LINEAR BOUNDED AUTOMATA (Module IV)

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Linear Bounded Automata

A linear bounded automaton is a multi-track non-deterministic Turing machine with a tape of some bounded finite length.

Length = function (Length of the initial input string, constant c)Here,

Memory information ≤ c × Input information

- The computation is restricted to the constant bounded area.
- The input alphabet contains two special symbols which serve as **left end markers and right end markers** which mean the transitions neither move to the left of the left end marker nor to the right of the right end marker of the tape.

A linear bounded automaton can be defined as an 8-tuple (Q, X, Σ , q_0 , ML, MR, δ , F) where

Q is a finite set of states

X is the tape alphabet

\(\) is the input alphabet

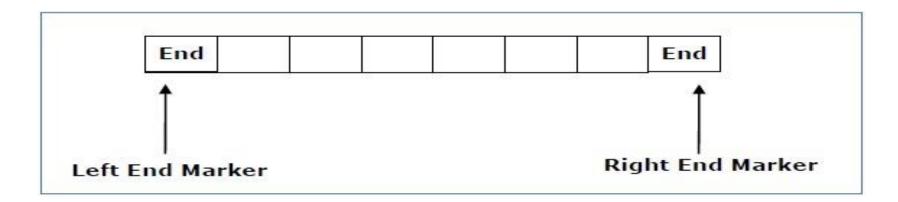
q₀ is the initial state

M_I is the left end marker

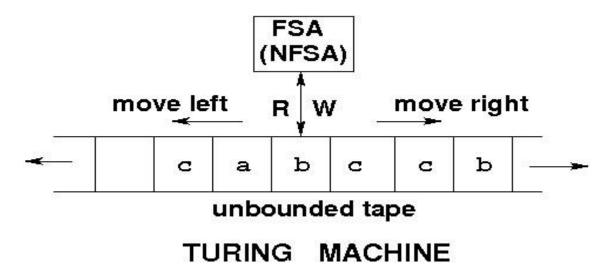
 M_R is the right end marker where $M_R \neq M_L$

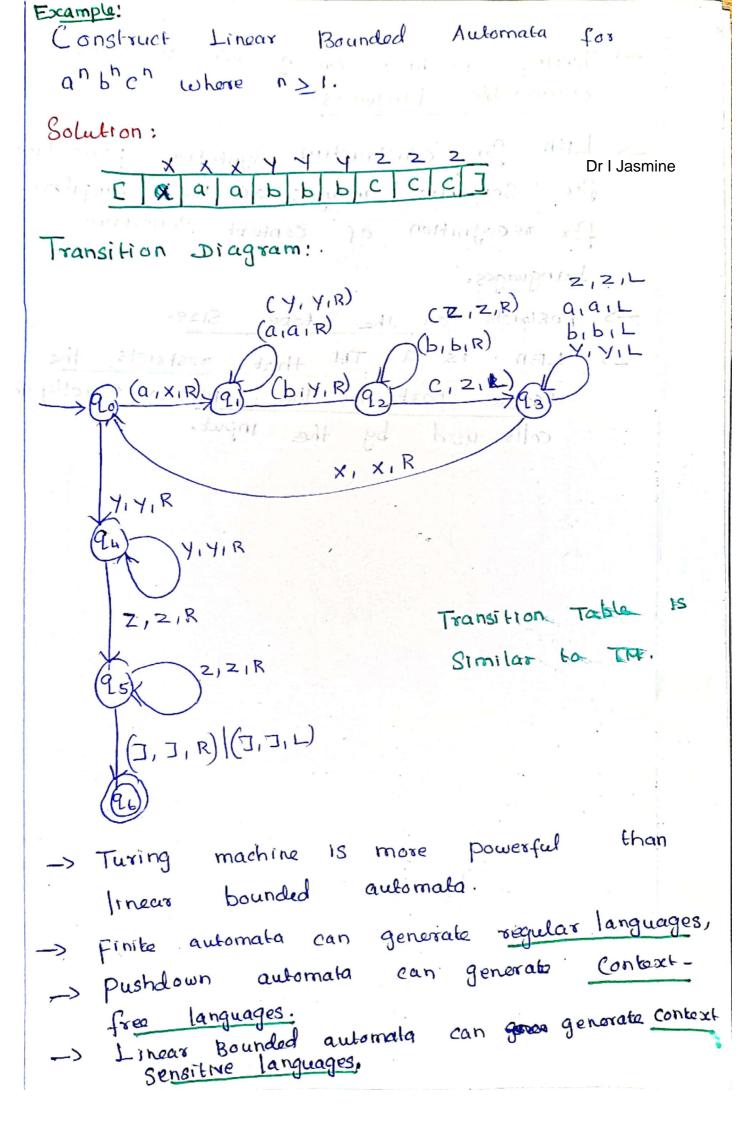
δ is a transition function which maps each pair (state, tape symbol) to (state, tape symbol, Constant 'c') where c can be 0 or +1 or -1

F is the set of final states



 A deterministic linear bounded automaton is always context-sensitive and the linear bounded automaton with empty language is undecidable.





-> Turing machines can generate recursively enumerable languages.

-> LBA for constructing syntatic parse trees
for secognition of context sensitive

Languages.

-> Rostricts on the tape 812e.

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-> LBA 13 a TM that restricts the usable part of the tape to exactly the cells used by the input.

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