## Inducing Game Rules from Varying Quality Game Play

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March 19, 2020

#### Abstract

General Game Playing (GGP) is a framework in which an artificial intelligence program is required to play a variety of games successfully. The framework includes repositories of game descriptions written in a logic programming language. These game descriptions can be used separately as an interesting Inductive Logic Programming problem.

The Inductive General Game Playing (IGGP) problem challenges ILP systems to learn these GGP game rules by watching the game being played. This challenge motivates investigation into ways to better induce general rules from specific examples and provides a good test bed for existing ILP systems.

Existing work on IGGP has always assumed that the game player being observed makes random moves. This is not representative of how a human learns to play a game, to learn to play chess we watch someone who is playing to win.

To address this limitation, in this paper we analyse the effect of using optimal verses random gameplay traces as well as the effect of varying the number of traces in the training set.

We use the General Game Playing competition winner in 2014, Sancho, to generate optimal game traces for a large number of games and the systems, Metagol, Aleph and ILASP are trained and tested with combinations of optimal and random data.

### Introduction

General Game Playing (GGP) is a framework in which artificial intelligence programs are required to play a large number of games successfully.[7]. Games in the framework range in both number of players and complexity; from the single player Eight Puzzle to the six player Chinese Checkers, and from the relatively simple Rock Paper Scissors to Chess[1]. Every year the GGP competition takes place, aiming to find the best GGP program. A consequence of this framework is large repositories of game descriptions written in the Game Description Language (GDL), a logic programming language built for describing games as state machines[13]. A logic programming language being any that is mainly based on formal logic, such as Prolog.

These GDL game descriptions form the basis for the Inductive General Game Playing (IGGP) problem. The task is an inversion of the GGP problem. Rather than taking game rules using them to play the game in IGGP the aim is to learn the rules from observations of gameplay. Cropper et al. define the IGGP problem in their 2019 paper; given a set of gameplay observations the goal is to induce the rules of the game[5]. The games are those from the GGP competition problem set, specified in GDL, meaning they are widely varied in complexity.

A way of solving this problem is Inductive Logic Programming (ILP).

In Inductive Logic Programming, machine learning systems are tasked with learning logic programs given some background knowledge and a set of values for which the programs are true or false. The ILP system will derive a hypothesis, a logic program that when combined with the background knowledge, entails all of the positive and none of the negative examples[12]. In the IGGP paper it is shown that the problem is hard for current ILP systems, with on average only 40% of the rules being learned by the best performing systems.

In this paper we try and increase the success rate of the ILP systems. We

use the IGGP framework to evaluate the ability of ILP agents to correctly induce the rules of a game given different sets of gameplay observations - optimal gameplay verses random gameplay as well as combinations of the two. So far no work has been done on how the observations given as training data to the systems affect their ability to learn games. It is not obvious whether random or optimal gameplay would be best. When learning the rules of chess would a human rather watch moves being made randomly, or a match between two grandmasters? It is not an easy question to answer. Both situations will result in a restricted view of the game, with curtain situations never occurring in each one. It is possible watching both types of play viewed together would give a better understanding which is why we test on a mixture of optimal and random traces together.

- Q1 Does varying the quality of game traces influence the ability for learners to solve the IGGP problems? Specifically, does the quality of game place affect predictive accuracy?
- Q2 Does varying the amount of game traces influence the ability for learners to solve the IGGP problems? Specifically, does the quality of game place affect predictive accuracy?
- Q3 Can we improve the performance of a learner by mixing the quality of traces?

We would also expect models trained on the same distribution as they are tested on to perform best since it is generally accepted that the accuracy of a model increases the closer the test data distribution is to the training data distribution.[11]. However, Gonzales [8] suggests that a system trained on a different domain to the one it is tested can outperform a system trained and tested on the same domain. In his 2010 paper Ben-David shows that training data taken from multiple different domains can in fact give lower error on testing data that training data taken from any single domain, including the testing domain [2]. It is not clear in our case what selection of training data will result in the most effect learning.

In machine learning the probably approximately correct learning framework (PAC) is a mathematical model that can be applied to machine learning systems for analysis. It gives an upper bound on the number of training examples needed for a learner to probably (with probability at least  $[1 - \delta]$ ) learn a hypothesis that is approximately (within error  $\epsilon$ )correct[11]. In section ? we test to find the number of examples beyond which the reduction in  $\epsilon$  and  $\delta$  becomes negligible.

We will train the ILP systems on different combinations of optimal and random data as well as testing them on each individually. In section ? we test to find the optimal selection of distributions to take our training data from.

#### Contributions

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 $\bullet$  We experimentally explore the research questions (Q1, Q2, and Q3). Our results show ...

### Related work

AC: Use paragraphs

### 2.1 GGP

General game playing is a framework for evaluating an agenet's general intelligence accross a wide range of tasks [5]. The idea is that agents are able to accept declarative descriptions of arbitrary games at run time and are able to use such descriptions to play those games effectively. All the games are finite, discrete, deterministic multi-player games of complete information. The games can vary in number of players, dimensions and complexy. For example games such as rock paper scissons have 0 dimensions and only 10 rules in the given GGP ruleset, more complex games such as checkers has 52 rules and is 2 dimensional. There are also single player games such as eight-puzzle or Fizz Buzz. The agents play games selected from an exsisting database which lists their descriptions described in the Game Description Language (GDL). GDL is a language based on first order logic described in section 2.1.1. The list of games along with their descriptions is available online[1]. Matches in the GGP framework take place through an online framework called the Gamemaster. The agents connect to an online Game Manager (part of the Gamemaster) service which arbitrates indavidual matches. The connecting agents send several pieces of information to the Game Manger incliding a game that is already know by the manager and the required number of players. The match can then be started by pressing a start button, when pressed the agents will recive a Start message including the role of the player (e.g. black or white in chess) and a description on the game in GDL. The Game Manger communicates all instructions to the players through HTTP[7]. In 2005 an annual International General Game Playing Competition (IGGPC)

was set up which still runs to this day. Each year hopeful participants pit GGP agents against one another to find the best one. The competitors take part in a series of rounds of increasing complexity. The agent that wins the most games in these rounds is the winner. The 2014 winner Sancho is used in this paper to generate optimal game traces for the IGGP task. The Game description language is used in the IGGP as an example of the rules to learn and the rules to use when generating examples. It is a type of logical programming language.

#### 2.1.1 GDL

Game definition language is the formal language used in the GGP competition to specify the rules of the games.[13] It is suited to specifying game rules because:

- It is a purely declarative language
- It has restrictions to ensure that all questions of logical entailment are decidable
- There are some reserved words, which tailor the language to the task of defining games

These descriptions define the games in terms of a set of true facts capturing the information needed to give the following predicates:

- The initial game state
- The goal state
- The terminal state

In addition, logical rules are used to describe the following:

- The legal moves for a given player and state
- The next state for a given player state and move
- The termination and goal conditions

In the IGGP problem the tasks are the rules for the goal, the next state, the legal states and the terminal state.

### 2.2 Logical Programming

Logic programming is a programming paradigm based on formal logic. Programs are made up of facts and rules. Rules are made up of two parts: the head and the body. They can be read as logical implications where the conjunction of all the elements in the body imply the head. The syntax is different for different logical programming languages but the head is usually written before the body. A fact is simply a rule without a body, that is, a statement that is taken as true. The compiler takes queries and returns weather they are true of false. If there are free variables in the query the compiler assigns them values for which the query is true. Logical programming is good for symbolic non-numeric computation. It is well suited to solving problems that involve well defined objects and relations between them, such as a GGP game.

### 2.3 Prolog

Prolog is one of the most popular and well used logical programming languages. The syntax in Prolog for rules and facts is fairly intuitive. A simple fact might be son(bob, alice).. This tells us that the atom alice relates the atom bob in the son relation. A rule is written parent(X,Y) := son(Y,X) which means if we have the relation son of Y and X then X relates Y in parent. The symbol: is the same as reverse implication ( $\Leftarrow$ ). We can query a program in the Prolog environment. If we typed in "parent(alice,bob)." then it would return true because we have son(bob, alice) and the rule which tells us that if X is a son of Y else then Y is a parent of X so alice is a parent of bob. Predicates can be conjoined with the comma "," (e.g. a(X,Y,Z) := b(X,Y), c(Z,Y)). Disjunction is expressed with the semicolon";" (e.g. a(X,Y,Z) := b(X,Y); c(Z,Y)).

Prolog answers queries using a process know as proof seach and unification. The unification process is similar to logical unification, two terms unify if they are the same term or if they contain variables that can be instantiated with terms in such a way that the new terms are equal. For example the terms  $\mathtt{name}(\mathtt{bob})$ . and  $\mathtt{name}(\mathtt{X})$ . will unify with  $\mathtt{X} = \mathtt{bob}$ . The Prolog ISO defines the Herbrand Algorithm for unification [9]. A Prolog query is a set of goals. The Prolog System decides weather they are satisfiable or not. When a query is asked the of the Prolog system it executes something similar to the following algorithm. The exact implementations vary among Prolog systems however they are all roughly this algorithm.

the program listing is wearched for a term to unify with. The listing is

```
Algorithm 1: Execute Prolog Goals
 Input: A list of goals GoalList = G_1, G_2, ..., G_M
 Function execute (Program, GoalList, Success):
     if empty(GoalList) then
        Success \leftarrow true
     end
     while not empty(GoalList) do
         Goal \leftarrow head(GoalList);
         OtherGoals \leftarrow tail(GoalList);
         Satisfied \leftarrow false;
         while not Satisfied and there are more clauses in the
          program do
            Let next clause in the program be H \vdash B_1, ..., B_n;
            Construct a variant of this clause H' \vdash B'_1, ..., B'_n;
            match(Goal, H', MatchOK, Instant);
            if MatchOK then
                NewGoals \leftarrow append([B'_1,...,B'_n],OtherGoals);
                NewGoals \leftarrow substitue(Instant, NewGoals);
                execute(Program, NewGoals, Satisfied);
            end
         end
         Success \leftarrow Satisfied;
     end
```

searched in the order it is written in. When Prolog finds a matching rule it then attempts to sequentially unify the terms of the body using the same method. If the rule has no body then the variables are assigned and the terms unify. If Prolog fails to unify two terms then it backtracks, assignes differnt values to [3]

### 2.4 IGGP

The IGGP problem uses the GDL game descriptions from the GGP framework to construct the IGGP dataset containing 50 games. The IGGP problem itself is described in section ??. A machanism is also provided by Andrew Cropper et al. [5] to turn the GGP games into new IGGP tasks. This mechanism has been modified and added to to generate optimal traces.

#### 2.5 ILP

Inductive logic programming is a form of machine learning that uses logic programming to represent examples, background knowledge, and learned programs[6]. To learn the program is supplied with positive examples, negative examples and the background knowledge. In the general inductive setting we are provided with three languages.

- $\mathcal{L}_O$ : the language of observations (positive and negative examples)
- $\mathcal{L}_B$ : the language of background knowledge
- $\mathcal{L}_H$ : the language of hypotheses

The general inductive problem is as follows: given a consistent set of examples or observations  $O \subseteq \mathcal{L}_O$  and consistent background knowledge  $B \subseteq \mathcal{L}_B$  find an hypothesis  $H \in \mathcal{L}_H$  such that

$$B \wedge H \models O$$

[12] That is that the generated hypothesis and the background knowledge imply the positive examples and not imply the negative examples.

Example task (grandparent relation)

ILP systems generally regard ILP as a search problem. The search space is the set of well formed hypothesis. Often a set of inference rules are applied to the starting hypothesis, the new hypothesise are then pruned and expanded according to how often  $B \wedge H \models O$  in the observations O. When all the observations are implied then a correct program has been found.

### 2.6 Machine learning

Machine learning theory contains a lot of work on the general setting of the problem we are considering, specifically, it covers how the relation between the training data and the testing data can affect the accuracy of the hypothesis as well as how the amount of training data can affect the result. In this paper

(in PAC) The cardinality of this example set must be at most a polynomial function of the size of the vocabulary used in constructing hypotheses.

More examples allows for longer programs therefore better programs?

### 2.6.1 Metagol?

?Should I add a section on how metagol works here as an example?

To evaluate ILP systems the learning tasks they are given need to be small enough that it is feasible for them to learn the programs, but also diverse enough that we test the full capability of the system. The Inductive general game playing framework provides this. To understand IGGP we first need to explain the General game playing framework.

## Logical setting

- What precisely is the problem? - You can paraphrase a lot from the IGGP paper

The IGGP problem is defined in Andrew Croppers paper Inductive general game playing [5]. Much like the problem of ILP 2.5 the problem setting consists of examples about the truth or falsity of a formula F and a hypothesis H which covers F if H entails F. We assume the languages of:

- $\bullet$   $\mathscr{E}$  the language of examples (observations)
- $\mathcal{B}$  the language of background knowledge
- $\mathcal{H}$  the language of hypotheses

Each of these languages can be see as a subset of those described in 2.5. All the predicates involved in the task are taken from the GDL descriptions of games in the Stanford GGP\* library. A lot of the atoms in the descriptions are not function-free, that is, they are nested predicates. For example true(count(9)). We flatten all of these to single, non nested predicates, i.e. true\_count(9). This is done because some ILP systems do not support function symbols. We can therefore assume that both & and B are function-free. The language of hypotheses & can be assumed to consist of datalog programs with stratified negation as described here[14]. Stratified negation is not necessary but in practice allows significantly more concise programs, and thus often makes the learning task computationally easier. We first define an IGGP input the use it to define the IGGP problem. An IGGP input needs to capture the idea of an observation about a game. The input is based on the general input for the Logical induction problem of section 2.5 since this is a sub problem of it.

The IGGP Input: An input  $\Delta$  is a set of triples  $\{(B_i, E_i^+, E_i^-)\}_m^{i=1}$  where

- $B_i \subset \mathcal{B}$  represents background knowledge
- $E_i^+ \subseteq \mathscr{E}$  and  $E_i^- \subseteq \mathscr{E}$  represent positive and negative examples respectively

An IGGP input forms the IGGP problem:

The IGGP Problem: Given an IGGP input  $\Delta$ , the IGGP problem is to return a hypothesis  $H \in \mathscr{H}$  such that for all  $(B_i, E_i^+, E_i^-) \in \Delta$  it holds that  $H \cup B_i \vDash E_i^+$  and  $H \cup B_i \nvDash E_i$ 

#### Problem Setting

Let the accuracy of a set I of ILP systems in problem setting be defined as the mean of the percentage accuracy of each of them when tested on a given set of examples.

- Given an set of ILP systems I, for what selection of game traces  $\Delta$  combined from an optimal gameplay distribution and a random gameplay distribution are the systems most accurate when solving the IGGP problem. The accuracy of I when solving the IGGP problem will be tested by evaluating each  $i \in I$  with data taken from a 50/50 combination of both optimal and random distributions.

We will also be investigating how the size of the trace set  $|\Delta|$  affects the accuracy of the hypothesis H.

#### 3.0.1 Sancho

To generate optimal gameplay we decided that the best approach was to use a previous winner of the GGP competition. Since the aim of the competition is to find the program that performs best at the set of games used in the IGGP problem we conclude that there is no better way to generate a comprehensive set of game traces. The winner of the 2014 GGP competition 'Sancho' was selected <sup>1</sup> since they are the most recent winner to have published their code <sup>2</sup>. Some small modifications to the information logged by the game server are made but otherwise the code is unchanged.

The core of Sancho is the Monte Carlo tree search (MCTS) algorithm.

#### Monte Carlo Tree Search

Given a game state the basic MCTS will return the most promising next move. The algorithm achives this by simulating random playouts of the

<sup>&</sup>lt;sup>1</sup>http://ggp.stanford.edu/iggpc/winners.php accessed on 12/03/2020

<sup>&</sup>lt;sup>2</sup>http://sanchoggp.github.io/sancho-ggp/ access 12/03/2020

game many times. The technique was developed for computer Go but has since been applied to play a wide range of games effectively including board games and video games[15][4].

The use of random simulation to evaluate game states is a powerful tool. Information about the game such as heuristics for evaluating non terminal states are not needed at all, the rules of the game are enough on their own. In the case of the GGP problem this is ideal. The rules of the game are only revealed shortly before the game is played meaning deriving effective heuristics is a hard problem.

In our case all games being played are sequential, finite and discrete so we only need to consider MCTS for this case.

The MCTS is a tree search algorithm, the tree being searched is the game tree. A game tree being a tree made up of nodes representing states of the game and the children of each node being all the states that can be moved to from that state. The leaves of this tree are the terminal states. The search is a sequence of traversals of the game tree. A traversal is a path that starts at the root node and continues down until it reaches a node that has at least one unvisited child. One of these unvisited children are then chosen to be the start state for a simulation of the rest of the game. The simulation chooses random moves, playing the game out to a terminal state. The result of the simulation is propagated back from the node it started at all the way to the root node to update statistics attached to each node. These statistics are used to choose future paths to traverse so more promising moves are investigated more.

Exactly how the tree search algorithm chooses new nodes to simulate is where the complexity lies. The node chosen by traversal of the tree should strike a balance between exploiting promising nodes and exploring nodes with few simulations. The Upper Confidence Bound for trees algorithm was designed to do exactly this[10]. The algorithm chooses a child node for each state based on the UTC formula. The formulae can be written as:

$$UTC(v_i) = \frac{Q(v_i)}{N(v_i)} + c\sqrt{\frac{ln(N(v_i))}{N(v_i)}}$$

Where we have:

- $v_i$  is node v after move i
- Q(v) is the number of winning simulations that have taken place below it.
- N(v) is the number of simulations that have taken place below it

The function is the sum of two components.  $\frac{Q(v_i)}{N(v_i)}$  is called the exploitation component. It is a ratio of winning to loosing simulation that resulted from making this move. This encourages traversal of promising nodes that have a high win rate. However, if only this factor was used the algorithm would quickly find a winning node and only explore that hence we need a second component: the exploration factor.  $c\sqrt{\frac{\ln(N(v_i))}{N(v_i)}}$  favours nodes that have not yet been explored. The value c is a constant that balances the two components, it can be adjusted depending on the use case.

Sancho makes a few modifications to MCTS<sup>3</sup>, the main one being adding of heuristics. The tree is also replaced with a more general graph which allows for transitions between lines of play without duplication of states. There are other modifications to increase efficiency.

In GGP matches a period of time before the match is given in which to do pre match calculations. Sancho uses this period to derive basic heuristics and optimise the value of c in the formulae. A heuristic should take a game state or move and return a value based on how promising that state or move is in relation to the goal state. To take the heuristics into account when choosing the next state to explore each node (state) v is seeded with a heuristic value on creation.

The identification of possible heuristics takes place in two stages. The first of these is static analysis of the game rules. This static analysis identifies possible heuristics that can be applied to the current game. These include things like piece capture: if certain rules indicate capture of a piece these can be selected against, or numeric quantity detection: a number in the state acts as a heuristic (like number of coins a player has). The second stage is simulation of the game. Many (possibly tens of thousands) of full simulations of the game are run. After the simulations are complete a correlation coefficient is calculated between each candidate heuristic's observed values and the eventual score achieved in the game. Those heuristics that show correlation above a fixed threshold are then enabled, and will be used during play.

 $<sup>^3 \</sup>rm https://sanchoggp.blogspot.com/2014/05/what-is-sancho.html accessed on <math display="inline">15/03/2020$ 

## **Implementations**

### 4.1 Eight Puzzle

To write an optimal player for Eight Puzzle I decide to use an approach based on best first search.

#### 4.1.1 Best first search

Best first search generates a graph of the game states, expanding the most promising nodes first according to a heuristic for how close to the goal state the current state is. When the goal node is reached the search gives the path from the root node to it, i.e. the list of moves made to get to the goal. Since Best first search is a well know and often used algorithm in Prolog we decided to use a standard implementation from Bratko[3]. This decision was made to avoid needless mistakes and inefficiency that would have come with a new independent implementation of this algorithm.

Best First search is a generic algorithm that needs certain problem specific predicates to be implemented for it to function. These predicates are:

- The successor predicate s(Node, Node1, Cost): This predicate is true if there is an arc costing Cost between state Node and state Node1. In Eight Puzzle we set the cost of all arcs to be 1.
- The Goal predicate goal (Node): This is true if the state Node is a goal state.
- The heuristic h(Node, H): This is a relation that relates a state Node to a value H that is a heuristic for how close to the state is to the goal state.

To represent the board I decided to use a 9 element list with b representing the blank tile INSERT IMAGE OF GOAL BOARD

```
[1,2,3,4,5,6,7,8,b]
```

. I defined some useful helper predicates to start with. The Manhatten distance between two tiles between two tiles is the difference in X coordinates plus the difference in Y coordinates, I wrote a predicate to calculate this relation:

```
mandist(X1-Y1, X2-Y2,D) :-
    diff(X1, X2,Dx),
    diff(Y1, Y2,Dy),
    D is Dx+Dy.
```

I also wrote predicates to calculate the coordinates of a tile from its position in the list, one that relates two tiles if they are next to each other on the grid and one that gives the position in the list of the blank tile.

#### Successor

To define the successor predicate I decided the best way to do it was to relate two boards by swapping the blank tile with its neighbours. I first wrote a function swap:

```
swap(I,J,L1,L3) :-
    same_length(L1,L3),
    append(BeforeI,[AtI|PastI],L1),
    append(BeforeI,[AtJ|PastI],L2),
    append(BeforeJ,[AtJ|PastJ],L2),
    append(BeforeJ,[AtI|PastJ],L3),
    length(BeforeI,I),
    length(BeforeJ,J).
```

This function only uses built in predicates that all do the obvious thing. In the first two appends this rule takes L1 and replaces the element I with the element J to make list L2 then in the second two takes list L2 and replaces J with I to make list L3. The variable AtI and AtJ will be instantiated with the elements at index I and J because of the restriction on the length of the sublist before them at the end of the rule. The successor function is defined as:

```
s(B1,B2,1) :-
member(N,B1),
```

```
nth0(NI,B1,N),
blank_pos(B1,BI),
neighbor(BI,NI),
swap(BI,NI,B1,B2).
```

Since all arc (move) costs in my version of Eight Puzzle are 1 the predicate is only true when Cost is 1. member(N,B1) is true when N is a member of B1. Here it restricts N to values on the board. nth0(NI,B1,N) sets NI to the index of the value N in the list B1, i.e. it tells us the position on the board of the value we are considering. blank\_pos(B1,BI) sets BI to the position on the board of the blank. neighbor(BI,NI) is true if the positions BI and NI are next to each other on the board. If all of these are true we then instantiate B2 with the new board by swapping the blank and its neighbour.

#### Goal

Defining the goal was actually very simple, I added the predicate:

```
goal([1,2,3,4,5,6,7,8,b])
```

This is all that is needed to define the goal state.

#### Heuristic

A good heuristic for Eight Puzzle needs to capture an estimate of the distance of the current state from the goal. I initially decided to use the sum of the Manhattan distances of each tile from its position in the goal state. To calculate this I wrote a predicate totdist(L,Goal,N) which relates a board L, a goal Goal and the sum of the Manhattan distances N.

```
totdist(L,Goal,N) :-
    totdist1(L,Goal,N,0).

totdist1([],_Goal,0,_Pos).

totdist1([b|T],Goal,N,Pos) :-
    !,Pos1 is Pos + 1,
    totdist1(T,Goal,N,Pos1).

totdist1([V|T],Goal,N,Pos) :-
    nth0(I,Goal,V),
    coord(Pos,X1-Y1),
    coord(I,X2-Y2),
```

```
mandist(X1-Y1,X2-Y2,D),
Pos1 is Pos + 1,
totdist1(T,Goal,N1,Pos1),
N is N1 + D.
```

My implementation of this predicate iterates through the list working out the Manhattan distance for each tile. When iterating through the list the way to keep track of how far through it you are (the position on the board you are currently looking at) is use a separate variable as a counter. Here I used Pos. After testing my program with this heuristic I found that it took about 6 seconds to come up with the first solution. I decided that I could probably do better if I had another heuristic that I combined with the Manhattan distance. I decided to use the number of tiles that were out of the row and/or column they were in. I wrote a predicate outofpos(B,Goal,N,Pos) that relates the board and the goal to the number of tiles out column plus the number of tiles out of row.

```
outofpos([],_Goal,0,_Pos).
outofpos([b|T],Goal,N,Pos) :- !,
    Pos1 is Pos + 1,
    outofpos(T,Goal,N,Pos1).
outofpos([V|T],Goal,N,Pos) :-
    coord(Pos,X1-Y1),
    nth0(PosG, Goal, V),
    coord(PosG, X2-Y2),
    Pos1 is Pos + 1,
    outofpos(T,Goal,N1,Pos1),
         xandy(X1 = X2, Y1 = Y2) \rightarrow
         N is N1
        xory(X1 = X2, Y1 = Y2) \rightarrow
        N is N1 + 1
        N is N1 + 2
    ).
xandy(X,Y) := X,Y.
xory(X,Y) := X;Y.
```

This predicate iterates through the board and for each tile looks at the place it is on the goal board and checks which coordinates match. If both match then it is the right column and row so the total remains the same, otherwise it is incremented. I eventually found that the heuristic the worked the fastest was Manhattan distance total + 2 \* number of tiles out of position. I wrote this predicate for the heuristic:

```
h(B,H) :-
   goal(Goal),
   totdist(B,Goal,D),
   outofpos(B,Goal,N,0),
   H is D + (2*N).
```

### 4.2 Noughts and crosses

# **Experiments**

- 5.1 Optimal Play vs Random with fixed sample size
- 5.2 Mixed datasets (50/50 random and optimal)
- 5.2.1 Results

### Random training and testing

metagol

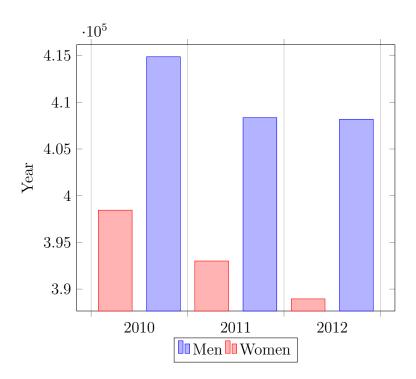
metagoi				
game	next	goal	legal	terminal
alquerque	51	50	50	100
asylum	50	50	50	100
$battle\_of\_numbers$	54	50	50	100
breakthrough	50	50	50	100
buttons_and_lights	50	50	50	100
duikoshi	50	50	50	100
$eight\_puzzle$	55	50	50	100
freeforall	54	50	50	100
hexforthree	51	50	50	100
horseshoe	57	50	50	100
hunter	52	50	50	100
$knights\_tour$	50	50	50	100
kono	52	50	50	100
pentago	51	50	50	50
pilgrimage	50	50	50	100
rainbow	50	50	50	100
sudoku	50	50	50	50
sukoshi	50	50	50	100
$untwisty\_corridor$	62	50	50	100
1 1				

aleph

game	next	goal	legal	terminal
alquerque	56	95	50	100
asylum	52	99	50	100
battle_of_numbers	68	98	76	100
breakthrough	51	75	50	100
buttons_and_lights	100	75 75	50	100
duikoshi	90	83	90	100
eight_puzzle	96	99	84	100
freeforall	98	98	76	100
hexforthree	67	99	72	100
horseshoe	93	99 97	97	100
hunter	95	98	99	100
knights_tour	98	99	99	100
kono	94	98	95	100
pentago	98	83	98	100
pentago pilgrimage	98 97	98	83	100
rainbow	93	96 75	70	100
sudoku	95	75 75	50	100
sukoshi	93 87	75 75	90	100
untwisty_corridor	100	75 75	50	100
	100	10	- 50	100
specialised ilasp				
specialised ilasp	next	goal	legal	terminal
game	next 46	goal 90	legal 99	terminal
game alquerque	46	90	99	100
game alquerque asylum	46 69	90 99	99 83	100 100
game alquerque asylum battle_of_numbers	46 69 54	90 99 97	99 83 88	100 100 100
game alquerque asylum battle_of_numbers breakthrough	46 69 54 56	90 99 97 50	99 83 88 100	100 100 100 50
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights	46 69 54 56 100	90 99 97 50 50	99 83 88 100 0	100 100 100 50 100
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi	46 69 54 56 100 85	90 99 97 50 50 66	99 83 88 100 0 60	100 100 100 50 100
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi eight_puzzle	46 69 54 56 100 85 96	90 99 97 50 50 66 98	99 83 88 100 0 60 75	100 100 100 50 100 100
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi eight_puzzle freeforall	46 69 54 56 100 85 96 76	90 99 97 50 50 66 98 71	99 83 88 100 0 60 75 85	100 100 100 50 100 100 100
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi eight_puzzle freeforall hexforthree	46 69 54 56 100 85 96 76 83	90 99 97 50 50 66 98 71 100	99 83 88 100 0 60 75 85 100	100 100 100 50 100 100 100 100
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi eight_puzzle freeforall hexforthree horseshoe	46 69 54 56 100 85 96 76 83 93	90 99 97 50 50 66 98 71 100 95	99 83 88 100 0 60 75 85 100 96	100 100 100 50 100 100 100 100 100
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi eight_puzzle freeforall hexforthree horseshoe hunter	46 69 54 56 100 85 96 76 83 93	90 99 97 50 50 66 98 71 100 95	99 83 88 100 0 60 75 85 100 96 83	100 100 100 50 100 100 100 100 100
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi eight_puzzle freeforall hexforthree horseshoe hunter knights_tour	46 69 54 56 100 85 96 76 83 93 71 83	90 99 97 50 50 66 98 71 100 95 99	99 83 88 100 0 60 75 85 100 96 83 95	100 100 100 50 100 100 100 100 100 100
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi eight_puzzle freeforall hexforthree horseshoe hunter knights_tour kono	46 69 54 56 100 85 96 76 83 93 71 83 83	90 99 97 50 50 66 98 71 100 95 99 98	99 83 88 100 0 60 75 85 100 96 83 95	100 100 100 50 100 100 100 100 100 100 1
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi eight_puzzle freeforall hexforthree horseshoe hunter knights_tour kono pentago	46 69 54 56 100 85 96 76 83 93 71 83 83 99	90 99 97 50 50 66 98 71 100 95 99 98 97 42	99 83 88 100 0 60 75 85 100 96 83 95 98 83	100 100 100 50 100 100 100 100 100 100 1
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi eight_puzzle freeforall hexforthree horseshoe hunter knights_tour kono pentago pilgrimage	46 69 54 56 100 85 96 76 83 93 71 83 83	90 99 97 50 50 66 98 71 100 95 99 98 97 42 96	99 83 88 100 0 60 75 85 100 96 83 95 98 83 89	100 100 100 50 100 100 100 100 100 100 1
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi eight_puzzle freeforall hexforthree horseshoe hunter knights_tour kono pentago	46 69 54 56 100 85 96 76 83 93 71 83 83 99	90 99 97 50 50 66 98 71 100 95 99 98 97 42 96 50	99 83 88 100 0 60 75 85 100 96 83 95 98 83 89 27	100 100 100 50 100 100 100 100 100 100 1
game alquerque asylum battle_of_numbers breakthrough buttons_and_lights duikoshi eight_puzzle freeforall hexforthree horseshoe hunter knights_tour kono pentago pilgrimage rainbow	46 69 54 56 100 85 96 76 83 93 71 83 83 99 92	90 99 97 50 50 66 98 71 100 95 99 98 97 42 96	99 83 88 100 0 60 75 85 100 96 83 95 98 83 89	100 100 100 50 100 100 100 100 100 100 1

### 5.2.2 Optimal res

game	predicate	metagol	aleph	specialised_ilasp
alquerque	next	51	90	47
alquerque	goal	50	95	90
alquerque	legal	50	50	99
alquerque	terminal	100	100	100
asylum	next	50	52	49
asylum	goal	50	99	99
asylum	legal	50	50	99
asylum	terminal	100	100	100
$battle\_of\_numbers$	next	53	65	75
$battle\_of\_numbers$	goal	50	98	97
$battle\_of\_numbers$	legal	50	76	88
$battle\_of\_numbers$	terminal	100	100	100
breakthrough	next	50	90	90
breakthrough	goal	50	75	50
breakthrough	legal	50	50	87
breakthrough	terminal	100	100	65
buttons_and_lights	next	50	100	100
buttons_and_lights	goal	50	75	50
buttons_and_lights	legal	50	50	0
buttons_and_lights	terminal	100	100	100
duikoshi	next	50	90	87
duikoshi	goal	50	83	66
duikoshi	legal	50	91	69
duikoshi	terminal	100	100	100
eight_puzzle	next	55	96	88
eight_puzzle	goal	50	99	50
eight_puzzle	legal	50	84	68
eight_puzzle	terminal	100	100	100
freeforall	next	55	98	74
freeforall	goal	50	98	71
freeforall	legal	50	76	97
freeforall	terminal	100	100	100
hexforthree	next	51	67	83
hexforthree	goal	50	99	100
hexforthree	legal	50	71	100
hexforthree	terminal	100	100	100
horseshoe	next	57	93	93
horseshoe	goal	50	97	95
horseshoe	legal	50	98	100
horseshoe	terminal	$100 \\ 22_{52}$	100	100
hunter	next	$22_{52}^{00}$	95	74
hunter	goal	50	98	99
hunter	legal	50	99	-100
hunter	terminal	100	100	100
knights_tour	next	50	93	50
knights_tour	goal	50 50	99	99
knights_tour	goar legal	50 50	99 99	99



5.3 Optimal Play vs Random with varying sample size

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