

$$\rightarrow \Delta Q = \Delta U + \Delta W$$

\rightarrow State of entropy \uparrow

given 'n'

Relation Between Entropy and Probabilities

- "The Boltzmann formula shows the relationship between entropy and the number of ways the atoms or molecules of a thermodynamic system can be arranged."

$$S = k \log W$$

- "At equilibrium, the system will be in its most probable state and the entropy will be maximum."



Ludwig Boltzmann (1844 - 1906)

Statistical entropy:

$$S = -k \sum_i P_i \ln P_i$$

What is classical info:

\rightarrow Sun will rise on the east today.

\rightarrow We may have rainfall tonight at Chennai.

\rightarrow Statement 1 is a certain event (doesn't add to knowledge)

Information is all about probabilities

- Statement 3: One-third of the world's population is India+ China population.
- Statement 4: Two-third of the earth's surface is covered by water.
- Is there any difference between the information contents of these two statements?
- No! Given any person, our ignorance about his/her nationality is same as our ignorance about the character (viz. water or solid) of any given surface area of the Earth.
- So information is all about the probabilities of occurrences of different events: probability that a person (from all over the world) is (India + China) is $1/3$, probability that a surface area of the Earth is solid is equal to $1/3$, etc.

- Thus the amount of information about an event is the amount of ignorance (or, **uncertainty**) about that event.
- Ignorance increases with increase of the inverse of the probability p of occurrence of the event.
- Total amount of ignorance of two independent events is sum of the ignorance.
- So the amount of ignorance $H(p)$ should be an additive function of p .

Shannon entropy $H = - \sum p_i \log p_i$
 \Rightarrow We define $0 \log 0 = 0$

• **Example of Shannon's noiseless coding theorem**

Code 4 symbols $\{1, 2, 3, 4\}$ with probabilities $\{1/2, 1/4, 1/8, 1/8\}$.

Code without compression:

$$[1, 2, 3, 4] \xrightarrow{\text{without compr.}} [00, 01, 10, 11]$$

But, what happens if we use this code instead?

$$[1, 2, 3, 4] \xrightarrow{\text{with compr.}} [0, 10, 110, 111]$$

Average string length for the second code:

$$\langle \text{length} \rangle = \frac{1}{2} \cdot 1 + \frac{1}{4} \cdot 2 + \frac{1}{8} \cdot 3 + \frac{1}{8} \cdot 3 = \frac{7}{4} < 2$$

Note: $H\left(\frac{1}{2}, \frac{1}{4}, \frac{1}{8}, \frac{1}{8}\right) = -\frac{1}{2} \log_2 \frac{1}{2} - \frac{1}{4} \log_2 \frac{1}{4} - \frac{2}{8} \log_2 \frac{1}{8} = \frac{7}{4}$

Classical encodings don't work
necessarily work in quantum

| | | | |
|-----|-----|-----|-------|
| 0-2 | 2-4 | 4-5 | 10-11 |
|-----|-----|-----|-------|

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|-----|--|
| 2-5 | |
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| 5-10 | |
|------|--|

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| 10-13 | |
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