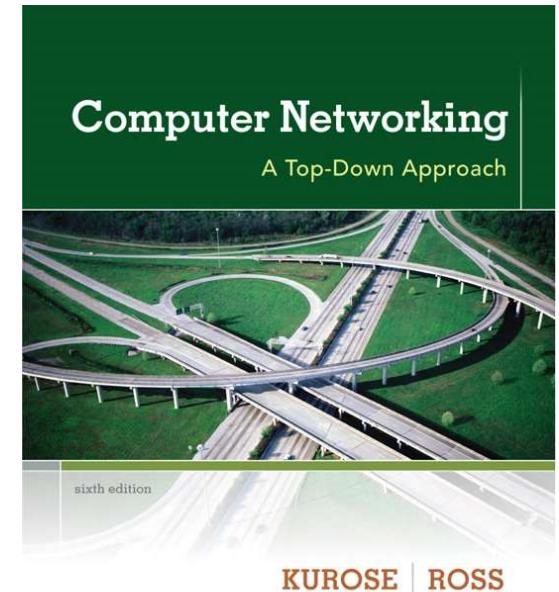


Chapter I Introduction



*Computer
Networking: A Top
Down Approach*
6th edition
Jim Kurose, Keith Ross
Addison-Wesley
March 2012

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Chapter 1: introduction

our goal:

- ❖ get “feel” and terminology
- ❖ more depth, detail
later in course
- ❖ approach:
 - use Internet as example

overview:

- ❖ what’s the Internet?
- ❖ what’s a protocol?
- ❖ network edge; hosts, access net, physical media
- ❖ network core: packet/circuit switching, Internet structure
- ❖ performance: loss, delay, throughput
- ❖ security
- ❖ protocol layers, service models
- ❖ history

Chapter I: roadmap

I.1 what *is* the Internet?

I.2 network edge

- end systems, access networks, links

I.3 network core

- packet switching, circuit switching, network structure

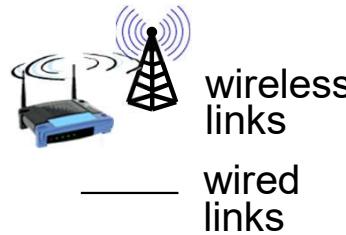
I.4 delay, loss, throughput in networks

I.5 protocol layers, service models

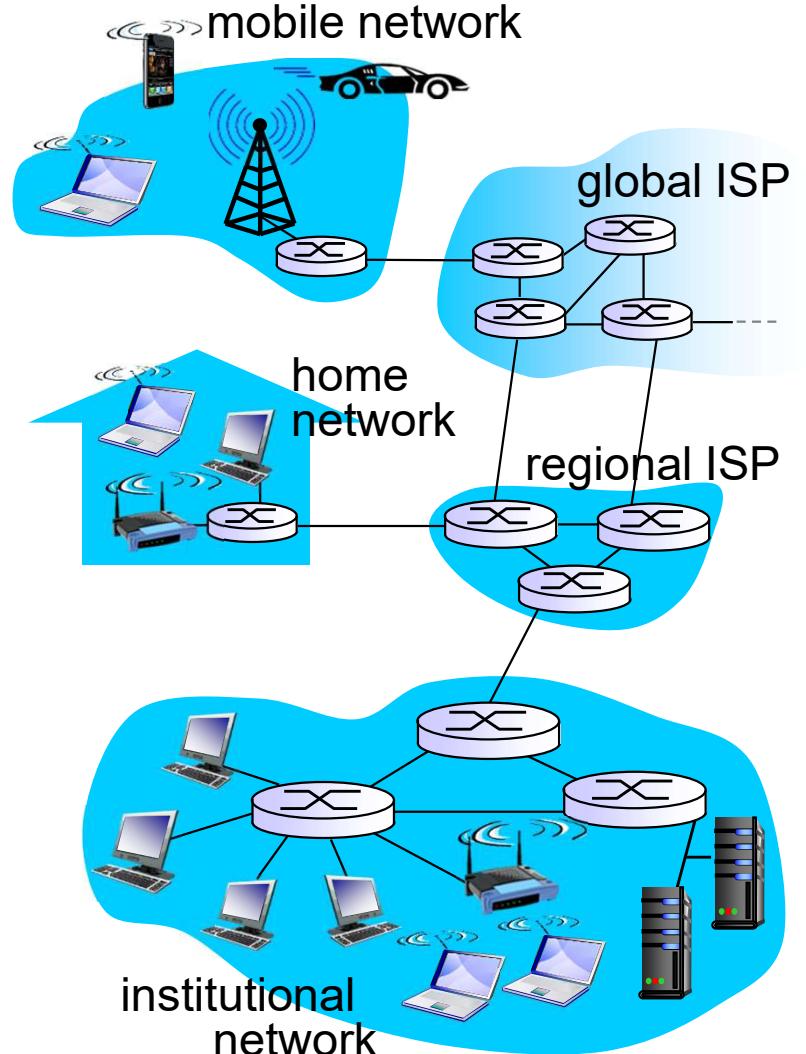
I.6 networks under attack: security

I.7 history

What's the Internet: “nuts and bolts” view



- ❖ millions of connected computing devices:
 - *hosts = end systems*
 - running *network apps*
- ❖ *communication links*
 - fiber, copper, radio, satellite
 - transmission rate: *bandwidth*
- ❖ *Packet switches*: forward packets (chunks of data)
 - *routers and switches*



“Fun” internet appliances



IP picture frame
<http://www.ceiva.com/>



Internet refrigerator



Web-enabled toaster +
weather forecaster



Tweet-a-watt:
monitor energy use



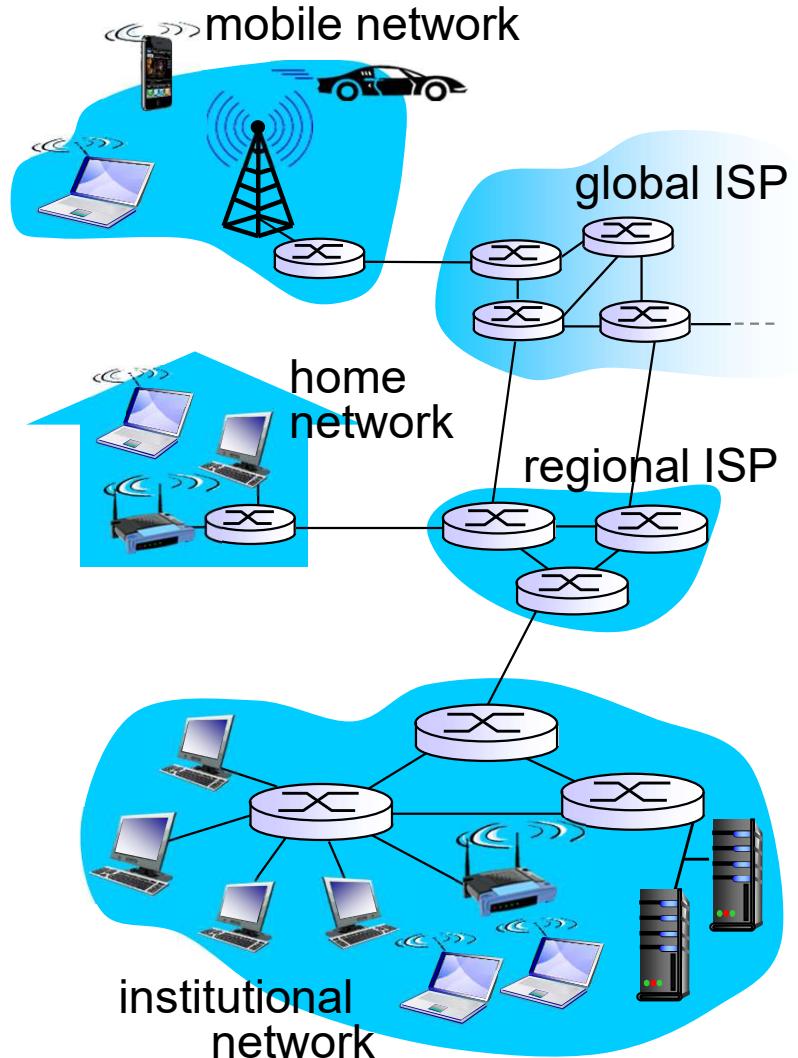
Slingbox: watch,
control cable TV remotely



Internet phones

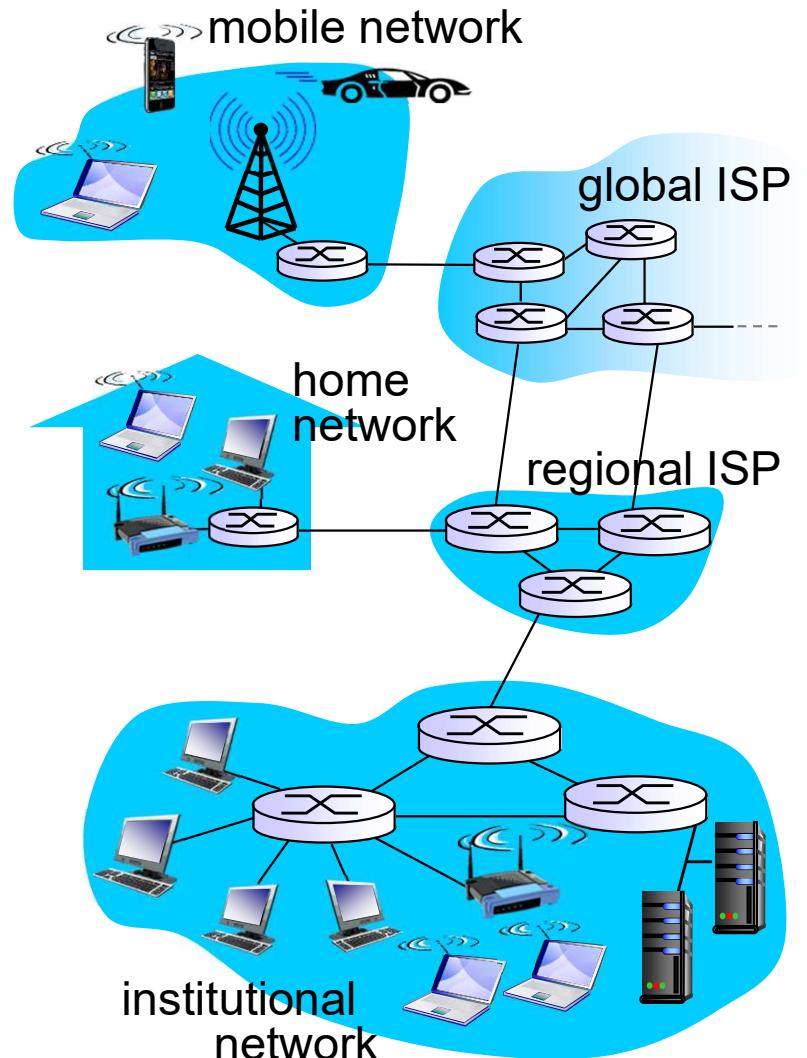
What's the Internet: “nuts and bolts” view

- ❖ *Internet: “network of networks”*
 - Interconnected ISPs
- ❖ *protocols* control sending, receiving of msgs
 - e.g., TCP, IP, 802.11, HTTP
- ❖ *Internet standards*
 - RFC: Request for comments
 - IETF: Internet Engineering Task Force



What's the Internet: a service view

- ❖ *Infrastructure that provides services to applications:*
 - Web, VoIP, email, games, e-commerce, social nets, ...
- ❖ *provides programming interface to apps*
 - hooks that allow sending and receiving app programs to “connect” to Internet
 - provides service options, analogous to postal service



What's a protocol?

human protocols:

- ❖ “what’s the time?”
- ❖ “I have a question”
- ❖ introductions

... specific msgs sent

... specific actions taken
when msgs received, or
other events

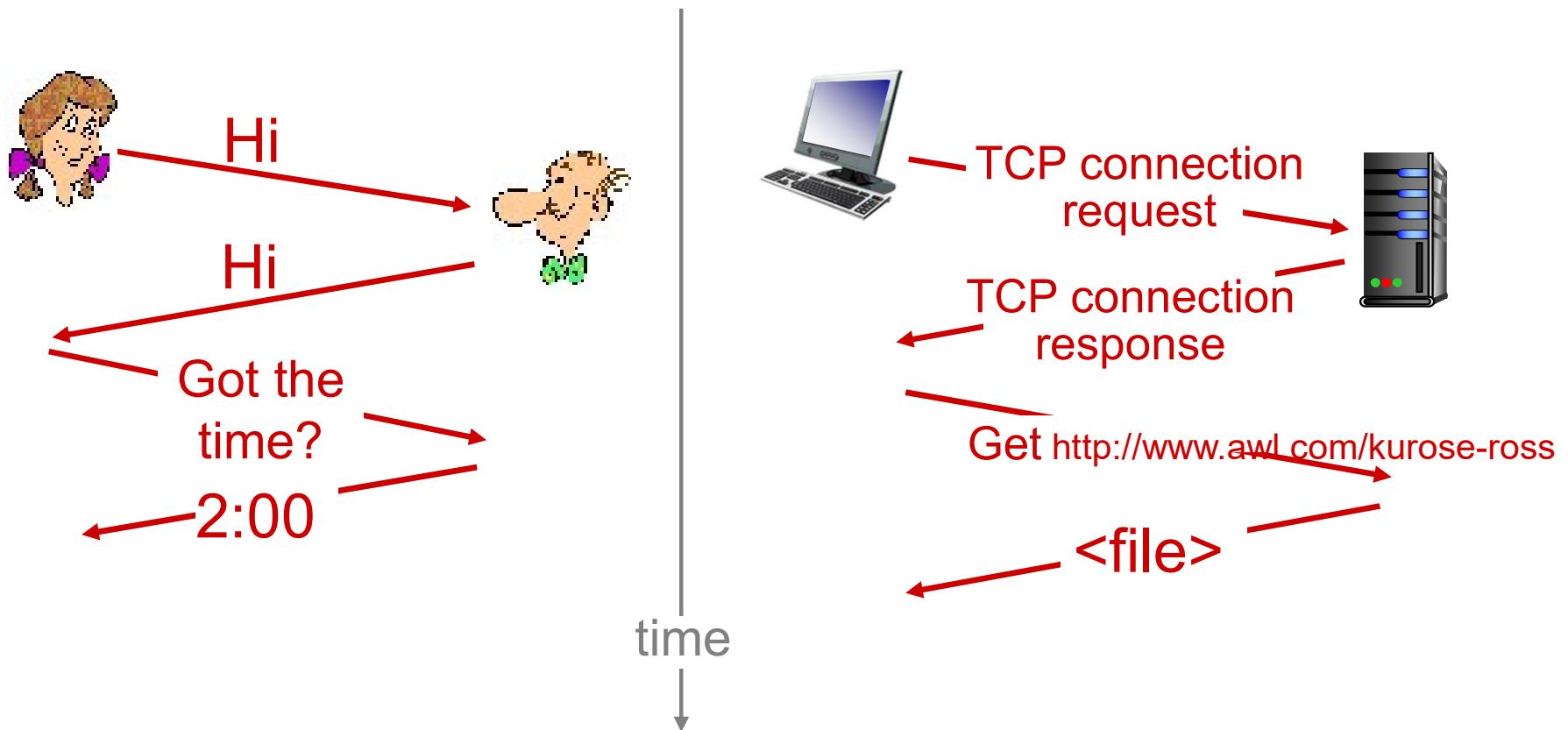
network protocols:

- ❖ machines rather than humans
- ❖ all communication activity in Internet governed by protocols

*protocols define format, order
of msgs sent and received
among network entities,
and actions taken on msg
transmission, receipt*

What's a protocol?

a human protocol and a computer network protocol:



Q: other human protocols?

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- packet switching, circuit switching, network structure

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I.5 protocol layers, service models

I.6 networks under attack: security

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A closer look at network structure:

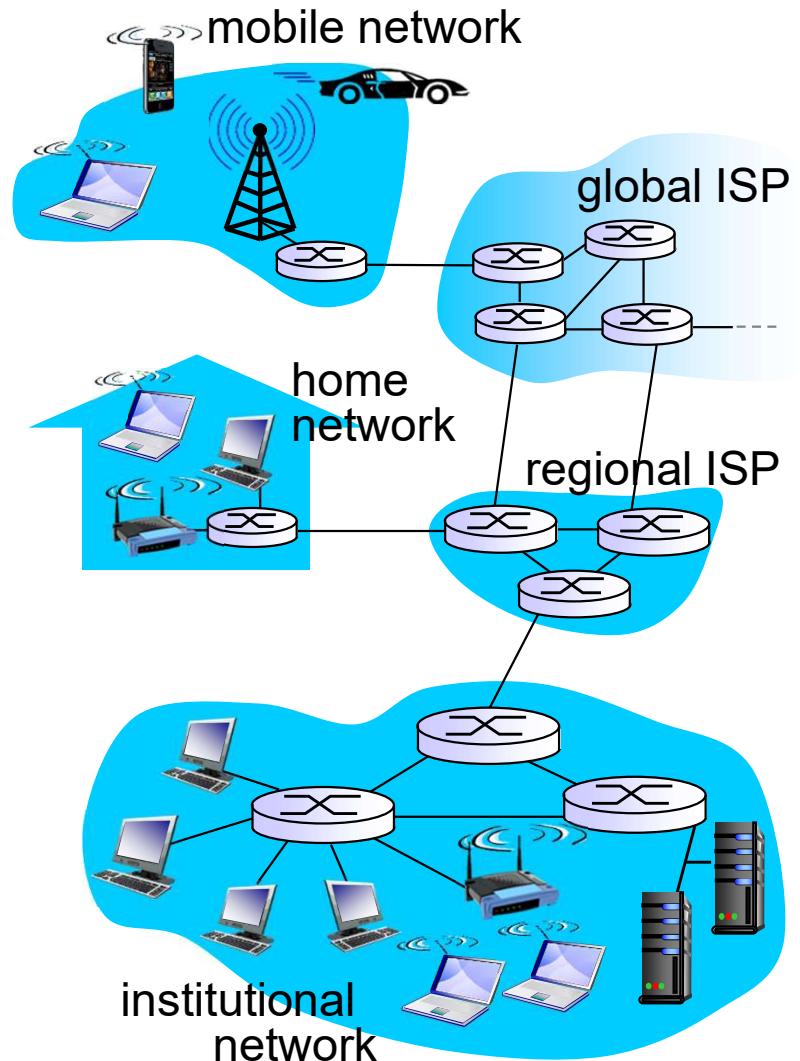
❖ *network edge:*

- hosts: clients and servers
- servers often in data centers

❖ *access networks, physical media:* wired, wireless communication links

❖ *network core:*

- interconnected routers
- network of networks



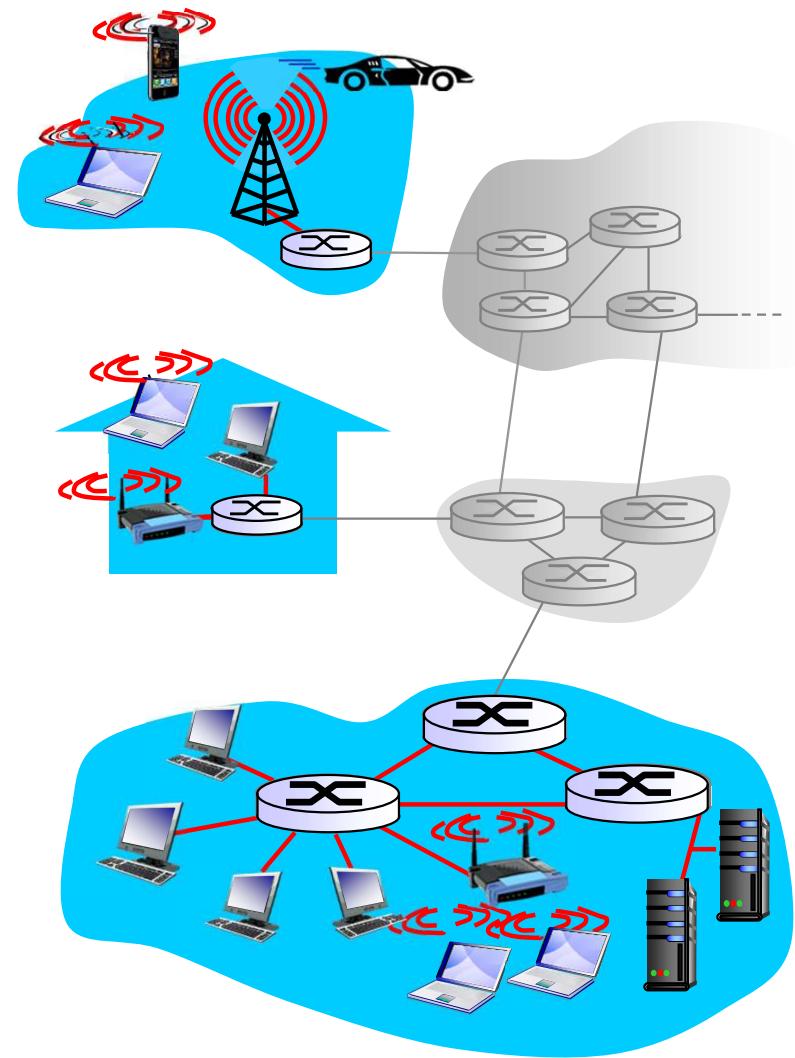
Access networks and physical media

Q: How to connect end systems to edge router?

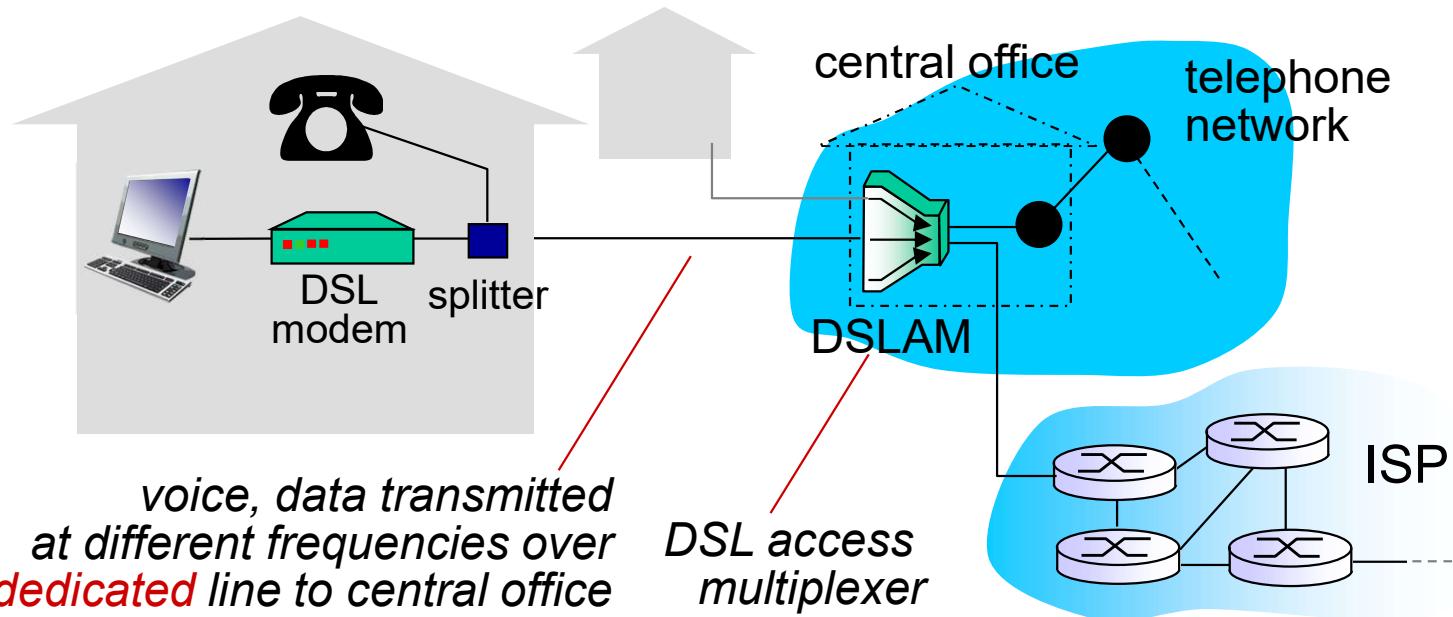
- ❖ residential access nets
- ❖ institutional access networks (school, company)
- ❖ mobile access networks

keep in mind:

- ❖ bandwidth (bits per second) of access network?
- ❖ shared or dedicated?

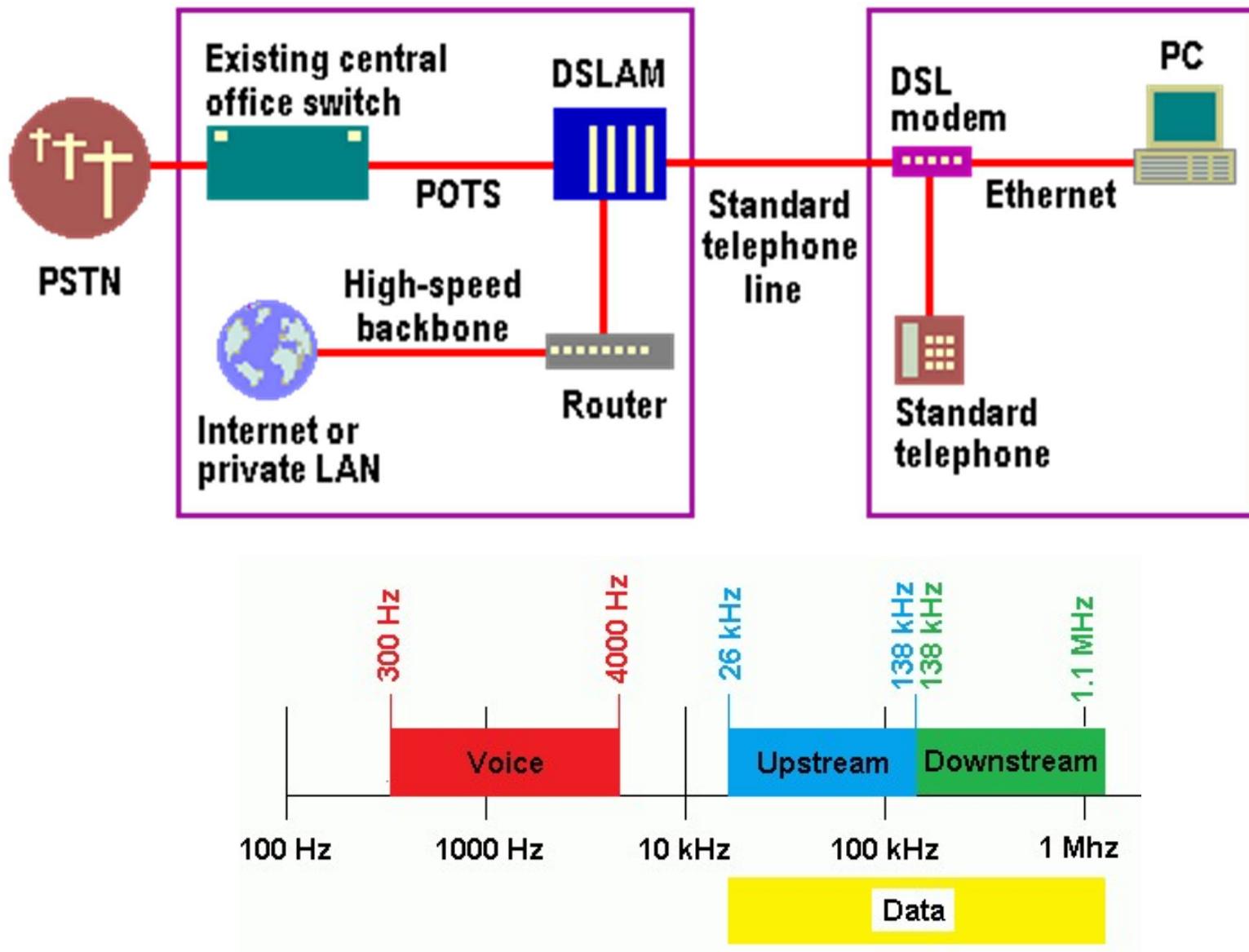


Access net: digital subscriber line (DSL)



- ❖ use **existing** telephone line to central office DSLAM
 - data over DSL phone line goes to Internet
 - voice over DSL phone line goes to telephone net
- ❖ < 2.5 Mbps upstream transmission rate (typically < 1 Mbps)
- ❖ < 24 Mbps downstream transmission rate (typically < 10 Mbps)

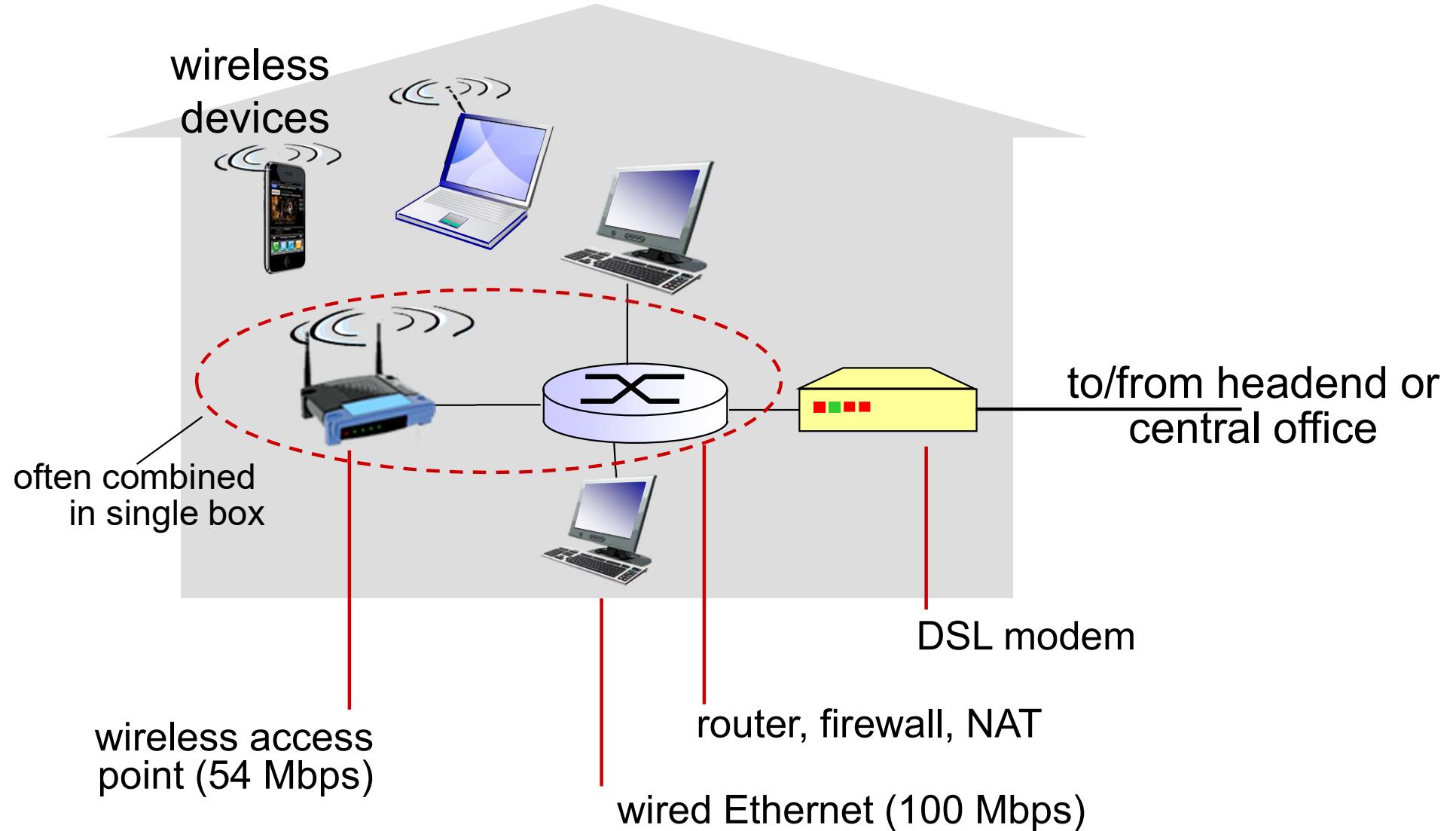
Access net: digital subscriber line (DSL)



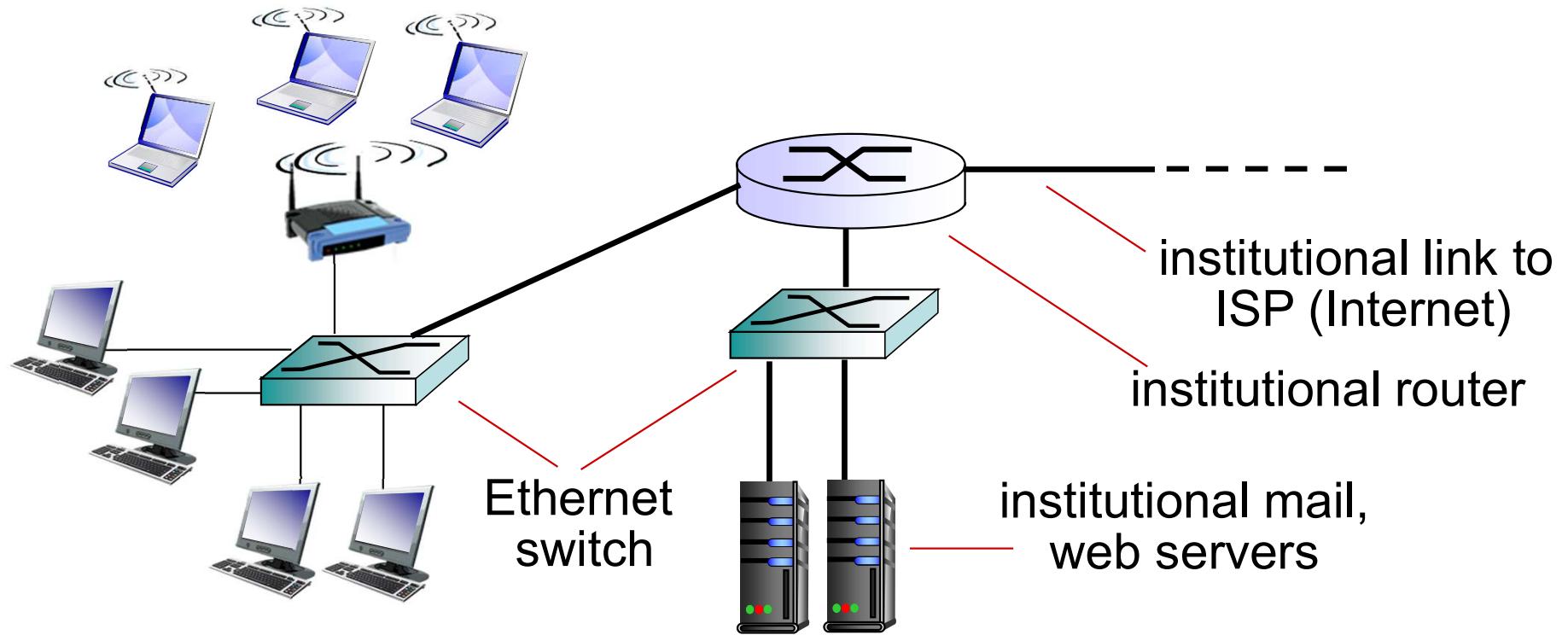
“DSLAM”



Access net: home network



Enterprise access networks (Ethernet)



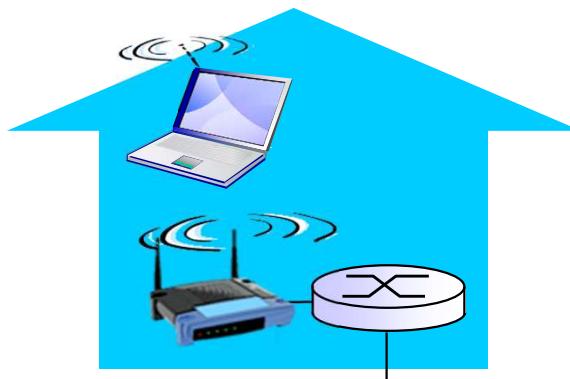
- ❖ typically used in companies, universities, etc
- ❖ 10 Mbps, 100Mbps, 1 Gbps, 10Gbps transmission rates
- ❖ today, end systems typically connect into Ethernet switch

Wireless access networks

- ❖ shared wireless access network connects end system to router
 - via base station aka “access point”

wireless LANs:

- within building (100 ft)
- 802.11b/g (WiFi): 11, 54 Mbps transmission rate



to Internet

wide-area wireless access

- provided by telco (cellular) operator, 10's km
- between 1 and 10 Mbps
- 3G, 4G: LTE

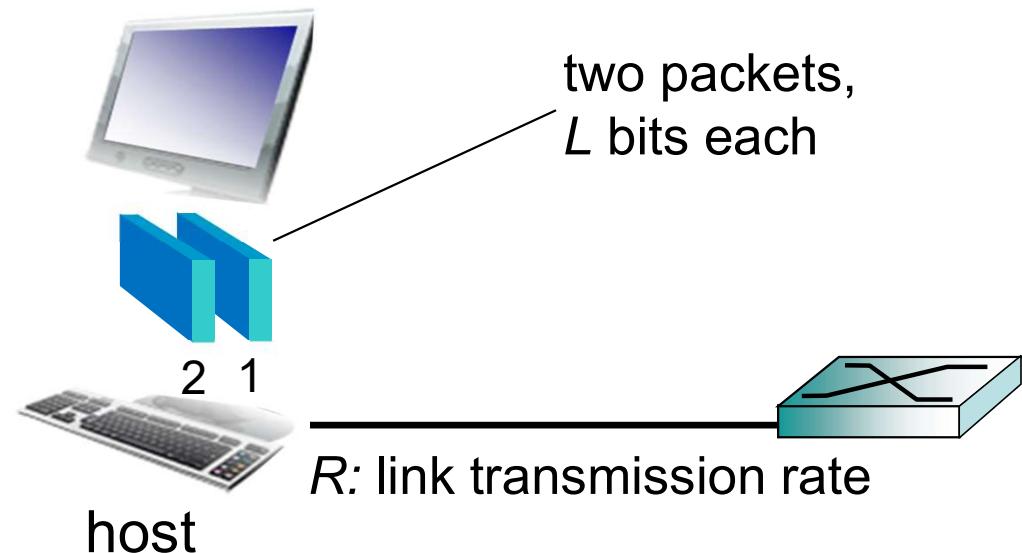


to Internet

Host: sends packets of data

host sending function:

- ❖ takes application message
- ❖ breaks into smaller chunks, known as *packets*, of length L bits
- ❖ transmits packet into access network at *transmission rate R*
 - link transmission rate, aka link *capacity*, aka *link bandwidth*



$$\text{packet transmission delay} = \frac{\text{time needed to transmit } L\text{-bit packet into link}}{R \text{ (bits/sec)}}$$

Physical media

- ❖ **bit:** propagates between transmitter/receiver pairs
- ❖ **physical link:** what lies between transmitter & receiver
- ❖ **guided media:**
 - signals propagate in solid media: copper, fiber, coax
- ❖ **unguided media:**
 - signals propagate freely, e.g., radio

twisted pair (TP)

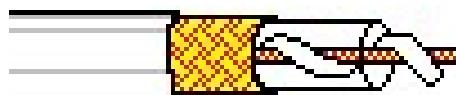
- ❖ two insulated copper wires
 - Category 5: 100 Mbps, 1 Gbps Ethernet
 - Category 6: 10Gbps



Physical media: coax, fiber

coaxial cable:

- ❖ two concentric copper conductors
- ❖ bidirectional
- ❖ broadband:
 - multiple channels on cable
 - HFC



fiber optic cable:

- ❖ glass fiber carrying light pulses, each pulse a bit
- ❖ high-speed operation:
 - high-speed point-to-point transmission (e.g., 10' s-100' s Gpbs transmission rate)
- ❖ low error rate:
 - repeaters spaced far apart
 - immune to electromagnetic noise



Physical media: radio

- ❖ signal carried in electromagnetic spectrum
- ❖ no physical “wire”
- ❖ bidirectional
- ❖ propagation environment effects:
 - reflection
 - obstruction by objects
 - interference

radio link types:

- ❖ **terrestrial microwave**
 - e.g. up to 45 Mbps channels
- ❖ **LAN (e.g., WiFi)**
 - 11Mbps, 54 Mbps
- ❖ **wide-area (e.g., cellular)**
 - 3G cellular: ~ few Mbps
- ❖ **satellite**
 - Kbps to 45Mbps channel (or multiple smaller channels)
 - 270 msec end-end delay
 - geosynchronous versus low altitude

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- packet switching, circuit switching, network structure

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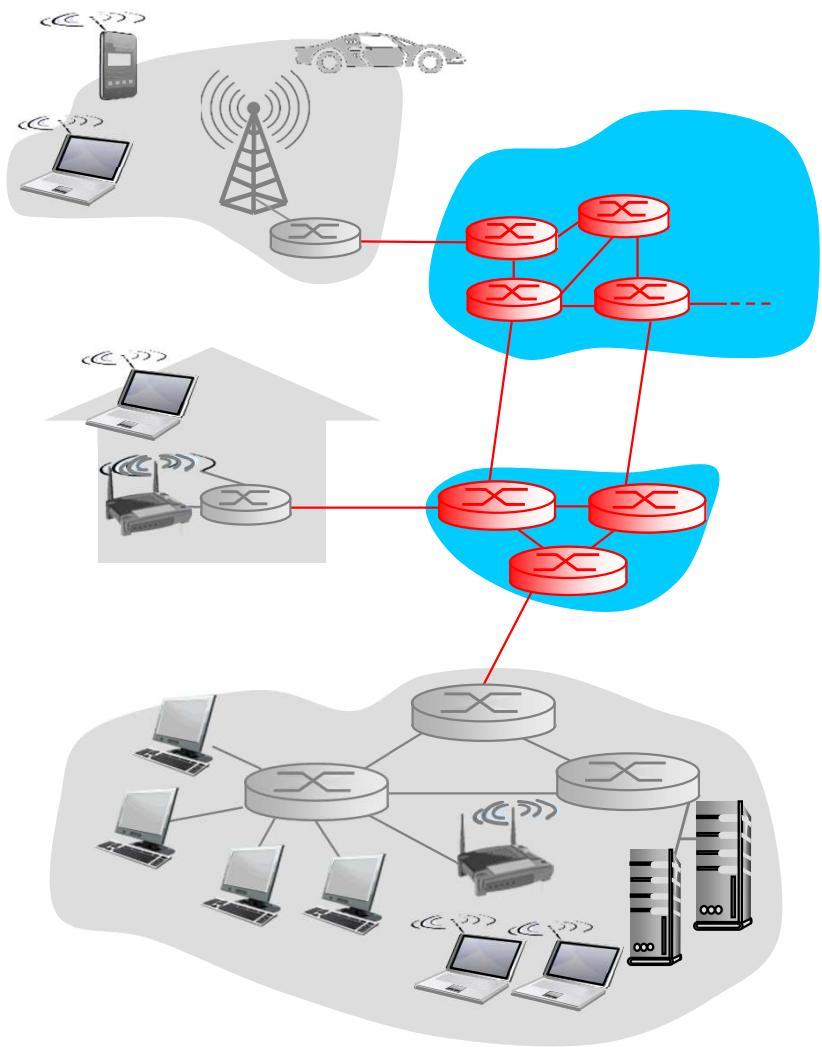
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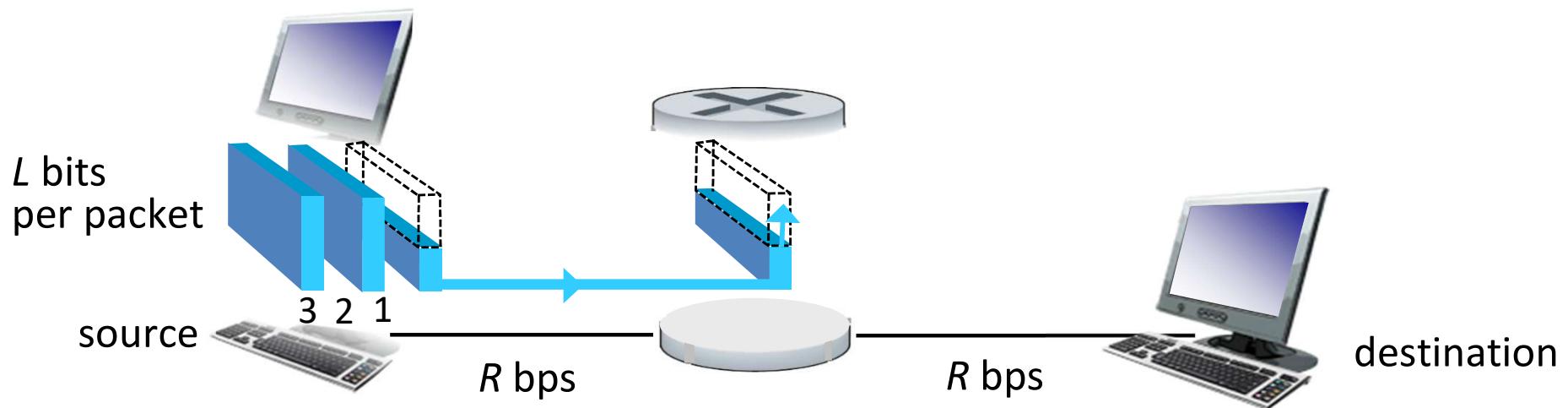
I.7 history

The network core

- ❖ mesh of interconnected routers
- ❖ **packet-switching: hosts break application-layer messages into packets**
 - forward packets from one router to the next, across links on path from source to destination
 - each packet transmitted at full link capacity



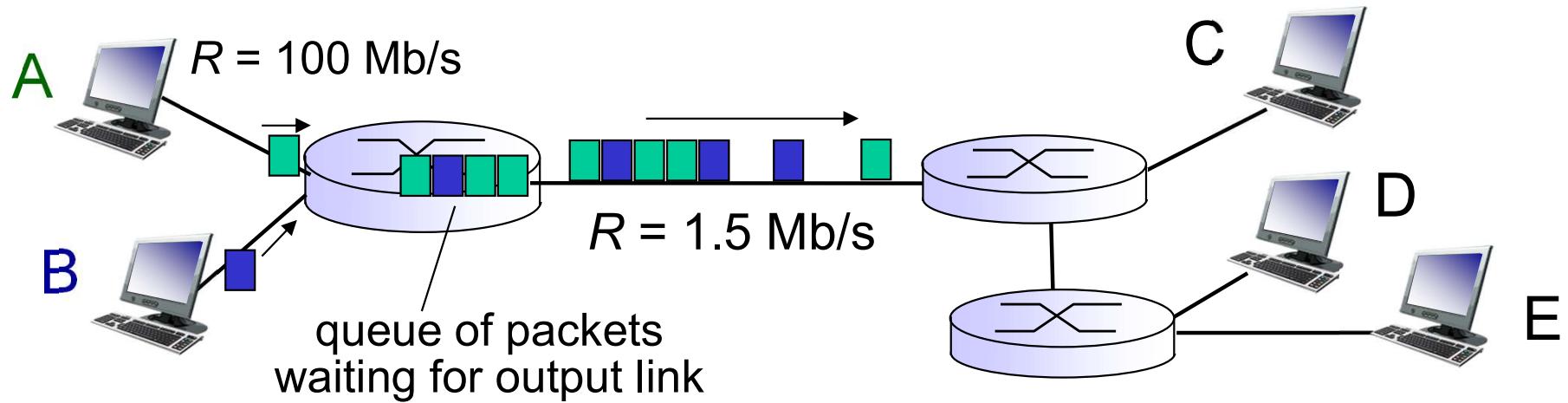
Packet-switching: store-and-forward



- ❖ takes L/R seconds to transmit (push out) L -bit packet into link at R bps
- ❖ *store and forward*: entire packet must arrive at router before it can be transmitted on next link
- ❖ end-end delay = $2L/R$ (assuming zero propagation delay)

- one-hop numerical example:*
- $L = 7.5$ Mbits
 - $R = 1.5$ Mbps
 - one-hop **transmission delay** = 5 sec
- } more on delay shortly ...

Packet Switching: queueing delay, loss



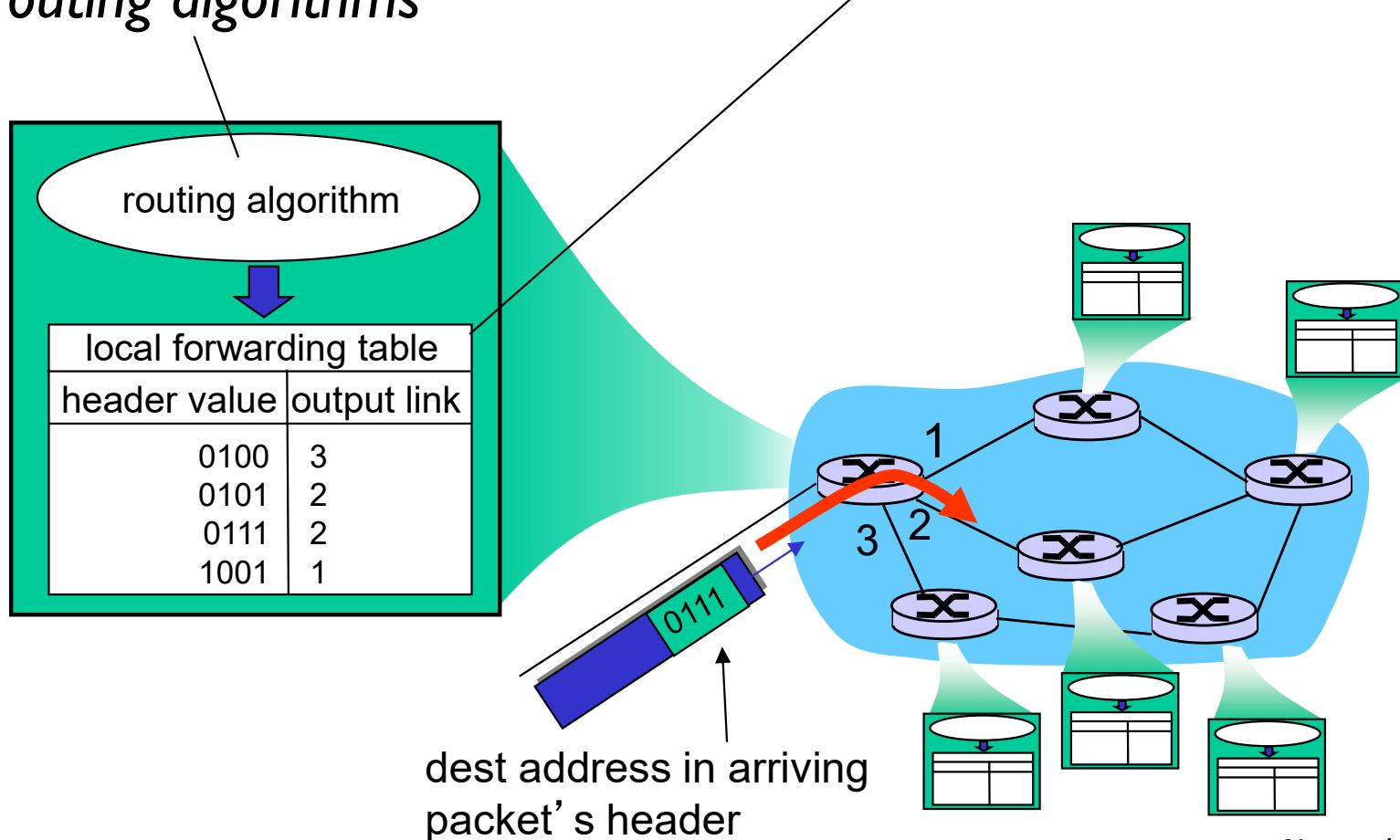
queuing and loss:

- ❖ If arrival rate (in bits) to link exceeds transmission rate of link for a period of time:
 - packets will queue, wait to be transmitted on link
 - packets can be dropped (lost) if memory (buffer) fills up

Two key network-core functions

routing: determines source-destination route taken by packets

- *routing algorithms*

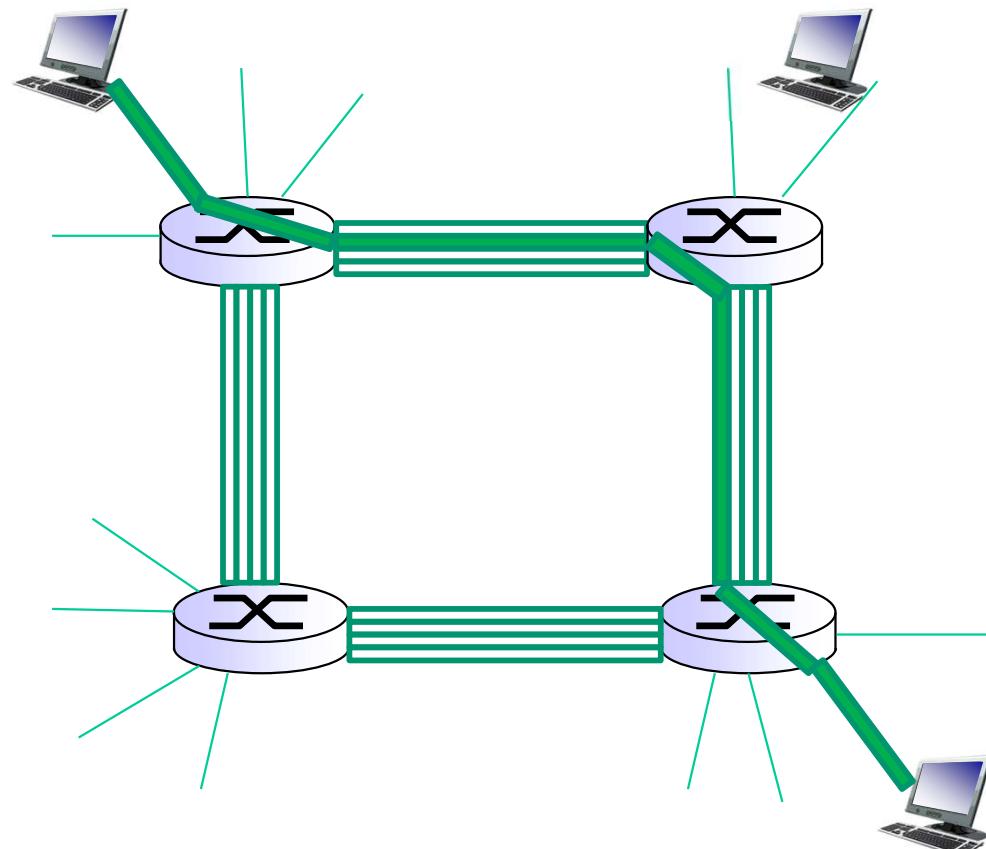


forwarding: move packets from router's input to appropriate router output

Alternative core: circuit switching

end-end resources allocated to, reserved for “call” between source & dest:

- ❖ In diagram, each link has four circuits.
 - call gets 2nd circuit in top link and 1st circuit in right link.
- ❖ dedicated resources: no sharing
 - circuit-like (guaranteed) performance
- ❖ circuit segment idle if not used by call (*no sharing*)
- ❖ Commonly used in traditional telephone networks



Exercise : Circuit switching

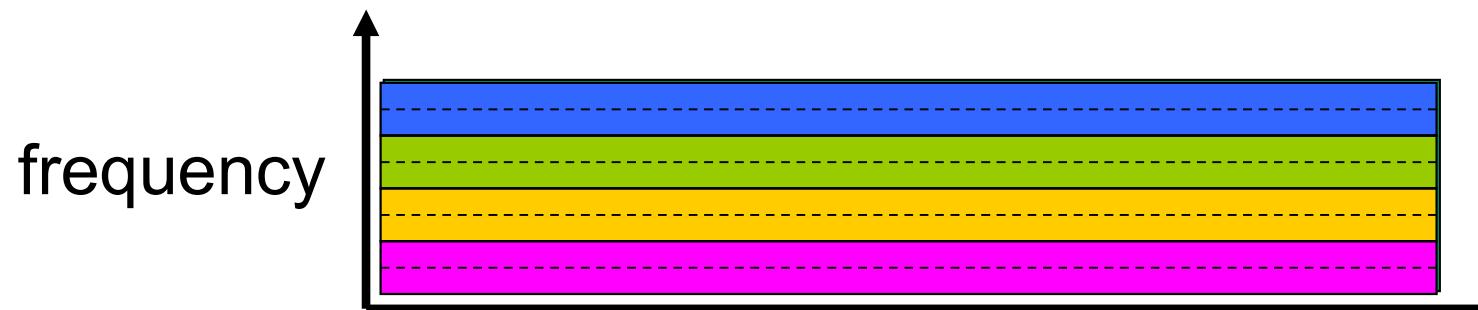
- ❖ จากรูปใน Slide ก่อนหน้า สมมติว่า Switch มีชื่อ A, B, C, D ตามเข็มนาฬิกา
- ❖ ให้ตอบคำถาม
 - มีกี่ Circuit ที่สามารถสร้างได้พร้อมกัน ให้ยกตัวอย่าง
 - สมมติว่าต้องการติดต่อจาก Switch A ไปยัง Switch C สามารถสร้างการติดต่อได้พร้อมกัน ได้กี่ Circuit
 - ถ้าต้องการสร้าง 4 Circuit จาก A -> C และอีก 4 จาก B -> D สามารถทำได้หรือไม่ อย่างไร

Circuit switching: old telephone



Circuit switching: FDM versus TDM

FDM

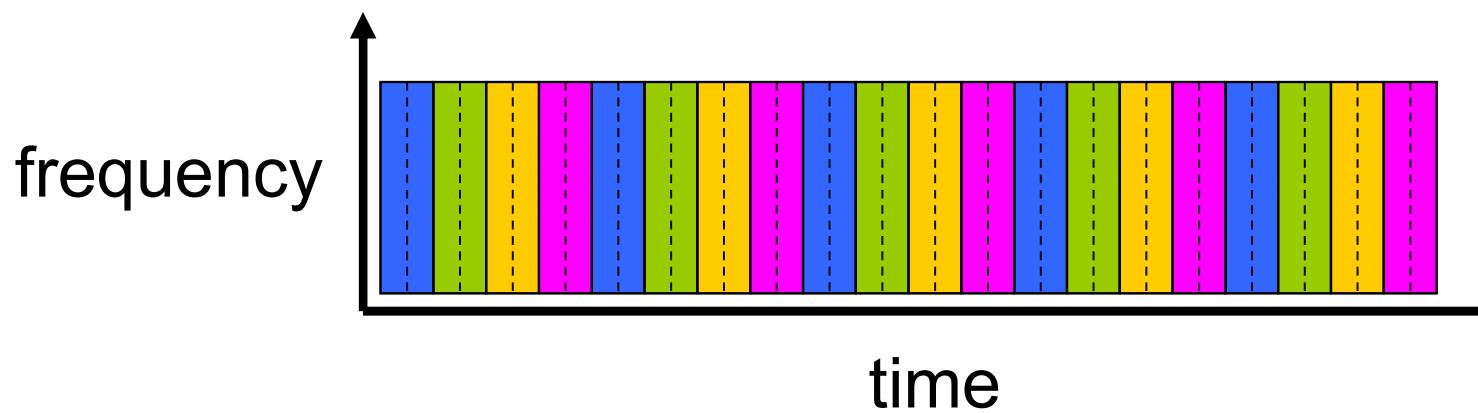


Example:

4 users



TDM

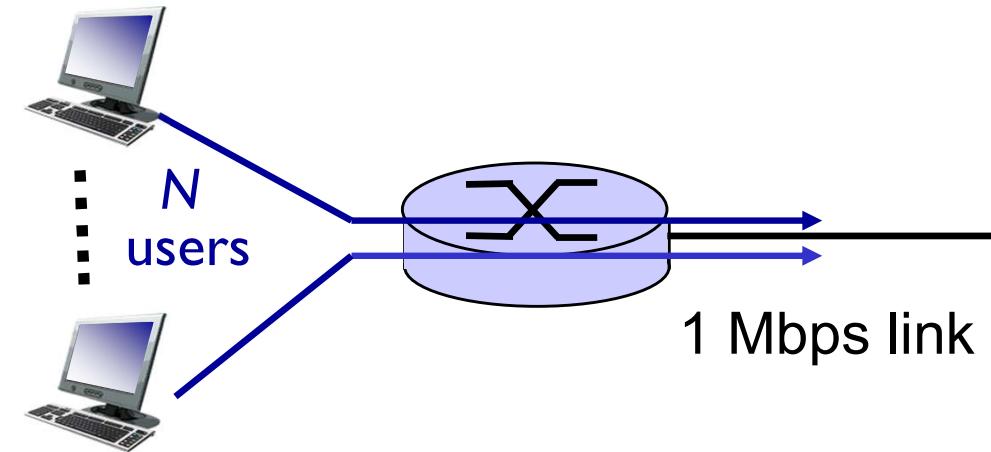


Packet switching versus circuit switching

packet switching allows more users to use network!

example:

- 1 Mb/s link
- each user:
 - 100 kb/s when “active”
 - active 10% of time



❖ *circuit-switching:*

- 10 users

❖ *packet switching:*

- with 35 users, probability > 10 active at same time is less than .0004 *

* Check out the online interactive exercises for more examples

Packet switching versus circuit switching

is packet switching a “slam dunk winner?”

- ❖ great for bursty data
 - resource sharing
 - simpler, no call setup
- ❖ **excessive congestion possible:** packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- ❖ **Q: How to provide circuit-like behavior?**
 - bandwidth guarantees needed for audio/video apps
 - still an unsolved problem (chapter 7)

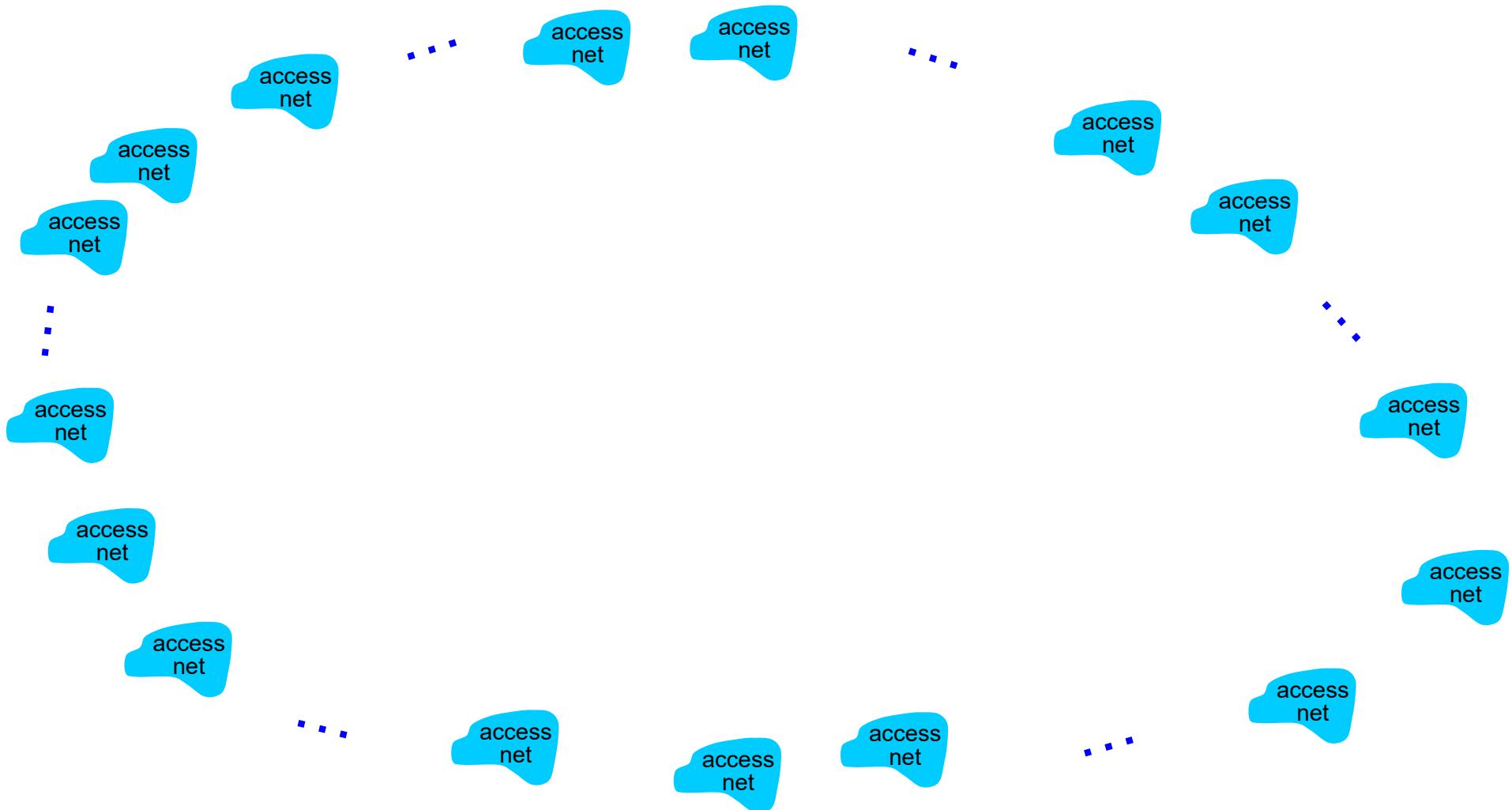
Q: human analogies of reserved resources (circuit switching)
versus on-demand allocation (packet-switching)?

Internet structure: network of networks

- ❖ End systems connect to Internet via **access ISPs** (Internet Service Providers)
 - Residential, company and university ISPs
- ❖ Access ISPs in turn must be interconnected.
 - ❖ So that any two hosts can send packets to each other
- ❖ Resulting network of networks is very complex
 - ❖ Evolution was driven by **economics** and **national policies**
- ❖ Let's take a stepwise approach to describe current Internet structure

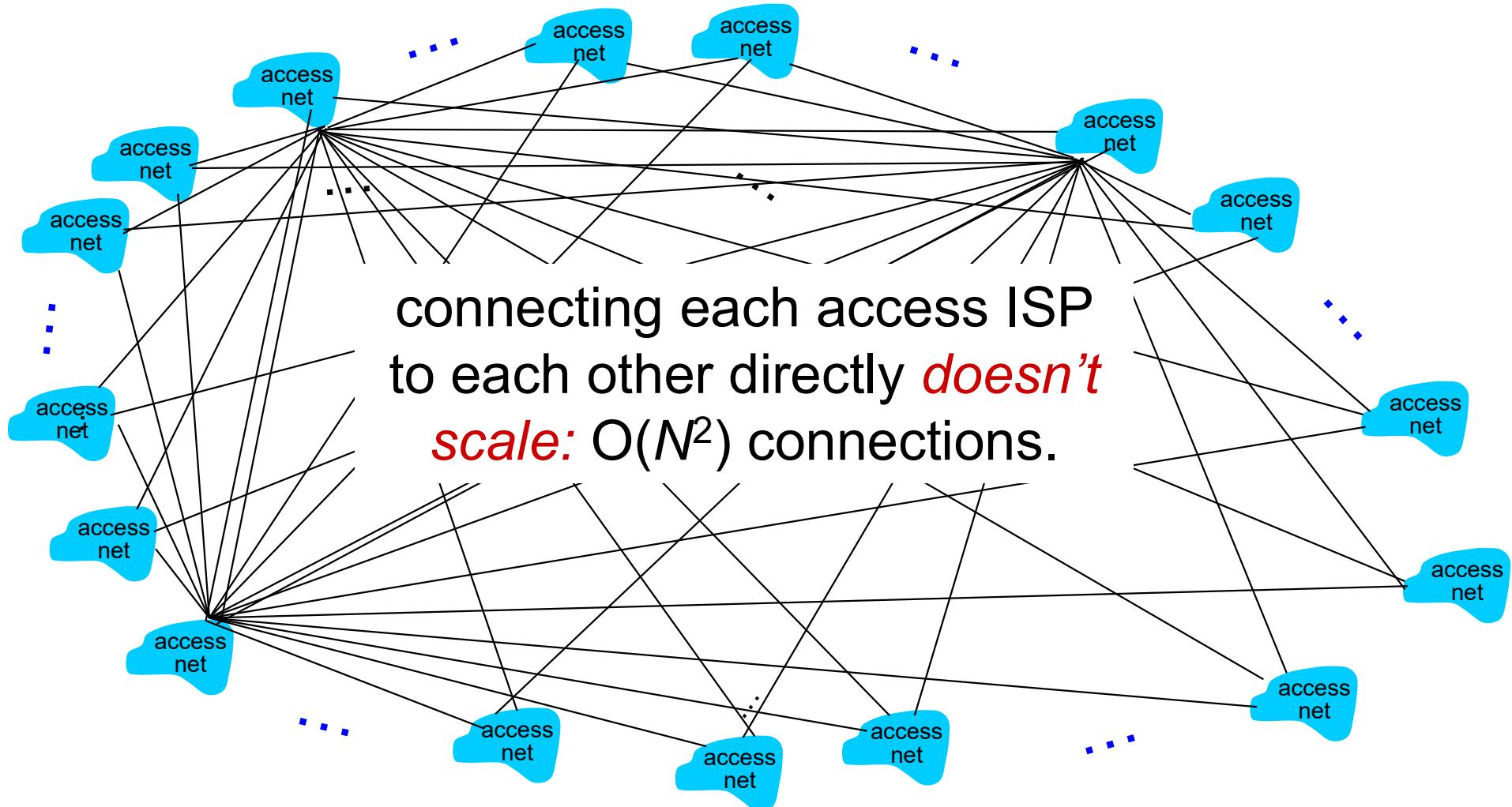
Internet structure: network of networks

Question: given *millions* of access ISPs, how to connect them together?



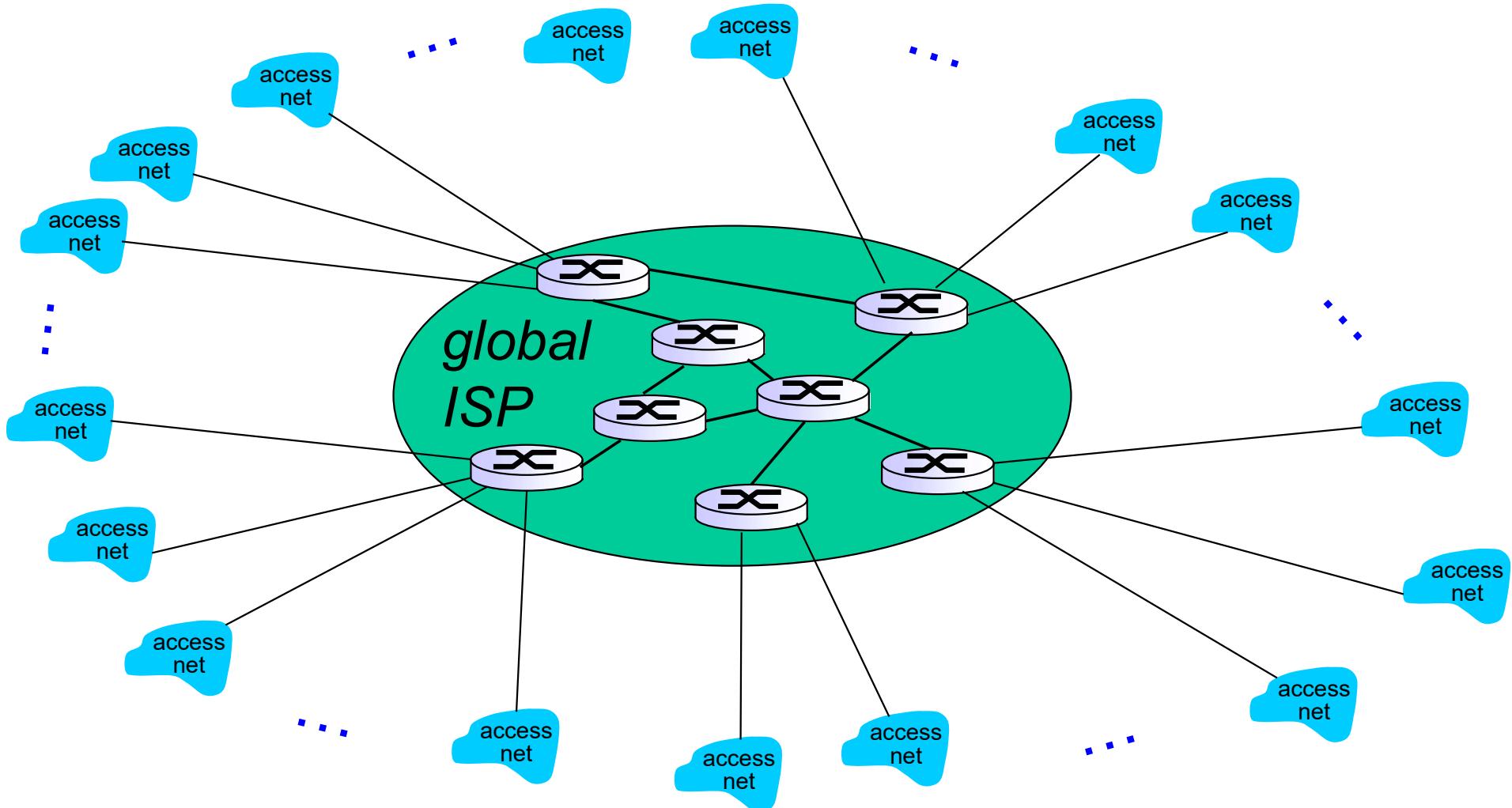
Internet structure: network of networks

Option: connect each access ISP to every other access ISP?



Internet structure: network of networks

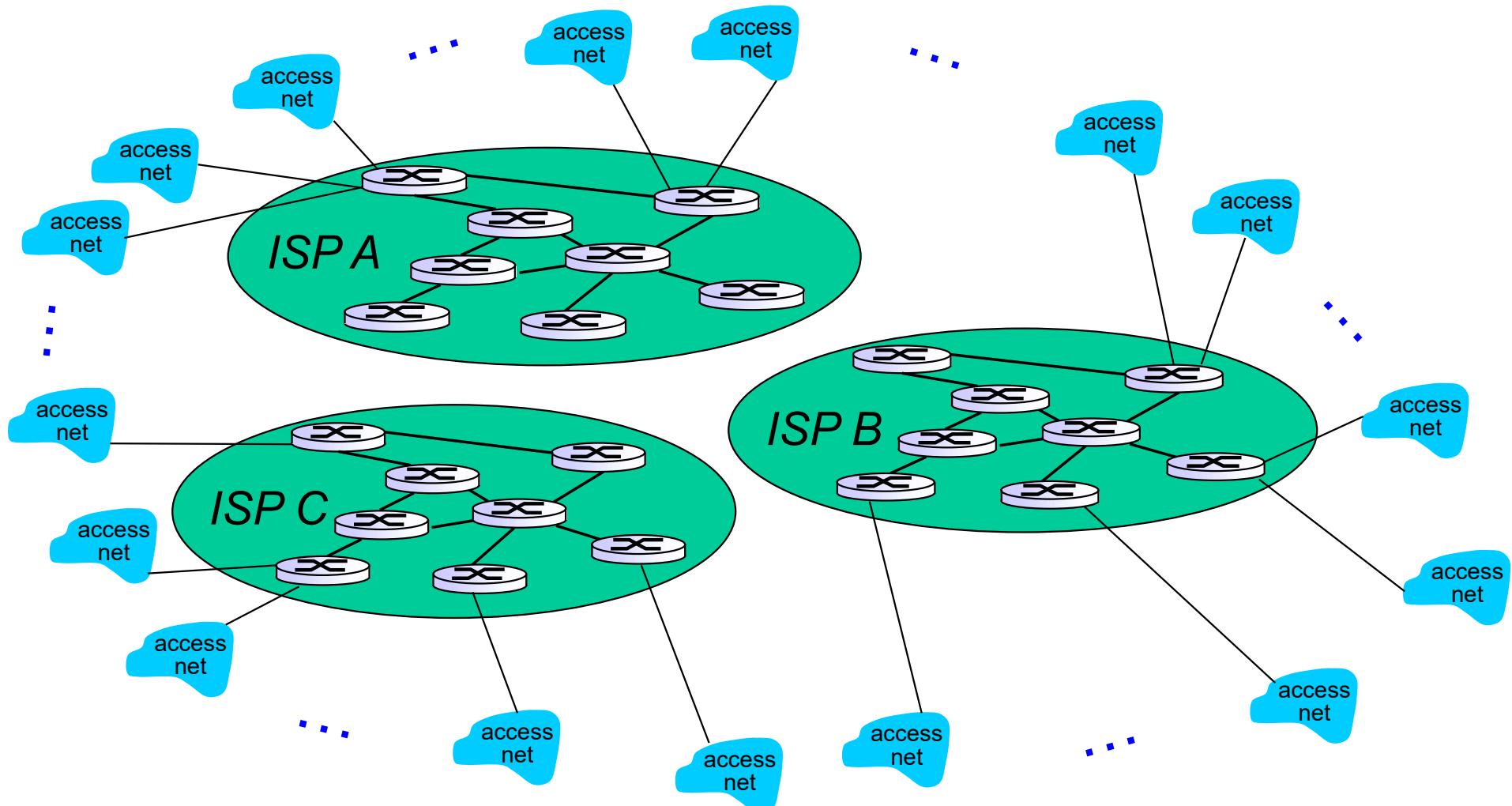
Option: connect each access ISP to a global transit ISP? Customer and provider ISPs have economic agreement.



Internet structure: network of networks

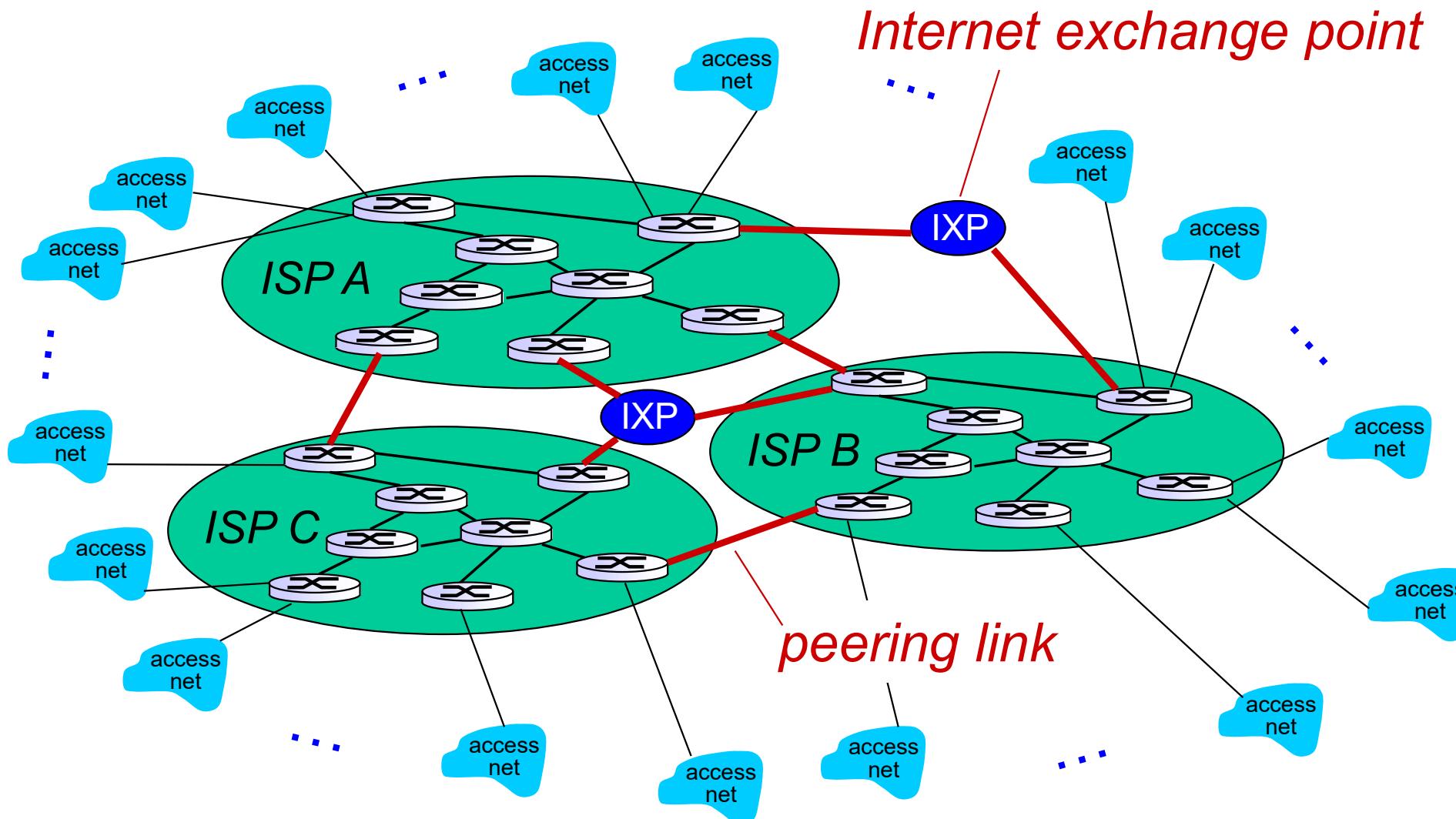
But if one global ISP is viable business, there will be competitors

....



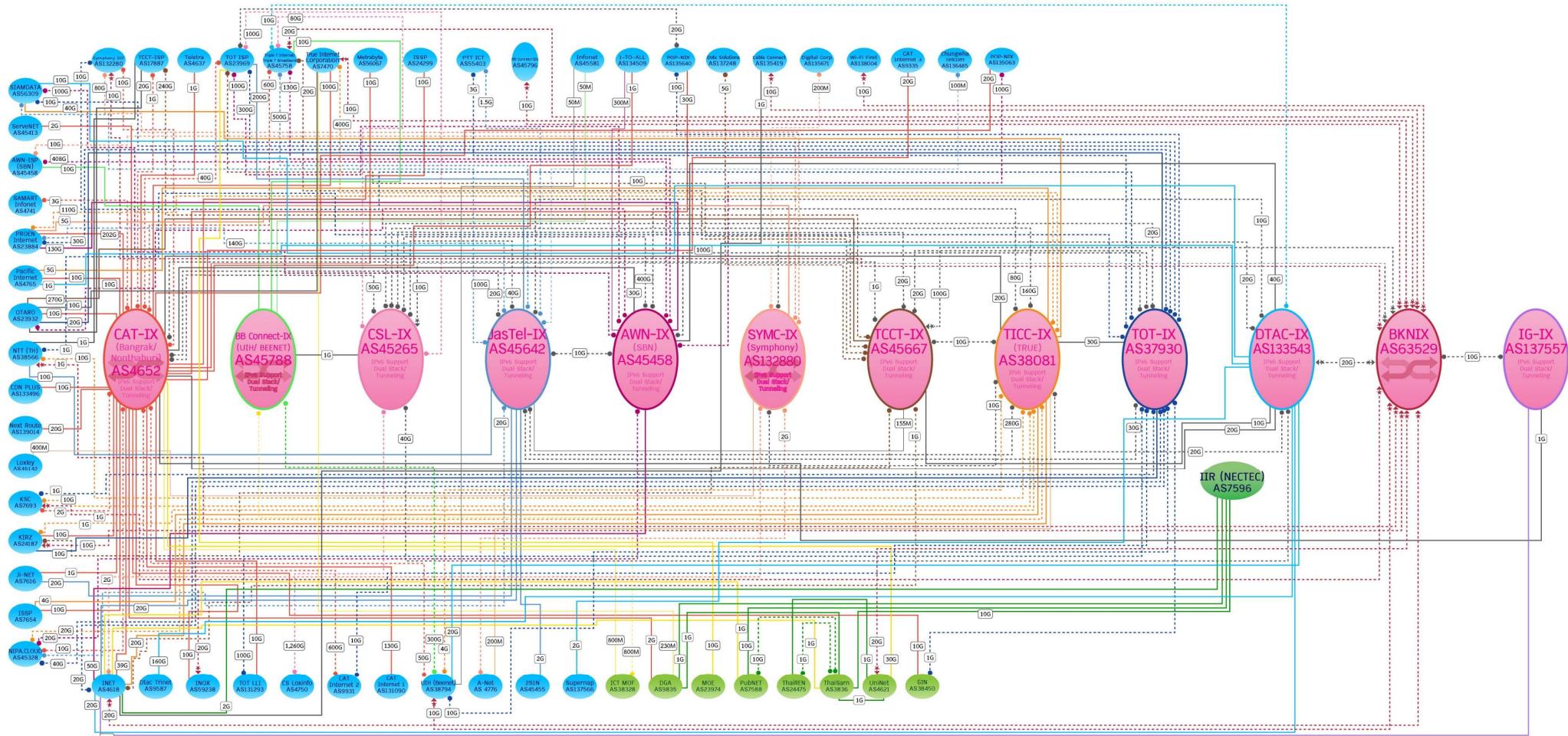
Internet structure: network of networks

But if one global ISP is viable business, there will be competitors
.... which must be interconnected



THAILAND DOMESTIC INTERNET EXCHANGE

แผนภาพการเชื่อมต่อเครือข่ายอินเทอร์เน็ตภายในประเทศไทย



Legend

Thailand Internet Exchange : IX	Layer 2 Exchange Point
Commercial Internet Provider	CAT Internet Exchange
Government/Research/Academic	BKNIX
	BB Connect Internet Exchange
	CS Lookin Internet Exchange
	ICT Internet Exchange
	TICC Internet Exchange
	ISP to ISP/Internet Exchange to Internet Exchange
	Government/Research/Academic
	TOT Internet Exchange
	Tot-IX Internet Exchange
	Icru Internet Exchange
	JasTel Internet Exchange
	AWN Internet Exchange
	SYMC Internet Exchange
	TCCT Internet Exchange
	TOT-ALL Internet Exchange
	Pop-IX Internet Exchange
	Digi Connect Internet Exchange
	Digital Corp Internet Exchange
	Wi-Fi First Internet Exchange
	CAT Internet Exchange
	Chiangmai Telecom Internet Exchange
	Chonburi Internet Exchange
	IIR (NECTEC) Internet Exchange
	Peering/ IPv6 Native
	Internet Protocol Version 6 Native
	Internet Protocol Version 6 Tunneling
	Internet Protocol Version 6 Dual Stack

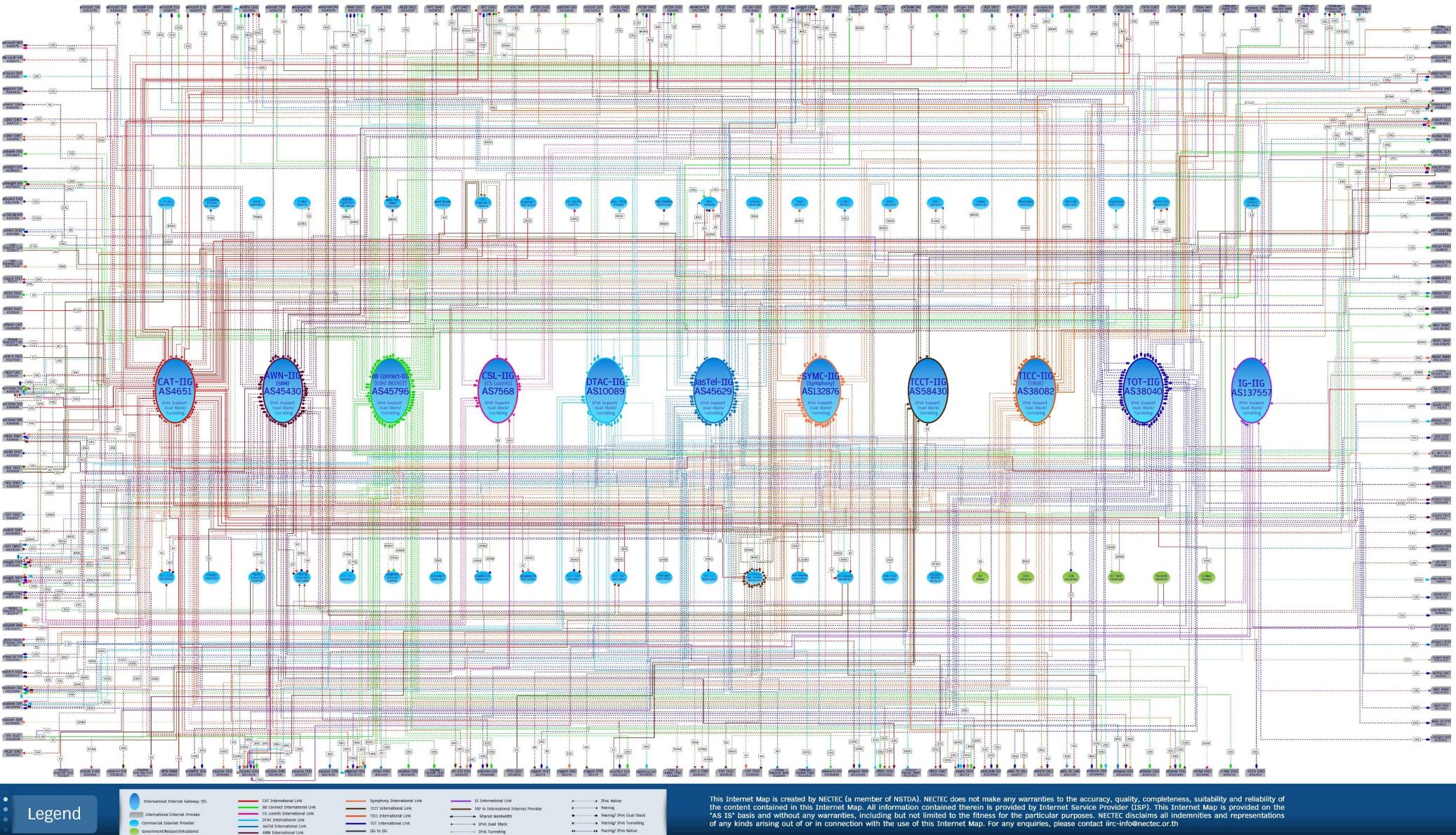
This Internet Map is created by NECTEC (a member of NSTDA). NECTEC does not make any warranties to the accuracy, quality, completeness, suitability and reliability of the content contained in this Internet Map. All information contained therein is provided by Internet Service Provider (ISP). This Internet Map is provided on the "AS IS" basis and without any warranties, including but not limited to the fitness for the particular purposes. NECTEC disclaims all indemnities and representations of any kinds arising out of or in connection with the use of this Internet Map. For any enquiries, please contact iirc-info@nectec.or.th

THAILAND INTERNATIONAL INTERNET GATEWAY

แผนภาพการเชื่อมต่อเครือข่ายอินเทอร์เน็ตระหว่างประเทศ

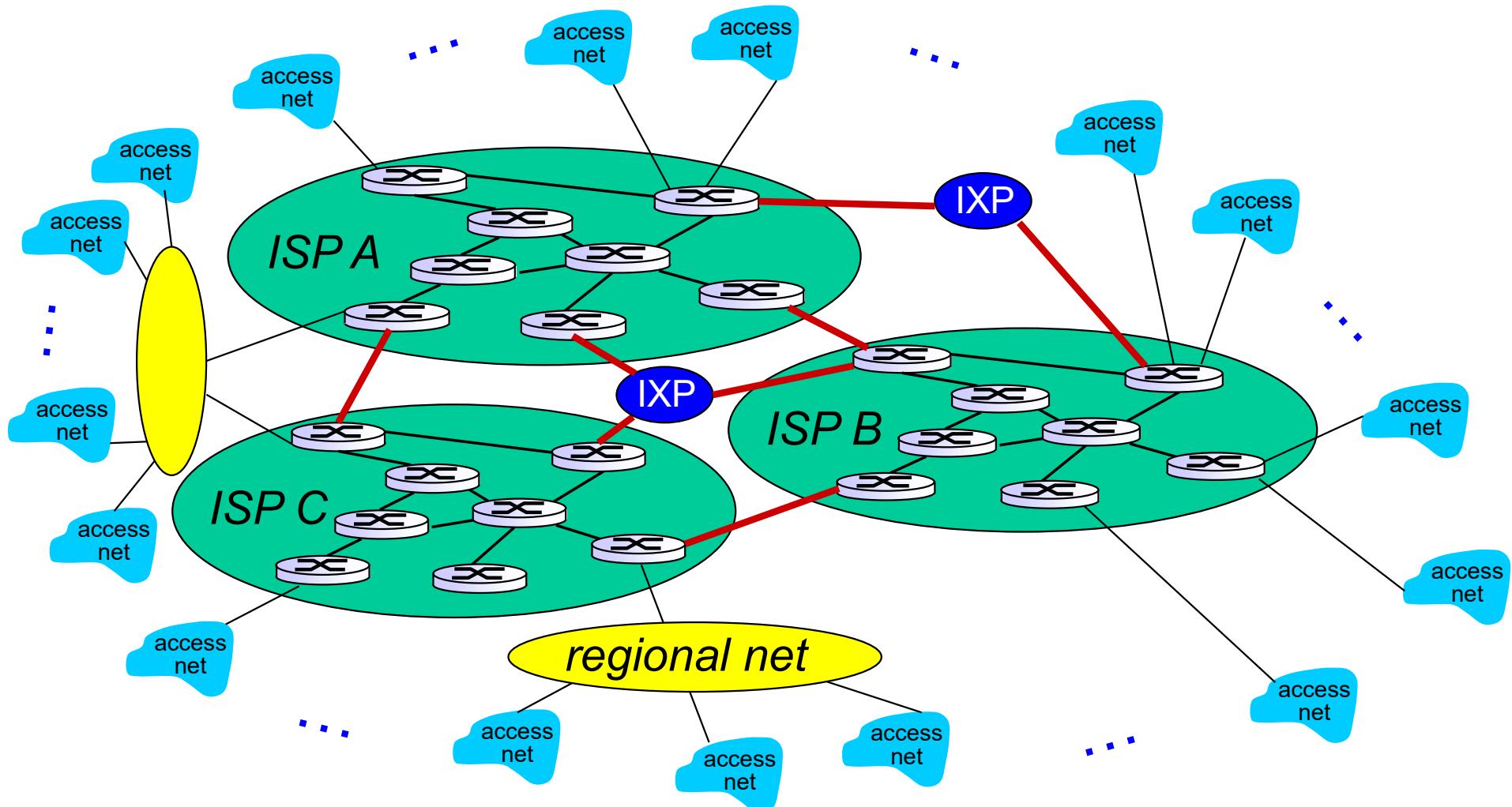


Total International Bandwidth
14,041.148 Gbps.
Last update: 07/12/2020



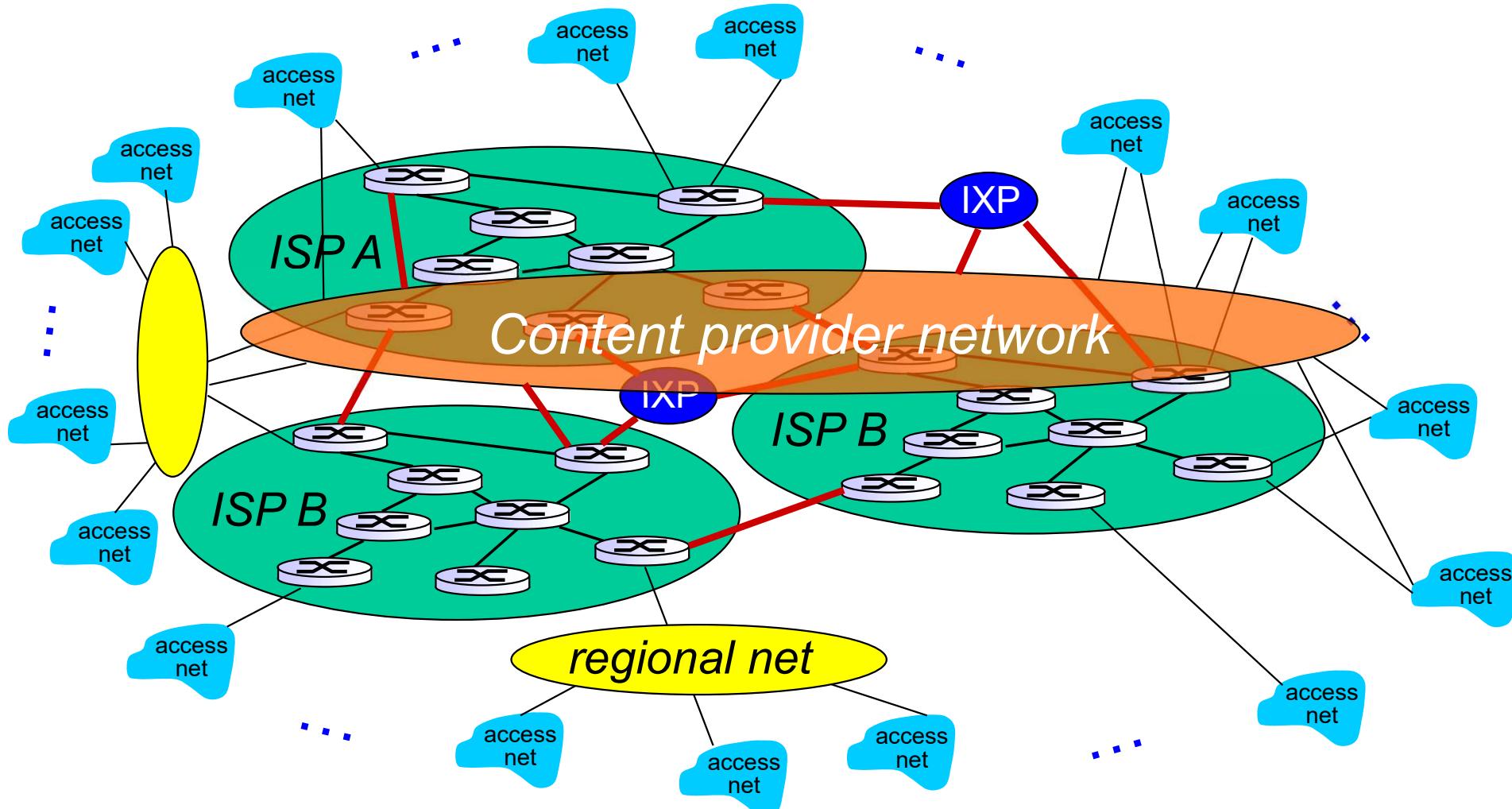
Internet structure: network of networks

... and regional networks may arise to connect access nets to ISPs

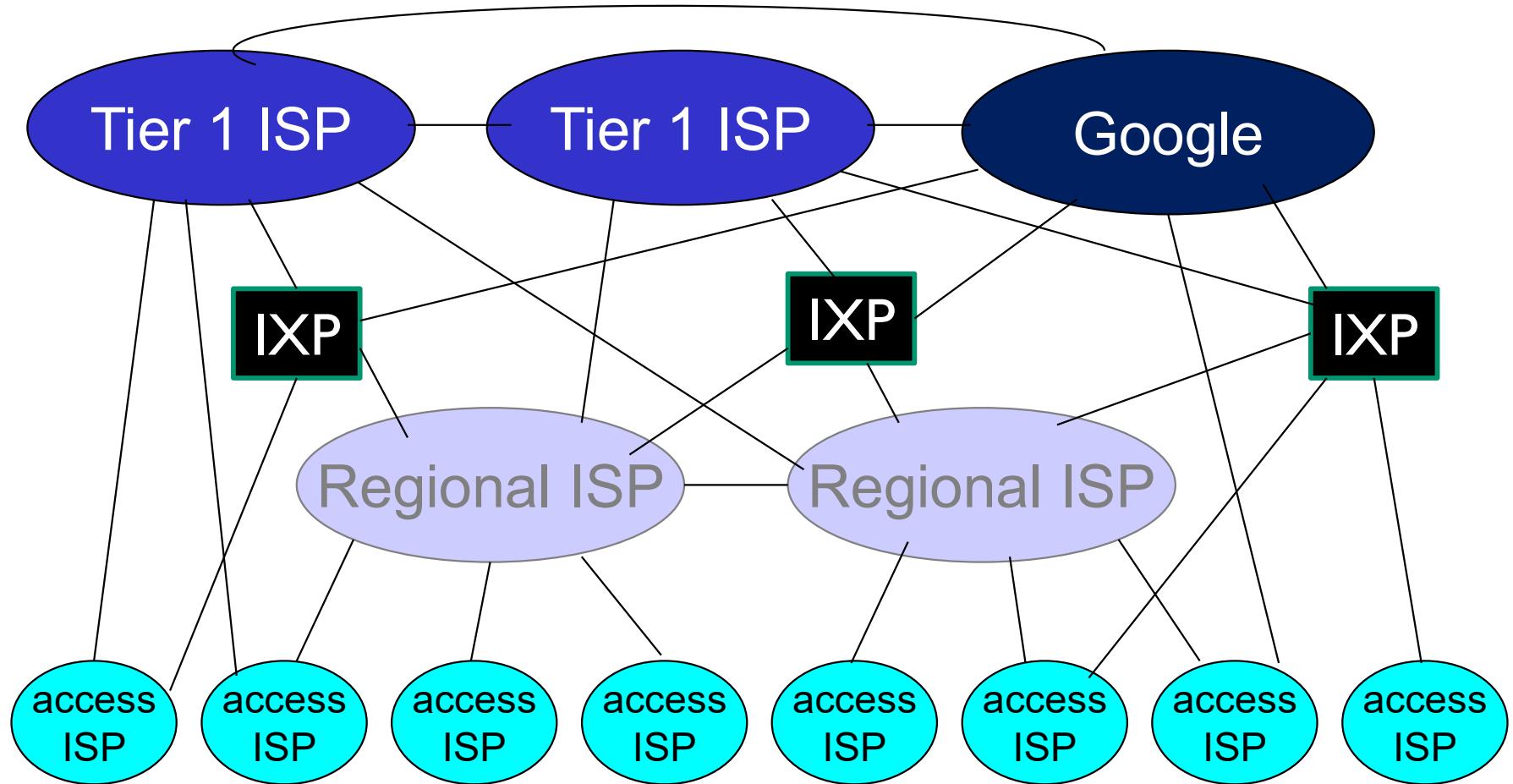


Internet structure: network of networks

... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users

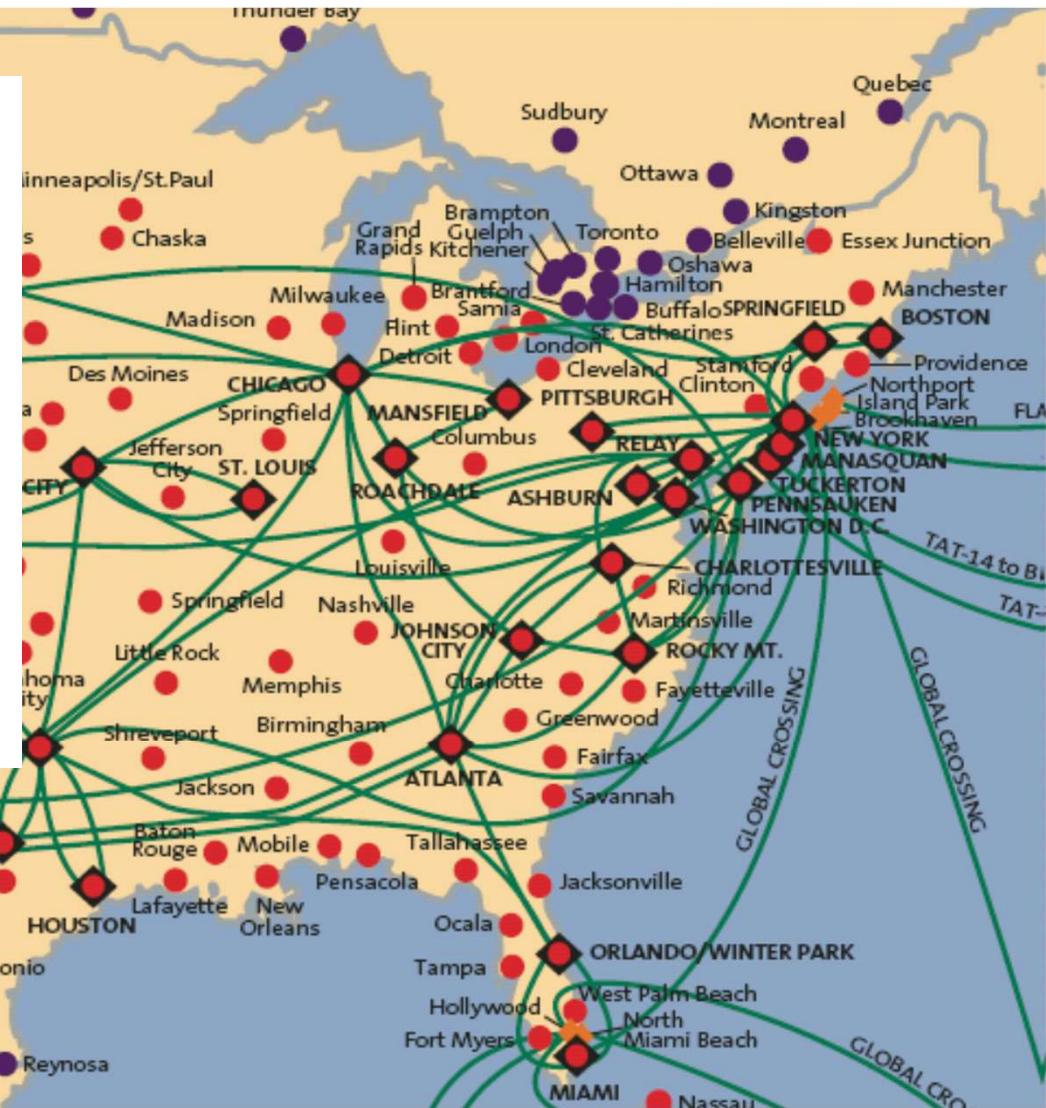
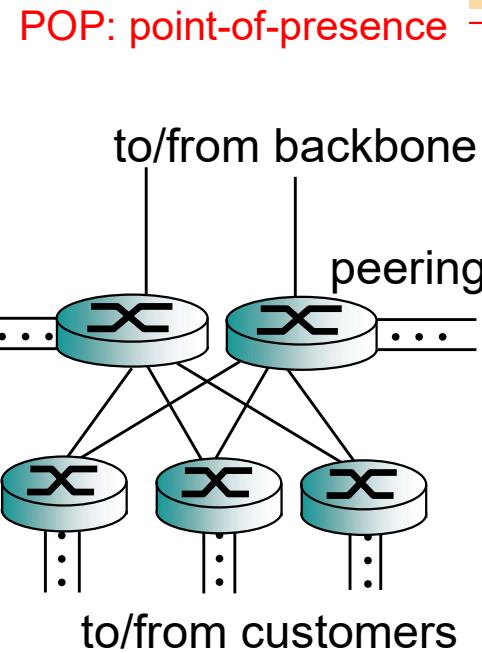


Internet structure: network of networks



- ❖ at center: small # of well-connected large networks
 - “tier-1” commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
 - **content provider network** (e.g, Google): private network that connects its data centers to Internet, often bypassing tier-1, regional ISPs

Tier-I ISP: e.g., Sprint



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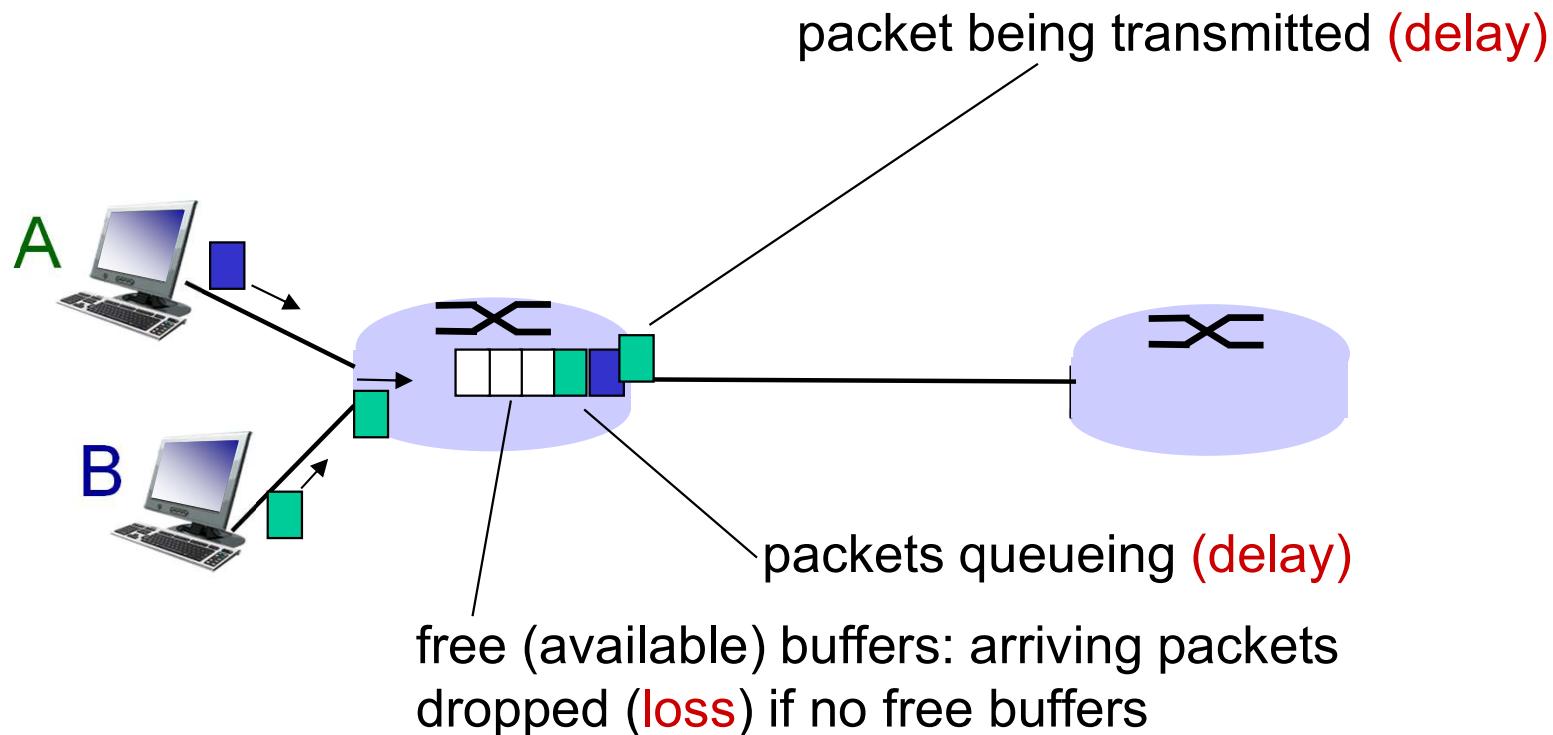
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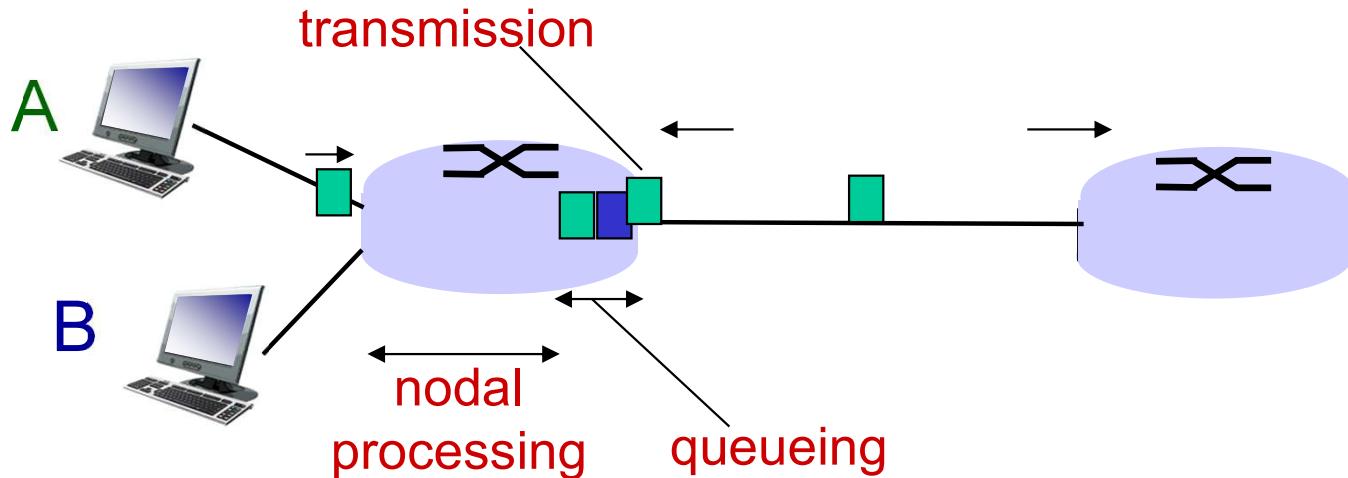
How do loss and delay occur?

packets queue in router buffers

- ❖ packet arrival rate to link (temporarily) exceeds output link capacity
- ❖ packets queue, wait for turn



Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

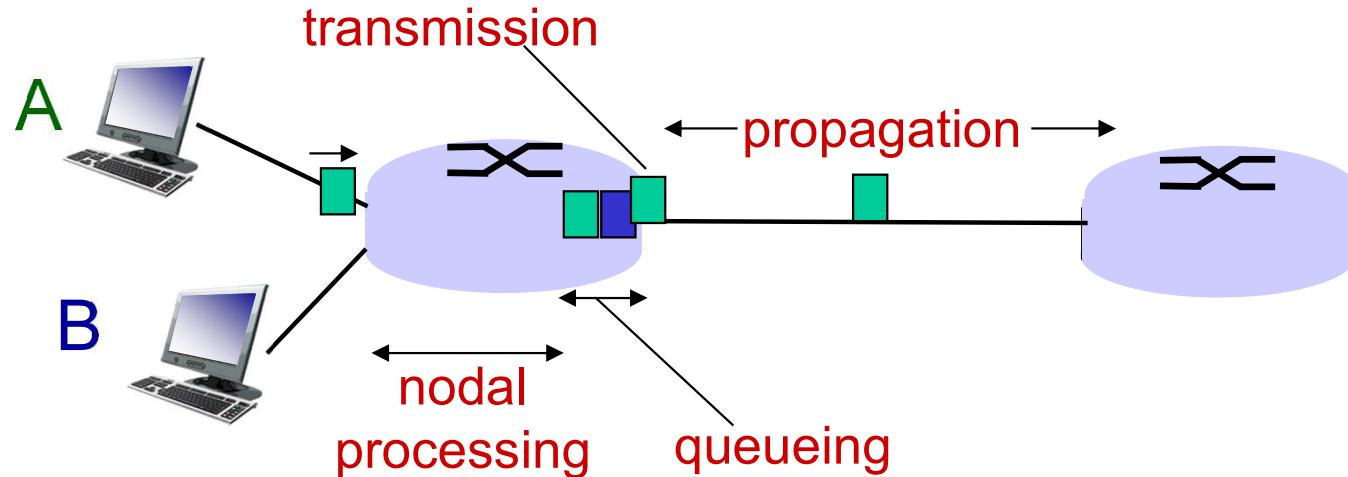
d_{proc} : nodal processing

- check bit errors
- determine output link
- typically < msec

d_{queue} : queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

d_{trans} : transmission delay:

- L : packet length (bits)
- R : link bandwidth (bps)
- $d_{\text{trans}} = L/R$

d_{trans} and d_{prop}
very different

d_{prop} : propagation delay:

- d : length of physical link
- s : propagation speed in medium ($\sim 2 \times 10^8 \text{ m/sec}$)
- $d_{\text{prop}} = d/s$

* Check out the Java applet for an interactive animation on trans vs. prop delay

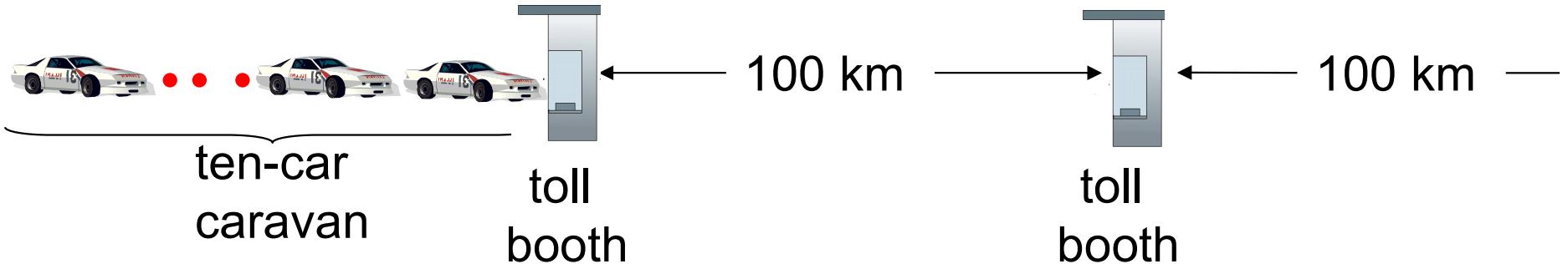
Exercise : packet delay |

- ❖ เข้าไปที่เว็บ

https://media.pearsoncmg.com/aw/ecs_kurose_computer-network_7/cw/content/interactiveanimations/transmission-vs-propogation-delay/transmission-propagation-delay-ch1/index.html

- ❖ ให้ปรับค่า rate, propagation delay, packet size เพื่อสร้างสถานการณ์ที่ผู้ส่งสามารถส่งข้อมูลได้หมด ก่อนที่บิตแรกของข้อมูลจะไปถึงผู้รับ มีปัจจัยใดบ้าง อธิบาย

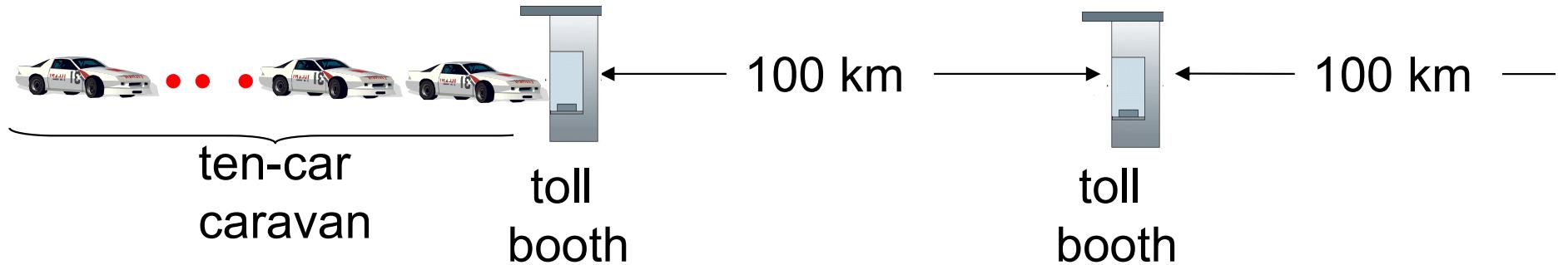
Caravan analogy



- ❖ cars “propagate” at 100 km/hr
- ❖ toll booth takes 12 sec to service car (bit transmission time)
- ❖ car~bit; caravan ~ packet
- ❖ **Q: How long until caravan is lined up before 2nd toll booth?**

- time to “push” entire caravan through toll booth onto highway = $12*10 = 120$ sec
- time for last car to propagate from 1st to 2nd toll both:
 $100\text{km}/(100\text{km/hr}) = 1\text{ hr}$
- **A: 62 minutes**

Caravan analogy (more)



- ❖ suppose cars now “propagate” at 1000 km/hr
- ❖ and suppose toll booth now takes one min to service a car
- ❖ **Q: Will cars arrive to 2nd booth before all cars serviced at first booth?**
 - **A: Yes!** after 7 min, 1st car arrives at second booth; three cars still at 1st booth.

Exercise : packet delay 2

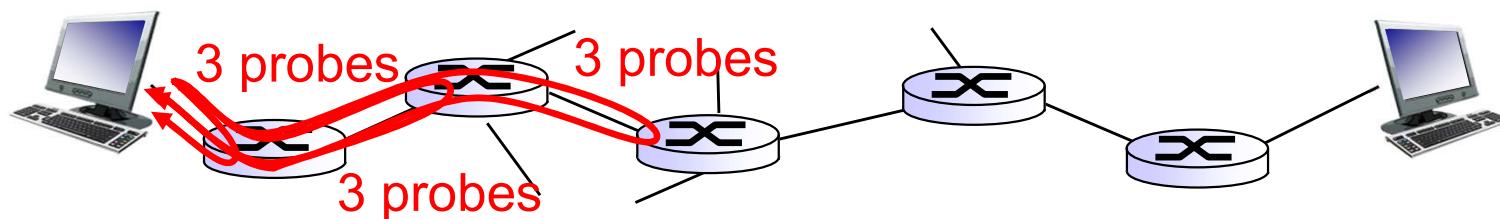
- ❖ กำหนดข้อมูล Packet หนึ่ง ส่งจาก Source Host ไปยัง Destination Host ให้บอก delay ต่างๆ ที่เกิดขึ้นในจุดต่างๆ จากปลายถึงปลาย (end to end delay) และให้บอกว่า delay ได้เป็นแบบคงที่ หรือ เปลี่ยนแปลง

Exercise : packet delay 3

- ❖ กำหนดให้ Link หนึ่งความยาว 2,500 กม. มี propagation speed = 2.5×10^8 m/s และ transmission rate 2 Mbps กำหนดให้ข้อมูลยาว 1000 ไบต์ ให้หาว่า
 - ใช้เวลาเท่าไรข้อมูลทั้ง 1000 ไบต์จึงส่งผ่าน Link ข้างต้นได้หมด
 - ค่า delay ขึ้นกับความยาว Packet หรือไม่
 - ค่า delay ขึ้นกับ transmission rate หรือไม่

“Real” Internet delays and routes

- ❖ what do “real” Internet delay & loss look like?
- ❖ `traceroute` program: provides delay measurement from source to router along end-end Internet path towards destination. For all i :
 - sends three packets that will reach router i on path towards destination
 - router i will return packets to sender
 - sender times interval between transmission and reply.



“Real” Internet delays, routes

traceroute: gaia.cs.umass.edu to www.eurecom.fr

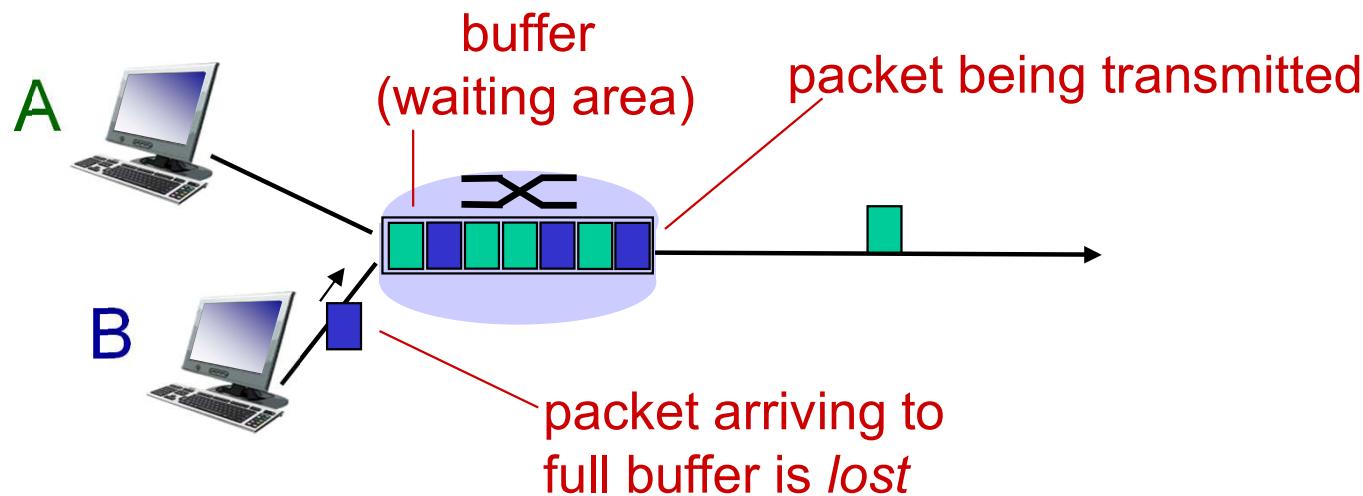
3 delay measurements from
gaia.cs.umass.edu to cs-gw.cs.umass.edu

1	cs-gw (128.119.240.254)	1 ms	1 ms	2 ms	
2	border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145)	1 ms	1 ms	2 ms	
3	cht-vbns.gw.umass.edu (128.119.3.130)	6 ms	5 ms	5 ms	
4	jn1-at1-0-0-19.wor.vbns.net (204.147.132.129)	16 ms	11 ms	13 ms	
5	jn1-so7-0-0-0.wae.vbns.net (204.147.136.136)	21 ms	18 ms	18 ms	
6	abilene-vbns.abilene.ucaid.edu (198.32.11.9)	22 ms	18 ms	22 ms	
7	nycm-wash.abilene.ucaid.edu (198.32.8.46)	22 ms	22 ms	22 ms	
8	62.40.103.253 (62.40.103.253)	104 ms	109 ms	106 ms	trans-oceanic link
9	de2-1.de1.de.geant.net (62.40.96.129)	109 ms	102 ms	104 ms	
10	de.fr1.fr.geant.net (62.40.96.50)	113 ms	121 ms	114 ms	
11	renater-gw.fr1.fr.geant.net (62.40.103.54)	112 ms	114 ms	112 ms	
12	nio-n2.cssi.renater.fr (193.51.206.13)	111 ms	114 ms	116 ms	
13	nice.cssi.renater.fr (195.220.98.102)	123 ms	125 ms	124 ms	
14	r3t2-nice.cssi.renater.fr (195.220.98.110)	126 ms	126 ms	124 ms	
15	eurecom-valbonne.r3t2.ft.net (193.48.50.54)	135 ms	128 ms	133 ms	
16	194.214.211.25 (194.214.211.25)	126 ms	128 ms	126 ms	
17	***				
18	***	*	means no response (probe lost, router not replying)		
19	fantasia.eurecom.fr (193.55.113.142)	132 ms	128 ms	136 ms	

* Do some traceroutes from exotic countries at www.traceroute.org

Packet loss

- ❖ queue (aka buffer) preceding link in buffer has finite capacity
- ❖ packet arriving to full queue dropped (aka lost)
- ❖ lost packet may be retransmitted by previous node, by source end system, or not at all



* Check out the Java applet for an interactive animation on queuing and loss

Exercise : packet loss

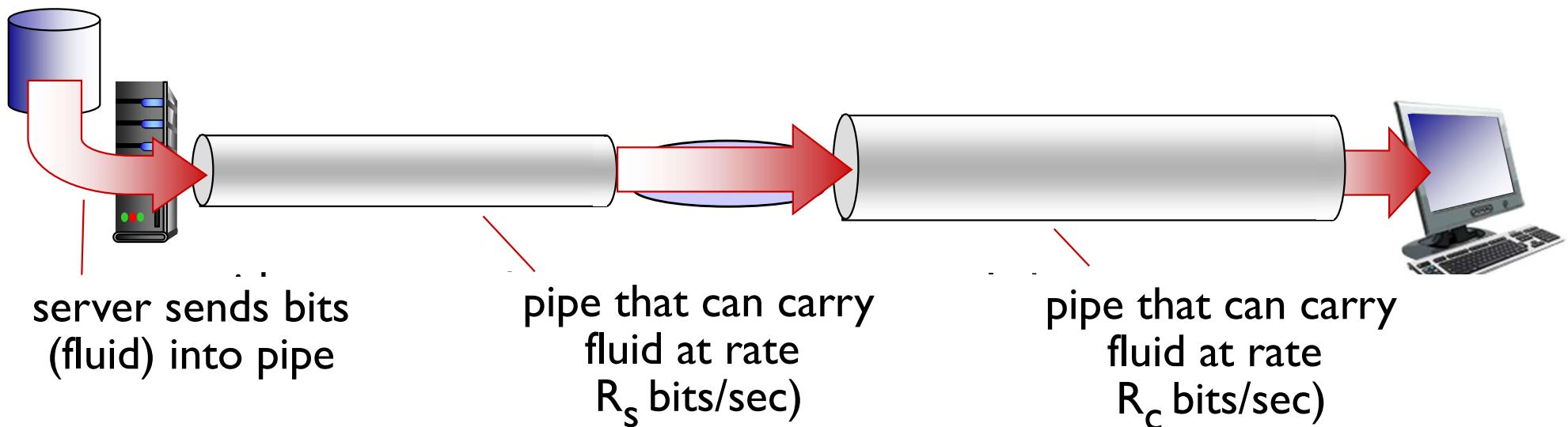
- ❖ เข้าไปที่เว็บ

https://media.pearsoncmg.com/aw/ecs_kurose_computer-network_7/cw/content/interactiveanimations/queuing-loss-applet/index.html

- ❖ เมื่อใช้ค่า Emission rate ที่มากที่สุด และ Transmission Rate ที่น้อยที่สุด เป็นเวลานานเท่าไรที่ Packet จะถูก Drop
- ❖ อธิบายว่า Packet Drop เกิดจากอะไร

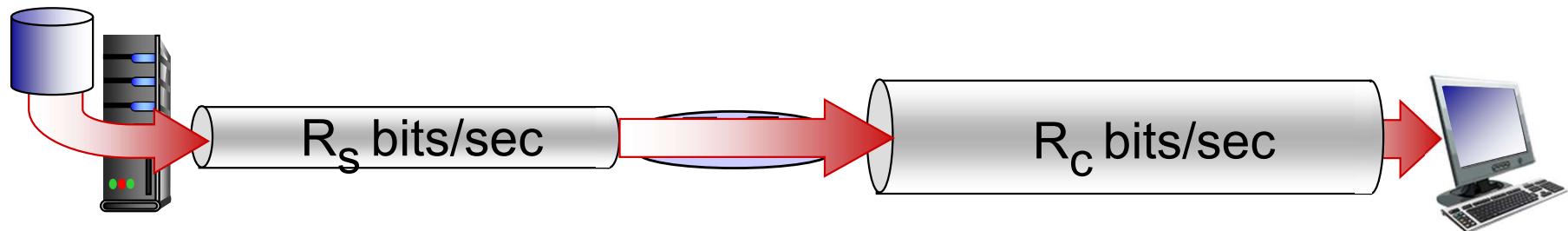
Throughput

- ❖ **throughput:** rate (bits/time unit) at which bits transferred between sender/receiver
 - *instantaneous:* rate at given point in time
 - *average:* rate over longer period of time

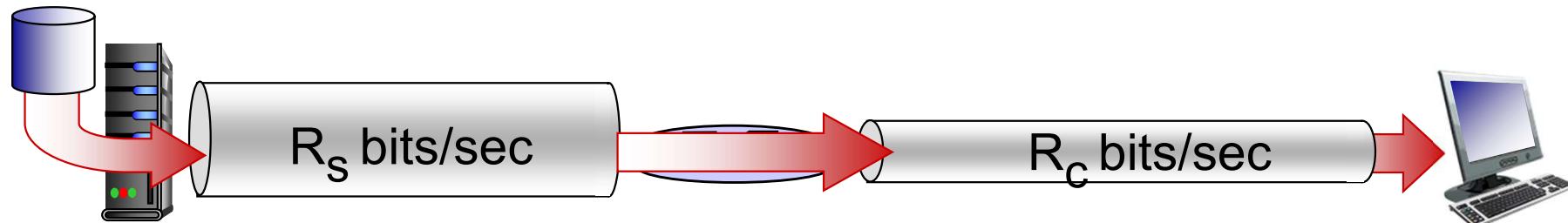


Throughput (more)

- ❖ $R_s < R_c$ What is average end-end throughput?



- ❖ $R_s > R_c$ What is average end-end throughput?



bottleneck link

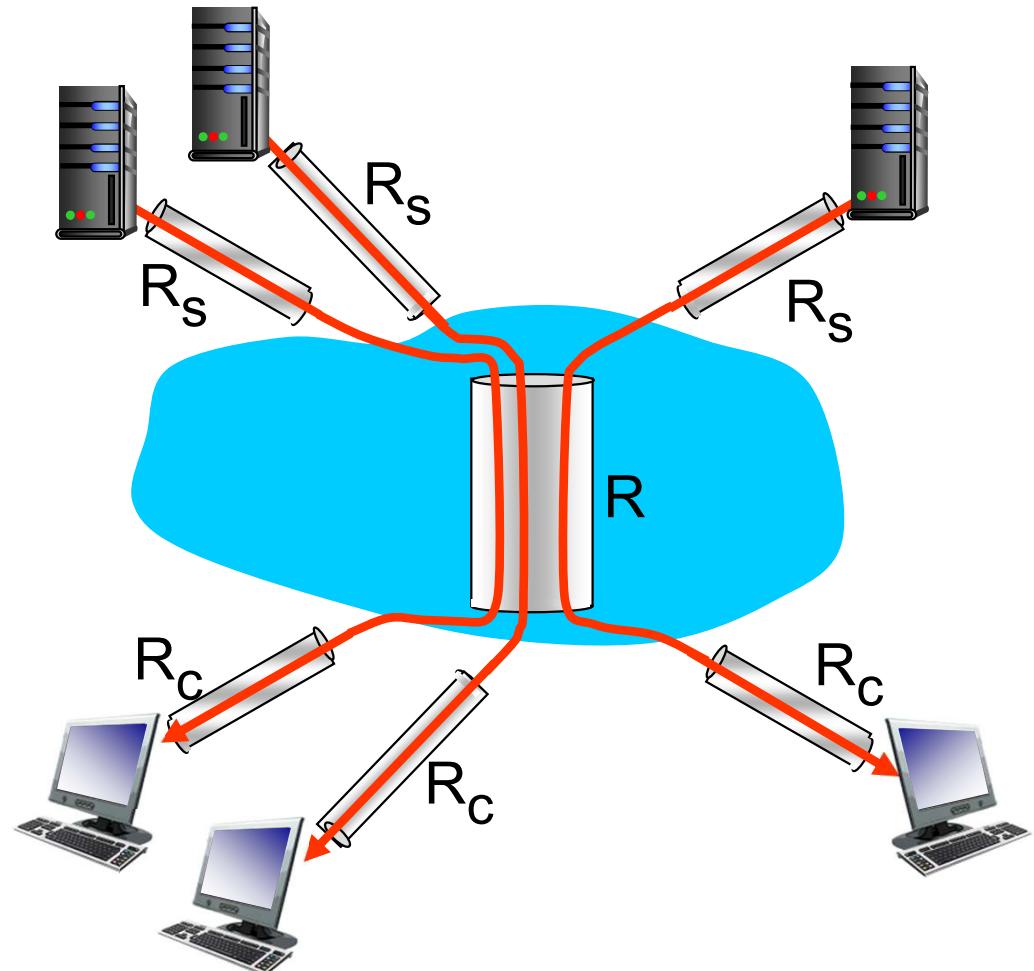
link on end-end path that constrains end-end throughput

Exercise : Throughput

- ❖ กำหนดให้ Host A ต้องการส่งไฟล์ไปยัง Host B โดยมี 3 เส้นทางที่ความเร็วต่างกัน คือ 500 kbps, 2 Mbps, 1 Mbps
- ❖ ให้หา Throughput ของการส่งไฟล์ โดยกำหนดให้ไฟล์ขนาด 4 MB และหาว่าจะใช้เวลาเท่าไรจึงจะส่งถึง Host B ในแต่ละเส้นทาง

Throughput: Internet scenario

- ❖ per-connection end-end throughput:
 $\min(R_c, R_s, R/10)$
- ❖ in practice: R_c or R_s is often bottleneck



10 connections (fairly) share
backbone bottleneck link R bits/sec

Exercise : message segment

- ❖ เข้าไปที่เว็บ

https://media.pearsoncmg.com/aw/ecs_kurose_compenetwork_7/cw/content/interactiveanimations/message-segmentation/index.html

- ❖ ทดลองกำหนดค่า Package Size, Message Size และ อธิบายการทำงาน

Chapter I: roadmap

I.1 what *is* the Internet?

I.2 network edge

- end systems, access networks, links

I.3 network core

- packet switching, circuit switching, network structure

I.4 delay, loss, throughput in networks

I.5 protocol layers, service models

I.6 networks under attack: security

I.7 history

Protocol “layers”

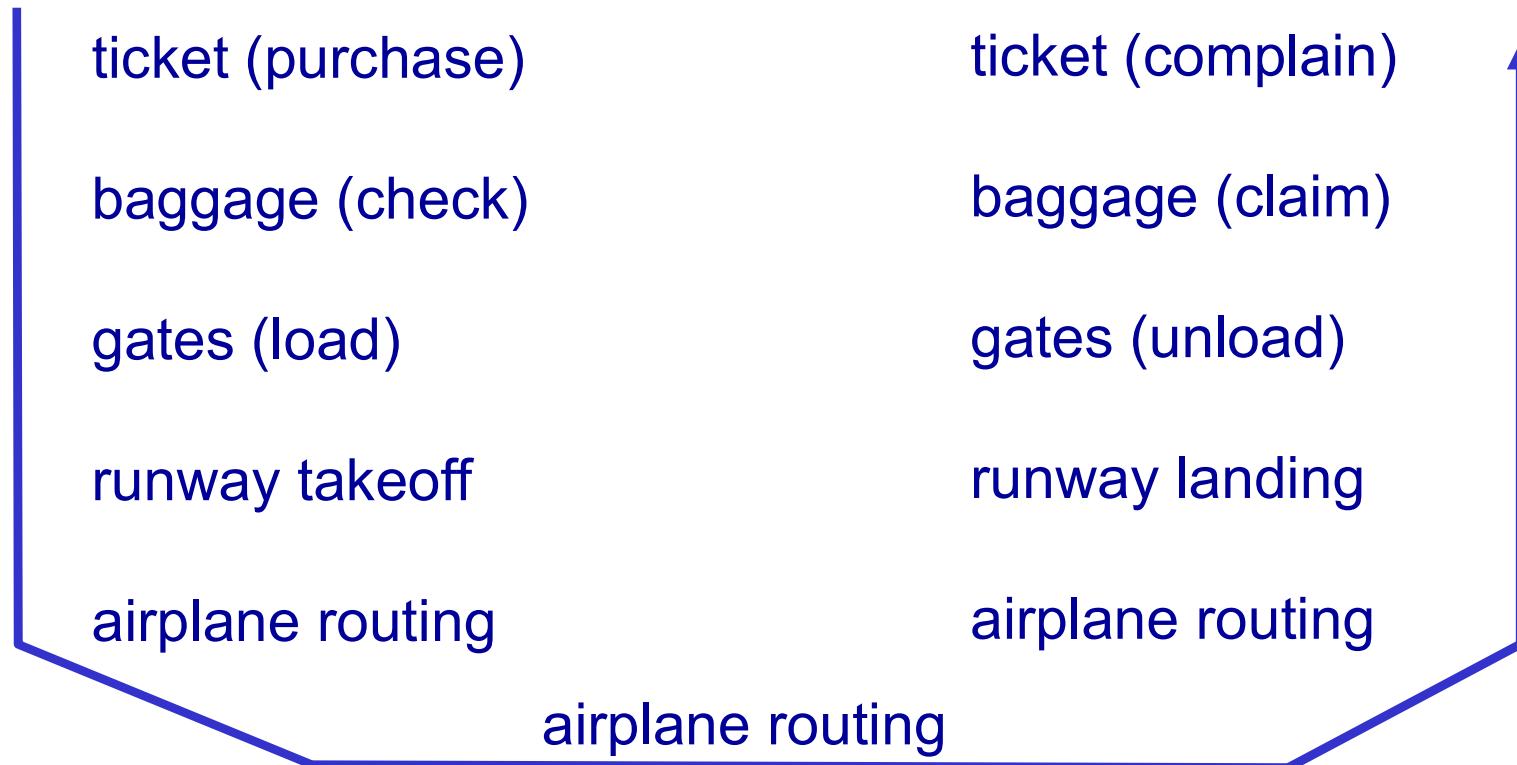
*Networks are complex,
with many “pieces”:*

- hosts
- routers
- links of various media
- applications
- protocols
- hardware, software

Question:
is there any hope of
organizing structure of
network?

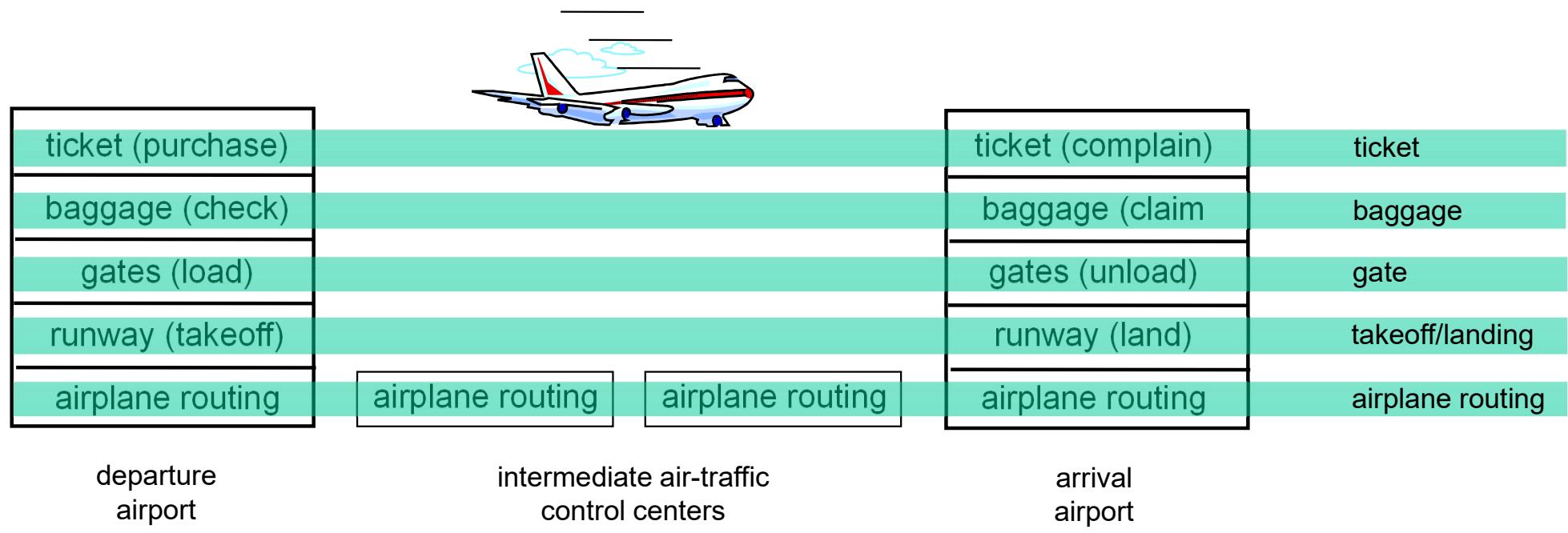
.... or at least our
discussion of networks?

Organization of air travel



- ❖ a series of steps

Layering of airline functionality



layers: each layer implements a service

- via its own internal-layer actions
 - relying on services provided by layer below

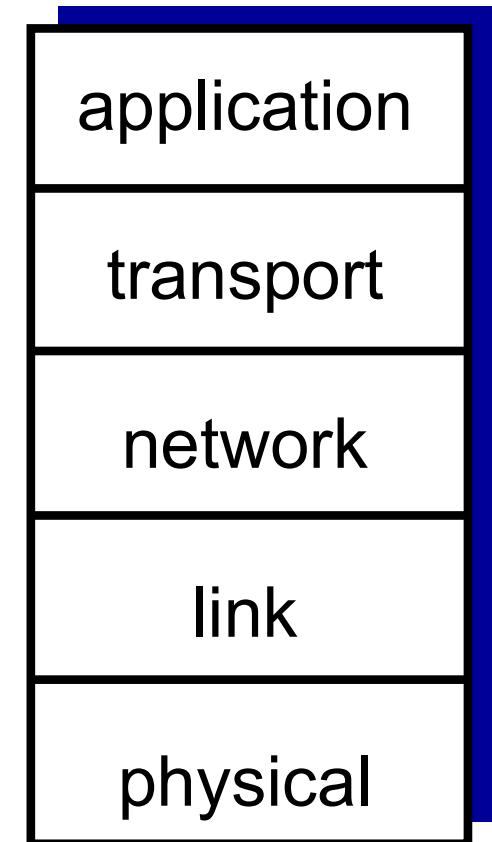
Why layering?

dealing with complex systems:

- ❖ explicit structure allows identification, relationship of complex system's pieces
 - layered *reference model* for discussion
- ❖ modularization eases maintenance, updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system
- ❖ layering considered harmful?

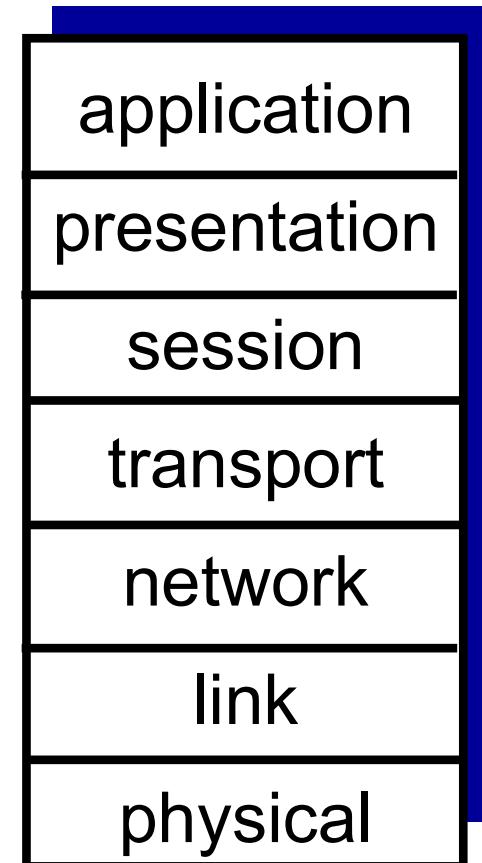
Internet protocol stack

- ❖ *application*: supporting network applications
 - FTP, SMTP, HTTP
- ❖ *transport*: process-process data transfer
 - TCP, UDP
- ❖ *network*: routing of datagrams from source to destination
 - IP, routing protocols
- ❖ *link*: data transfer between neighboring network elements
 - Ethernet, 802.111 (WiFi), PPP
- ❖ *physical*: bits “on the wire”



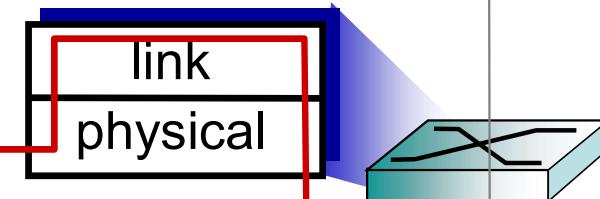
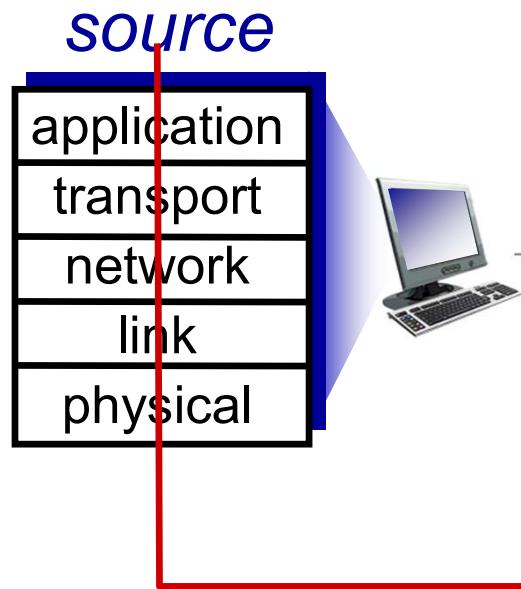
ISO/OSI reference model

- ❖ *presentation*: allow applications to interpret meaning of data, e.g., encryption, compression, machine-specific conventions
- ❖ *session*: synchronization, checkpointing, recovery of data exchange
- ❖ Internet stack “missing” these layers!
 - these services, *if needed*, must be implemented in application
 - needed?

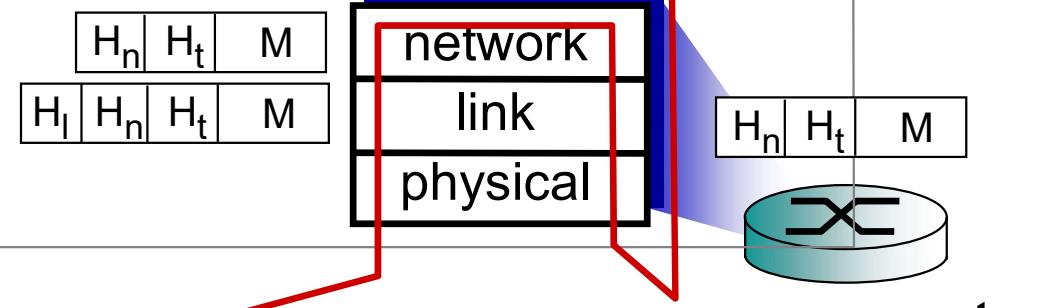
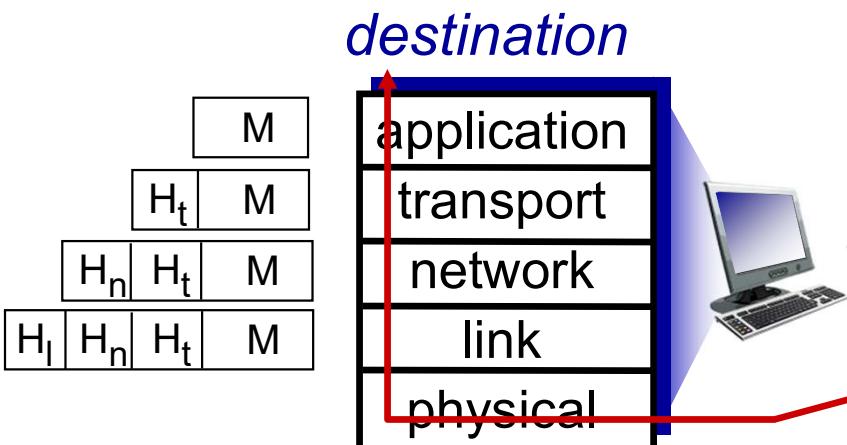


Encapsulation

message	M
segment	H _t M
datagram	H _n H _t M
frame	H _l H _n H _t M



switch



router

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Network security

- ❖ field of network security:
 - how bad guys can attack computer networks
 - how we can defend networks against attacks
 - how to design architectures that are immune to attacks
- ❖ Internet not originally designed with (much) security in mind
 - *original vision:* “a group of mutually trusting users attached to a transparent network” ☺
 - Internet protocol designers playing “catch-up”
 - security considerations in all layers!

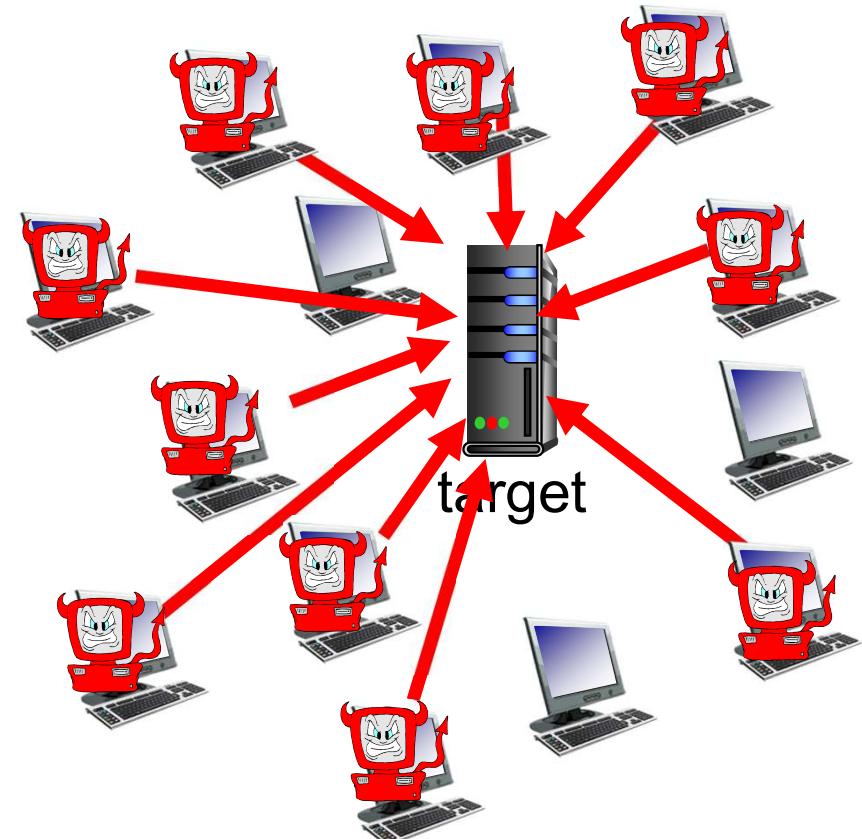
Bad guys: put malware into hosts via Internet

- ❖ malware can get in host from:
 - **virus**: self-replicating infection by receiving/executing object (e.g., e-mail attachment)
 - **worm**: self-replicating infection by passively receiving object that gets itself executed
- ❖ **spyware malware** can record keystrokes, web sites visited, upload info to collection site
- ❖ infected host can be enrolled in **botnet**, used for spam, DDoS attacks

Bad guys: attack server, network infrastructure

Denial of Service (DoS): attackers make resources (server, bandwidth) unavailable to legitimate traffic by overwhelming resource with bogus traffic

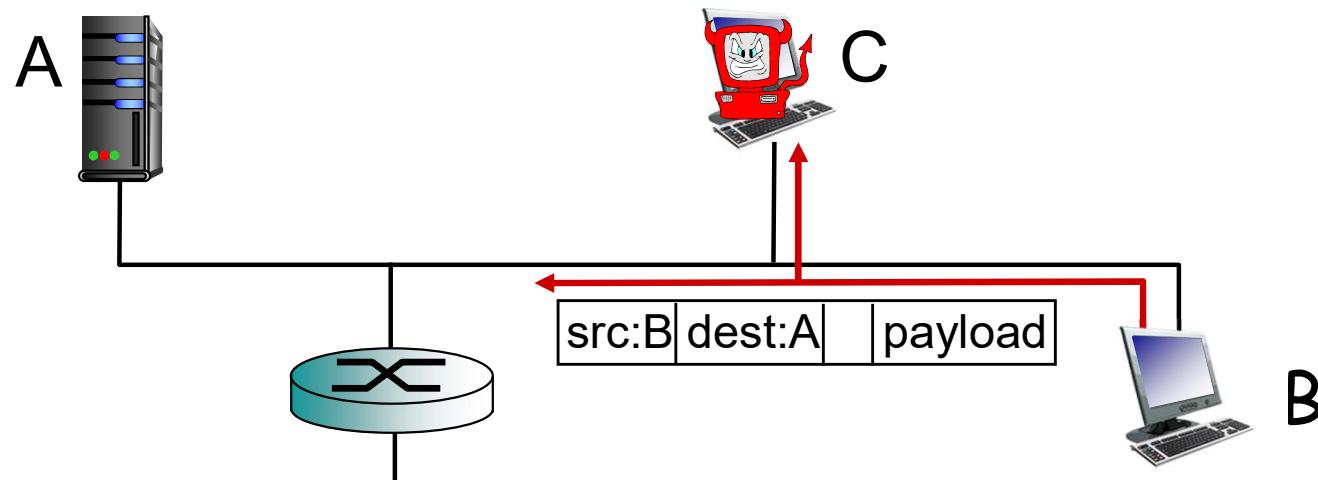
1. select target
2. break into hosts around the network (see botnet)
3. send packets to target from compromised hosts



Bad guys can sniff packets

packet “sniffing”:

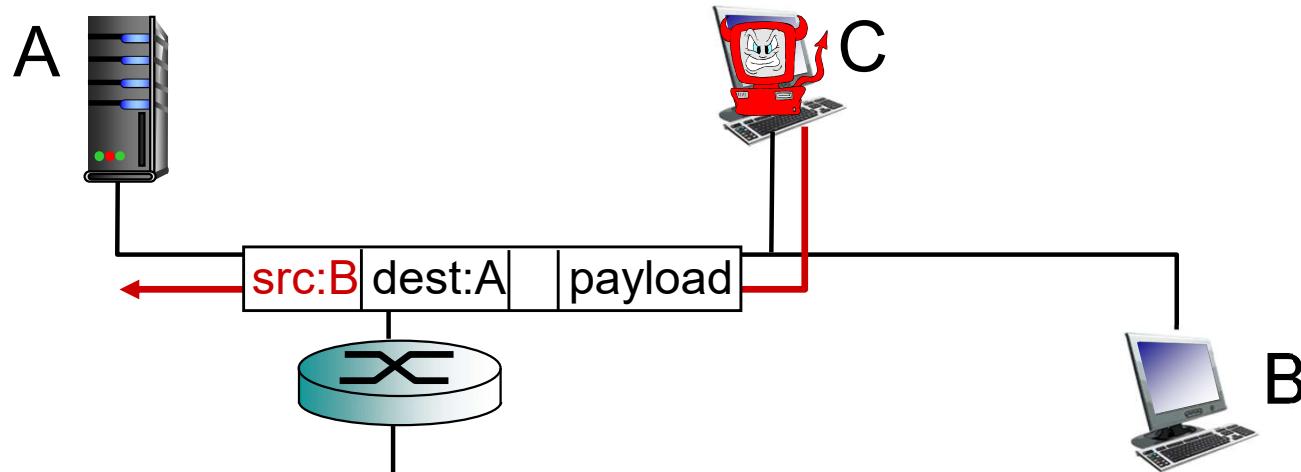
- broadcast media (shared ethernet, wireless)
- promiscuous network interface reads/records all packets (e.g., including passwords!) passing by



- ❖ wireshark software used for end-of-chapter labs is a (free) packet-sniffer

Bad guys can use fake addresses

IP spoofing: send packet with false source address



... lots more on security (throughout, Chapter 8)

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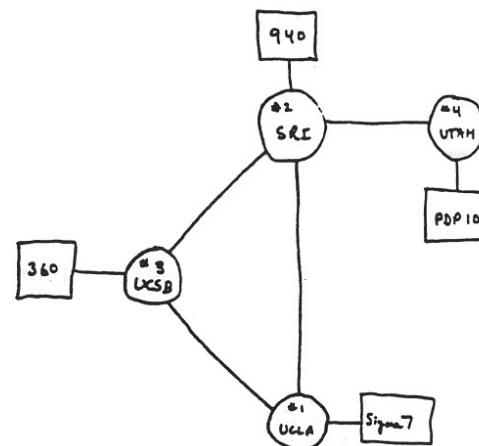
I.6 networks under attack: security

I.7 history

Internet history

1961-1972: Early packet-switching principles

- ❖ 1961: Kleinrock - queueing theory shows effectiveness of packet-switching
- ❖ 1964: Baran - packet-switching in military nets
- ❖ 1967: ARPAnet conceived by Advanced Research Projects Agency
- ❖ 1969: first ARPAnet node operational
- ❖ 1972:
 - ARPAnet public demo
 - NCP (Network Control Protocol) first host-host protocol
 - first e-mail program
 - ARPAnet has 15 nodes



THE ARPA NETWORK

Internet history

1972-1980: Internetworking, new and proprietary nets

- ❖ 1970: ALOHAnet satellite network in Hawaii
- ❖ 1974: Cerf and Kahn - architecture for interconnecting networks
- ❖ 1976: Ethernet at Xerox PARC
- ❖ late 70' s: proprietary architectures: DECnet, SNA, XNA
- ❖ late 70' s: switching fixed length packets (ATM precursor)
- ❖ 1979: ARPAnet has 200 nodes

Cerf and Kahn's
internetworking principles:

- minimalism, autonomy - no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

define today's Internet architecture

Internet history

1980-1990: new protocols, a proliferation of networks

- ❖ 1983: deployment of TCP/IP
- ❖ 1982: smtp e-mail protocol defined
- ❖ 1983: DNS defined for name-to-IP-address translation
- ❖ 1985: ftp protocol defined
- ❖ 1988: TCP congestion control
- ❖ new national networks: Csnets, BITnet, NSFnet, Minitel
- ❖ 100,000 hosts connected to confederation of networks

Internet history

1990, 2000's: commercialization, the Web, new apps

- ❖ early 1990's: ARPAnet decommissioned
- ❖ 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- ❖ early 1990s: Web
 - hypertext [Bush 1945, Nelson 1960's]
 - HTML, HTTP: Berners-Lee
 - 1994: Mosaic, later Netscape
 - late 1990's:
commercialization of the Web

- late 1990's – 2000's:
 - ❖ more killer apps: instant messaging, P2P file sharing
 - ❖ network security to forefront
 - ❖ est. 50 million host, 100 million+ users
 - ❖ backbone links running at Gbps

Internet history

2005-present

- ❖ ~750 million hosts
 - Smartphones and tablets
- ❖ Aggressive deployment of broadband access
- ❖ Increasing ubiquity of high-speed wireless access
- ❖ Emergence of online social networks:
 - Facebook: soon one billion users
- ❖ Service providers (Google, Microsoft) create their own networks
 - Bypass Internet, providing “instantaneous” access to search, email, etc.
- ❖ E-commerce, universities, enterprises running their services in “cloud” (eg, Amazon EC2)

Introduction: summary

covered a “ton” of material!

- ❖ Internet overview
- ❖ what’s a protocol?
- ❖ network edge, core, access network
 - packet-switching versus circuit-switching
 - Internet structure
- ❖ performance: loss, delay, throughput
- ❖ layering, service models
- ❖ security
- ❖ history

you now have:

- ❖ context, overview, “feel” of networking
- ❖ more depth, detail to follow!