

Shor's Algorithm - In Depth Analysis

To Factor an odd integer N (Let's choose 15) :

1. Determine if the number n is a prime
2. Or a even number,
3. or an integer power of a prime number 2.
4. If it is 2 we will not use Shor's algorithm.
5. There are efficient classical methods.

6. Choose an integer q such that $N^2 < q < 2N^2$

7. Choose a random integer x such that $\text{GCD}(x, N) = 1$

at least 1 period
or sth like that

let's pick 256 \rightarrow 8 qubits

and
enough
qubits to
represent

- An important result from Number Theory:

$F(a) = x^a \bmod N$ is a periodic function in 'a'

- Choose $N = 15$ and $x = 7$ and we get the following:

We need
to find the
period

$$7^0 \bmod 15 = 1$$

$$7^1 \bmod 15 = 7$$

$$7^2 \bmod 15 = 4$$

$$7^3 \bmod 15 = 13$$

$$7^4 \bmod 15 = 1$$

⋮

Shor's Algorithm - In Depth Analysis

4. Create two quantum registers (these registers must also be entangled so that the collapse of the input register corresponds to the collapse of the output register)

- Input register: must contain enough qubits to represent numbers as large as $q-1$. up to 255, so we need 8 qubits
- Output register: must contain enough qubits to represent numbers as large as $N-1$. up to 14, so we need 4 qubits

5. Load the input register with an equally weighted superposition of all integers from 0 to $q-1$. **0 to 255**
6. Load the output register with all zeros.

The total state of the system at this point will be:

$$\frac{1}{\sqrt{256}} \sum_{a=0}^{255} |a, 000\rangle$$

Input
Register

Output
Register

Note: the comma here denotes that the registers are entangled

7. Apply the transformation $x^a \bmod N$ to each number in the input register, storing the result of each computation in the output register.

Input Register	$7^a \bmod 15$	Output Register
$ 0\rangle$	$7^0 \bmod 15$	1
$ 1\rangle$	$7^1 \bmod 15$	7
$ 2\rangle$	$7^2 \bmod 15$	4
$ 3\rangle$	$7^3 \bmod 15$	13
$ 4\rangle$	$7^4 \bmod 15$	1
$ 5\rangle$	$7^5 \bmod 15$	7
$ 6\rangle$	$7^6 \bmod 15$	4
$ 7\rangle$	$7^7 \bmod 15$	13

⋮

8. Now take a measurement on the output register. This will collapse the superposition to represent **just one** of the results of the transformation, let's call this value c .

Our output register will collapse to represent one of the following:

$$|1\rangle, |4\rangle, |7\rangle, \text{ or } |13\rangle$$

For sake of example, let's choose $|1\rangle$

Now things really get interesting !

9. Since the two registers are entangled, measuring the output register will have the effect of partially collapsing the input register into an **equal superposition** of each state between 0 and $q-1$ that yielded c (the value of the collapsed output register.)

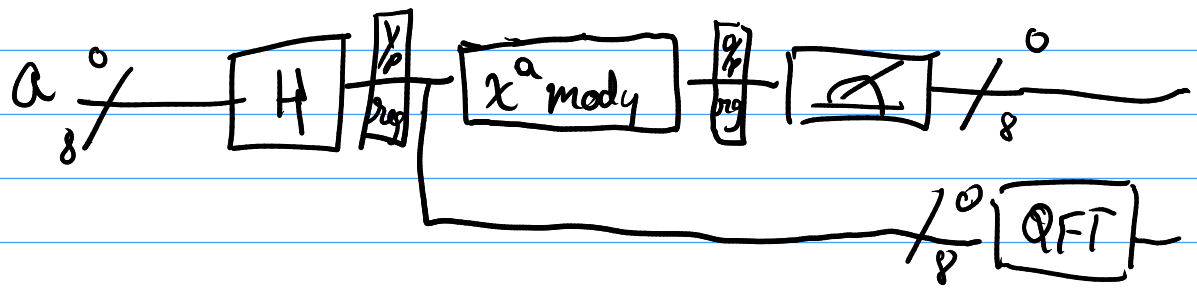
Since the output register collapsed to $|1\rangle$, the input register will partially collapse to:

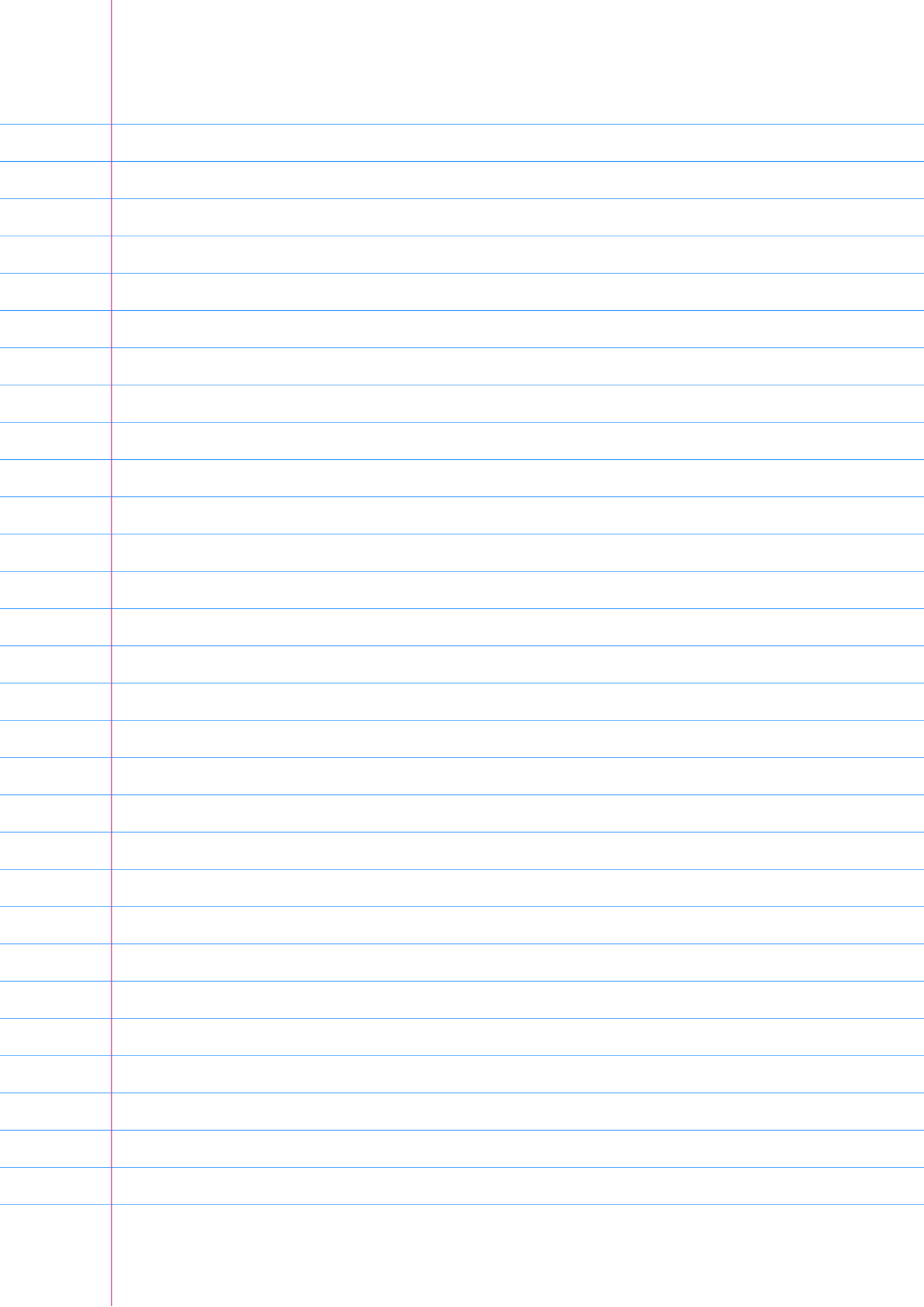
$$\frac{1}{\sqrt{64}} |0\rangle + \frac{1}{\sqrt{64}} |4\rangle + \frac{1}{\sqrt{64}} |8\rangle + \frac{1}{\sqrt{64}} |12\rangle, \dots$$

The probabilities in this case are $\frac{1}{64}$ since our register is now in an equal superposition of 64 values (0, 4, 8, ... 252)

We now apply the Quantum Fourier transform on the partially collapsed input register. The Fourier transform has the effect of taking a state $|a\rangle$ and transforming it into a state given by:

$$\frac{1}{\sqrt{q}} \sum_{c=0}^{q-1} |c\rangle * e^{2\pi i a c / q}$$





$$| \chi \rangle = \frac{1}{\sqrt{256}} [|0,1\rangle + |1,4\rangle + |2,7\rangle + |3,13\rangle + |4,1\rangle + |5,4\rangle + |6,7\rangle + |7,13\rangle + |8,1\rangle + |9,13\rangle + |10,1\rangle + |11,13\rangle + |12,7\rangle + |13,13\rangle + |14,1\rangle + |15,4\rangle]$$

Shor's Algorithm - The Factors :)

10. Now that we have the period, the factors of N can be determined by taking the greatest common divisor of N with respect to $x^{(P/2)} + 1$ and $x^{(P/2)} - 1$. The idea here is that this computation will be done on a classical computer.

We compute:

$$\text{Gcd}(7^{4/2} + 1, 15) = 5$$

$$\text{Gcd}(7^{4/2} - 1, 15) = 3$$

We have successfully factored 15!

$$\begin{aligned} a^P &= 1 \pmod{N} \\ (a^P - 1) &= 0 \pmod{N} \\ \underline{\underline{(a^{P/2} - 1)(a^{P/2} + 1)}} &= 0 \pmod{N} \end{aligned}$$

Particle
e gun
→

$|s_1\rangle$
 $|s_2\rangle$

wave
wave
particle
duality
"probability"
go to where
it's striking

$|e^- \text{ only} \rangle$
→

\angle
 \angle
 \angle
 \angle
 \angle
 \angle

will pinpoint it out,
where it's striking