Orthogonal bases

A collection of vectors $\{\boldsymbol{v}_1, \boldsymbol{v}_2, \dots, \boldsymbol{v}_N\}$ in a finite dimensional vector space \mathcal{S} is called an **orthogonal basis** if

- 1. span($\{\boldsymbol{v}_1, \boldsymbol{v}_2, \dots, \boldsymbol{v}_N\}$) = \mathcal{S} ,
- 2. $\mathbf{v}_j \perp \mathbf{v}_k$ (i.e. $\langle \mathbf{v}_j, \mathbf{v}_k \rangle = 0$) for all $j \neq k$.

If in addition the vectors are normalized (under the induced norm),

$$\|\boldsymbol{v}_n\| = 1$$
, for $n = 1, \dots, N$,

we will call it an **orthonormal basis** or **orthobasis**.

A note on infinite dimensions

In infinite dimensions, we need to be a little more careful with what we mean by "span". Traditionally, the span is defined as the set of all possible linear combinations of *finitely many* elements of \mathcal{S} . Thus, if $\mathcal{B} = \{v_n\}_{n \in \mathbb{Z}}$ is an infinite sequence of orthogonal vectors in a Hilbert space \mathcal{S} , it is an orthobasis if the *closure* of span(\mathcal{B}) is \mathcal{S} ; this is written

$$\operatorname{cl}\operatorname{Span}\left(\{\boldsymbol{v}_n\}_n\right)=\mathcal{S}.$$

We don't need to get into too much, but basically this means that every vector in S can be approximated arbitrarily well by a finite linear combination of vectors in S.

Here is an example which illustrates the point: Let x(t) be any function on [0,1] which is not a polynomial — say $x(t) = \sin(2\pi t)$. Let $\mathcal{B} = \{1, t, t^2, t^3, \ldots\}$; the span (set of a finite linear combinations of elements) of \mathcal{B} is all polynomials on [0,1]. So $\mathbf{x} \notin \text{span}(\mathcal{B})$. But x(t) can be approximated arbitrarily well by elements in \mathcal{B} (using higher and higher order polynomials) so $\mathbf{x} \in \text{cl Span}(\mathcal{B})$).

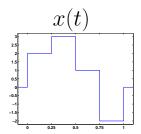
Examples.

1. $\mathcal{S} = \mathbb{R}^2$, equipped with the standard inner product

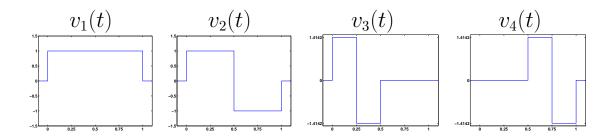
$$oldsymbol{v}_1 = rac{1}{\sqrt{2}} egin{bmatrix} 1 \ 1 \end{bmatrix}, \qquad oldsymbol{v}_2 = rac{1}{\sqrt{2}} egin{bmatrix} 1 \ -1 \end{bmatrix}$$

2. S = space of piecewise constant functions on [0, 1/4), [1/4, 1/2), [1/2, 3/4), [3/4, 1]

Example signal:



The following four functions form an orthobasis for this space



3. Fourier series

$$\left\{v_k(t) = \frac{1}{\sqrt{2\pi}}e^{jkt}, k \in \mathbb{Z}\right\}$$
 is an orthobasis for $L_2([0, 2\pi])$

(with the standard inner product).

Let's quickly check the orthogonality:

$$\left\langle \frac{1}{\sqrt{2\pi}} e^{jk_1 t}, \frac{1}{\sqrt{2\pi}} e^{jk_2 t} \right\rangle = \frac{1}{2\pi} \int_0^{2\pi} e^{j(k_1 - k_2)t} dt$$
$$= \begin{cases} 1, & k_1 = k_2 \\ 0, & k_1 \neq k_2 \end{cases}.$$

It is also true that the closure of span($\{(2\pi)^{-1/2}e^{jkt}\}_{k=-\infty}^{\infty}$) is $L_2([0,2\pi])$. The proof of this is a bit involved; if you are interested, see Chapter 5 of Young's *Introduction to Hilbert Space*.

4. Sampling

 $B_{\pi/T}(\mathbb{R})$ = real-valued functions which are bandlimited to π/T , equipped with the standard inner product. The set of functions

$$\left\{ v_n(t) = \sqrt{T} \frac{\sin(\pi(t - nT)/T)}{\pi(t - nT)}, \ n \in \mathbb{Z} \right\}$$

is an orthobasis for $B_{\pi/T}(\mathbb{R})$. (Notice that we have a slightly different normalization than when we looked at the sampling theorem — we have a \sqrt{T} out front instead of T.)

Check the orthogonality:

$$\left\langle \sqrt{T} \frac{\sin(\pi(t - n_1 T)/T)}{\pi(t - n_1 T)}, \sqrt{T} \frac{\sin(\pi(t - n_2 T)/T)}{\pi(t - n_2 T)} \right\rangle$$

$$= \frac{1}{2\pi} \int_{-\pi/T}^{\pi/T} T e^{-j\Omega T n_1} e^{j\Omega T n_2} d\Omega \qquad \text{(Parseval)}$$

$$= \frac{T}{2\pi} \int_{-\pi/T}^{\pi/T} e^{j\Omega T (n_1 - n_2)} d\Omega$$

$$= \begin{cases} 1, & n_1 = n_2 \\ 0, & n_1 \neq n_2 \end{cases}.$$

That the (closure of the) span of this set is $B_{\pi/T}(\mathbb{R})$ is essentially the content of the Shannon-Nyquist sampling theorem.

Again, sampling $x(t) \in B_{\pi/T}(\mathbb{R})$ is equivalent to a Fourier Series analysis of $X(j\Omega)$ on $[-\pi/T, \pi/T]$.

5. Legendre Polynomials

Define

$$p_0(t) = 1, \quad p_1(t) = t,$$

and then for $n = 1, 2, \dots$

$$p_{n+1}(t) = \frac{2n+1}{n+1} t p_n(t) - \frac{n}{n+1} p_{n-1}(t),$$

and so

$$p_{2}(t) = \frac{1}{2}(3t^{2} - 1)$$

$$p_{3}(t) = \frac{1}{2}(5t^{3} - 3t)$$

$$p_{4}(t) = \frac{1}{8}(35t^{4} - 30x^{2} + 3)$$

$$\vdots \text{ etc.}$$

These $p_n(t)$ are called *Legendre polynomials*, and if we renormalize them, taking

$$v_n(t) = \sqrt{\frac{2n+1}{2}} \, p_n(t),$$

then $v_0(t), \ldots, v_N(t)$ are an orthobasis for polynomials of degree N on [-1, 1].

Computing approximations with the Legendre basis is far more stable than computing the approximation in the standard basis.

Linear approximation and orthobases

Let's return to our linear approximation problem: Given $x \in \mathcal{S}$, we want to find the closest point in a subspace \mathcal{T} .

Suppose we have an orthobasis $\{\boldsymbol{v}_1,\ldots,\boldsymbol{v}_N\}$ for \mathcal{T} . Then solving this problem is easy. Here's why: we know the solution is

$$\hat{\boldsymbol{x}} = a_1 \boldsymbol{v}_1 + a_2 \boldsymbol{v}_2 + \dots + a_N \boldsymbol{v}_N \tag{1}$$

where the a_n are given by

$$\begin{bmatrix} a_1 \\ a_2 \\ \vdots \\ a_N \end{bmatrix} = \boldsymbol{G}^{-1}\boldsymbol{b}, \quad \text{with } \boldsymbol{G} = \begin{bmatrix} \langle \boldsymbol{v}_1, \boldsymbol{v}_1 \rangle & \cdots & \langle \boldsymbol{v}_N, \boldsymbol{v}_1 \rangle \\ \langle \boldsymbol{v}_1, \boldsymbol{v}_2 \rangle & \cdots & \langle \boldsymbol{v}_N, \boldsymbol{v}_2 \rangle \\ \vdots & & & \\ \langle \boldsymbol{v}_1, \boldsymbol{v}_N \rangle & \cdots & \langle \boldsymbol{v}_N, \boldsymbol{v}_N \rangle \end{bmatrix}, \ \boldsymbol{b} = \begin{bmatrix} \langle \boldsymbol{x}, \boldsymbol{v}_1 \rangle \\ \langle \boldsymbol{x}, \boldsymbol{v}_2 \rangle \\ \vdots \\ \langle \boldsymbol{x}, \boldsymbol{v}_N \rangle \end{bmatrix}$$

Now since $\langle \boldsymbol{v}_n, \boldsymbol{v}_k \rangle = 1$ if n = k and 0 otherwise, $\boldsymbol{G} = \mathbf{I}$ (the identity matrix), and so $\boldsymbol{G}^{-1} = \mathbf{I}$ as well, and

$$\begin{bmatrix} a_1 \\ a_2 \\ \vdots \\ a_N \end{bmatrix} = \begin{bmatrix} \langle \boldsymbol{x}, \boldsymbol{v}_1 \rangle \\ \langle \boldsymbol{x}, \boldsymbol{v}_2 \rangle \\ \vdots \\ \langle \boldsymbol{x}, \boldsymbol{v}_N \rangle \end{bmatrix}. \tag{2}$$

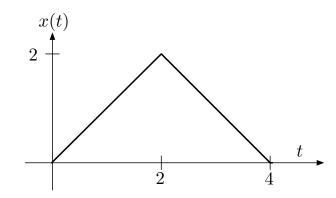
So calculating the closest point is as easy as computing N inner products — no matrix inversion necessary.

Combining the expressions (1) and (2) gives us the compact expression

$$\hat{oldsymbol{x}} = \sum_{n=1}^N \langle oldsymbol{x}, oldsymbol{v}_n
angle oldsymbol{v}_n.$$

Example. Suppose $x(t) \in L_2([0,4])$ is

$$x(t) = \begin{cases} t & 0 \le t \le 2\\ 4 - t & 2 \le t \le 4 \end{cases}$$



Let \mathcal{T} = piecewise constant functions on [0,1), [1,2), [2,3), [3,4].

Find the closest point in \mathcal{T} to \boldsymbol{x} . A good orthobasis to use is

$$v_n(t) = \begin{cases} 1 & (n-1) \le t \le n \\ 0 & \text{otherwise} \end{cases}, \quad n = 1, 2, 3, 4.$$

Orthobasis expansions

The orthogonality principle (easily) gives us an expression for the **expansion coefficients** of a vector in an orthobasis.

Suppose a finite dimensional space S has an orthobasis $\{v_1, \ldots, v_n\}$. Given any $x \in S$, the closest point in S to x is x itself (of course). This gives us the following **reproducing formula**:

$$oldsymbol{x} = \sum_{n=1}^N \langle oldsymbol{x}, oldsymbol{v}_n
angle oldsymbol{v}_n, \quad ext{for all } oldsymbol{x} \in \mathcal{S}.$$

In infinite dimensions, if S has an orthobasis $\{v_n\}_{n=-\infty}^{\infty}$ and $x \in S$ obeys

$$\sum_{n=-\infty}^{\infty} |\langle oldsymbol{x}, oldsymbol{v}_n
angle|^2 \ < \ \infty,$$

then we can write

$$oldsymbol{x} = \sum_{n=-\infty}^{\infty} \langle oldsymbol{x}, oldsymbol{v}_n
angle oldsymbol{v}_n.$$

(We need the sequence of expansion coefficients to be square-summable to make sure the sum of vectors above converges to something.)

In other words, $\boldsymbol{x} \in \mathcal{S}$ is captured without loss by the discrete list of numbers

$$\ldots, \langle \boldsymbol{x}, \boldsymbol{v}_{-1} \rangle, \langle \boldsymbol{x}, \boldsymbol{v}_{0} \rangle, \langle \boldsymbol{x}, \boldsymbol{v}_{1} \rangle, \ldots$$

An orthobasis gives us a natural way to discretize vectors in \mathcal{S} through a set of expansion coefficients. Moreover, there is a straightforward and explicit way to compute these expansion coefficients — you simply take an inner product with the corresponding basis vector.

Example: Sampling a bandlimited function.

 $B_{\pi/T}$ = space of bandlimited signals equipped with the standard inner product. We have seen already that

$$v_n(t) = \sqrt{T} \frac{\sin(\pi(t - nT)/T)}{\pi(t - nT)}, \quad n \in \mathbb{Z}$$

is an orthobasis for $B_{\pi/T}$. This means that any $\boldsymbol{x} \in B_{\pi/T}$ can be written

$$oldsymbol{x} = \sum_{n=-\infty}^{\infty} \langle oldsymbol{x}, oldsymbol{v}_n
angle oldsymbol{v}_n.$$

What are the $\langle \boldsymbol{x}, \boldsymbol{v}_n \rangle$?

$$\langle \boldsymbol{x}, \boldsymbol{v}_n \rangle = \left\langle x(t) , \sqrt{T} \frac{\sin(\pi(t - nT)/T)}{\pi(t - nT)} \right\rangle$$
$$= \frac{1}{2\pi} \int_{-\pi/T}^{\pi/T} X(j\Omega) \sqrt{T} e^{jn\Omega T} d\Omega$$
$$= \sqrt{T} x(nT),$$

which is simply a sample scaled by \sqrt{T} . So the reproducing formula is just a restatement of the sampling theorem:

$$x(t) = \sum_{n = -\infty}^{\infty} \langle \boldsymbol{x}, \boldsymbol{v}_n \rangle \boldsymbol{v}_n$$

$$= \sum_{n = -\infty}^{\infty} \sqrt{T} x(nT) \frac{\sqrt{T} \sin(\pi(t - nT)/T)}{\pi(t - nT)}$$

$$= \sum_{n = -\infty}^{\infty} x(nT) g_T(t - nT).$$

The moral of the story is that we can recreate a vector in a Hilbert space from the sequence of numbers $\{\langle \boldsymbol{x}, \boldsymbol{v}_n \rangle\}$. We can think of every different orthobasis for \mathcal{S} as a different **transform**, and the $\{\langle \boldsymbol{x}, \boldsymbol{v}_n \rangle\}$ as **transform coefficients**.

Next we will see that our notions of **distance** and **angle** also carry over to this discrete space.

Parseval's Theorem

One handy fact (and a fact we have used many times in this course already) about the Fourier transform is that it is **energy preserving**,

$$||x(t)||_2^2 = \int_{-\infty}^{\infty} |x(t)|^2 dt = \frac{1}{2\pi} \int_{-\infty}^{\infty} |X(j\Omega)|^2 d\Omega = \frac{1}{2\pi} ||X(j\Omega)||_2^2,$$

and more generally, it preserves the L_2 inner product:

$$\langle x(t), y(t) \rangle = \int_{-\infty}^{\infty} x(t) \overline{y(t)} dt = \frac{1}{2\pi} \int_{-\infty}^{\infty} X(j\Omega) \overline{Y(j\Omega)} d\Omega$$
$$= \frac{1}{2\pi} \langle X(j\Omega), Y(j\Omega) \rangle.$$

It is not not too hard to show that something very similar is true for any orthobasis expansion. Let S be a Hilbert space with inner product $\langle \cdot, \cdot \rangle_S$ which induces norm $\|\cdot\|_S$. Let $\{v_k\}_{k\in\Gamma}$ be an orthobasis for S. Then for every $x, y \in S$,

$$\langle oldsymbol{x}, oldsymbol{y}
angle_S = \sum_{k \in \Gamma} lpha_k \overline{eta_k},$$

where

$$\alpha_k = \langle \boldsymbol{x}, \boldsymbol{v}_k \rangle_S, \qquad \beta_k = \langle \boldsymbol{y}, \boldsymbol{v}_k \rangle_S.$$

You can think of the $\{\alpha_k\}$ as the transform coefficients of \boldsymbol{x} and the $\{\beta_k\}$ as the transform coefficients of \boldsymbol{y} . So we have

$$egin{aligned} \langle oldsymbol{x}, oldsymbol{y}
angle_S &= \langle oldsymbol{lpha}, oldsymbol{eta}
angle_{\ell_2}, \ \|oldsymbol{x}\|_S^2 &= \|oldsymbol{lpha}\|_2^2. \end{aligned}$$

¹We are using Γ to be an arbitrary index set here; it can be either finite, e.g. $\Gamma = 1, 2, ..., N$, or infinite, e.g. $\Gamma = \mathbb{Z}$.

 \Rightarrow An orthobasis makes every Hilbert space **equivalent** to ℓ_2 .

All of the geometry (lengths, angles) maps into standard Euclidean geometry in coefficient space. As you can imagine, this is a pretty useful fact.

Proof of Parseval. With $\alpha_k = \langle \boldsymbol{x}, \boldsymbol{v}_k \rangle$ and $\beta_k = \langle \boldsymbol{y}, \boldsymbol{v}_k \rangle$, we can write

$$oldsymbol{x} = \sum_{k \in \Gamma} lpha_k \, oldsymbol{v}_k, \quad ext{and} \quad oldsymbol{y} = \sum_{k \in \Gamma} eta_k \, oldsymbol{v}_k,$$

and so

$$\langle \boldsymbol{x}, \boldsymbol{y} \rangle_{S} = \left\langle \sum_{k \in \Gamma} \alpha_{k} \boldsymbol{v}_{k}, \sum_{\ell \in \Gamma} \beta_{\ell} \boldsymbol{v}_{\ell} \right\rangle_{S}$$

$$= \sum_{k \in \Gamma} \alpha_{k} \left\langle \boldsymbol{v}_{k}, \sum_{\ell \in \Gamma} \beta_{\ell} \boldsymbol{v}_{\ell} \right\rangle_{S}$$

$$= \sum_{k \in \Gamma} \sum_{\ell \in \Gamma} \alpha_{k} \overline{\beta_{\ell}} \langle \boldsymbol{v}_{k}, \boldsymbol{v}_{\ell} \rangle_{S}.$$

For a fixed value of k, only one term in the inner sum above will be nonzero, as $\langle \boldsymbol{v}_k, \boldsymbol{v}_\ell \rangle = 0$ unless $\ell = k$. Thus

$$\langle oldsymbol{x}, oldsymbol{y}
angle_S = \sum_{k \in \Gamma} lpha_k \overline{eta_k}.$$

A straightforward consequence of the result above is that distances in S under the induced norm are equivalent to Euclidean (ℓ_2) distances in coefficient space:

$$\|oldsymbol{x} - oldsymbol{y}\|_S = \|oldsymbol{lpha} - oldsymbol{eta}\|_2 = \left(\sum_{k \in \Gamma} (lpha_k - eta_k)^2
ight)^{1/2}.$$

Thus changing the value of an orthobasis expansion coefficient by an amount ϵ will change the signal by an amount (as measured in $\|\cdot\|_S$) ϵ .

To be more precise about this, suppose \boldsymbol{x} has transform coefficients $\{\alpha_k = \langle \boldsymbol{x}, \boldsymbol{v}_k \rangle_S\}$. If I perturb one of them, say at location k_0 , by setting

$$\tilde{\alpha}_k = \begin{cases} \alpha_{k_0} + \epsilon & k = k_0 \\ \alpha_k & k \neq k_0 \end{cases},$$

and then synthesizing

$$\tilde{m{x}} = \sum_{k \in \Gamma} \tilde{lpha}_k m{v}_k,$$

we will have

$$\|\boldsymbol{x} - \tilde{\boldsymbol{x}}\|_S = \epsilon.$$

Notice that while the error is localized to one expansion coefficient, it could effect the entire reconstruction, but its net effect will still be ϵ .

Here is another example. Suppose I sample a signal $x_c(t)$ which is bandlimited to π/T at a rate T, producing the sample sequence $x[n] = x_c(nT)$. Each of these samples gets perturbed by a (possibly different) amount $\epsilon[n]$:

$$\tilde{x}[n] = x[n] + \epsilon[n].$$

We resynthesize the signal using sinc interpolation:

$$\tilde{x}_c(t) = \sum_{n=-\infty}^{\infty} \tilde{x}[n]h_T(t-nT),$$

and the difference between this signal and the "true" signal is

$$x_c(t) - \tilde{x}_c(t) = \sum_{n = -\infty}^{\infty} (x[n] - \tilde{x}[n]) h_T(t - nT)$$
$$= \sum_{n = -\infty}^{\infty} \sqrt{T} (x[n] - \tilde{x}[n]) h_T(t - nT) / \sqrt{T}.$$

Since the $\{h_T(t-nT)/\sqrt{T}\}_{n\in\mathbb{Z}}$ are an orthobasis for $B_{\pi/T}$, we know

$$||x_{c}(t) - \tilde{x}_{c}(t)||_{L_{2}}^{2} = \int |x_{c}(t) - \tilde{x}_{c}(t)|^{2} dt$$

$$= \sum_{n = -\infty}^{\infty} |\sqrt{T}(x[n] - \tilde{x}[n])|^{2}$$

$$= T \sum_{n = -\infty}^{\infty} |\epsilon[n]|^{2}$$

The upshot of this is that as we change each sample, we know exactly what the net effect will be on the reconstruction error.