The Kernel Abstraction

อดีต

Single task system

ปัจจุบัน

Multitasking system

Activity #1

• การเปลี่ยนแปลงจากระบบแบบ single task ไปเป็นระบบ แบบ Multitask ... สิ่งที่จะต้องมีการปรับแต่งหรือเพิ่มเติมเข้า มาใน OS ได้แก่...

What does an OS do...

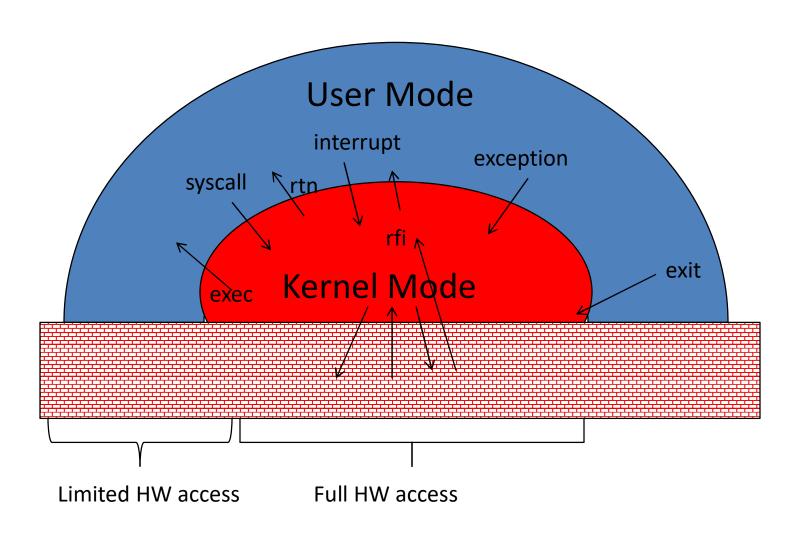
Hiding Complexity

 Kernel is the part of the OS that running all the time on the computer

UNIX Structure

User Mode		Applications	(the users)			
		Standard Libs	shells and commands ompilers and interpreters system libraries			
		system-call interface to the kernel				
Kernel Mode	Kernel	signals terminal handling character I/O system terminal drivers	file system swapping block I/O system disk and tape drivers	CPU scheduling page replacement demand paging virtual memory		
		kernel interface to the hardware				
Hardware		terminal controllers terminals	device controllers disks and tapes	memory controllers physical memory		

User/Kernel (Privileged) Mode



One of the major goals of OS is...

- Protecting Process and the Kernel
 - Running multiple programs
 - Keep them from interfering with the OS kernel
 - Keep them from interfering with each other

Activity #2: Protection: WHY?

เวลา 10 นาที

การ protect Process และ Kernel ทำให้เกิด impact อะไรกับระบบบ้าง และยังต้อง protect อะไรอีกบ้าง เพื่ออะไร

Protection: How? (HW/SW)

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Hardware Support: Dual-Mode Operation

- Kernel mode
 - Execution with the full privileges of the hardware
 - Read/write to any memory, access any I/O device, read/write any disk sector, send/read any packet
- User mode
 - Limited privileges
 - Only those granted by the operating system kernel
- On the x86, mode stored in EFLAGS register
- On the MIPS, mode in the status register

Hardware Support: Dual-Mode Operation

- Privileged instructions
 - Available to kernel
 - Not available to user code
- Limits on memory accesses
 - To prevent user code from overwriting the kernel
- Timer
 - To regain control from a user program in a loop
- Safe way to switch from user mode to kernel mode, and vice versa

Privileged instructions

Examples?

 What should happen if a user program attempts to execute a privileged instruction?

User Mode

- Application program
 - Running in process

Virtual Machine:VM

- Software emulation of an abstract machine
 - Give programs illusion they own the machine
 - Make it look like HW has feature you want
- 2 types of VM
 - Process VM
 - Supports the execution of a single program (one of the basic function of the OS)
 - System VM
 - Supports the execution of an entire OS and its applications

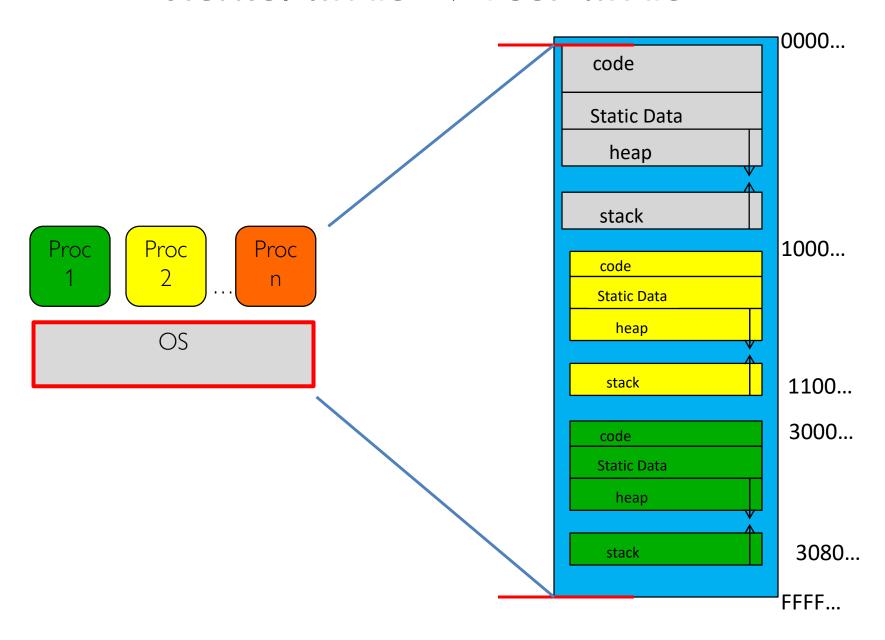
Process VMs

• GOAL:

- Provide an isolation to a program

- Portability (Program)

Kernel mode & User mode

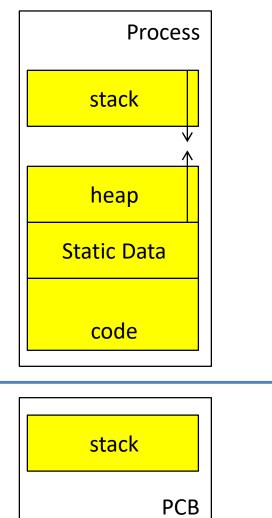


Process Abstraction

- Process: an *instance* of a program, running with limited rights
 - Thread: a sequence of instructions within a process
 - Potentially many threads per process (for now 1:1)
 - Address space: set of rights of a process
 - Memory that the process can access
 - Other permissions the process has (e.g., which system calls it can make, what files it can access)

Process

- 2 parts
 - PCB in kernel
 - Others in user



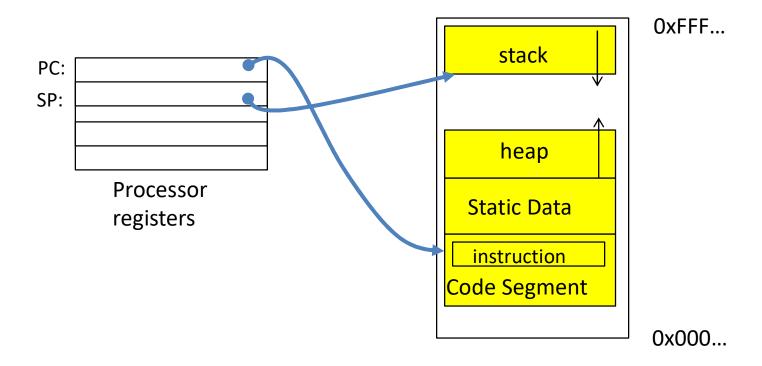
User

Kernel

Process Control Block: PCB

- Kernel represents each process as a process control block (PCB)
 - Status (running, ready, blocked, ...)
 - Registers, SP, ... (when not running)
 - Process ID (PID), User, Executable, Priority, ...
 - Execution time, ...
 - Memory space, translation tables, ...
- Kernel Scheduler maintains a data structure containing the PCBs
- Scheduling algorithm selects the next one to run

Address Space: In a Picture



Break