

NPTEL MOOC

PROGRAMMING, DATA STRUCTURES AND ALGORITHMS IN PYTHON

Week 8, Lecture 2

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Grid Paths

(5,10)

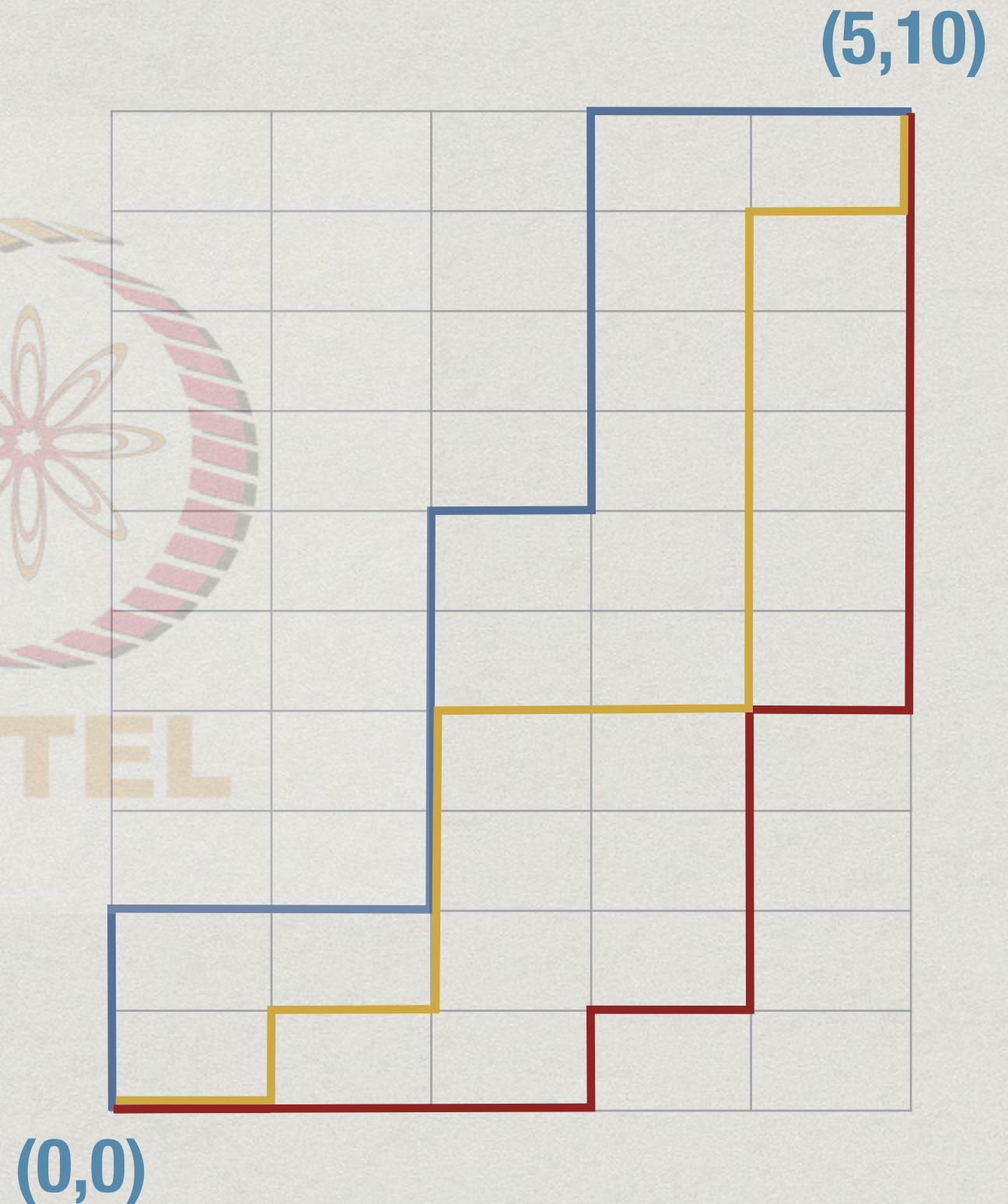
- * Roads arranged in a rectangular grid
- * Can only go up or right
- * How many different routes from (0,0) to (m,n)?



(0,0)

Grid Paths

- * Roads arranged in a rectangular grid
- * Can only go up or right
- * How many different routes from $(0,0)$ to (m,n) ?

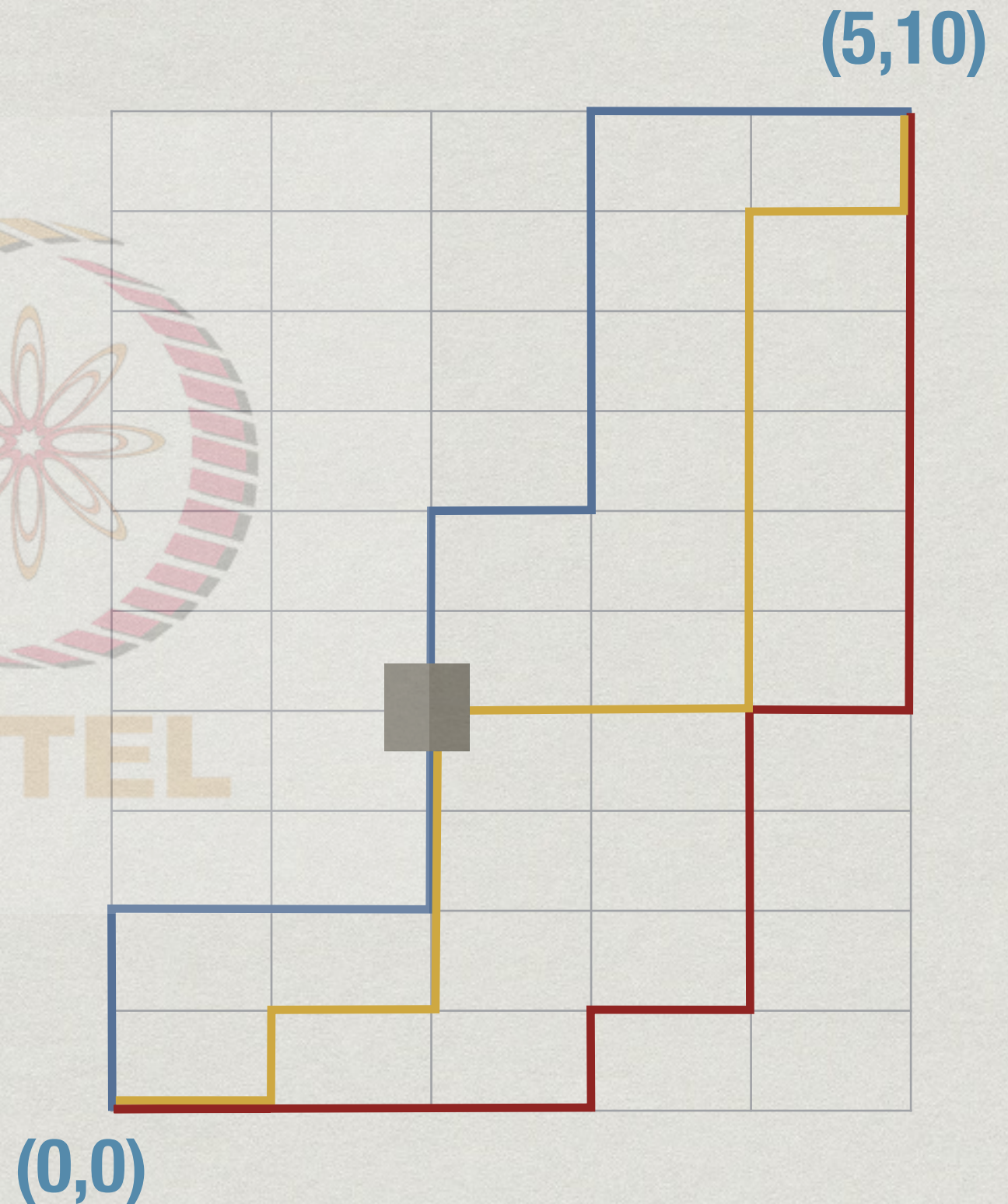


Combinatorial solution

- * Every path from (0,0) to (5,10) has 15 segments
 - * In general $m+n$ segments from (0,0) to (m,n)
- * Of these exactly 5 are right moves, 10 are up moves
- * Fix the positions of the 5 right moves among the overall 15 positions
 - * $15 \text{ choose } 5 = (15!)/(10!)(5!) = 3003$
 - * Same as 15 choose 10: fix the 10 up moves

Holes

- * What if an intersection is blocked?
- * (2,4), for example
- * Paths through (2,4) need to be discarded
- * Two of our earlier examples are invalid paths

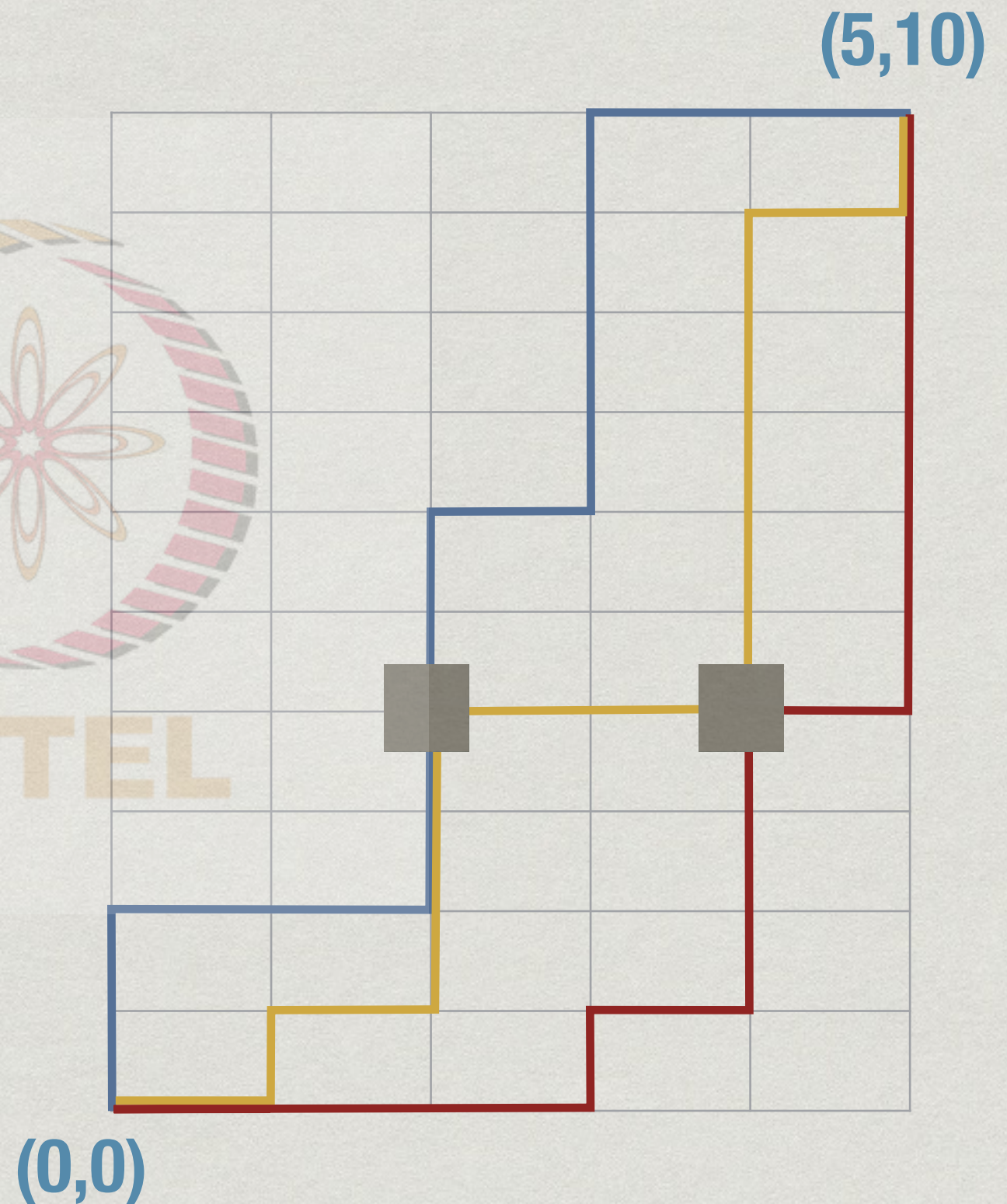


Combinatorial solution

- * Every path through (2,4) goes from (0,0) to (2,4) and then from (2,4) to (5,10)
- * Count these separately:
 - * $(4+2) \text{ choose } 2 = 15$
 - * $(6+3) \text{ choose } 3 = 84$
- * Multiply to get all paths through (2,4): 1260
- * Subtract from 15 choose 5 = 3003 to get valid paths that avoid (2,4): 1743

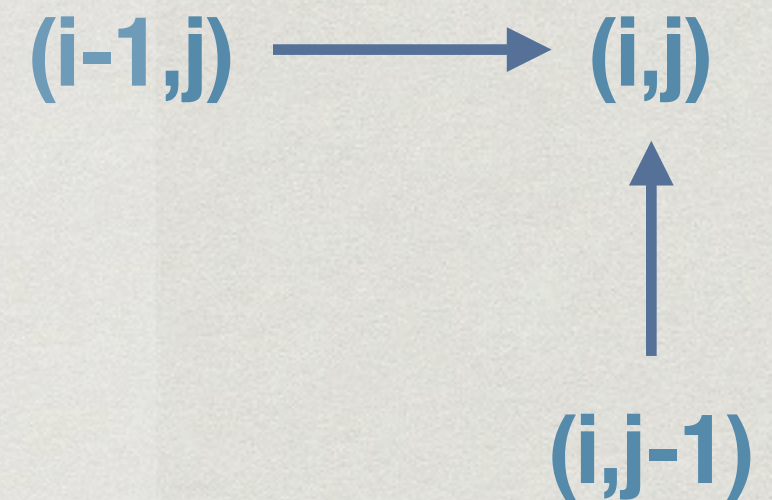
Holes

- * What if two intersections are blocked?
- * Subtract paths through $(2,4)$, $(4,4)$
 - * Some paths are counted twice!
- * Add back paths through both holes
- * Inclusion-exclusion: messy



Inductive formulation

- * How can a path reach (i,j)
 - * Move up from $(i,j-1)$
 - * Move right from $(i-1,j)$
- * Every path to these neighbours extends in a unique way to (i,j)



Inductive formulation

- * $\text{Paths}(i,j)$: Number of paths from $(0,0)$ to (i,j)
- * $\text{Paths}(i,j) = \text{Paths}(i-1,j) + \text{Paths}(i,j-1)$
- * Boundary cases
 - * $\text{Paths}(i,0) = \text{Paths}(i-1,0)$ # Bottom row
 - * $\text{Paths}(0,j) = \text{Paths}(0,j-1)$ # Left column
 - * $\text{Paths}(0,0) = 1$ # Base case

Dealing with holes

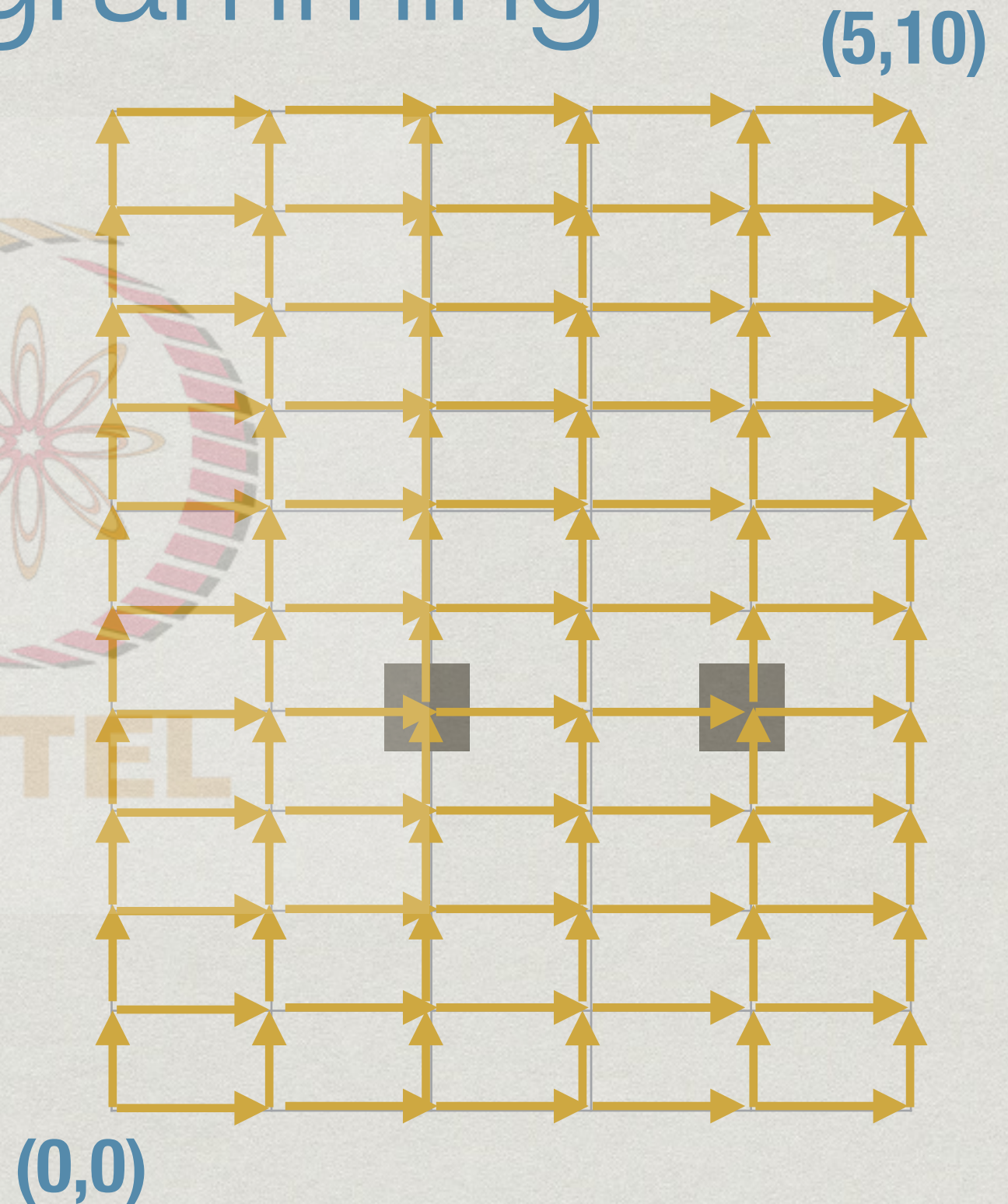
- * $\text{Paths}(i,j) = 0$, if there is a hole at (i,j)
- * $\text{Paths}(i,j) = \text{Paths}(i-1,j) + \text{Paths}(i,j-1)$, otherwise
- * Boundary cases
 - * $\text{Paths}(i,0) = \text{Paths}(i-1,0)$ # Bottom row
 - * $\text{Paths}(0,j) = \text{Paths}(0,j-1)$ # Left column
 - * $\text{Paths}(0,0) = 1$ # Base case

Computing Paths(i,j)

- * Naive recursion will recompute multiple times
 - * Paths(5,10) requires Paths(4,10) and Paths(5,9)
 - * Both Paths(4,10) and Paths(5,9) require Paths(4,9)
- * Use memoization ...
- * ... or compute the subproblems directly in a suitable way

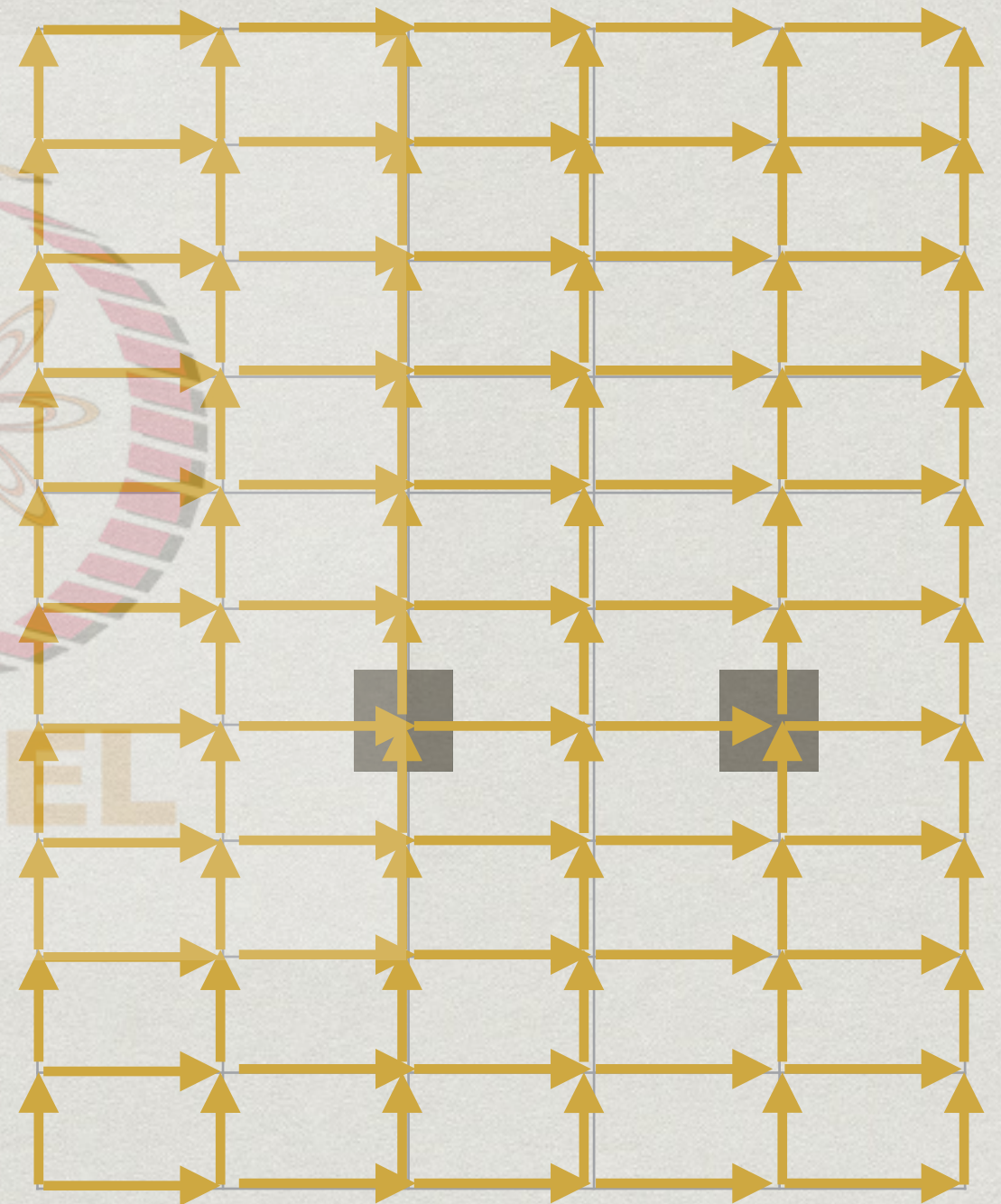
Dynamic programming

- * Identify dependency structure
- * Paths(0,0) has no dependencies
- * Start at (0,0)



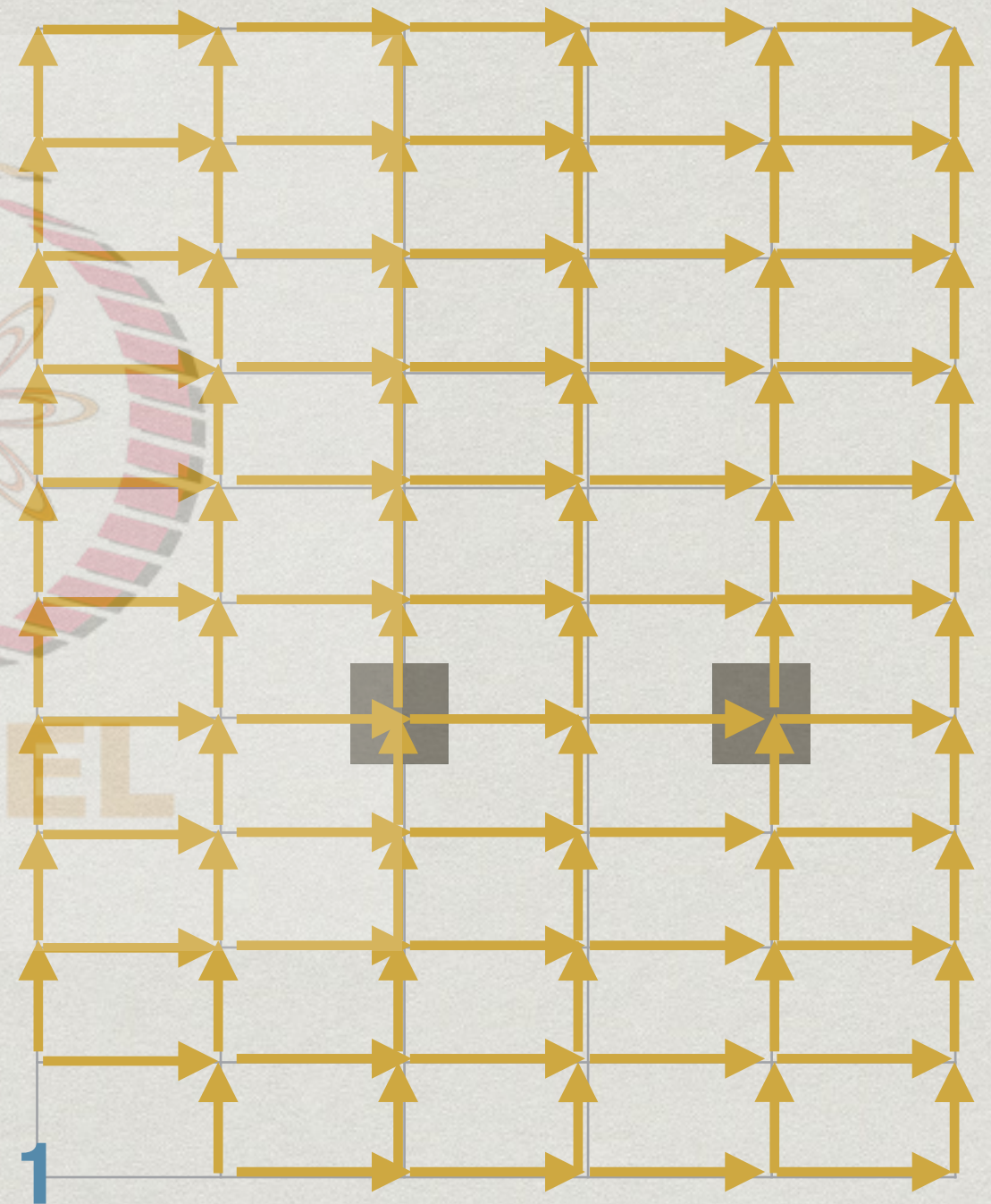
Dynamic programming

- * Start at (0,0)
- * Fill row by row



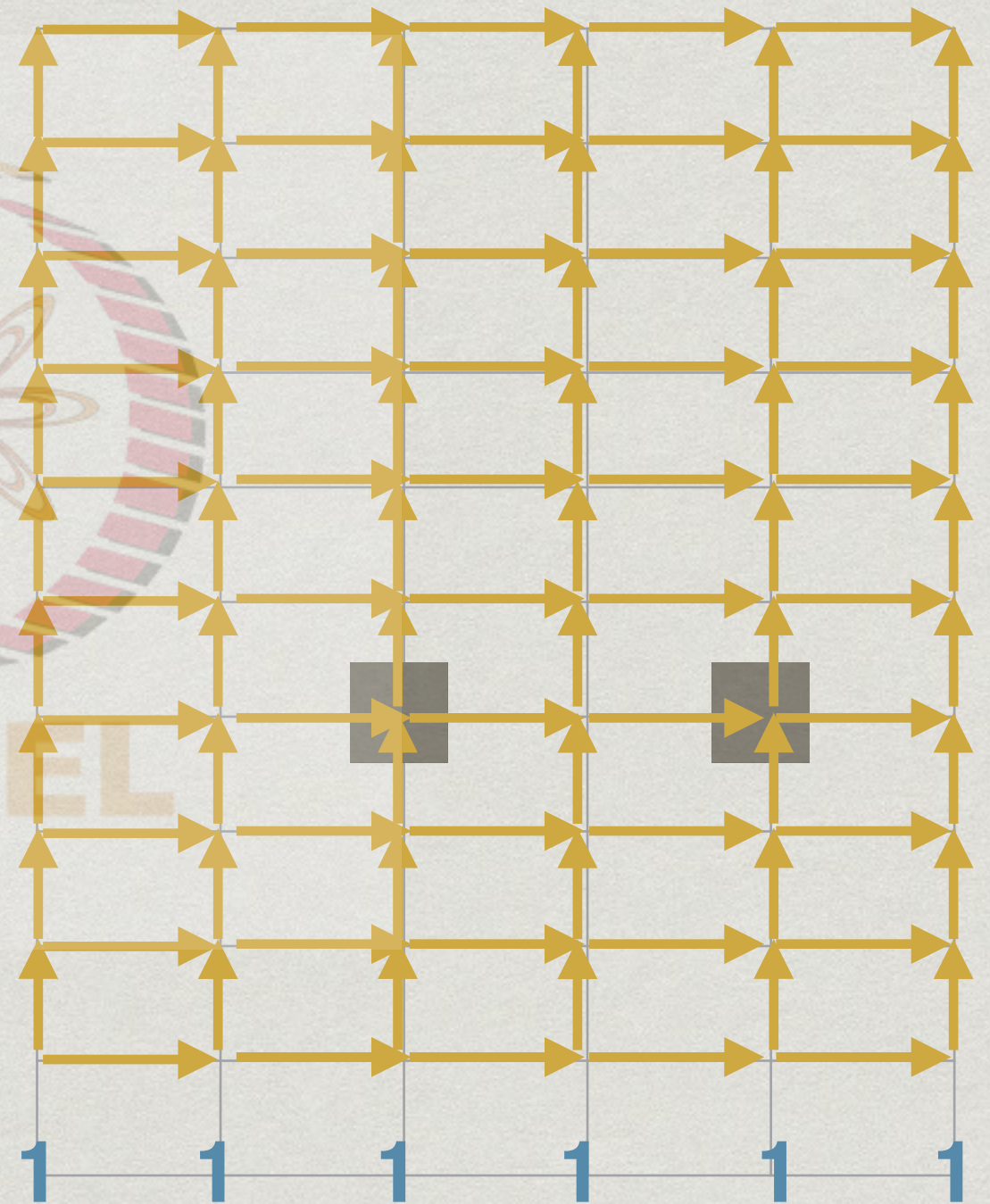
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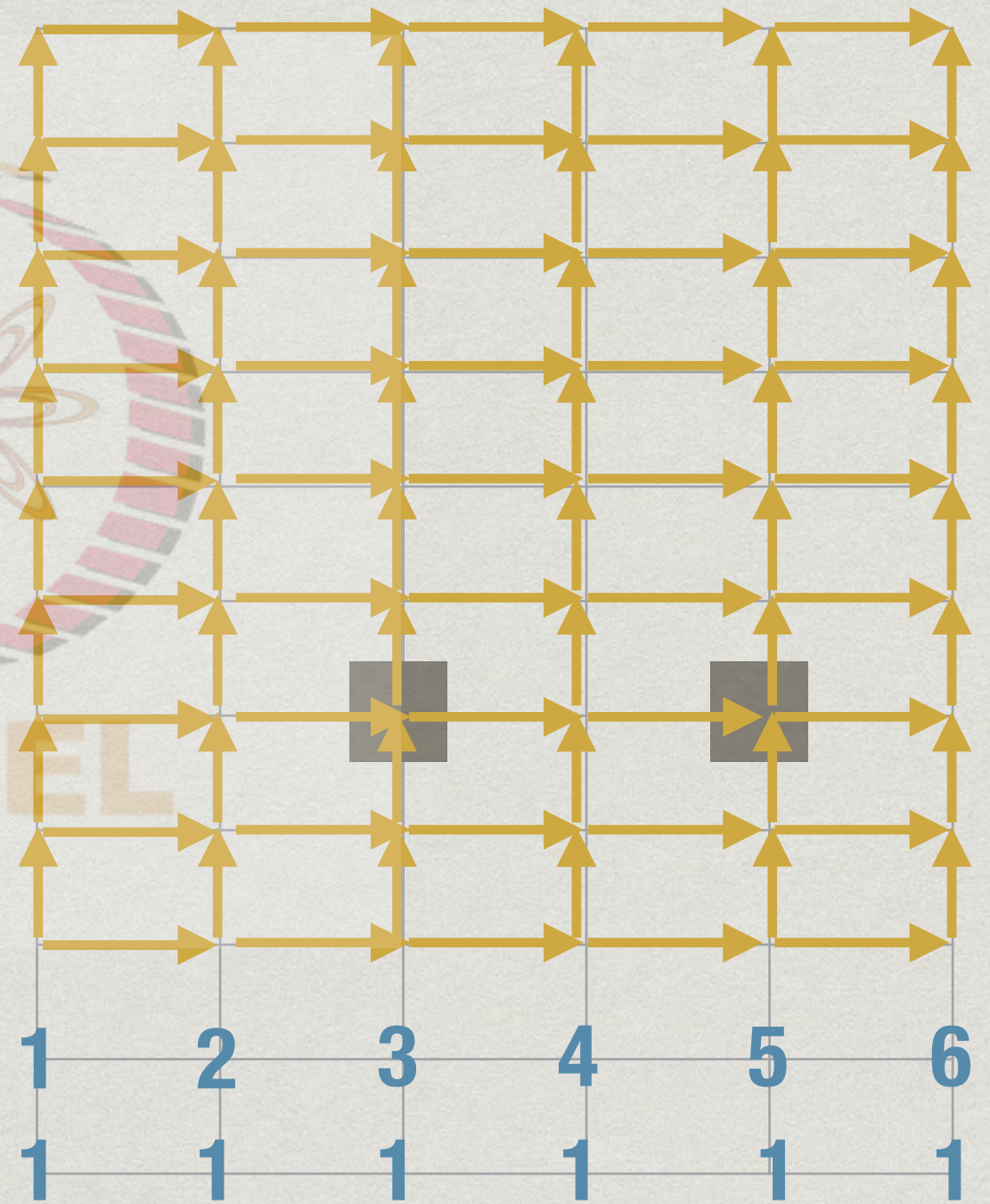
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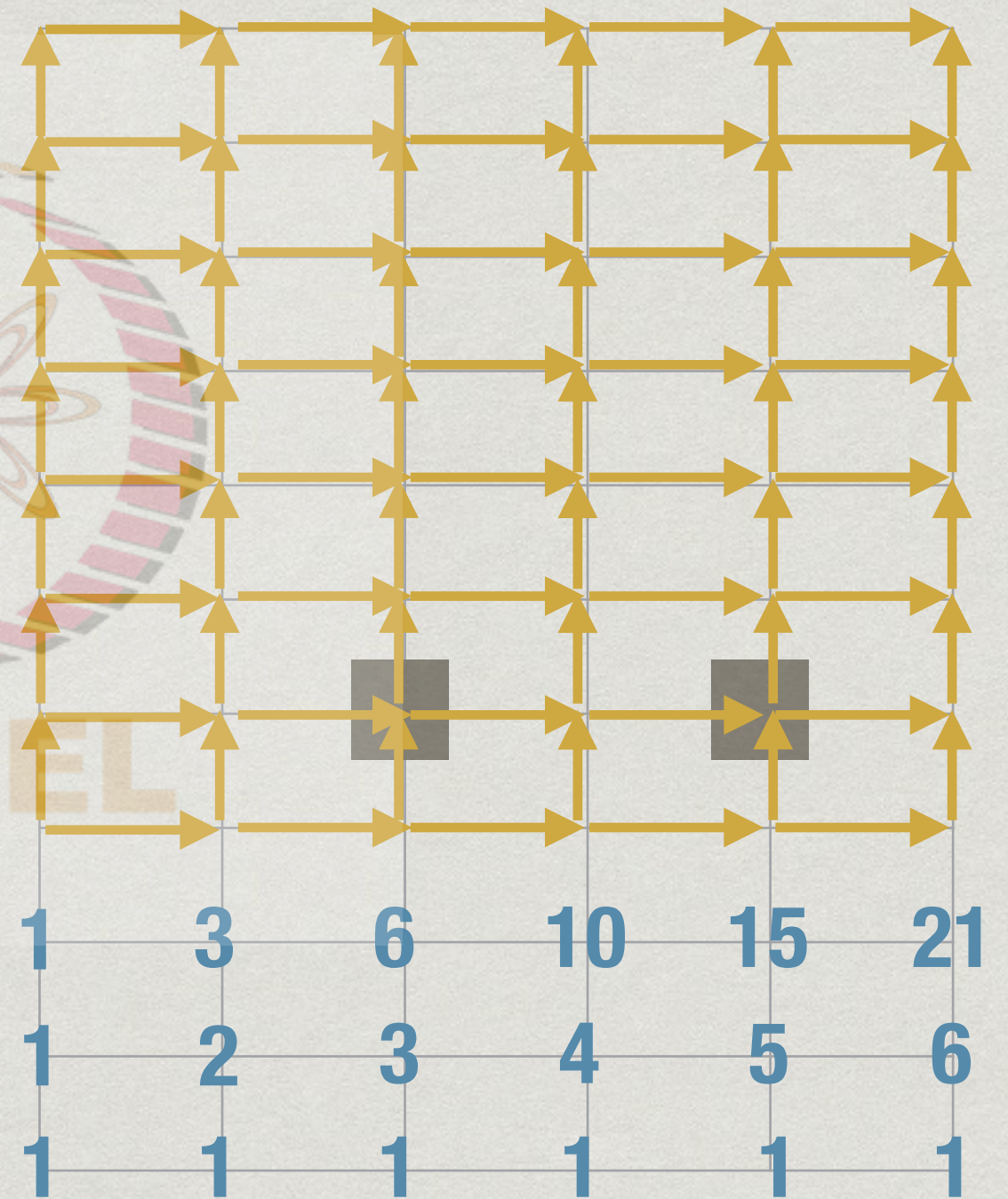
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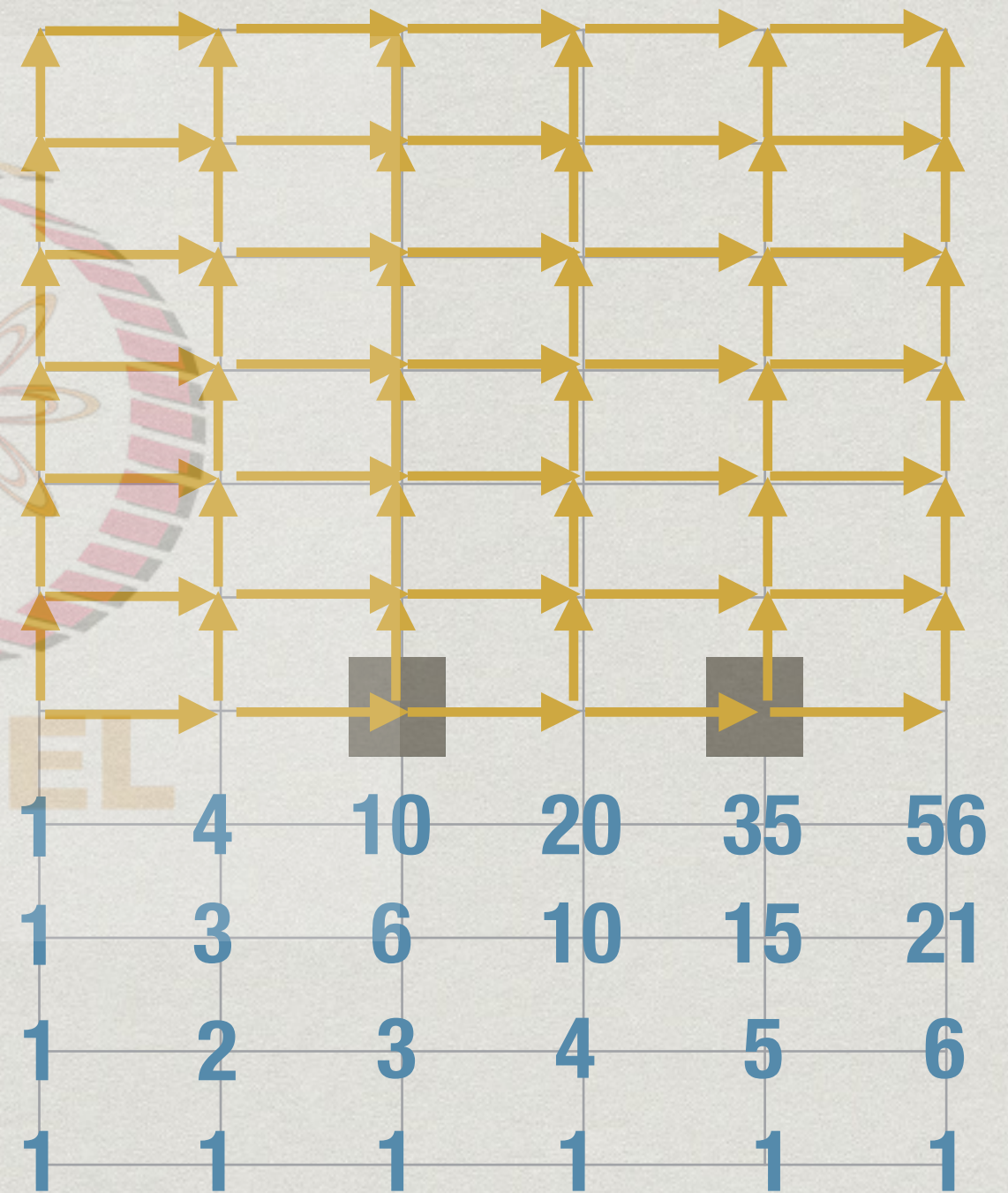
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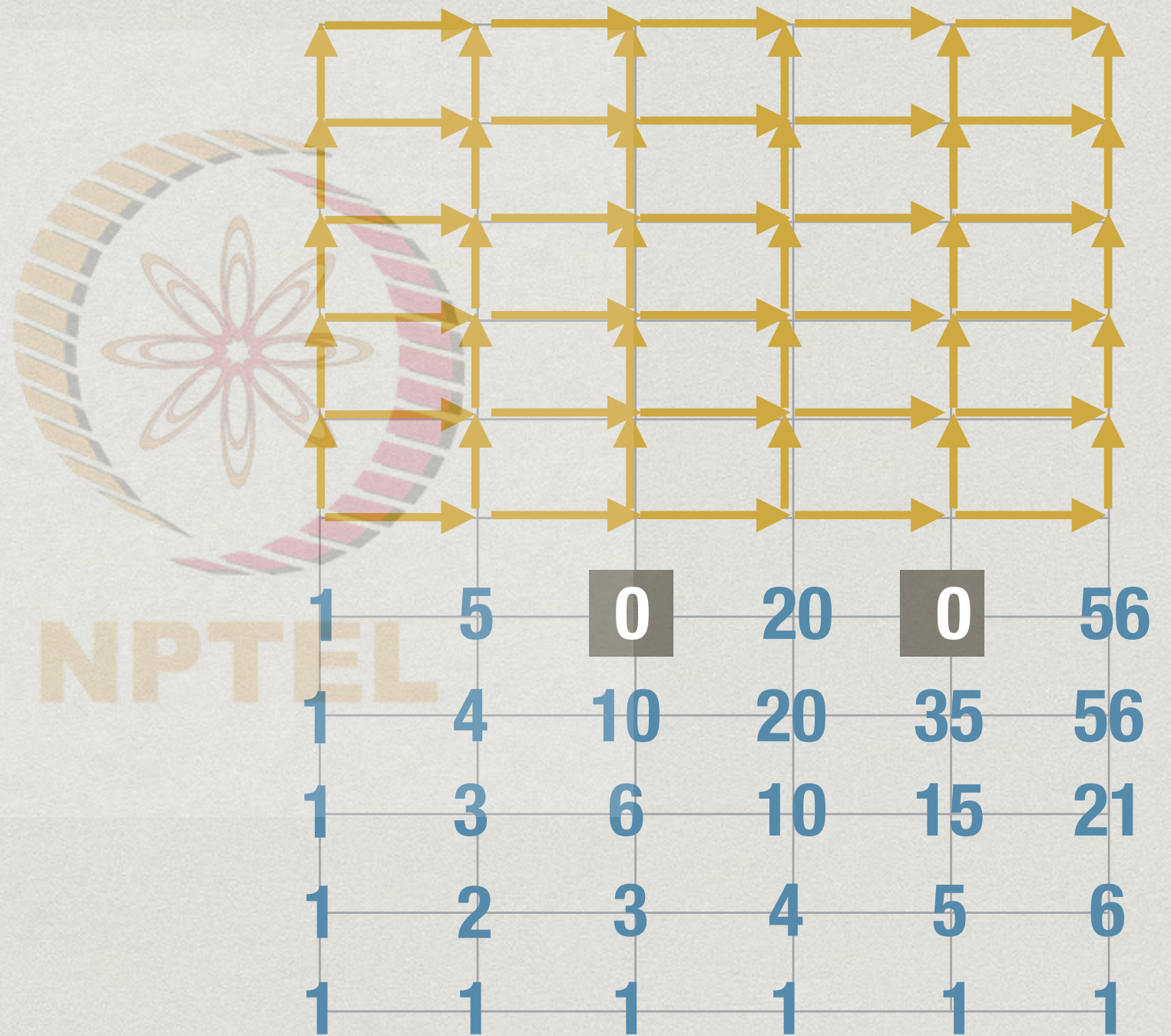
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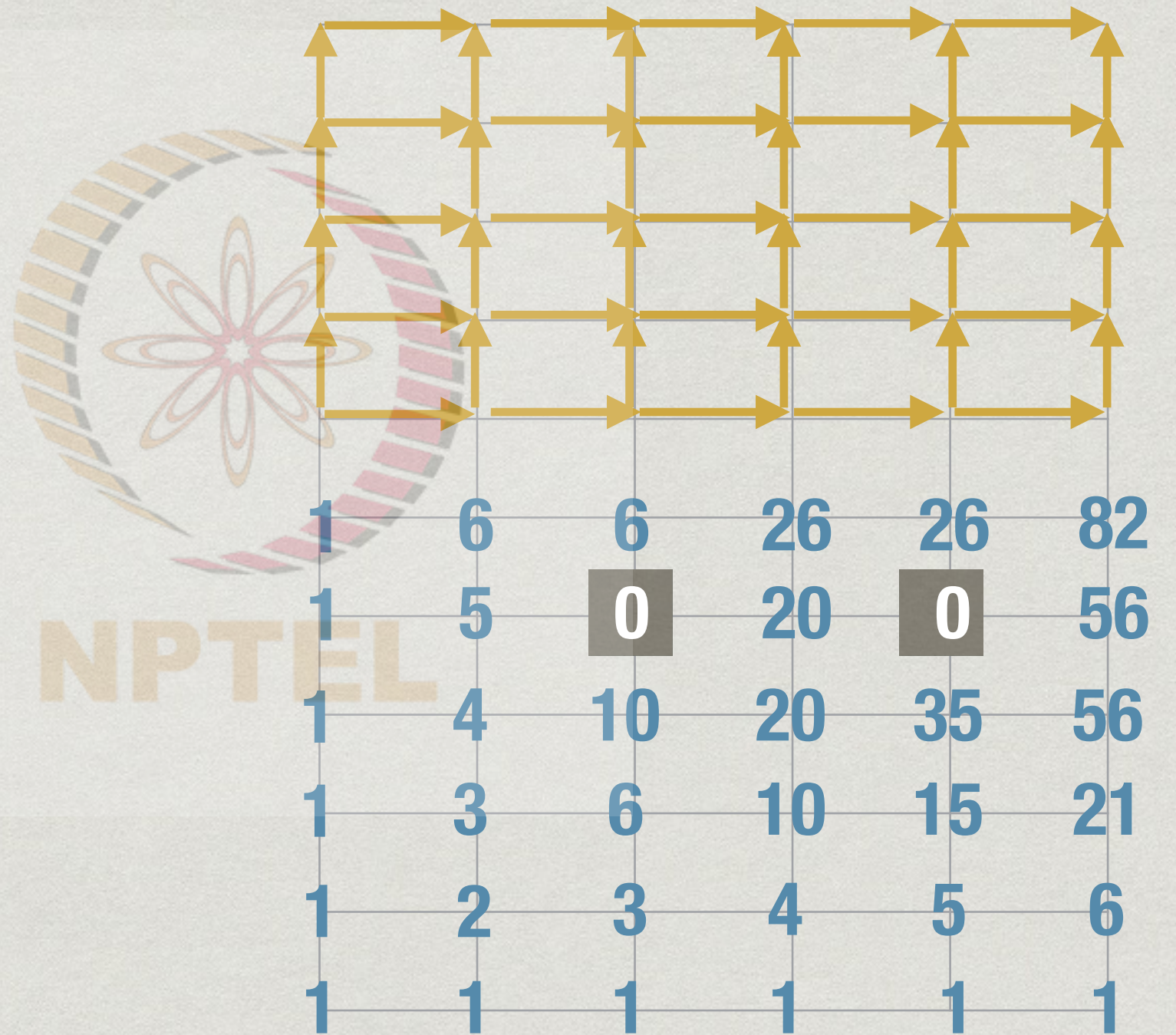
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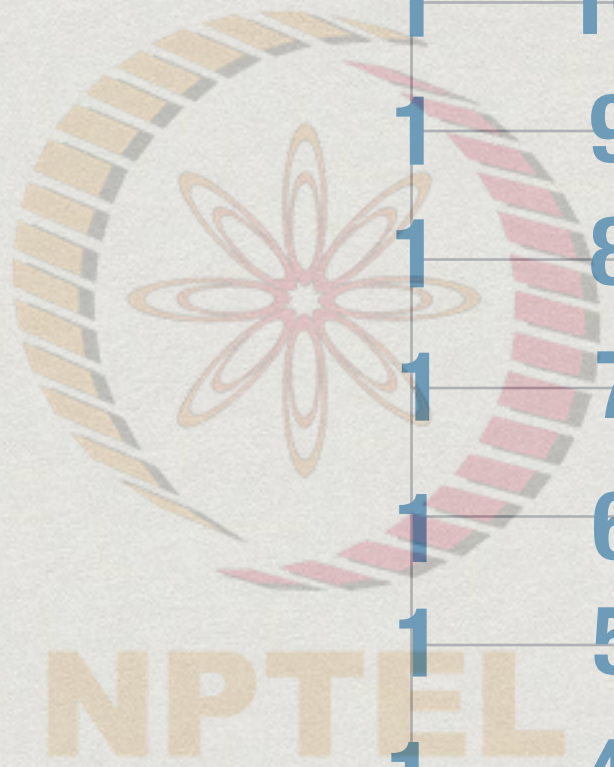
Dynamic programming

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Dynamic programming

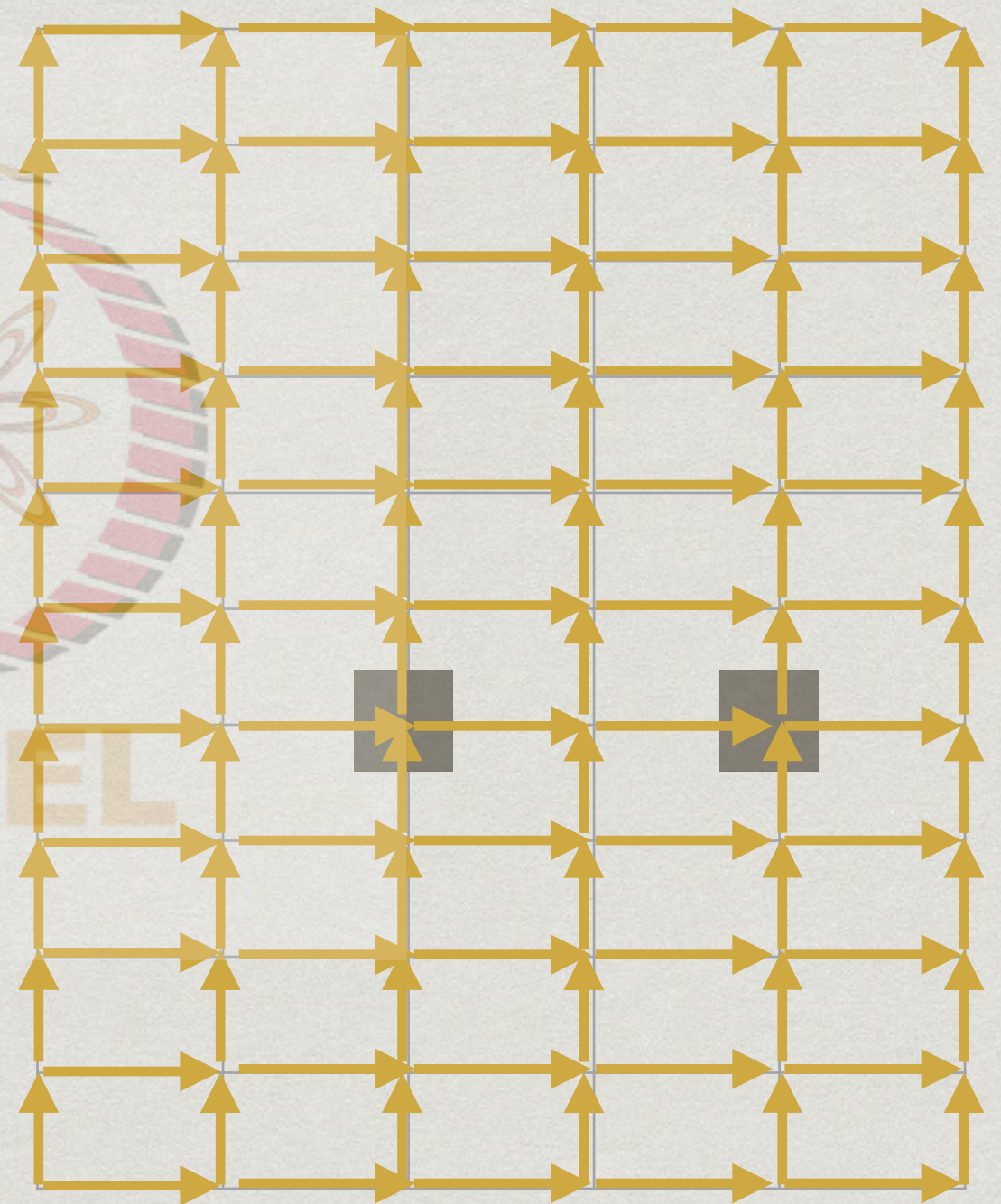
- * Start at (0,0)
- * Fill row by row



1	11	51	181	526	1363
1	10	40	130	345	837
1	9	30	90	215	492
1	8	21	60	125	272
1	7	13	39	65	147
1	6	6	26	26	82
1	5	0	20	0	56
1	4	10	20	35	56
1	3	6	10	15	21
1	2	3	4	5	6
1	1	1	1	1	1

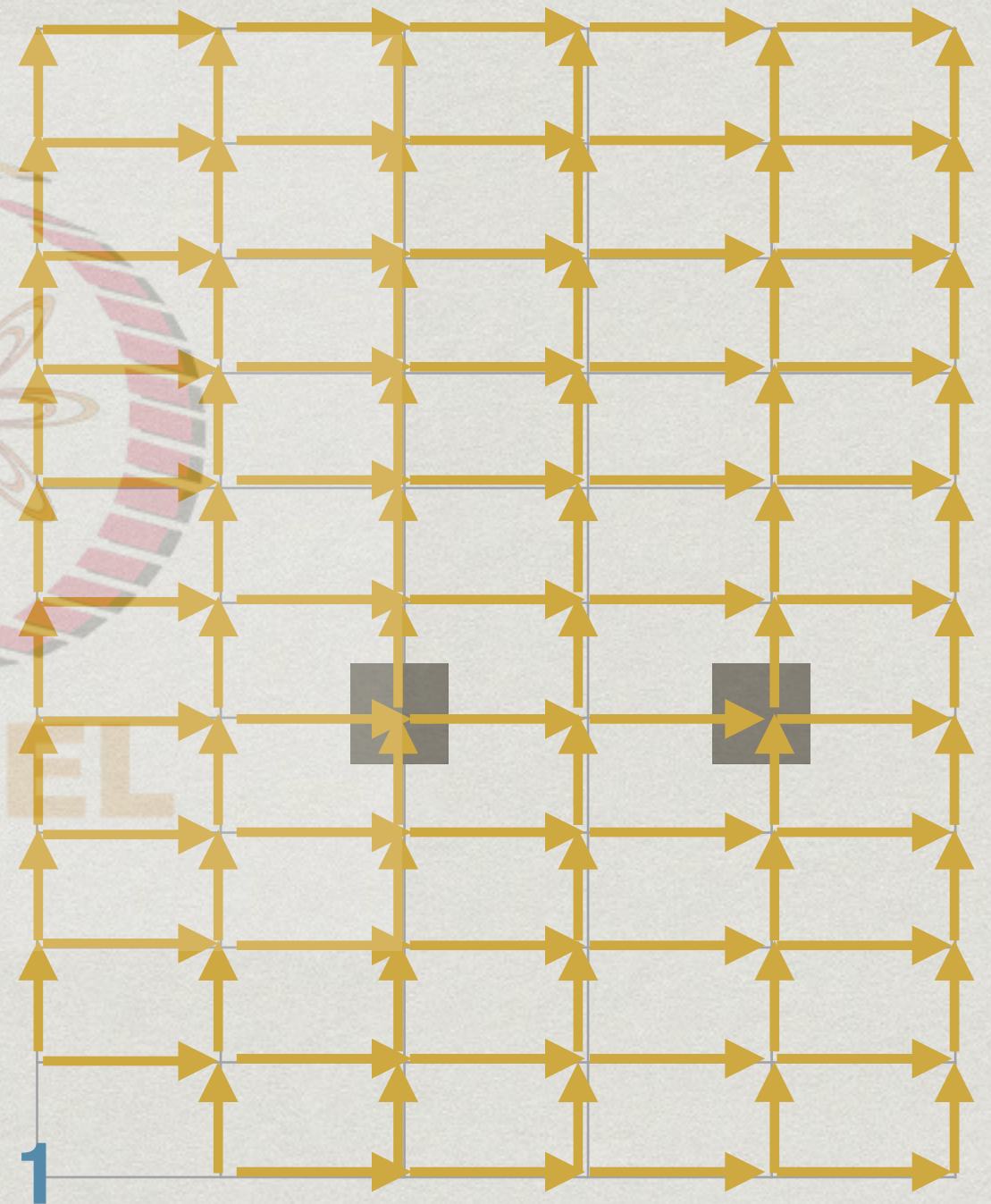
Dynamic programming

- * Start at (0,0)
- * Fill by column



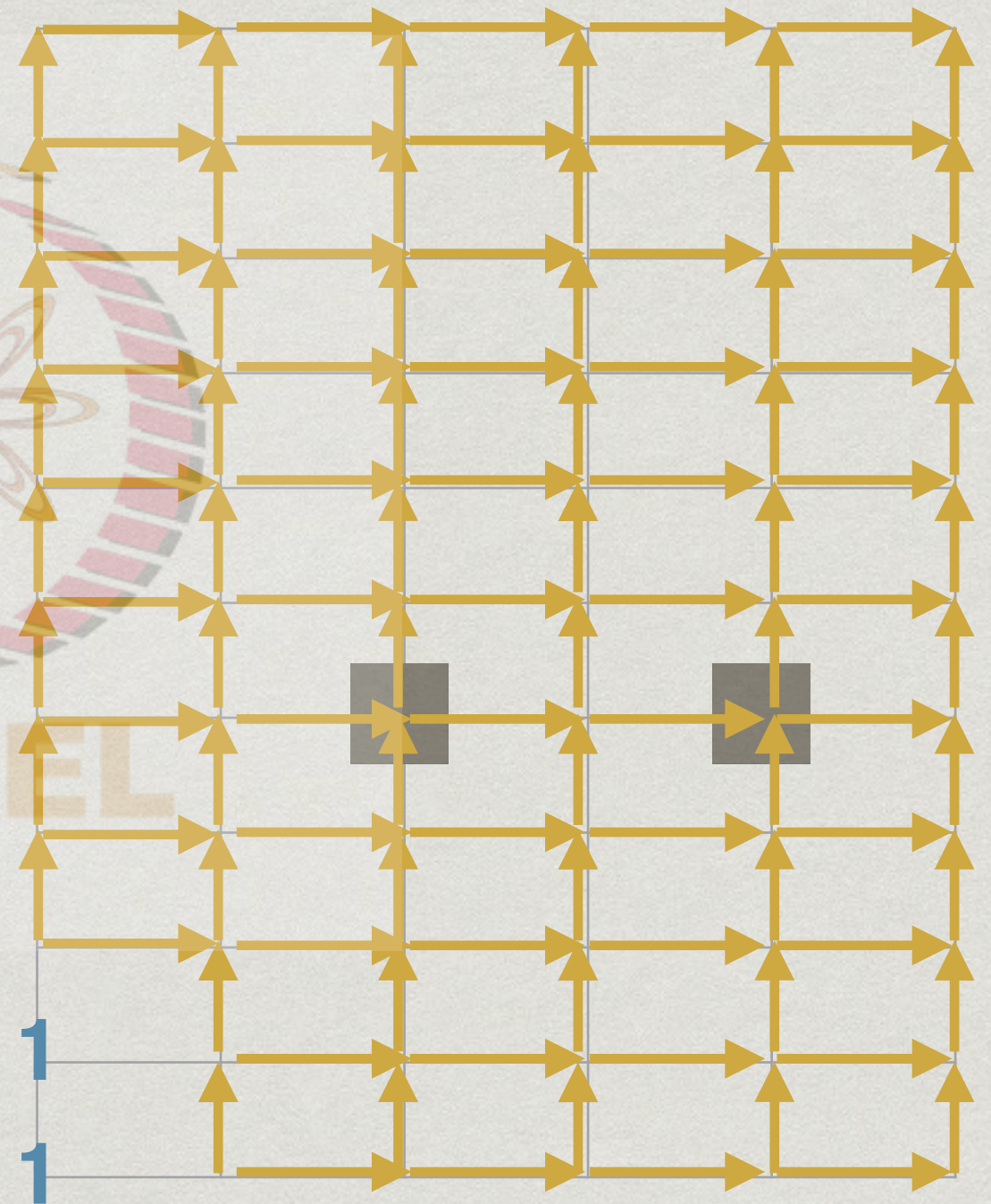
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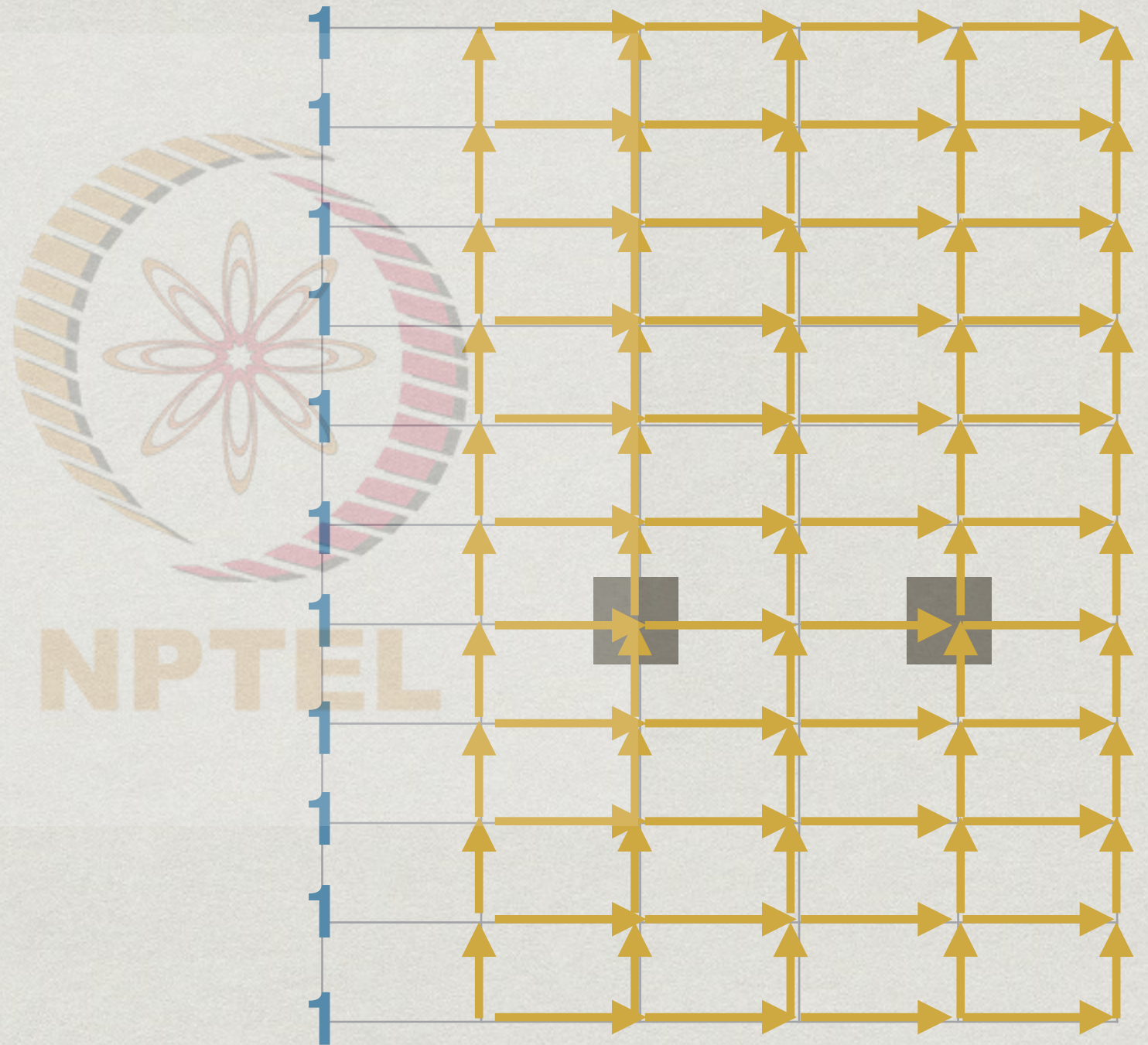
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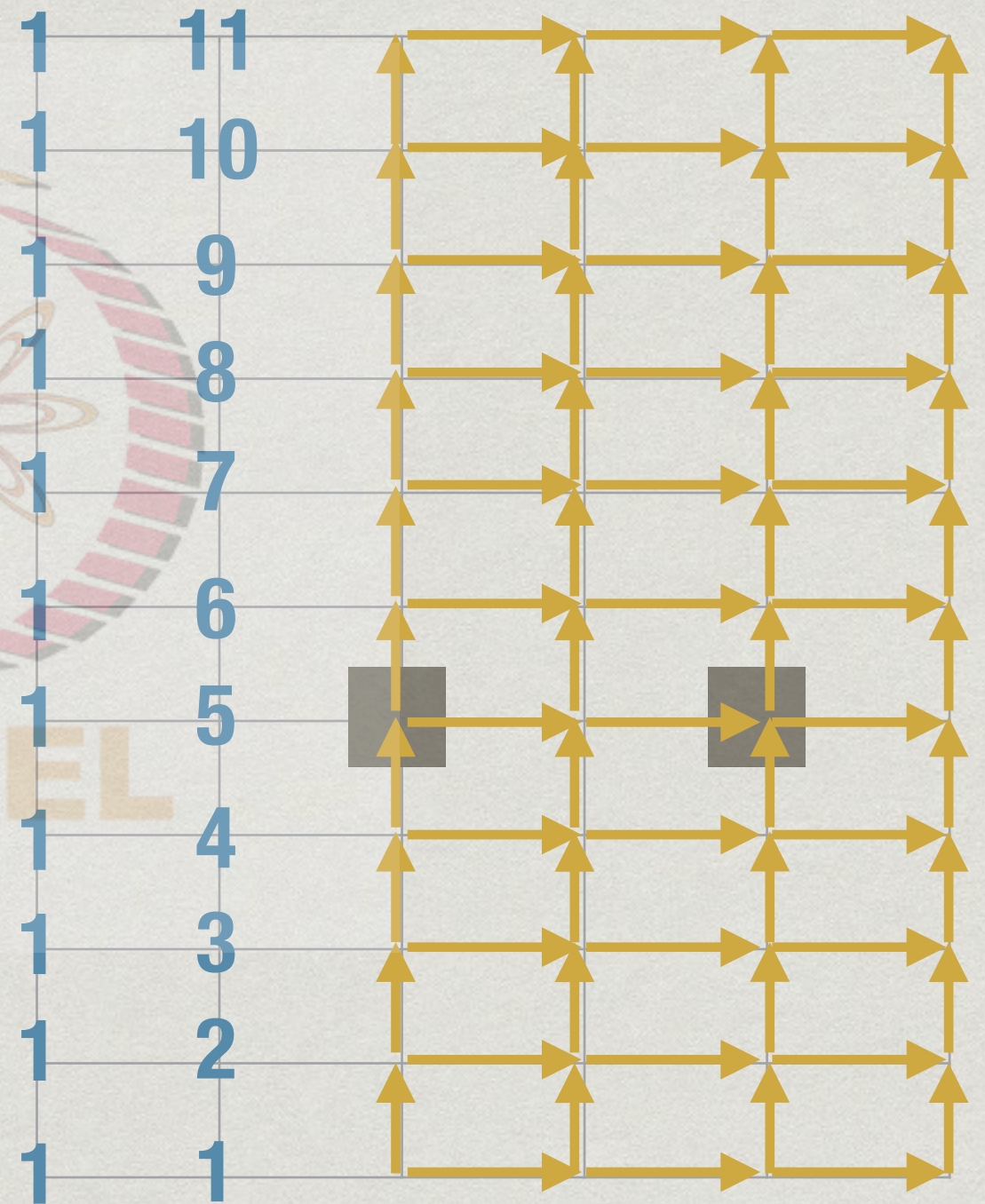
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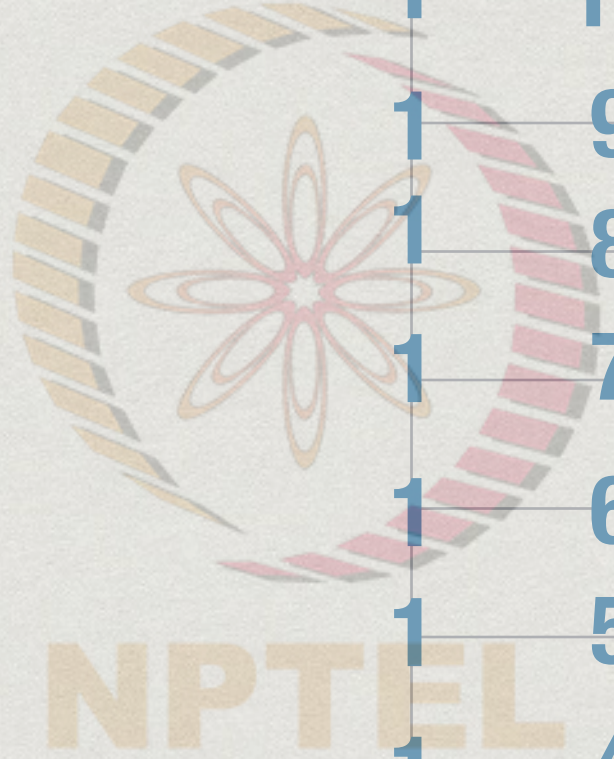
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Dynamic programming

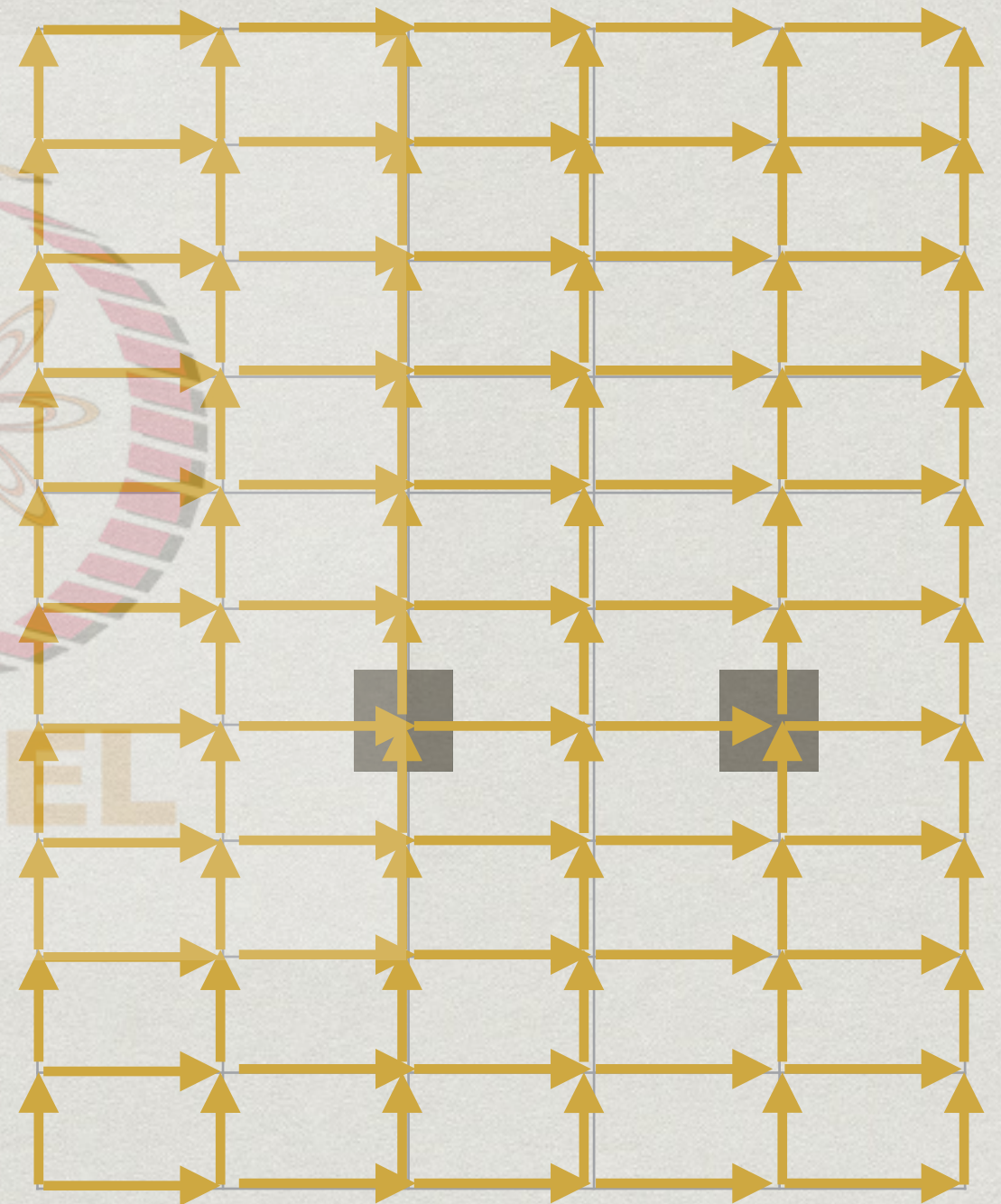
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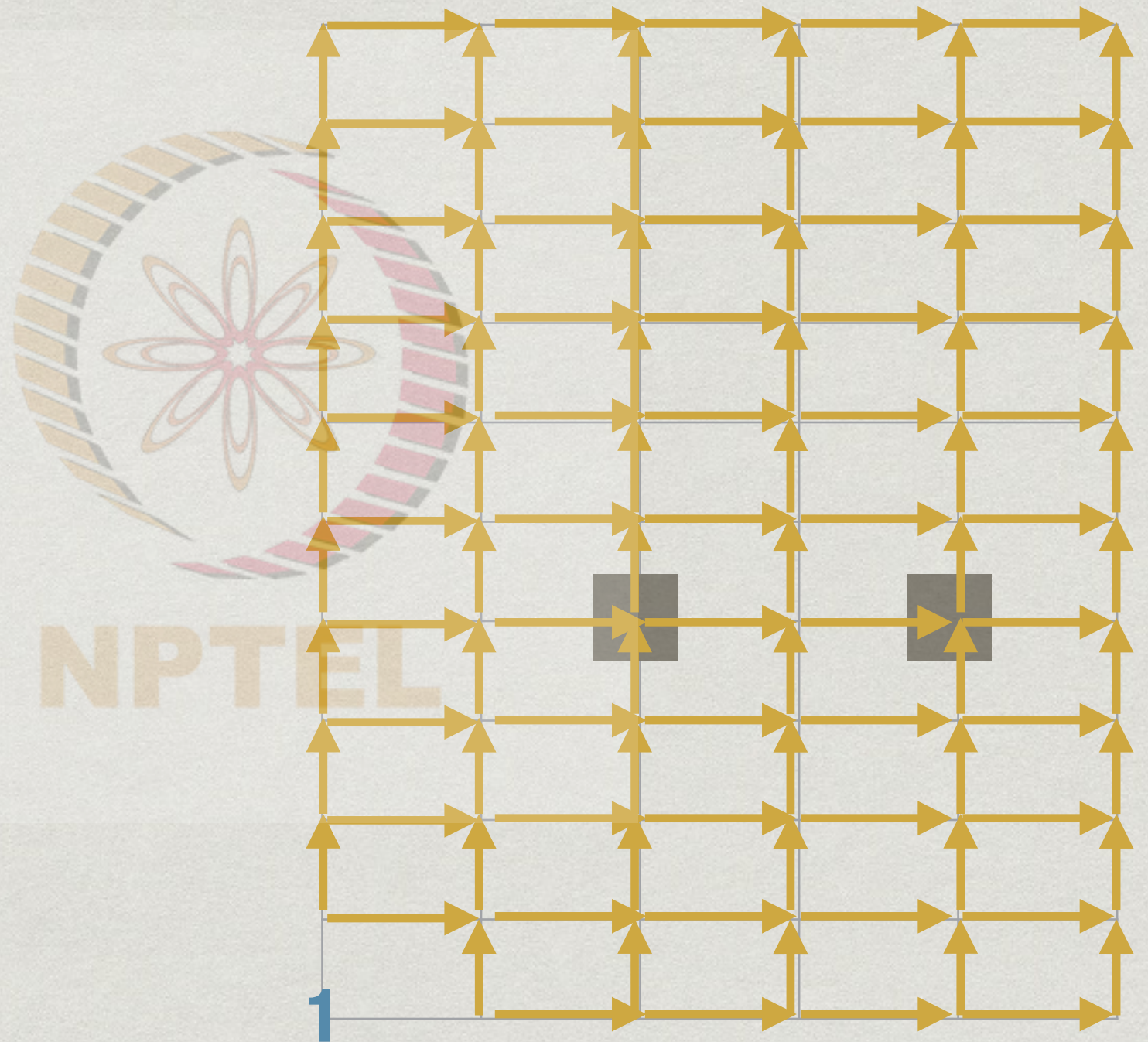
Dynamic programming

- * Start at (0,0)
- * Fill by diagonal



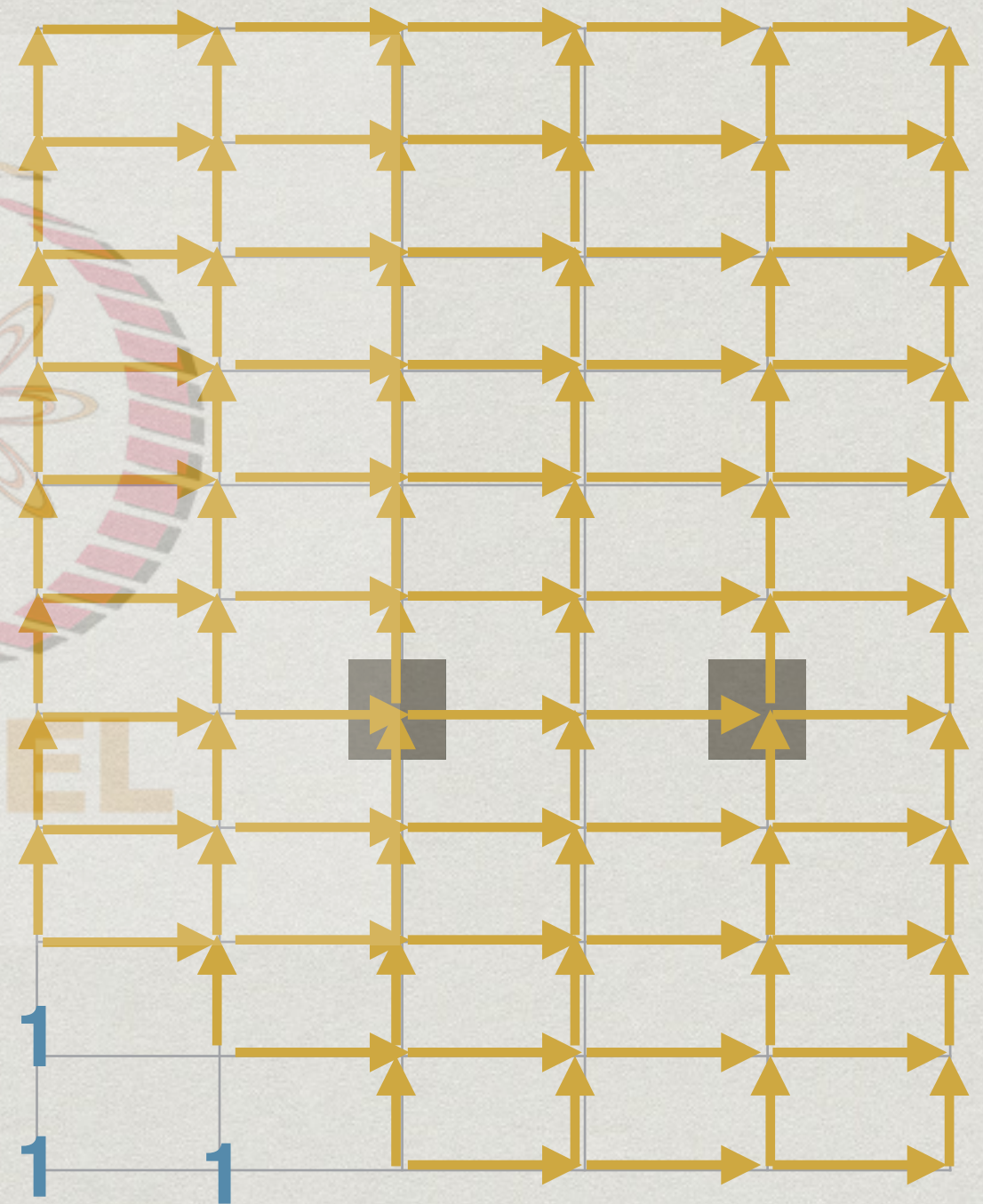
Dynamic programming

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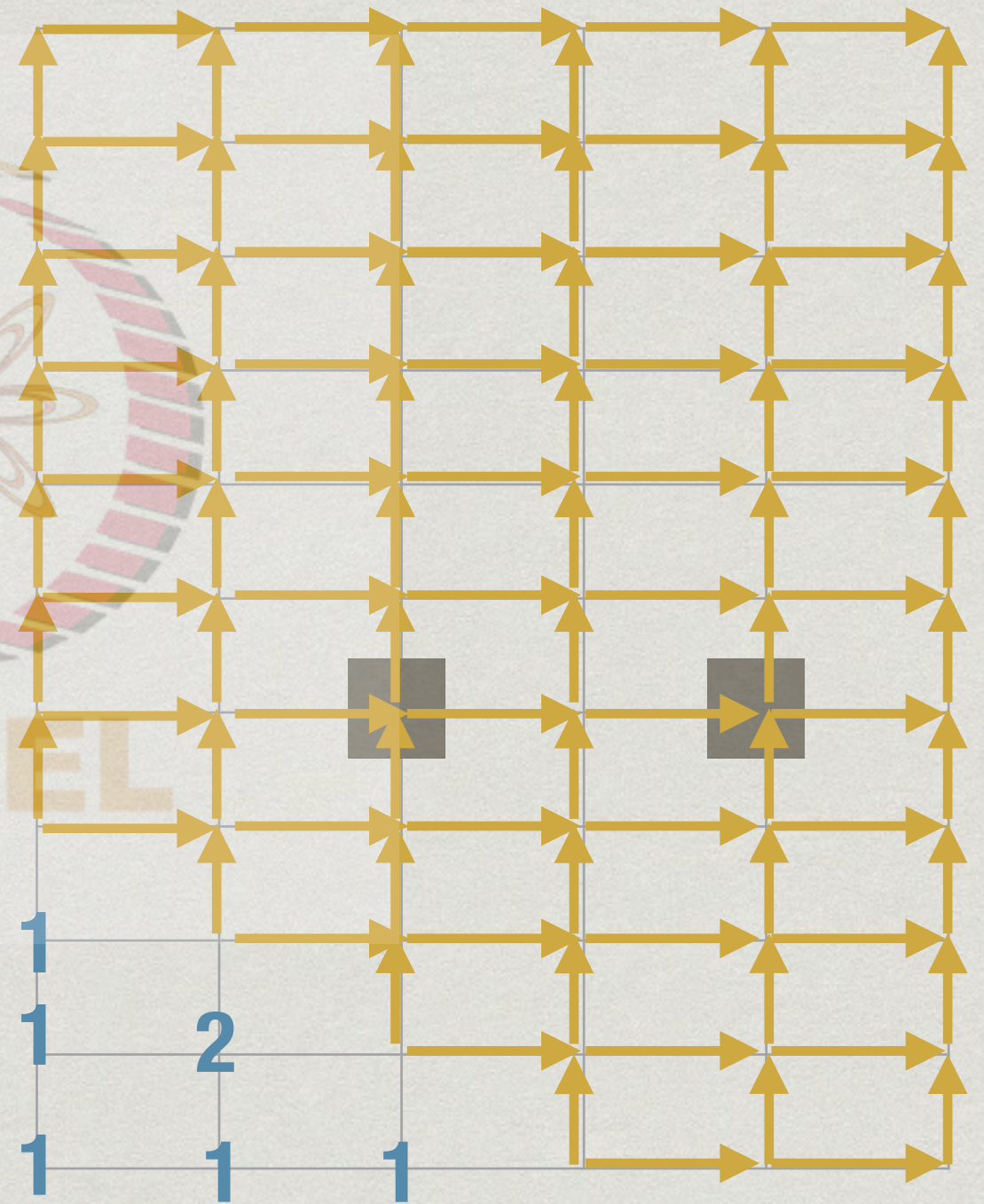
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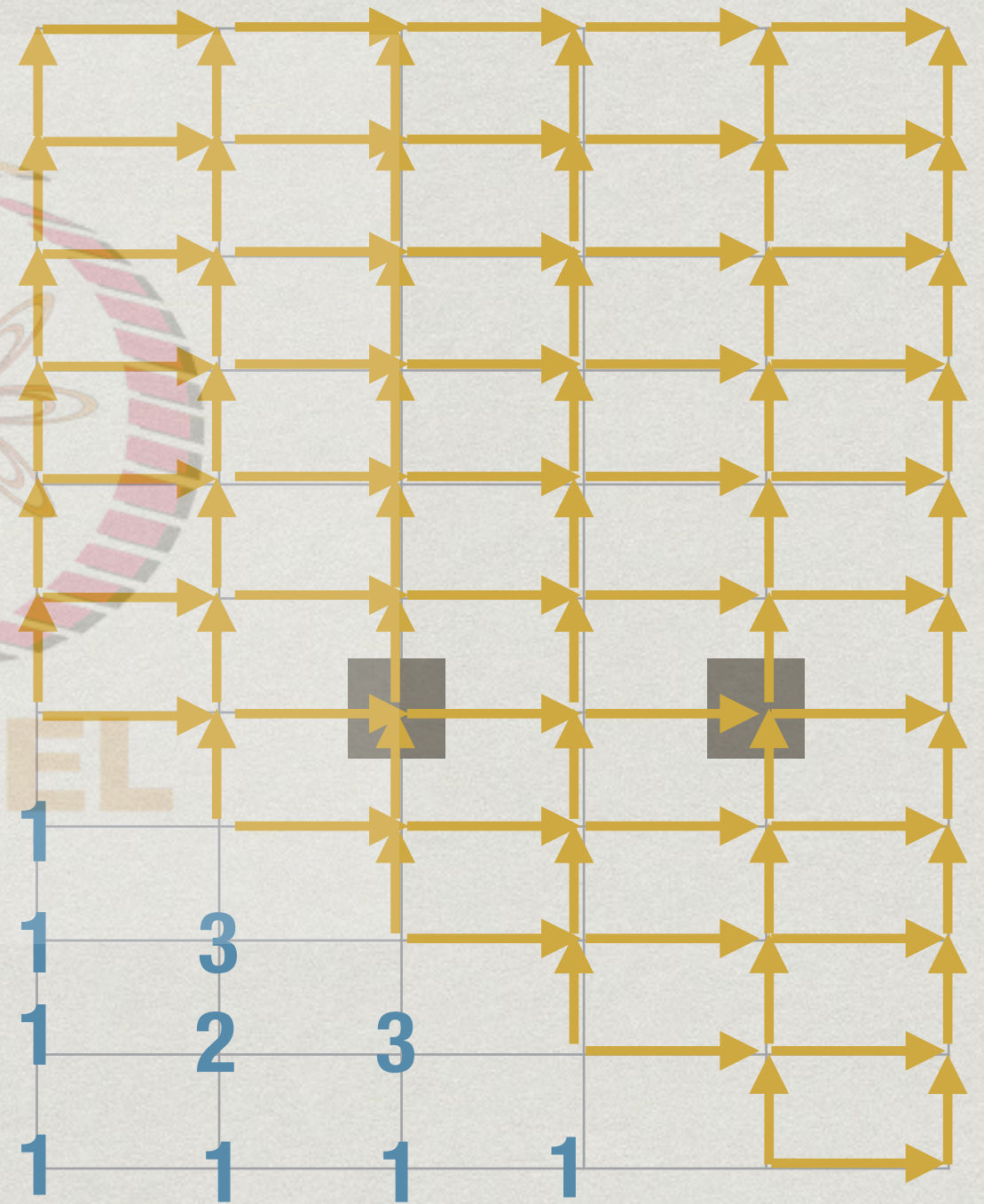
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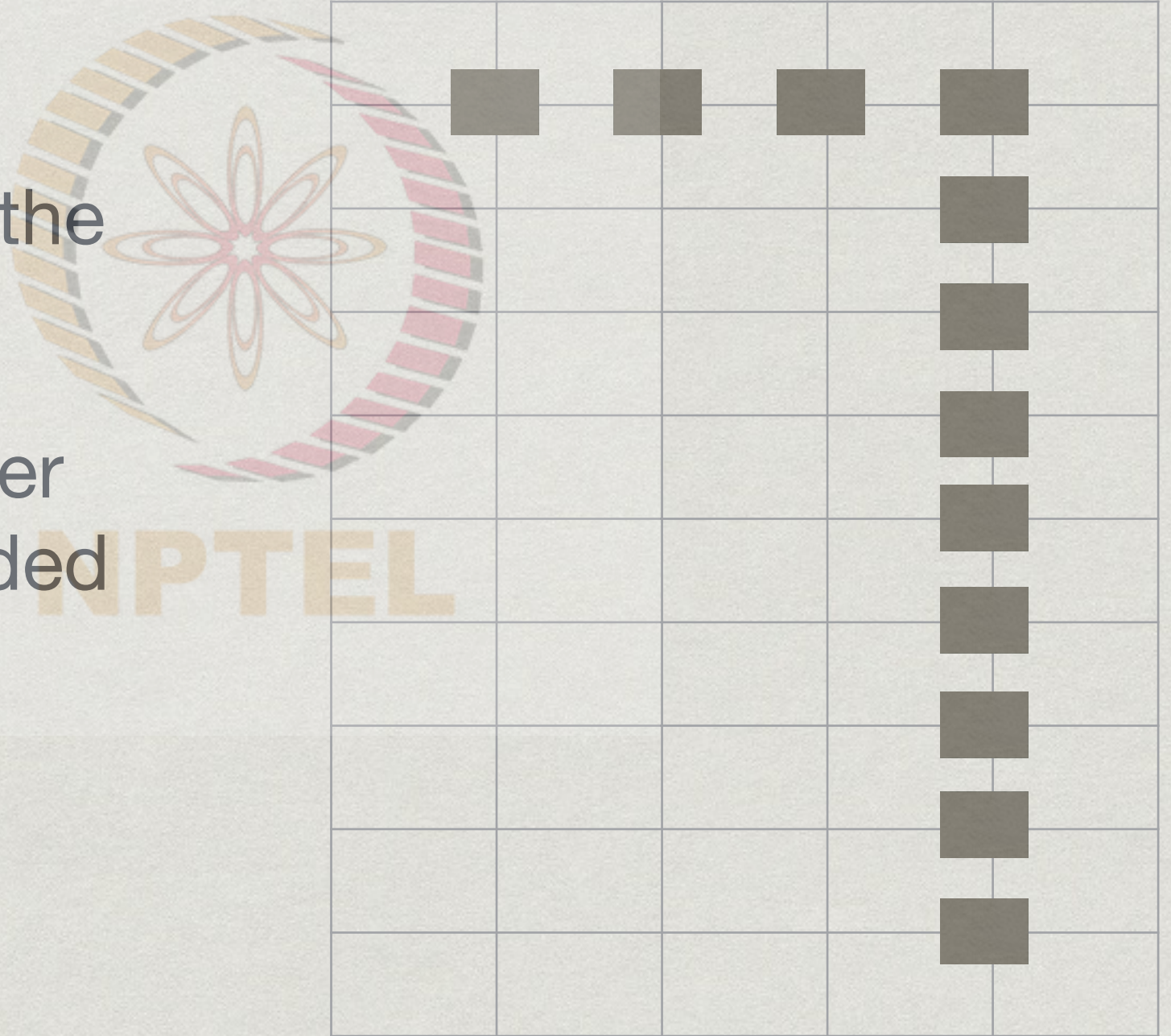
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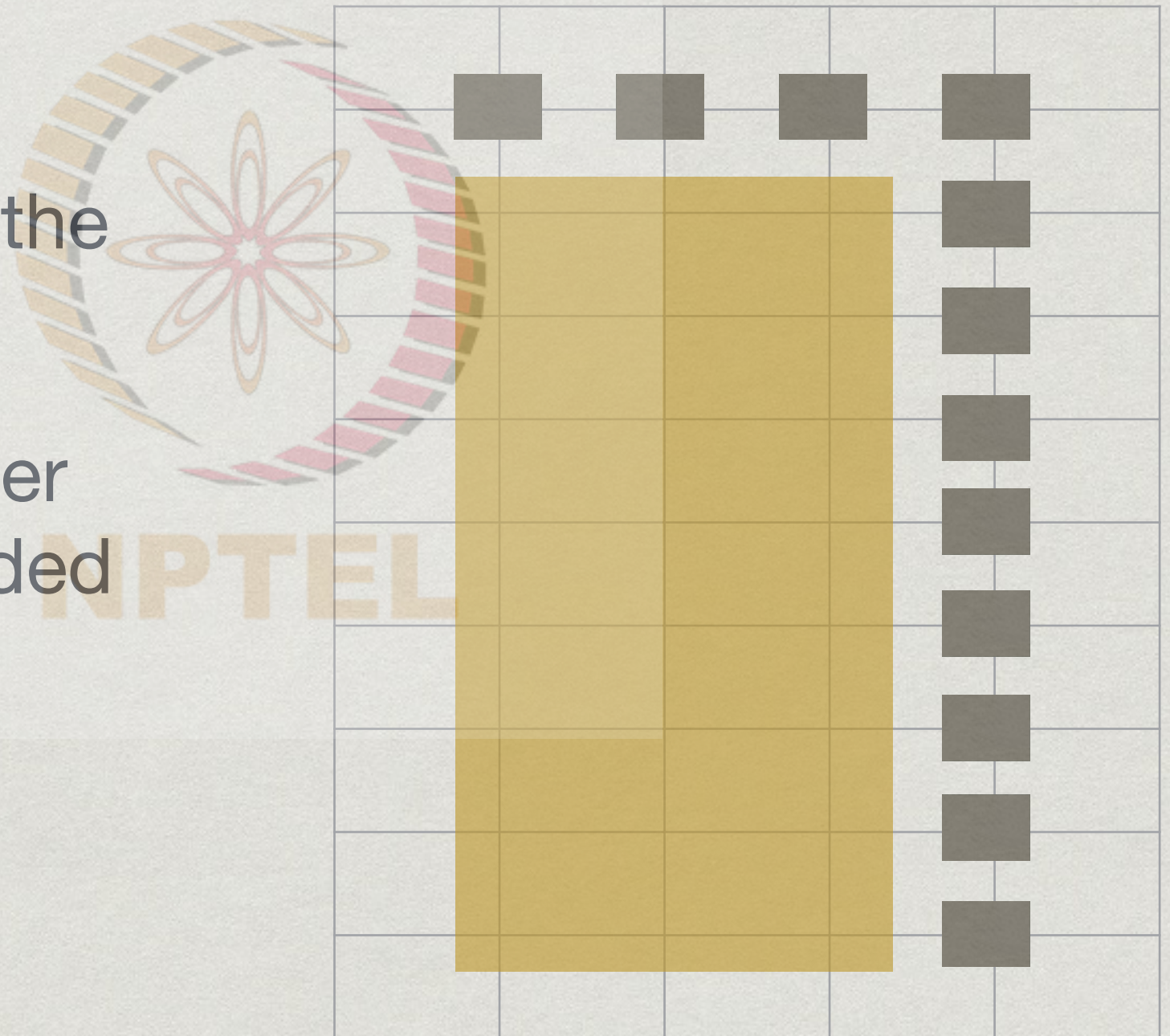
Memoization vs dynamic programming

- * Holes just inside the border
- * Memoization never explores the shaded region



Memoization vs dynamic programming

- * Holes just inside the border
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Memoization vs dynamic programming

- * Memo table has $O(m+n)$ entries
- * Dynamic programming blindly fills all $O(mn)$ entries
- * Iteration vs recursion
— “wasteful”
dynamic programming is still better, in general

