

# **Content**

- 1. Zoned Namespace SSD
- 2. Motivation
- 3. LSM-ZGC Design
- 4. Evaluation
- 5. Conclusions
- Appendix

#### 1. Zoned Namespace SSD

#### What is ZNS SSD?

# LBA space LBA space Zone 1 Zone 2 Zone 3 NAND NAND

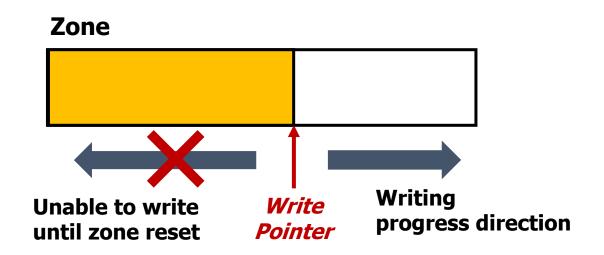
#### ✓ Benefits

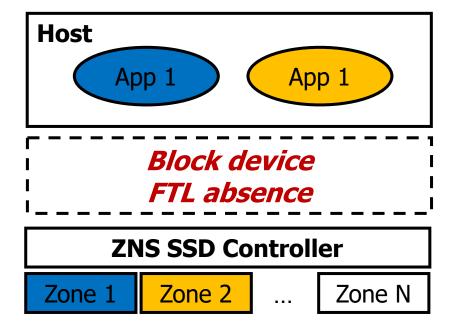
- Better performance and WAF by distributing different workloads into different zones
- Better isolation (IO Determinism)
- Reduce DRAM usage and Over-provisioning area in SSDs

#### 1. Zoned Namespace SSD

# What are the issues of ZNS SSD?

- ✓ Sequential write constraint: writes need to be conducted in a sequential manner, like the SMR drives.
- ✓ Host needs to control zones directly such as zone open, close, reset and zone garbage collection.



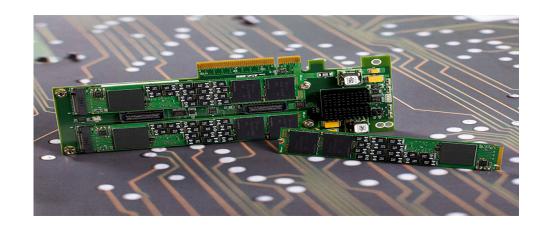


Sequential write constraint

Host Needs to handle zone controls

# How much is the Zone Garbage Collection (hereafter ZGC) overhead?

- ✓ Using real ZNS SSD prototype
- ✓ Zone size: 1GB (note that the typical segment size in LFS is 2MB)

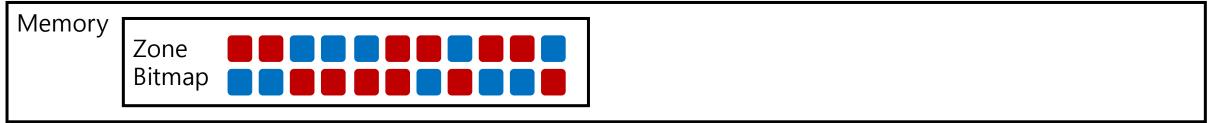


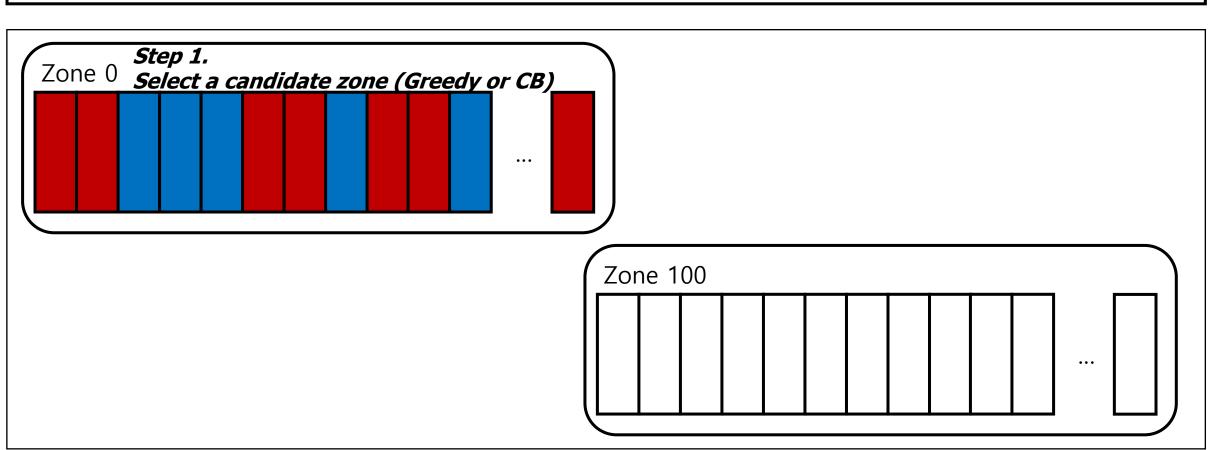
SK Hynix Prototype ZNS SSD

Table 1: ZNS SSD prototype information

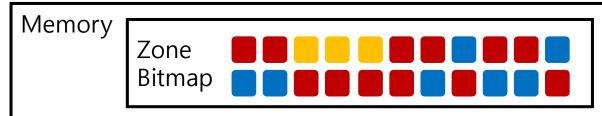
Item	Specification
SSD Capacity	1TB
Size of a Zone	1GB
Number of Zones	1024
Interface	PCIe Gen3
Protocol	NVMe 1.2.1

# **Basic Zone Garbage Collection (Basic\_ZGC)**

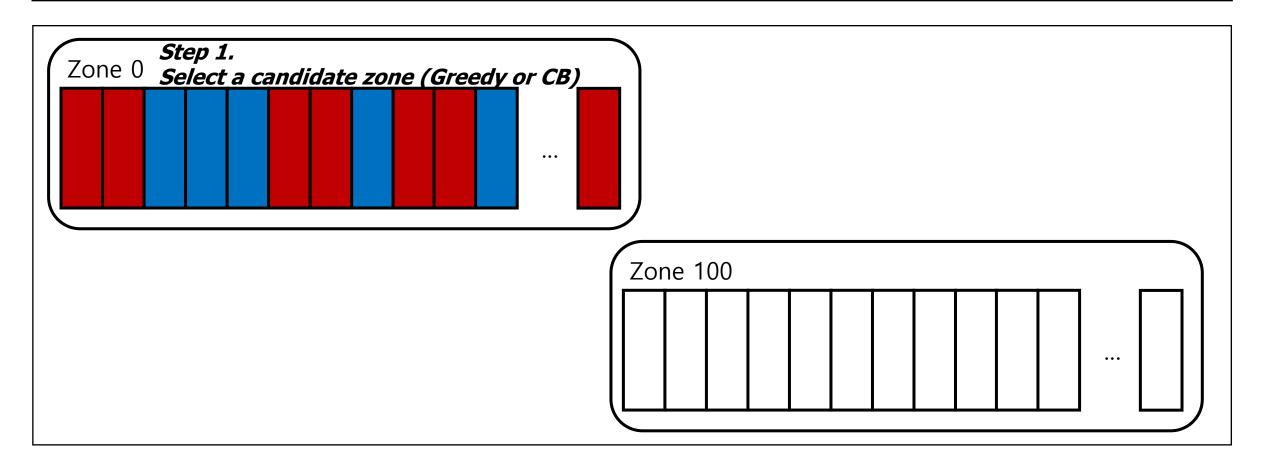




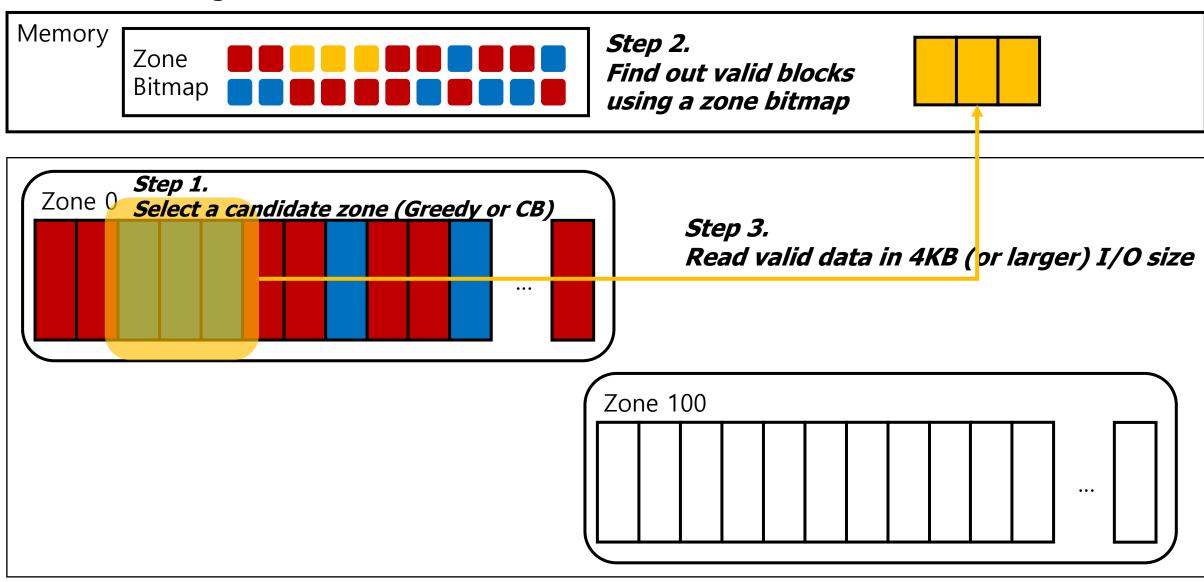
# **Basic Zone Garbage Collection**



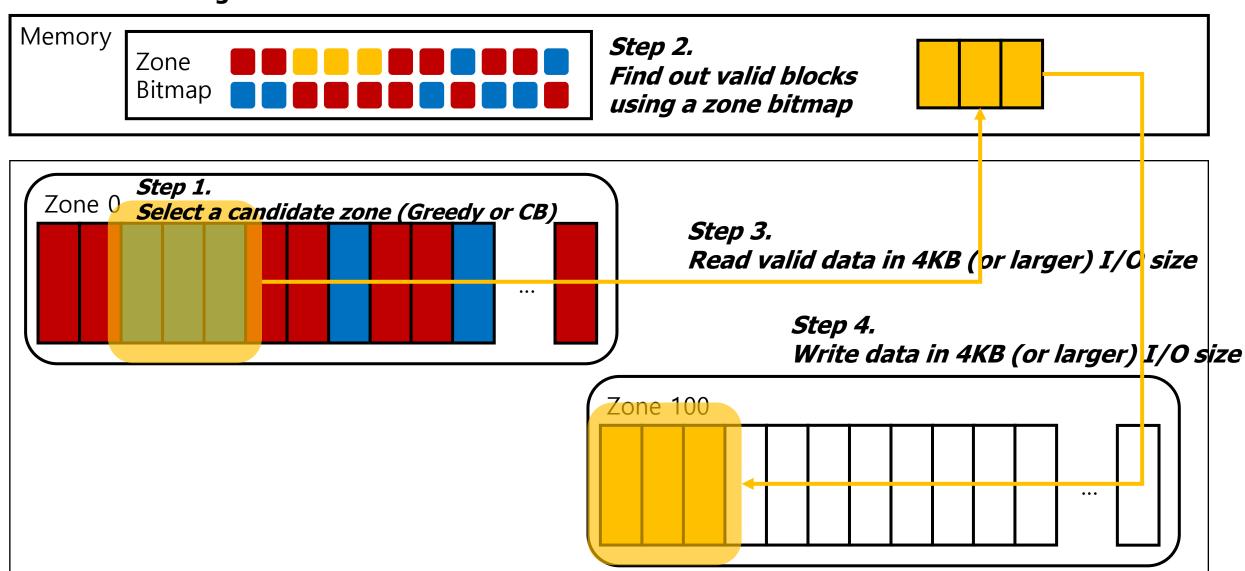
Step 2.
Find out valid blocks
using a zone bitmap



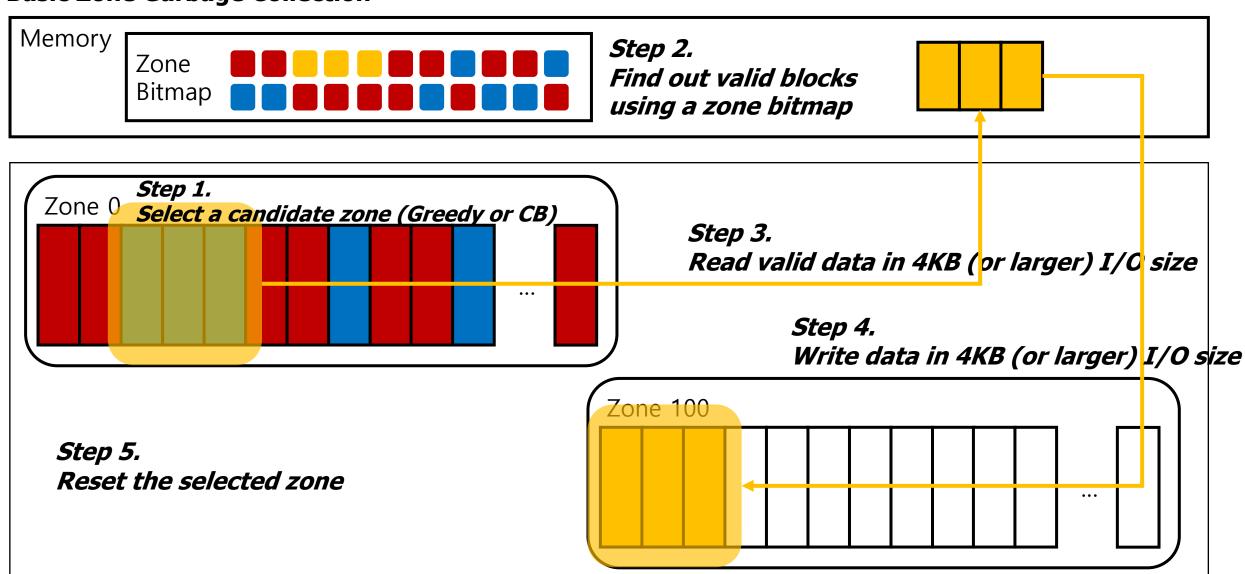
# **Basic Zone Garbage Collection**



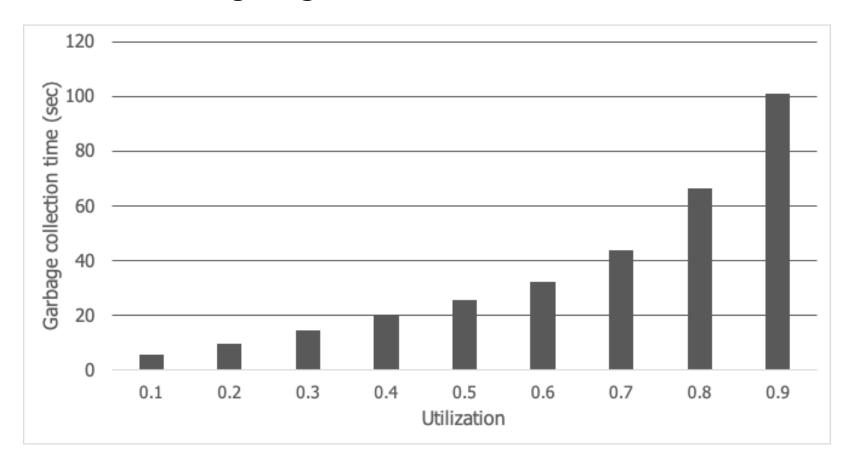
# **Basic Zone Garbage Collection**



#### **Basic Zone Garbage Collection**



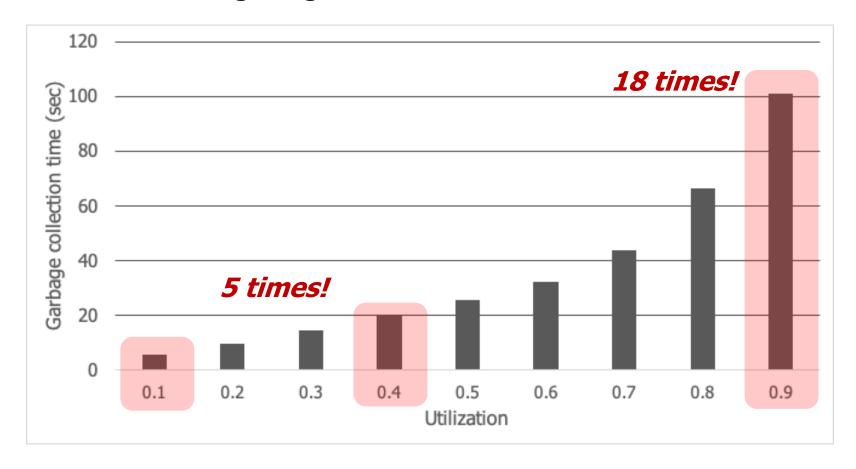
# **Observation 1: Zone garbage collection overhead**



• Zone: 1GB

• Block: 4KB

#### **Observation 1: Zone garbage collection overhead**



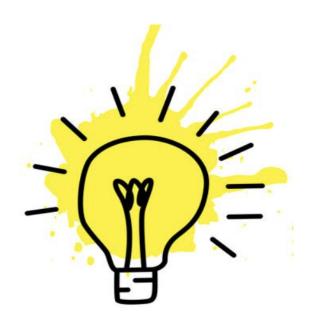
• Zone: 1GB

• Block: 4KB

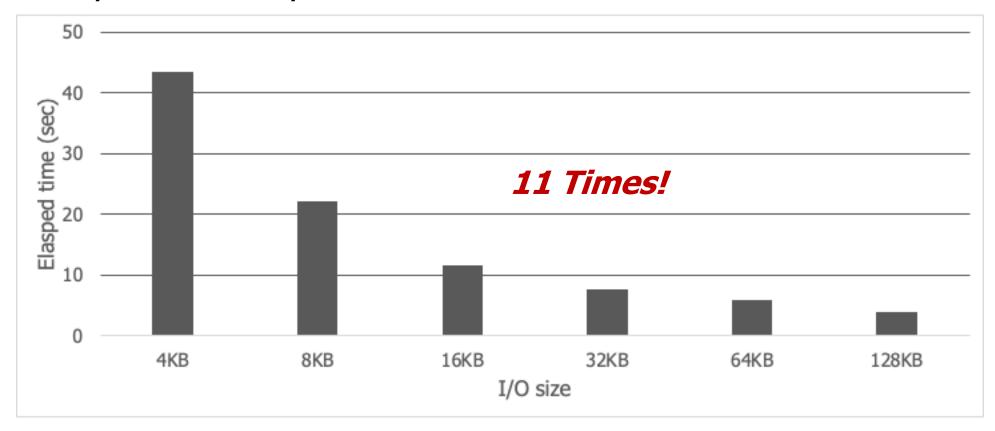
Motivation 1: reducing utilization of a candidate zone is indispensable

#### **Observation 2: I/O size for Read/Write**

- Another feature of ZNS SSD
  - ✓ A zone is, in general, mapped into multiple channels/ways.
- Then, how about read/write data in a larger I/O size (e.g. 128KB)?



## **Observation 2: I/O size for Read/Write**

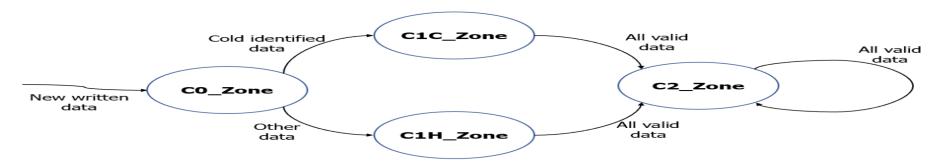


■ Motivation 2: accessing in a larger I/O size is beneficial in ZNS SSDs

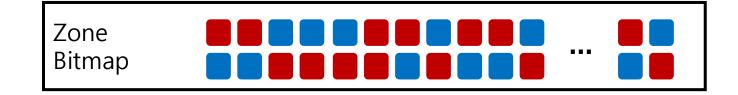
# So, Our ideas are

- 1) Make the utilization of a candidate zone low
- 2) Access data in a larger I/O size
- How to access data in a larger I/O size?
  - ✓ The coexistence of valid and invalid data makes it difficult
  - ✓ Read not only valid but also invalid data in a larger I/O size
- How to make the utilization of a candidate zone low?
  - √ Traditional hot/cold separation is not applicable in ZNS SSDs since zone is quite big
  - ✓ Employ the segment concept for finer-grained hot/cold separation

- Two management units
  - ✓ Zone: for garbage collection vs. Segment: for hot/cold separation
  - ✓ A zone is divided into multiple segments (1GB vs. 2MB in this study)
- Segment state and transition rule (refer to our paper for details)
  - ✓ New data → C0
  - ✓ During ZGC, survived data from CO
    - Data in a high utilized segment ( > threshold<sub>cold</sub>): cold → C1C
    - Others: hot (or unknown) → C1H
    - Reasoning: spatial locality, also observed in previous studies such as F2FS (FAST'15),
       Multi-stream (FAST'19), Key-range locality (FAST'20)
  - ✓ During ZGC, survived data from C1C or C1H (second survived data)→ C2



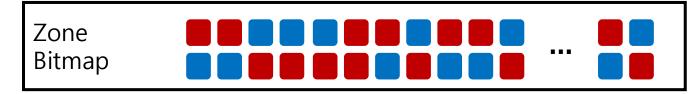
# LSM(Log Structured Merge) Zone GC

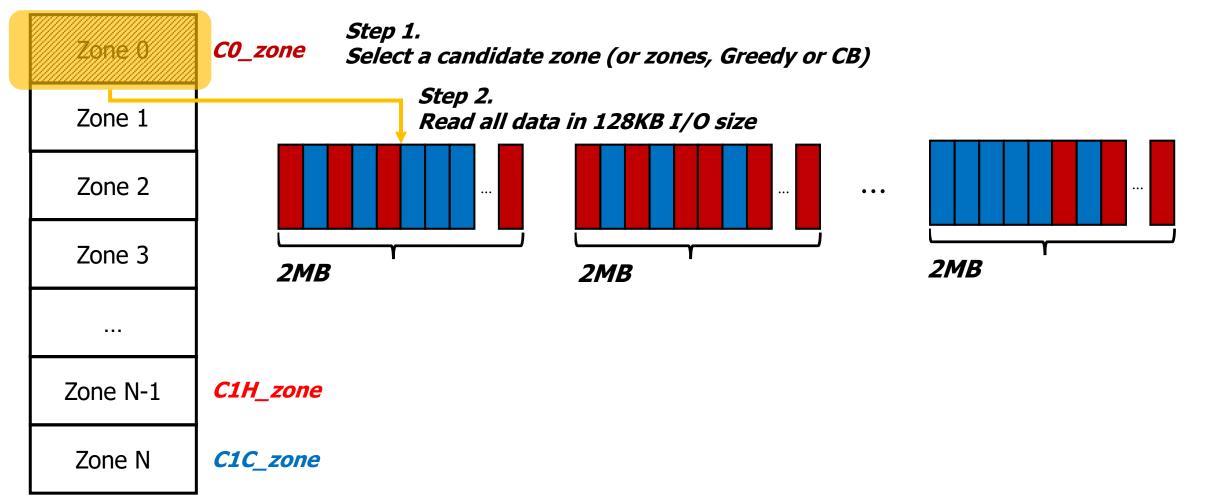


Step 1. Zone 0 CO\_zone Select a candidate zone (or zones, Greedy or CB) Zone 1 Zone 2 Zone 3 Zone N-1 C1H\_zone Zone N C1C\_zone

18

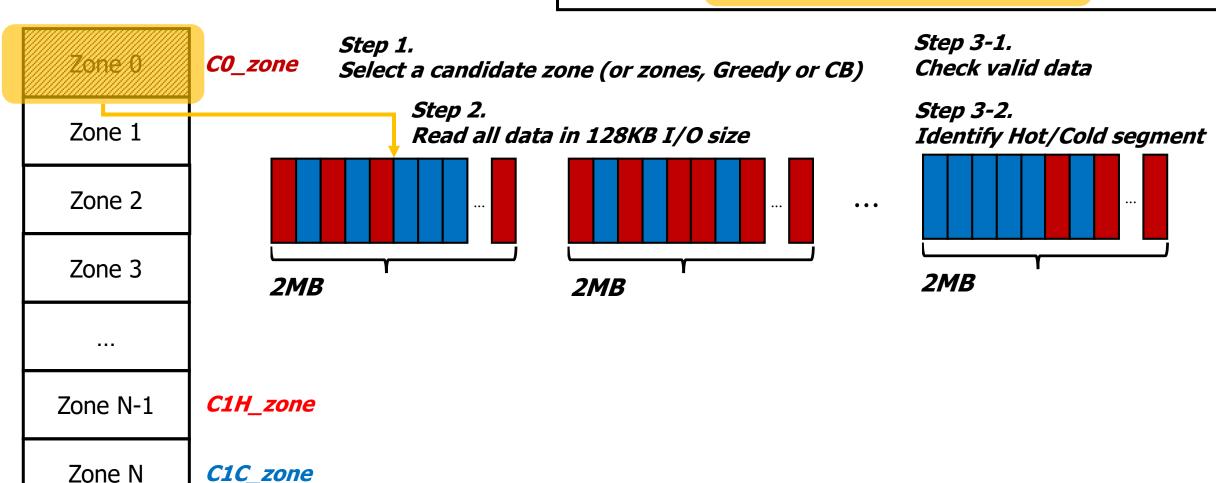
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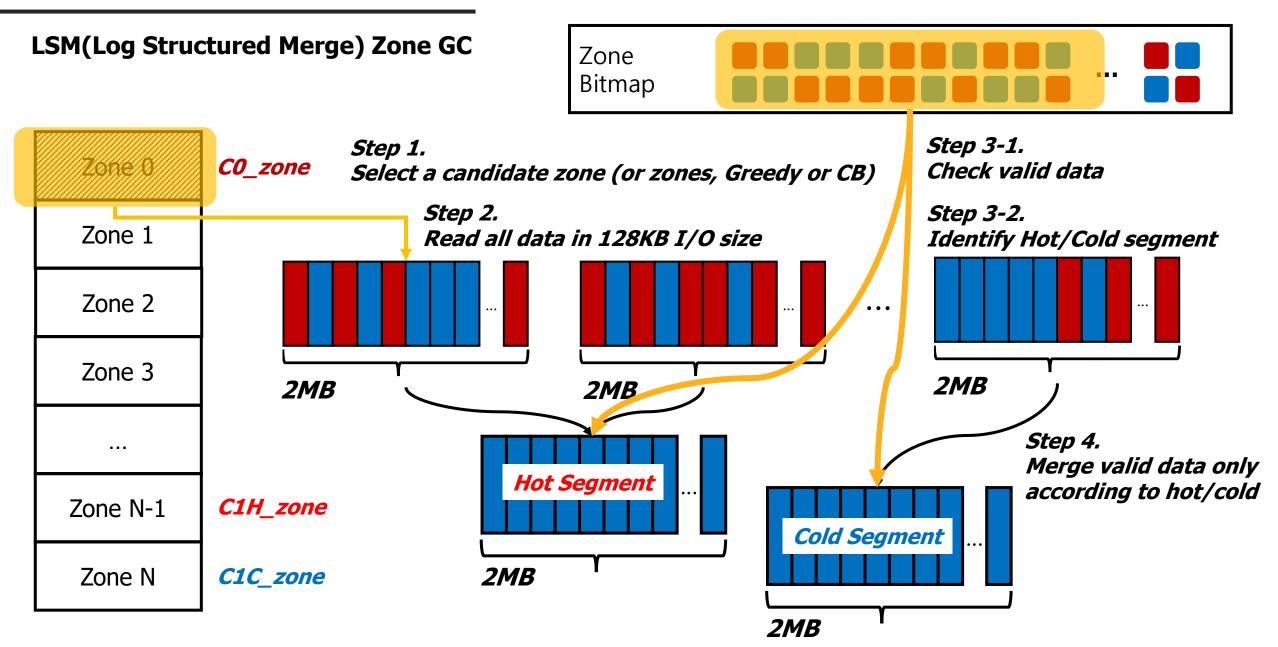


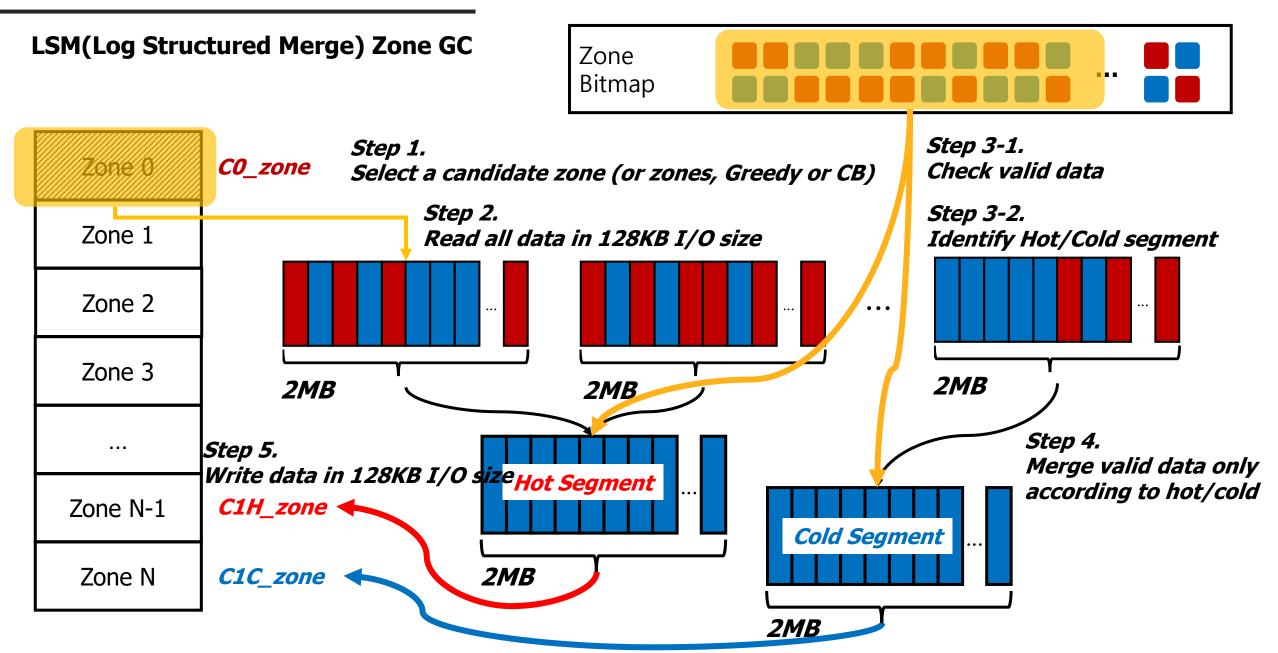


#### LSM(Log Structured Merge) Zone GC



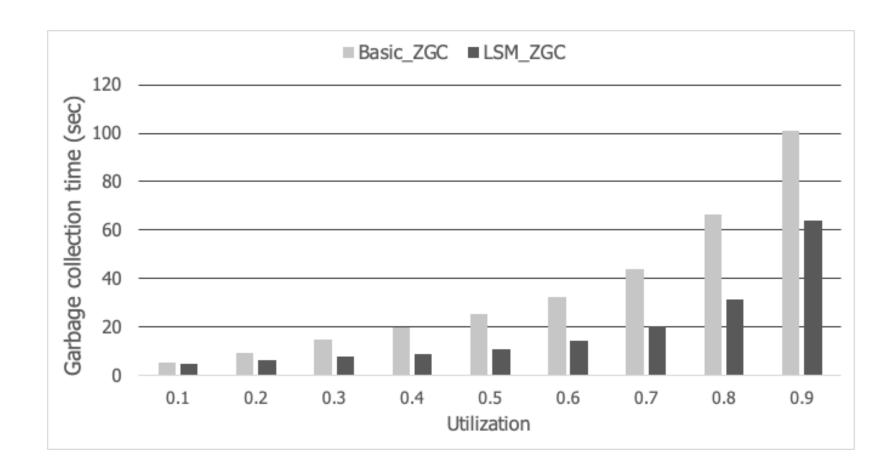






#### 4. Evaluation

#### Garbage collection overhead: uniform update pattern



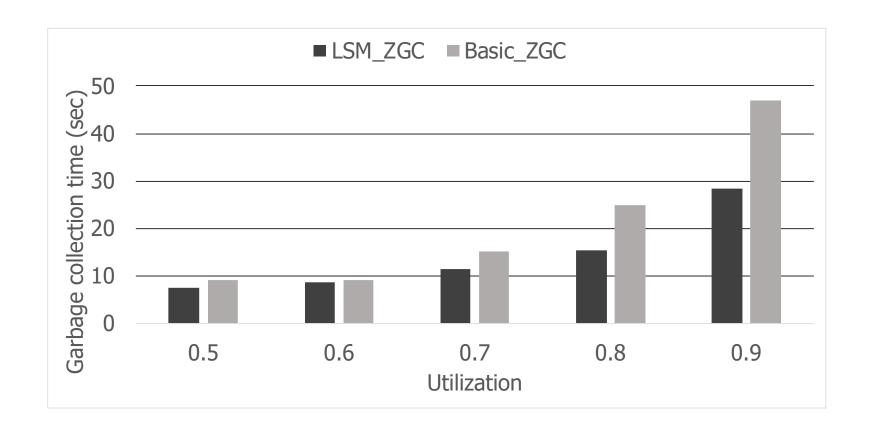
# Average of 1.9 times Max of 2.3 times

#### Experimental environment

- Intel Core i7 (8 core)
- 16GB DRAM
- 1TB ZNS SSD
- Size of Zone: 1GB

#### 4. Evaluation

#### Garbage collection overhead: skewed update pattern



# Average of 1.4 times Max of 1.6 times

#### **Parameters**

- Workload: 70/30 hot/cold ratio
- Threahold<sub>cold</sub>: 0.8
- average utilization: x-axis

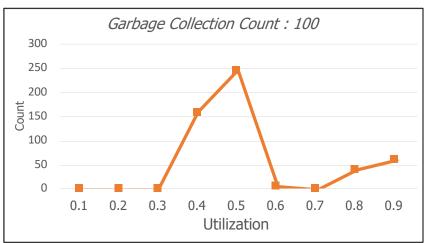
#### Experimental environment

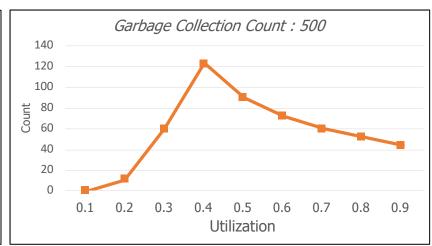
- Intel Core i7 (8 core)
- 16GB DRAM
- 1TB ZNS SSD
- Size of Zone: 1GB

#### 4. Evaluation

#### **Hot/Cold Separation**

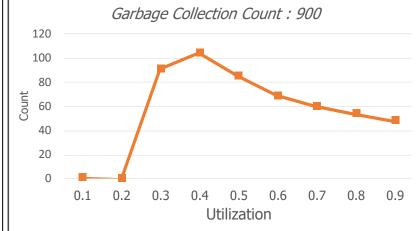
# ✓ Without hot/cold separation



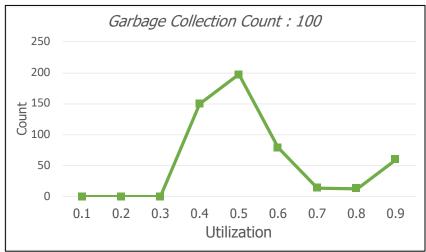


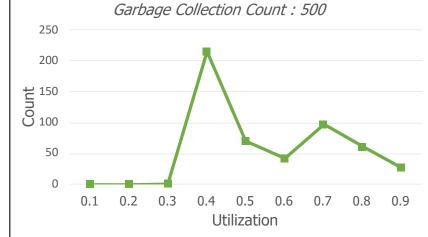


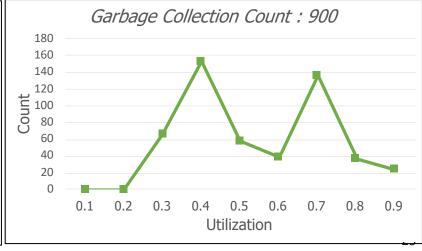
- Workload: 70/30 hot/cold ratio
- Threahold<sub>cold</sub>: 0.8
- Average utilization: 0.6



#### ✓ With hot/cold separation







# Our contributions

- Observation: a zone garbage collection really matters
- Proposal: a new LSM-style zone garbage collection scheme
- Evaluation: real implementation based results

# Future work

- We are currently extending F2FS on our ZNS SSD prototype
- Also, evaluating LSM ZGC under diverse workloads with different hot /cold ratio, data size, initial placement and classification policies



12th USENIX Workshop on Hot Topic in Storage and File System (HotStorage 20), 2020

Thank You!

Questions?

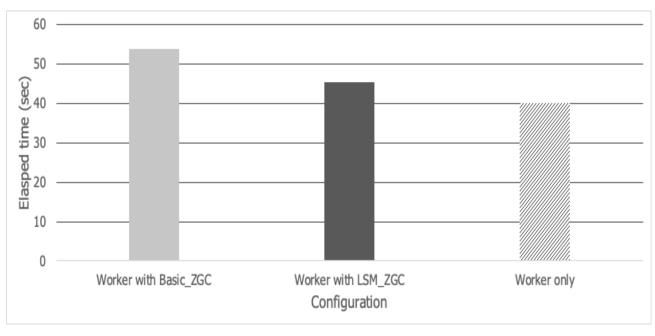
2022.12.15 张茂荣

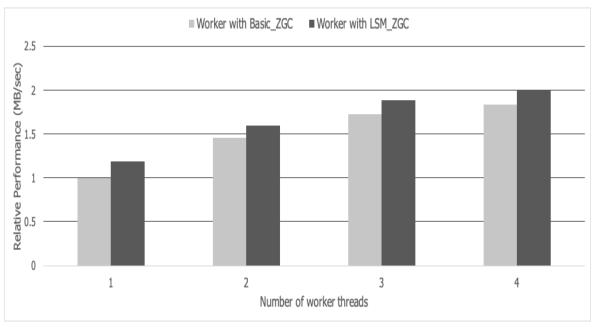
华中科技大学



## 6. Appendix

#### Performance comparison using multi-thread & Scalability





Worker only: 40

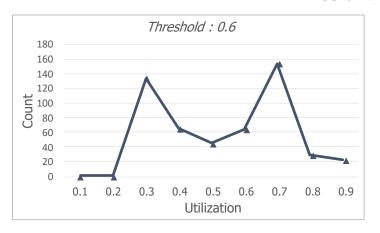
With LSM\_ZGC: 45

With Basic\_ZGC: 53

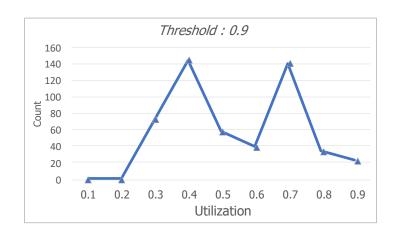
Non-linear Scalability

#### **Sensitive Analysis: various parameters**

# ✓ Effect of threshold<sub>cold</sub> (initial utilization: 0.6)







# ✓ Effect of initial utilization of a zone (threshold<sub>cold</sub>: 0.8)

