact n )  
    (if  
      (= n 0)  
      1  
      (* n (fact (- n 1)))))
```

```
(define  
  (fact-new n )  
  (fact-tail n 1)  
)
```

```
(define  
  (fact-tail n prod)  
    (if  
      (= n 0)  
      prod  
      (fact-tail (- n 1) (* n prod) )  
    )  
)
```

```
(define  
  (fact n )  
    (if  
      (= n 0)  
      1  
      (* n (fact (- n 1)))))
```

```
(define  
  (fact-new n )  
  (fact-tail n 1)  
)
```

```
(define  
  (fact-tail n prod)  
  (if  
    (= n 0)  
    prod  
    (fact-tail (- n 1) (* n prod) )  
  )  
)
```

```
(define  
  (fact n )  
  (if  
    (= n 0)  
    1  
    (* n (fact (- n 1)))))
```

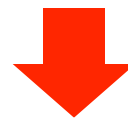
(fact-new n) and (fact n)
are the same, but proving any two
functions are equivalent or not is
undecidable !

To decide their equivalence between $(fact_new\ n)$ and $(fact\ n)$ is a lot more direct and involved than the simple example of $f(x) = x + 1$.

The **general answer** is that to decide in finite amount of time that any two functions are equivalent or not is **a task too hard** for a modern computer (Turing machine) to finish, so-called **undecidable**.

To decide their equivalence between $(fact_new\ n)$ and $(fact\ n)$ is a lot more direct and involved than the simple example of $f(x) = x + 1$.

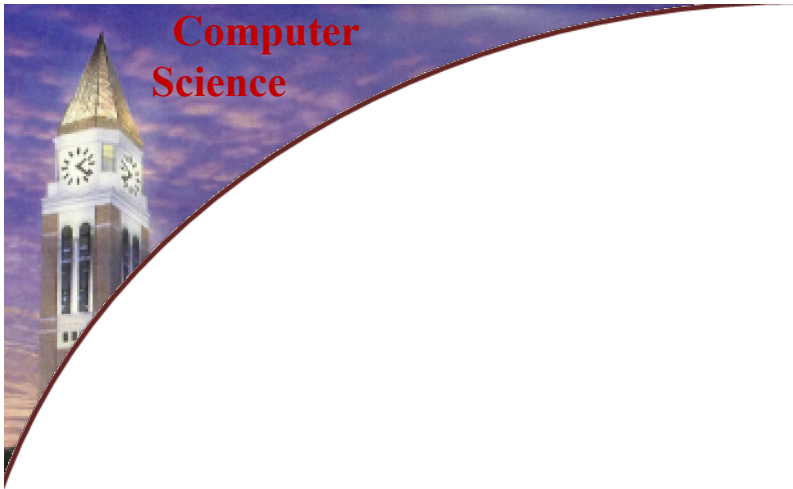
The **general answer** is that to decide in finite amount of time that any two functions are equivalent or not is **a task too hard** for a modern computer (Turing machine) to finish, so-called **undecidable**.



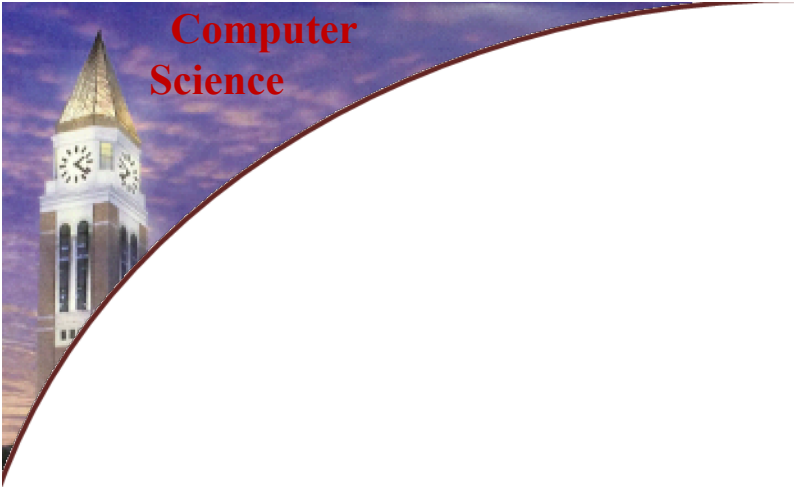
as a result

In Scheme language, the language designer just let the comparison of any functions to return a negative result, even for the simplest case i.e.,

```
(equal? (lambda (x) (+ x 1)) (lambda (x) (+ x 1))) = #f
```



HW02



HW02

Read the test file first !

Carpet

(carpet 3)

```
\ ( + + + + + + + )
    (+ % % % % % + )
    (+ % + + + % + )
    (+ % + % + % + )
    (+ % + + + % + )
    (+ % % % % % + )
    (+ + + + + + + ) )
```

(carpet 4)

```
\ ( (% % % % % % % % %)
    (% + + + + + + + %)
    (% + % % % % % + %)
    (% + % + + + % + %)
    (% + % + + + % + %)
    (% + % + % % % % + %)
    (% + + + + + + + %)
    (% % % % % % % % %) )
```

(carpet 4) =

step-1: for-each list in (carpet 3) expand it by adding \% to the beginning and end of it

↓

```
\ ( (% + + + + + + + %)
    (% + % % % % % + %)
    (% + % + + + % + %)
    (% + % + % + % + %)
    (% + % + + + % + %)
    (% + % % % % % + %)
    (% + + + + + + + %) )
```

↓

← **step-2:** add \% (% % % % % % % %) to the beginning and the end of the result returned by **step-1**

Carpet

(carpet 3)

```
\ ( + + + + + + + )
    (+ % % % % % + )
    (+ % + + + % + )
    (+ % + % + % + )
    (+ % + + + % + )
    (+ % % % % % + )
    (+ + + + + + + ) )
```

(carpet 4) =

step-1: for-each list in (carpet 3) expand it by adding \% to the beginning and end of it

↓

```
\ ( (% + + + + + + + %)
    (% + % % % % % + %)
    (% + % + + + % + %)
    (% + % + % + % + %)
    (% + % + + + % + %)
    (% + % % % % % + %)
    (% + + + + + + + %) )
```

(carpet 4)

```
\ ( (% % % % % % % %)
    (% + + + + + + + %)
    (% + % % % % % + %)
    (% + % + + + % + %)
    (% + % + + + % + %)
    (% + % + % + % + %)
    (% + % + + + + + %)
    (% % % % % % % %) )
```

← **step-2:** add \% (% % % % % % % %) to the beginning and the end of the result returned by **step-1**

Pascal Triangle

(define (pascal n) ...)

n = 1 \ ((1))

n = 2 \ ((1)
 (1 1))

n = 3 \ ((1)
 (1 1)
 (1 2 1))

n = 4 \ ((1)
 (1 1)
 (1 2 1)
 (1 3 3 1))


n = 3

<-> \ ((1) (1 1))

<-> \ ((1) (1 1) (1 2 1))

(pascal 4) = inserting '(1 3 3 1) to the end of
(pascal 3)

<-> \ ((1) (1 1) (1 2 1) (1 3 3 1))

Pascal Triangle

```
(define ( pascal n) ... )
```

```
n = 3      ` ( (1)
              (1 1)
              (1 2 1) )
```

```
n = 4      ` ( (1)
              (1 1)
              (1 2 1)
              (1 3 3 1) )
```

get the last element from (pascal 3)

generate ` (1 3 3 1)

insert it into (pascal 3)



(pascal 4) = inserting '(1 3 3 1) to the end of
(pascal 3)

Pascal Triangle

- Generate $\backslash (1 \ 3 \ 3 \ 1)$ from $\backslash (1 \ 2 \ 1)$

$$\begin{array}{ccccccc} & & & + & & + & \\ & & 1 & \rightarrow & 2 & \rightarrow & 1 \\ & & & \downarrow & & \downarrow & \\ & & & 3 & & 3 & \end{array}$$

Pascal Triangle

- Generate $\backslash (1 \ 3 \ 3 \ 1)$ from $\backslash (1 \ 2 \ 1)$

$$\begin{array}{cccc} \backslash & 1 & \xrightarrow{+} & 2 & \xrightarrow{+} & 1 & \backslash \\ & & \searrow & \downarrow & \searrow & \downarrow & \\ & & & 3 & & 3 & \backslash \end{array}$$

We first generate the core sub-list $\backslash (3 \ 3)$

Pascal Triangle

- Generate $\backslash (1 \ 3 \ 3 \ 1)$ from $\backslash (1 \ 2 \ 1)$

$$\begin{array}{ccccccc} & & + & & + & & \\ \backslash & (& 1 & \xrightarrow{\quad} & 2 & \xrightarrow{\quad} & 1 &) \\ & & \downarrow & & \downarrow & & \\ \backslash & (& 1 & & 3 & & 3 & & 1 &) \\ & & \uparrow & & & & & & \uparrow & \end{array}$$

At the very end, we just insert these two 1's into the beginning and end of $\backslash (3 \ 3)$

Pascal Triangle

```
(define ( pascal n) ... )
```

```
n = 3      ` ( (1)
              (1 1)
              (1 2 1) )
```

```
n = 4      ` ( (1)
              (1 1)
              (1 2 1)
              (1 3 3 1) )
```

get the last element from (pascal 3)

generate \downarrow ` (1 3 3 1)

insert it into \downarrow (pascal 3)



(pascal 4) = inserting '(1 3 3 1) to the end of
(pascal 3)

Pascal Triangle

- Generate $\backslash (1 \ 3 \ 3 \ 1)$ from $\backslash (1 \ 2 \ 1)$

$$\begin{array}{ccccccc} & & & + & & + & \\ & & 1 & \rightarrow & 2 & \rightarrow & 1 \\ & & & \downarrow & & \downarrow & \\ & 1 & & 3 & & 3 & & 1 \end{array}$$

At the very end, we just insert these two 1's into the beginning and end of $\backslash (3 \ 3)$

How to code it ?

Create-mapping

↓
' (I II III IV V) key-list
' (1 2 3 4 5) value-list
↑

Search-key: IV

Create-mapping

↓
' (I II III IV V) key-list
' (1 2 3 4 5) value-list
↑

Search-key: IV

Create-mapping

↓
'(II III IV V) key-list
'(2 3 4 5) value-list
↑

Search-key: IV

Create-mapping

↓
'(II III IV V) key-list
'(2 3 4 5) value-list
↑
Search-key: IV

Create-mapping

↓
'(III IV V) key-list
'(3 4 5) value-list
↑

Search-key: IV

Create-mapping

↓
'(III IV V) key-list

'(3 4 5) value-list
↑

Search-key: IV

Create-mapping

↓
'(IV V) key-list

'(4 5) value-list
↑

Search-key: IV

Create-mapping

'(IV V) key-list

'(4 5) value-list

Search-key: IV

Map Function

` (1 2 3)



` (1 8 27)

cubic each
number in a list

```
(map (lambda (x) (* x x x)) `(1 2 3) )
```