

Questions 1 through 33 form the written portion of the cybersecurity event. It counts for 50% of the event score. The questions will be a mixture of Cryptography and Web Architecture. Questions 34 and 35 form the hands-on cryptography questions. They count for 25% of the event score.

For the hands-on programming portion, go to https://www.hackerrank.com/bearso-2020-programming-hands-on (https://www.hackerrank.com/bearso-2020-programming-hands-on) and follow the instructions in this document: https://docs.google.com/document/d/1U5d6LedWSXWXi48io64GP_bopN5HOG3xugw2MgXlohs/edit?usp=sharing). That counts for 25% of the event score.

1. (2.00 pts) You've decided to encrypt a message with a substitution cipher. Instead of just scrambling the letters A-Z, you're going to scramble all the printable ASCII characters (of which there are 95), and your key is a random permutation. How many bits of security does your key have? A) 26 B) 95 C) 125 D) 259 E) 491 F) 9025

2. (1.00 pts) Which of the following are *hashing* algorithms? (You may select zero, one, or multiple.)

(Mark ALL correct answers)

✓ A) SHA256

□ B) Electronic Codebook
S by Leastonic Coccasion.
☐ C) Vigenère
✓ D) MD5
□ E) HTTPS
✓ F) Whirlpool
3. (6.00 pts) Consider the following encryption scheme:
Alice and Bob share a random secret password, PASS, that is 500 characters long. For Alice to sent a message M to Bob, she:
1. Divides M into chunks C_i of two characters each
2. Computes $H_i = SHA256(C_i + PASS)$ for each chunk, where + is string concatenation. 3. Alice sends these hashes to Bob, in order.
For Bob to read the message, he:
1. Loops through each hash, H_i , and:
2. tries each possible two-character string c until he finds $H_i=SHA256(c+PASS)$ 3. His final message is the concatenation of all his c 's.
Which of the following are problems with this scheme?
(Mark ALL correct answers)
A) Bob is likely to end up with a different string that what Alice meant to encode.
B) A 500 character password is not long enough for this type of scheme.
 C) SHA256 is one-way, which means Bob can't find the c's in any reasonable amount of time. ✓ D) This encryption has no diffusion, and patterns will show up in the output.
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6. (3.00 pts) is the process of creating a new shared secret password, that eavesdroppers will not be able to figure out based on your conversation.
Key exchange
7. (2.00 pts) Which of the following best describes the use of a <i>certificate</i> in cryptography, in particular for websites?
A) A certificate proves that the content has been encrypted, by the agreed upon protocol. This stops downgrade attacks.
O B) A certificate proves that the content is newly generated, with a verifiable timestamp. This stops replay attacks.
O C) A certificate provides the necessary information for the client to decrypt the data. Without it, eavesdroppers could read the traffic.
A certificate is a way for the server to prove they are who they say they are.
A certificate proves that the company has been vetted to use good practice in security, by a third-party like Digicert or the US government.
8. (4.00 pts) A attack is when an attacker has captured the unencrypted and encrypted versions of some messages, and now wants to figure out the secret key. A attack is when the attacker has sent their own message to be unencrypted, gotten it back decrypted form, and now wants to figure out the secret key. A attack is when the attacker steals secrets from the server using other information than the normal protocol such as timing, sound, or power usage. (All answers should be two words)
Known plaintext Chosen ciphertext Side channel

9. (7.00 pts)

The RC4 algorithm generates an arbitrarily long sequence of bytes, after being initialized with a key. This sequence of bytes can then be used as a stream cipher. It makes use of an array, S, that at all times contains a permutation of the numbers 0 to 255.

```
S := RandomPermutation([0,1,2 ... 255], key)
i := 0
j := 0
while True:
    i := (i + 1) mod 256
    j := (j + S[i]) mod 256
    swap values of S[i] and S[j]
    K := S[(S[i] + S[j]) mod 256]
    output K
```

The RC4 algorithm is widely regarded as a secure and trustworthy algorithm. Suppose the algorithm was changed to:

```
S := RandomPermutation([0,1,2 ... 255], key)
i := 0
j := 0
while True:
    i := (i + 1) mod 256
    j := (j + i) mod 256
    swap values of S[i] and S[j]
    K := S[(S[i] + S[j]) mod 256]
    output K
```

The change has been highlighted in red. Explain either why this would still be a trustworthy stream cipher, or why it would be vulnerable. Support your answer in terms of diffusion or confusion.

Expected Answer: Key point: the variable "j" no longer depends on the values in S. This means that S will undergo a *fixed* permutation every so often (every 512 loops), that does not depend on its current state, so the cipher will repeat quickly. 5 points for saying that it will repeat. 1 point for saying there is no longer diffusion of S into j. 1 point for saying that S will be permuted the same way every 512 loops. Incorrect answers would include things like, "It's still secure", "S stops being a permutation", "j will grow too quickly", "the key isn't used"...

10. (2.00 pts) The advent of post-quantum cryptography has made necessary the distinction between information-theoretic security and computational security. Which type of security is "more secure"?
A Information theoretic cocurity
A) Information-theoretic security
O B) Computational security
11. (2.00 pts) Of the two types of security listed in the previous question, one of them cannot be carried out with traditional internet communication. Which is that?
A) Information-theoretic security
O B) Computational security
49 (9.00 mts)
12. (2.00 pts) Give an example of how each of these two types of security were used in the year 2000 after modern cryptography, but before any quantum communications existed.
Expected Answer: Computational security: Anything involving hashing, RSA, public key, AES, etc most of the topics on this test. Information-theoretic security: One-time pads (and pretty much only those).
13. (9.00 pts) Consider the following encrypt function, written in pseudocode. Identify three flaws with this system. Each flaw will get you three points. You can assume that this key is used only for this one message, i.e. you don't need to worry about key reuse.
NOTE: You will use this same code again in the Hands-On portion, with a different question. You will be given a copy of the code again. This might be relevant if you're deciding
whether or not to attempt this problem.

```
import numpy
import random
#Get a random matrix, given a random number generator and a size.
#Fills it with random integers mod "modulo".
def Get Random Matrix(rand, modulo, size):
       mat = numpy.zeros((size,size))
        for i in range(size):
               for j in range(size):
                        mat[i,j] = numpy.floor(rand.random() * modulo)
        return mat
#Encrypt a message given the key.
\#Requirement: "message" is an array of bytes, with length a multiple of 4.
def encrypt(message, key):
        #The number of rows and columns in the matrix
        random.seed(key)
        output = []
        #Loop through every 4 bytes
        for block in range(0,len(message),n):
                matrix = Get_Random_Matrix(random, 256, n)
                vector = message[block : block+n]
                \#This is matrix multiplication, mod 256
                vector = numpy.dot( matrix, vector ) % 256
                output += list(vector)
        return output
```

Expected Answer: Intended flaws: * The matrices might not always be invertible. In this case, you can't recover the plaintext. * The series of bytes [0,0,0,0] will always be mapped to [0,0,0,0] again. * The python random number generator is not secure. ...and certainly any other new flaws they find should count! Things that are not flaws: * The key isn't hashed. * Someone might use a small key. * It maps the same set of bytes to the same set, every time, like ECB (it doesn't) * Linear operations are easy to undo * n=4 is too small to be secure

14. (2.00 pts)

Diffie-Hellman is often carried out using exponentiation of integers, but this has some people concerned that it could be broken using advanced number theory. So instead, some people use Diffie-Hellman based on a collection of points, and "adding" those points together. What is this set of points called?

Elliptic Curve

15. (4.00 pts)

"Post-quantum" cryptography refers to cryptographic protocols developed to address the rise of quantum computers. However, very few people use these protocols yet. What are the main reasons for their lack of adoption?

(Mark ALL correct answers)

- A) Physicists don't understand the laws of quantum mechanics well enough to know if the algorithms work.
- The protocols are several times slower than traditional protocols like RSA.
- O) Most people don't have a quantum computer yet, so they can't even run the algorithms!
- These algorithms are new, so they haven't been studied for as long, so people don't trust them yet.
- E) It takes a lot of work to upgrade these protocols that are core to the internet, and get everyone to switch over.
- F) The US Military issues export restrictions on cryptographic algorithms, which mean you can't use the new protocols outside of the US.

(Mark ALL correct a	answers)
□ B) n	
□ C) e	
✓ D) q	
✓ E) phi	
□ F) m	
17. (1.00 pts)	"Traditional" ciphers like substitution ciphers or Vigenere ciphers can have lots of keys, but are now considered a poor choice for secure communication. Why?
(Mark ALL correct a	
'	ip the same input to the same output over and over, so you can easily find repeated letters or words. re developed for English text, so there's no way to use them for other data like images or movies on a computer.
,	no confusion or diffusion in the encryption process, so you figure out the keys one part at a time.
_ ′	ers can brute-force all of the keys and find the plaintext easily.
18. (1.00 pts)	Which attack is most associated with session management and cookie abuse?
O A) Server-s	ide Template Injection
O B) Buffer O	
,	te Request Forgery
O D) Cross-si	
,	
19. (2.00 pts)	Identify all valid HTTP request verbs in the list below.
(Mark ALL correct a	
✓ A) PUT	answers)
☐ B) SEARCH	1
✓ C) TRACE	
✓ D) OPTION	s
✓ E) PATCH	
☐ F) LISTEN	
20. (1.00 pts)	Internet standards, which describe how internet-related processes should operate are documented by
A) Request	For Comments
'	ssion Control Protocol
O C) Internet	Protocol
	at Transfer Protocol
21. (2.00 pts)	An attacker is testing a website and tries the request
21. (2.00 pts)	
	http://somewebsite.com/?file=///etc/passwd
	What vulnerability is the attacker looking for?

O A) XML External Entities
B) Local File Inclusion
C) Cross-site Scripting
O D) SQL Injection
22. (2.00 pts) In the list below, identify both the character placed as a marker before a list of query parameters and the character that separates query parameters in a URL.
(Mark ALL correct answers)
☑ B) &
□ c):
☑ D) ?
□ E) %
□ F) \$
23. (2.00 pts) A cookie with the attribute is not accessible by JavaScript.
A) HttpOnly
○ B) Secure
C) Max-Age
D) Domain
24. (1.00 pts) The header of an HTTP request allows a server to identify information such as the application and operating system a request was made from.
User-Agent
User-Agent
User-Agent 25. (1.00 pts) CTF websites are infamous for overloading in the first hour of the event. In this case, you might encounter errors of the formxx while desperately refreshing the page.
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Dec Hx Oct Char 0 0 0000 NUL (null) 1 1 001 SOH (start of heading) 2 2 0002 STX (start of text) 3 3 003 ETX (end of text) 4 4 004 EOT (end of text) 5 5 005 ENO (enquiry) 6 6 006 ACK (acknowledge) 7 7 007 BEL (bell) 8 8 010 BS (backspace) 9 9 011 TAB (horizontal tab) 10 A 012 LF (NL line feed, new line 11 B 013 VT (vertical tab) 12 C 014 FF (NP form feed, new page 13 D 015 CR (carriage return) 14 E 016 SO (shift out) 15 F 017 SI (shift in) 16 10 020 DLE (data link escape) 17 11 021 DC1 (device control 1) 18 12 022 DC2 (device control 2) 19 13 023 DC3 (device control 3) 20 14 024 DC4 (device control 4) 21 15 025 NAK (negative acknowledge) 22 16 026 STM (synchronous idde) 23 17 027 ETB (end of trans. block) 24 18 030 CAN (cancel) 25 19 031 EM (end of medium) 26 1A 032 SUB (substitute) 27 1B 033 ESC (escape) 28 1C 034 FS (file separator) 30 1E 036 RS (record separator) 31 1F 037 US (unit separator)	43 2B 053 6#43; + 75 4B 113 6#75; K 107 6B 153 6#107; k
bearso%7B%25signs%3Dc	
27. (1.00 pts) What cookie attribute differentiates session Expected Answer: Exact name - Expires Description - expires	cookies and persistent cookies? Enter either its exact name or a description of it. ration date or similar
28. (3.00 pts) What is the purpose of a webpage's robots.	.txt file?
Expected Answer: Sample answer: It specifies how to infor	m the web crawlers about which areas of the website should not be processed or scanned.
29. (3.00 pts) Explain the difference between reflected an	nd persistent cross-site scripting attacks.
not stored in a database, but rather rendered immediately. U	o unsanitized data stored in a server's database; anyone who accesses the resource can be affected. Reflected XSS is sually, a client sends in unsanitized data and it's displayed in the browser; it doesn't affect many people outside the user. mall amount of people (usually just the client), feedback is immediate persistent - attack payload stored, affects anyone when affected resource is accessed

```
function f(s) {
  var res = "";
  for (i = 0; i < s.length; i++) {
    if (s[i] >= 'a' && s[i] <= 'z') {
      res += s[i].toUpperCase();
    } else if (s[i] >= 'A' && s[i] <= 'Z') {
      res += s.charCodeAt(i) - 'A'.charCodeAt();
    } else {
    res += s[i];
    }
  }
  return "bearso{" + res + "}";
}</pre>
```

What does f("bDErF_FA_FcEry") return?

bearso{B34R5_50_5C4RY}

31. (8.00 pts)

Consider an application with this snippet of code:

```
req = "SELECT * FROM users WHERE name='" + username + "' AND password='BEARSO'"
sqlserver.execute(req) # This line of code sends the string req to an SQL server as a request.
```

- a. (1 point) This app is vulnerable to a(n) _____.
- b. (3 points) To retrieve all elements of the table users what should the value of the variable username be?
- c. (4 points) If an attacker wanted to steal data from another table, let's say pocky-flavors, what should the value of the variable username be? (assume the pocky-flavors table has the same dimensions as users as to not trigger an error).

Please add a., b., and c., before each answer in your response.

Expected Answer: a. SQL injection b In general: 'OR {true expression} {comment delimiter ";" or "--"} ex. 'OR 1=1; c. In general: 'UNION SELECT * FROM pocky-flavors WHERE {true expression} {comment delimiter ";" or "--"} ex. 'UNION SELECT * FROM pocky-flavors WHERE 1=1-- Comment delimiters are necessary since otherwise AND password='BEARSO' would be included in the query, which would result in nothing being returned.

32. (6.00 pts) Consider an application with this snippet of code:

```
if ($_POST["secret"] == $SUPERSECRET) {
   admin_login();
} else {
   get_mad();
}
```

- a. (1 point) An attacker can exploit ______ in the conditional of the if statement to gain access to admin_login().
- b. (2 points) A value of secret that an attacker could supply is ______. (Assume \$SUPERSECRET is a string.)
- c. (3 points) Changing the program by one character to this program could fix its vulnerability. What should the programmer do?

Please add a., b., and c., before each answer in your response.

Expected Answer: a. Type juggling, ==, loose typing, or similar b. 0 c. Change the == in the conditional to === Please add a., b., and c., before each answer in your response.

```
POST /cookiejar.php HTTP/1.1
User-Agent: Mozilla/4.0 (compatible; MSIE5.01; Windows NT)
Host: www.bearsobakery.com
Content-Type: application/x-www-form-urlencoded

Accept-Language: en-us
Accept-Encoding: gzip, deflate
Connection: Keep-Alive
order=chocolatechip
```

Content-Length: 19

Cryptography Hands-On Portion

34. (24.00 pts)

Consider the following encrypt function, written in Python. It is identical to the system you saw in the Written portion. All blocks are available for copy-paste at https://pastebin.com/raw/BsUzPmdX (https://pastebin.com/raw/BsUzPmdX).

```
import numpy
import random
#Get a random matrix, given a random number generator and a size.
#Fills it with random integers mod "modulo".
def Get Random Matrix(rand, modulo, size):
        mat = numpy.zeros((size,size))
        for i in range(size):
               for j in range(size):
                        mat[i,j] = numpy.floor(rand.random() * modulo)
        return mat
\# Encrypt a message given the key.
#Requirement: "message" is an array of bytes, with length a multiple of 4.
def encrypt(message, key):
        \#The number of rows and columns in the matrix
        random.seed(key)
        output = []
        #Loop through every 4 bytes
        for block in range(0,len(message),n):
                matrix = Get_Random_Matrix(random, 256, n)
                vector = message[block : block+n]
                \#This is matrix multiplication, mod 256
                vector = numpy.dot( matrix, vector ) % 256
                output += list(vector)
        return output
```

Someone has used the *encrypt* function in the previous question to encrypt 8 bytes. Their secret key is the integer 28. The output (the ciphertext) is the array [5.0, 5.0, 43.0, 123.0, 68.0, 246.0, 201.0, 35.0]. Implement a decrypt function, to figure out:What are the first two numbers of the input (the plaintext)? You may wish to use the following code, for computing matrix inverses modulo p.

```
from numpy import matrix
from numpy import linalg
{\tt def\ modMatInv(A,p):}
                             # Finds the inverse of matrix A mod p
  n=len(A)
  A=matrix(A)
  adj=numpy.zeros(shape=(n,n))
  for i in range(0,n):
    for j in range(0,n):
      \label{eq:adj[i][j]=((-1)**(i+j)*int(round(linalg.det(minor(A,j,i)))))%p} adj[i][j]=((-1)**(i+j)*int(round(linalg.det(minor(A,j,i)))))%p
  return \ (modInv(int(round(linalg.det(A))),p)*adj)\%p
def modInv(a,p):
                             # Finds the inverse of a mod p, if it exists
  for i in range(1,p):
    if (i*a)%p==1:
      return i
  raise ValueError(str(a)+" has no inverse mod "+str(p))
def minor(A,i,j):
                      # Return matrix A with the ith row and jth column deleted
  A=numpy.array(A)
  \label{eq:minor} \verb|minor=numpy.zeros(shape=(len(A)-1,len(A)-1))| \\
  for s in range(0,len(minor)):
   if p==i:
      p=p+1
    q=0
    for t in range(0,len(minor)):
     if q==j:
        q=q+1
      minor[s][t]=A[p][q]
      q=q+1
    p=p+1
  return minor
```

You may wish to use https://repl.it/languages/python (https://repl.it/languages/python) or https://repl.it/languages/python3 (https://repl.it/languages/python3), they will have all the packages you need.

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35. (24.00 pts)

We've been intercepting transmissions from aliens for a while, and they use a special language with only ten characters: ACEFLMOPXY. By taking large samples, we have their digraph frequencies -- the frequency of every pair of letters when in this language. The matrix of frequencies is:

[2.71743254e-03, 2.95467253e-04, 1.40289673e-03, 7.20653007e-06,	
1.85226513e-02, 3.79145035e-03, 3.03941946e-04, 1.31030842e-03,	
2.08116563e-02, 4.08118653e-03],	
[5.86397974e-03, 5.00155945e-03, 1.76121250e-02, 2.49812154e-02,	
1.01608926e-03, 6.77517525e-04, 2.10792531e-02, 2.12760389e-03,	
6.94094301e-04, 2.74646235e-02],	
[1.30648789e-03, 4.30724718e-04, 1.81207876e-03, 4.00819280e-05,	
7.62017704e-03, 3.40079355e-03, 1.97508670e-02, 1.93630388e-02,	
2.13334334e-02, 8.90379568e-03],	
[6.03752785e-03, 6.79558987e-04, 5.78578167e-04, 1.89196884e-02,	
3.90869323e-02, 1.34144667e-06, 5.46029978e-05, 3.43034643e-02,	
1.57870130e-03, 4.48519966e-07],	
[1.08794700e-02, 5.48098533e-02, 6.69358702e-03, 1.74655811e-05,	
1.69510857e-06, 2.04915851e-02, 2.01802310e-04, 4.80953473e-02,	
5.05292049e-03, 1.95240540e-02],	
[6.32580292e-03, 2.71218352e-02, 2.67347435e-02, 3.00882655e-05,	
2.87534777e-04, 4.66718544e-05, 1.45967595e-02, 6.16627671e-04,	
5.45688335e-05, 1.27962918e-03],	
[1.33773698e-03, 1.47335003e-03, 1.02548815e-02, 3.98417829e-03,	
1.46758075e-02, 5.70272782e-04, 1.62947773e-03, 2.90758251e-02,	
8.14832182e-04, 1.91533197e-02],	
[1.86562469e-02, 3.37085700e-03, 9.25444336e-04, 3.81832477e-02,	
5.12699314e-02, 8.12671061e-03, 5.91072348e-04, 8.49518905e-04,	
1.74208748e-05, 1.37923656e-02],	
[2.84182925e-05, 1.14994712e-02, 5.69087979e-04, 5.01564480e-03,	
1.53687521e-04, 1.85763624e-02, 1.23625437e-02, 1.18727534e-05,	
1.13290886e-02, 1.64661646e-02],	
[9.10948547e-05, 1.83538397e-03, 1.73780557e-02, 1.00620275e-02,	
3.31332740e-02, 2.14115561e-02, 1.23993612e-02, 2.92084602e-05,	
1.43256257e-02, 6.74294919e-03]]	

You can also view these numbers at https://pastebin.com/raw/BsUzPmdX (https://pastebin.com/raw/0VnPNbcc) . This could be useful if you need to copy-paste the text and cannot copy-paste from the browser.

Recently they started using a substitution cipher. We need you to match the character frequencies to decode this block of text. It's alien, so don't expect it to make much sense in English...

ofpmaafcfcpoxyaofmfpoffmlcyaaaelxxxclycyaomelpofpoelelclclxlcooxlcoxxlcoofxlyclcfcpefcyaaaomefxcpoxlxlxcoemlcfcofpomfpelfaelclfpmmelxxcpoelcyaecelcyaefeta lfafemfmafclpmaefaaefecomaeclpofafpelfelcomecoefpmefclclxlyafcfclcfcoelyaofclxcyaapefmafmpelemlyfmlyaooefcoxlcpefclcoxcpoxlxyafecfaemcmaexcaoxlclfapoecoel pmelcfclxlxxcfpmclclxxcyapoxlxcaoxlfpmefcfcaefmapofpmyclcfpmefcoxcpoefcomemlclclpoxxlclyappoeclmafcomlclcoxoecpeycpoyaelcoefcpypmfcffmefapmlpmemclxyaoxlcl comclcyycpofafcypymmeloecofapmcfmafxlfclyafpoxlcmelcaefpomlxlxlyaofpmcaaopofcoelxlfmemfcycoxxcoelclxlcpcmelcomffcycpoeclfpomapmcoxcpolxlcpmefpoefmlcyyapoya ofmaaafxcoyclyaxcoeclypycoelfaapelcoxyfaaelyaofclfpoeclclxlpycoemlxlclfcpoemelfclfpelymfxxxlxlxlcfeffcoxcpoelxcmeclelfcpocoefclyapmaexyapmemlxcolcpefelcly apoyapmfapofcllxlcomfmeapmexlffxxmafefclxlcpelxcpofpefpmfpooxaoelxxcmlcmcoelcoemecoxcoxyapyyapefclxlxclxlxxlcomecfpmeclxxcfpmafpoemeclcpmaexcoefeyafcfcyafelyaelclclxlclxcapoelfmfpexxlxlpoxlxxxxxcfxcoxcoxxclycfafaapefapomafpypmcomlp

Once you've decrypted it, please encrypt the words "EPOXY FLAME FLEECE" and submit these as your answers.

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l			, !	

You're done with the test!