

Tugas 2

Analisis Algoritma



Disusun oleh :

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1. Pencarian nilai maksimal

```
#include <iostream>
using namespace std;

int main()
{
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++)
    {
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }

    int maks = x[0];
    int i = 1;
    while (i <= n)
    {
        if (x[i] > maks)
            maks = x[i];
        i++;
    }
    cout << "Maksimum Number : " << maks << endl;

    return 0;
}
```

Kompleksitas Waktu

$\text{maks} \leftarrow x_1$ 1 kali

$i \leftarrow 2$ 1 kali

$\text{maks} \leftarrow x_i$ n kali

$i \leftarrow i + 1$ n kali

$$T(n) = 1 + 1 + n + n = 2n + 2$$

2. Sequential Search

```
#include <iostream>
using namespace std;

int main()
{
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++)
    {
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }

    int y;
    cout << "Masukkan yang dicari : ";
    cin >> y;

    int i = 0;
    bool found = false;
    int idx;
    while ((i < n) && (!found))
    {
        if (x[i] == y)
            found = true;
        else
            i++;
    }
    if (found)
        idx = i+1;
    else
        idx = 0;

    cout << "Yang dicari berada di urutan : " << idx << endl;

    return 0;
}
```

Kompleksitas waktu

Best Case :

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow I$	1 kali

$$T_{\min}(n) = 1 + 1 + 1 + 1 = 4$$

Average Case :

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$i \leftarrow i + 1$	$\frac{1}{2} n$ kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow I$	1 kali

$$T_{\text{avg}}(n) = 1 + 1 + \frac{1}{2} n + 1 + 1 = \frac{1}{2} n + 4$$

Worst Case :

$i \leftarrow 1$	1 kali
$\text{found} \leftarrow \text{false}$	1 kali
$i \leftarrow i + 1$	n kali
$\text{found} \leftarrow \text{true}$	1 kali
$\text{idx} \leftarrow I$	1 kali

$$T_{\max}(n) = 1 + 1 + n + 1 + 1 = n + 4$$

3. Binary Search

```
#include <iostream>
using namespace std;
```

```
int main()
{
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++)
    {
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }

    int y;
    cout << "Masukkan yang dicari : ";
    cin >> y;
```

```

int i = 0;
int j = n-1;
bool found = false;
int idx;
int mid;
while ((i <= j) && (!found))
{
    mid = (i + j)/2;
    if (x[mid] == y)
        found = true;
    else
    {
        if (x[mid] < y)
            i = mid + 1;
        else
            j = mid - 1;
    }
}

if (found)
    idx = mid+1;
else
    idx = 0;

cout << "Yang dicari berada di urutan : " << idx << endl;

return 0;
}

```

Kompleksitas waktu

Best Case :

$i \leftarrow 1$	1 kali
$j \leftarrow n$	1 kali
$found \leftarrow false$	1 kali
$mid \leftarrow (i + j) \text{ div } 2$	1 kali
$found \leftarrow true$	1 kali
$Idx \leftarrow mid$	1 kali

$$T_{min}(n) = 1 + 1 + 1 + 1 + 1 + 1 = 6$$

Average Case :

$i \leftarrow 1$	1 kali
$j \leftarrow n$	1 kali

$\text{found} \leftarrow \text{false}$ 1 kali
 $\text{mid} \leftarrow (i + j) \text{ div } 2$ $\frac{1}{2}n + 1$ kali
 $i \leftarrow \text{mid} + 1 \text{ or } j \leftarrow \text{mid} - 1$ $\frac{1}{2}n$ kali
 $\text{found} \leftarrow \text{true}$ 1 kali
 $\text{Idx} \leftarrow \text{mid}$ 1 kali

$$T_{avg}(n) = 1 + 1 + 1 + \frac{1}{2}n + 1 + \frac{1}{2}n + 1 + 1 = n + 6$$

Worst Case :

$i \leftarrow 1$ 1 kali
 $j \leftarrow n$ 1 kali
 $\text{found} \leftarrow \text{false}$ 1 kali
 $\text{mid} \leftarrow (i + j) \text{ div } 2$ $n + 1$ kali
 $i \leftarrow \text{mid} + 1 \text{ or } j \leftarrow \text{mid} - 1$ n kali
 $\text{found} \leftarrow \text{true}$ 1 kali
 $\text{Idx} \leftarrow \text{mid}$ 1 kali

$$T_{max}(n) = 1 + 1 + 1 + n + 1 + n + 1 + 1 = 2n + 6$$

4. Insertion Sort

```
#include <iostream>
using namespace std;
```

```
int main()
{
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++)
    {
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }
    cout << "Data Sebelum di Sorting : ";
    for (int i = 0; i < n; i++)
        cout << x[i] << " ";
    cout << endl;

    int insert;
    int j;
    for (int i = 1; i < n; i++)
    {
        insert = x[i];
        j = i-1;
```

```

        while ((j >= 0) && (x[j] > insert))
        {
            x[j+1] = x[j];
            j--;
        }
        x[j+1] = insert;
    }

    cout << "Data setelah di Sorting : ";
    for (int i = 0; i < n; i++)
        cout << x[i] << " ";

    return 0;
}

```

Kompleksitas waktu

Best Case :

for i \leftarrow 2 to n do	1 kali
insert \leftarrow xi	n kali
j \leftarrow i	n kali
x[j] = insert	n kali
$T_{min}(n) = 1 + n + n + n = 3n + 1$	

Average Case :

for i \leftarrow 2 to n do	1 kali
insert \leftarrow xi	n kali
j \leftarrow I	n kali
x[j] \leftarrow x[j-1]	n * 1/2 n kali
j \leftarrow j-1	n * 1/2 n kali
x[j] = insert	n kali

$$T_{avg}(n) = 1 + n + n + \frac{1}{2} n^2 + \frac{1}{2} n^2 + n = n^2 + 3n + 1$$

Worst Case :

for i \leftarrow 2 to n do	1 kali
insert \leftarrow xi	n kali
j \leftarrow i	n kali
x[j] \leftarrow x[j-1]	n * n kali
j \leftarrow j-1	n * n kali
x[j] = insert	n kali

$$T_{max}(n) = 1 + n + n + n^2 + n^2 + n = 2n^2 + 3n + 1$$

5. Selection Sort

```
#include <iostream>
using namespace std;

int main()
{
    int n;
    int x[10];
    cout << "Masukkan Jumlah Data : ";
    cin >> n;
    for (int i = 0; i < n; i++)
    {
        cout << "Masukkan Data ke - " << i+1 << " : ";
        cin >> x[i];
    }
    cout << "Data Sebelum di Sorting : ";
    for (int i = 0; i < n; i++)
        cout << x[i] << " ";
    cout << endl;

    int imaks;
    int temp;
    for (int i = n-1; i >= 1; i--)
    {
        imaks = 0;
        for (int j = 1; j <= i; j++)
        {
            if (x[j] > x[imaks])
                imaks = j;
        }
        temp = x[i];
        x[i] = x[imaks];
        x[imaks] = temp;
    }

    cout << "Data setelah di Sorting : ";
    for (int i = 0; i < n; i++)
        cout << x[i] << " ";

    return 0;
}
```


Kompleksitas waktu

Best Case :

for i \leftarrow n down to 2 do	1 kali
i maks \leftarrow 1	n kali
for j \leftarrow 2 to i do	n kali
i maks \leftarrow j	n*1 kali
temp \leftarrow xi	n kali
xi \leftarrow xi maks	n kali
xi maks \leftarrow temp	n kali

$$T_{min}(n) = 1 + n + n + n * 1 + n + n + n = 6n + 1$$

Average Case :

for i \leftarrow n down to 2 do	1 kali
imaks \leftarrow 1	n kali
for j \leftarrow 2 to i do	n kali
imaks \leftarrow j	n * $\frac{1}{2}$ n kali
temp \leftarrow xi	n kali
xi \leftarrow xi maks	n kali
xi maks \leftarrow temp	n kali

$$T_{avg}(n) = 1 + n + n + \frac{1}{2}n^2 + n + n + n = \frac{1}{2}n^2 + 5n + 1$$

Worst Case :

for i \leftarrow n down to 2 do	1 kali
maks \leftarrow 1	n kali
for j \leftarrow 2 to i do	n kali
imaks \leftarrow j	n * n kali
temp \leftarrow xi	n kali
xi \leftarrow xi maks	n kali
ximaks \leftarrow temp	n kali

$$T_{max}(n) = 1 + n + n + n^2 + n + n + n = n^2 + 5n + 1$$