

$$\begin{array}{ll} S * T \rightarrow S & \text{(premise)} \\ S * T & \text{(premise)} \\ S & \text{(conclusion)} \end{array}$$

Specific arguments with similar forms have the same validity.

Therefore,  $p$  (AND elimination)

NOTE, order of symbols matter even if they don't affect validity

$p, q, z$

John is sick or John is bored  
John is not sick  
 $\therefore$  John is bored

$$\begin{array}{l} p \vee q \\ \sim p \\ \rightarrow q \end{array}$$

Example:

John is sick or John is bored.

John is not sick.

Therefore, John is bored.

One can symbolize this argument as the argument form:

(i)  $S \vee B$

$\sim S$

therefore, B

But this is talking in terms of specific statements of a specific argument. We can also show its general argument form, by saying that the argument in (i) HAS the form:

(ii)  $p \vee q$

$\sim p$

therefore, q

# Specific argument form

Specific argument form = "Specific argument form is an array of symbols that contains statement variables – each statement variable stands for a distinct **simple** statement – such that when statements are substituted for statement variables, the result is an argument."

## Test of validity

Premise true  $\Rightarrow$  conclusion must not be false  
for argument being valid

$p$	$q$	$p \vee q$	$\sim p$
1	1	1	0
1	0	1	0
0	1	1	1
0	0	0	1

There are three valid argument forms (argument formulas/equations) which are based on conditional/implication.

Any specific argument that is based on those argument forms shall also be valid.

These three argument forms are:

Modus ponens, modus tollens and hypothetical syllogism.

Affirm consequent by affirming antecedent  
Affirm 'q' by affirming 'p'

$p \rightarrow q$   
 $p$   
Therefore  $q$

$p$	$q$	Statement with connective ( $p \rightarrow q$ )
1	1	1
1	0	0
0	1	1
0	0	1

# Modus Tollens

Denying Antecedent by denying consequent

$p \rightarrow q$   
 $\sim q$   
 $\therefore \sim p$

$p \rightarrow q$   
 $\sim p$   
 $\sim q$

$p \rightarrow q$   
 $a$   
 $b$

Affirming  
the  
consequent

Hypothetical syllogisms

$\hookrightarrow$  Denying the antecedent