

# SIC/XE Assembler

## Names

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## **Overview**

The SIC machine has basic addressing, storing most memory addresses in hexadecimal integer format. Similar to most modern computing systems, the SIC architecture stores all data in binary and uses the two's complement to represent negative values at the machine level. Memory storage in SIC consists of 8-bit bytes, and all memory addresses in SIC are byte addresses. Any three consecutive bytes form a 24-bit 'word' value, addressed by the location of the lowest numbered byte in the word value. There is also a more complicated machine built on top of SIC called the Simplified Instruction Computer with Extra Equipment (SIC/XE). The XE expansion of SIC adds a 48-bit floating point data type, an additional memory addressing mode, and extra memory to the original machine. All SIC assembly code is upwards compatible with SIC/XE.

## **Goals**

The project is to implement **One Pass SIC/XE Assembler**. assembler, producing code for the absolute loader.

# **Requirements Specification**

- 1. Reading the source lines from a file.
- 2. A parser that is capable of handling source lines that are instructions, storage declaration, comments, and assembler directives.
- 3. The parser is to minimally be capable of decoding 2, 3 and 4-byte instructions.
- 4. The parser is to handle all storage directives (BYTE, WORD, RESW, and RESB).
- 5. The assembler should support:
  - a- EQU and ORG statements.
  - b- Simple expression evaluation. A simple expression includes simple (A B) operand arithmetic, where is one of +, -, \*, / and no spaces surround the operation, eg. A+B.
- 6. The output of the assembler should include (at least):
  - a- Intermediate file that contains the symbol table and a definition for any error.
  - b- Object-code file.

## Design

- → The user inserts the name of the file followed by (.txt) which contain the statements which need to be assembled, then:
  - ◆ The Main class calls the class ReadFromFile to read the statements one by one by for loop, in this loop the Parser is called to check that the statement is valid and doesn't have any syntax error.
    The Parser does that in two parts, first part by calling a function on it, and the second by calling Class DynamicTable which checks that the statement doesn't have any logical error. And if it's valid then it calls class Make Address to generate the Address as location counter of this line, then setting all (Location Counter, Label , Operation, Operant) in map SymbolMap by call class SymbolTable.
  - ◆ After reading the line and checking if this is a valid one, class <u>ReadFromFile</u> calls class <u>Controller</u> to generate the object code.
- → For the object code generation :
  - ◆ The <u>Controller</u> specifies which format is the instruction then calls the <u>formatchecker</u> that decides which pass to take from the three formats passess whether format 2 or 3 or 4
    - The 3 formats computes the first part of the object code by relatively the same way then for format 2 it finds the hexadecimal corresponding to each register used for format 3 and 4 takes the next step towards "find flags" class.
    - Find flags class specify which addressing mode the instruction uses and gets the corresponding flags combination (NIX) from "legalFlagsCombinationsClass" then performs some calculations to find the proper (BPE) part of the object code and gets it from "legalFlagsCombinationsClass".
    - After that the "<u>CalculateDisplacement</u>" does its work by calculating the displacement.
    - The controller combines the parts of the object code returned by each class and converts it to hexadecimal and stores the object code in the symbol table.
- → For the forward referencing part :
  - ◆ If the <u>CalculateDisplacement</u> class found an operand symbol that isn't in the symbol table(pre defined) it stores it in the "<u>LocationsTable</u>" singleton class and stores the address of the instruction which is missing this operand.
  - When the parser finds this symbol it gives it's address to "AddSymbol" class and it automatically completes the computation of the incomplete object codes.

### Main data structures

## I. Map (Hash map)

- Map to store the definition of data which follow the key word such as(WORD, BYTE, RESW, REWB, ORG, EQU).
- ➤ Map to store all Register names as key and his object code as parameters which Register are valid to use in Sic/XE.
- ➤ Map to store all Operation names as keys and his object code as parameters which Operation are valid to use in Sic/XE.

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#### II. Vector

- > Vector to store the Operant which is defined as follower reference .
- > Vector to split the operand which content of two parts one is Data and other is Register X or operand can be two register the vector store parts to check if the operant is valid .

## III. Array One Dimensional

- ➤ One dimension array in size seven elements to store the static label such as(WORD, BYTE, RESW, REWB, ORG, EQU) as String.
- ➤ One dimension array of size three to store the Label,Operation ,Operand.

# **Algorithms description**

- > First of all, the assembler reads the source lines from a file line by line and parsing it then sends it to the object code generator.
- Function **CheckWithRegex** in **Parsing** Class checks that the syntax of the line is correct, and splits the line to (Label Operation Operand), and then checks that all of them are valid, for example check the the operation belongs to **SIC/XE**.
- ➤ Class <u>DynamicTables</u> contains function <u>BuildDataTable</u> which check that statement don't have any logic error or any sentence error by many conditions such as:
  - Check that the Operand which after the keyword <u>"WORD"</u> the length less than 5 if the Operand is number, the length less than 8 and not contain

- white space if Operand is string and in two case must the length be more than zero.
- Check that the Operand which after the keyword <u>"BYTE"</u> the length less than 15 if the Operand is a hexadecimal number, is even number and contain character "X" and single quote, the length less than 16 and check that contain character "C" and single quote in the first and in the last and in two case must the length be more than zero.
- Check that the Operand which after the keyword <u>"RESB"</u> or <u>"RESW"</u> the length less than 5 number ,must be more than zero.
- Check that the Operand which after the keyword <u>"END"</u> or <u>"RESW"</u> the length be zero or in case "END" label not equal the label in the directive start.
- Check that the Operand which after the keyword <u>"START"</u> the length more than zero and less than 5 and the Operand must be a hexadecimal number.
- If the first character of Operand be '@','#','\*' must check that Operation not be any type of which deal with register and the operand size in case of '\*' must be one if the Operand after this symbol is word must check if have defined already if not store it to deal on it as forward reference if Operand is number the length must be less than 5 more than zero.

If Operand is two registers or register x and and word define or not must check that the Operation is at the end of character <u>'R'</u> or is equal <u>"STCH"</u>, "LDCH", and otherwise which deal on registers.

- Class <u>MakeAddress</u> contain function <u>LocationCounter</u> which return the (LOC) of every line as string of hexadecimal number that make by:
  - Define int variable <u>LOC</u> as static the initial value of it is the Operant which became after directive Start .
  - Then to calculate the value of location counter two any line take the value of LOC and ADD on it according to every situation :
    - ➤ If Operation is "RESB" then add Loc=Loc+Length of Operand
    - ➤ If Operation is "RESW" then add Loc=Loc+(Length of Operand) \*3
    - ➤ If Operation in format 4 then add LOC = 4 + LOC
    - ➤ If Operation is "WORD" then add Loc=Loc+3
    - ➤ If Operation is "BYTE" then add Loc=Loc+(Length of Operand-3) if the Operant is character if is hexadecimal then
      - Loc=Loc+(Length of Operand-3)/2.
    - Otherwise Loc=Loc+3 and for all that return the string equal the value of LOC at hexadecimal.

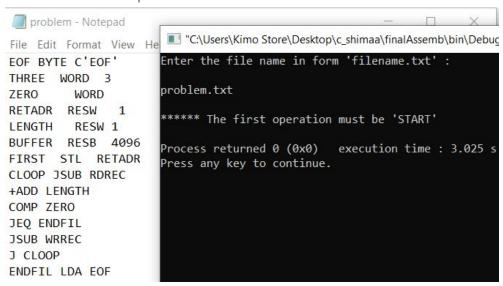
- ➤ If Operation is "ORG", then the next address will be the address of the operand. And if there is no operand, then it returns the default address.
- ➤ If Operation is "EQU", then the current address becomes equal to the value in the operand field or the address of the operand.

# **Assumptions**

> Owr program dealing with forward path technique for all Operand Not only in <u>JSUB</u> instructions but in any Operandand label except the Operand Which comes before Operation (Base) the Operandmust be defined before using it.

# Sample runs

First the parser check that the operation of the first instruction is "START" **Ex.1:** The first line operation isn't "Start"

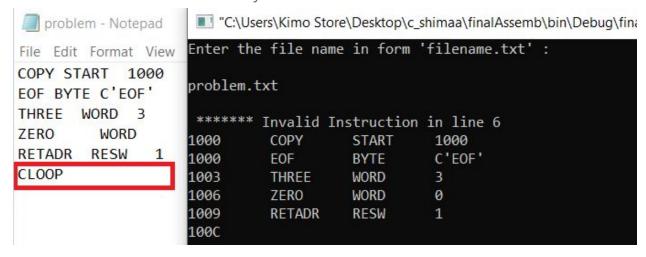


**Then** the parser checks that the instruction is a valid one with "regex". We can say the instruction is correct if:

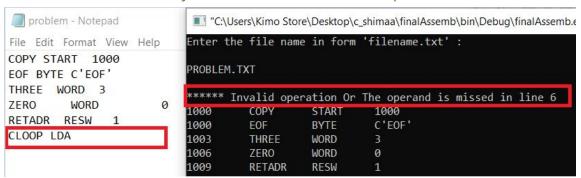
- It contains 3 words, which are label, operation and operand.
- It contains 2 words, operation and operand.
- It contains 1 word, like (RSUB, ORG, END).

So anything else will be considered as an invalid instruction. Let's try that.

**Ex.2:** The instruction is "CLOOP" only:



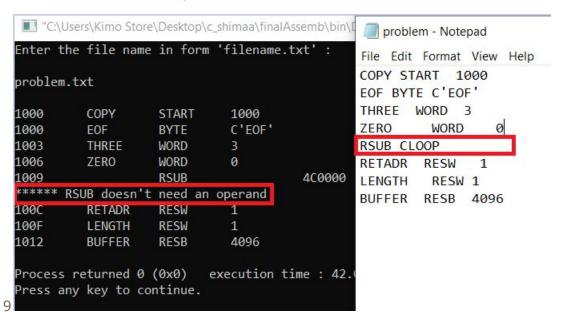
**Ex.3:** The instruction is only two words but it misses the operand.



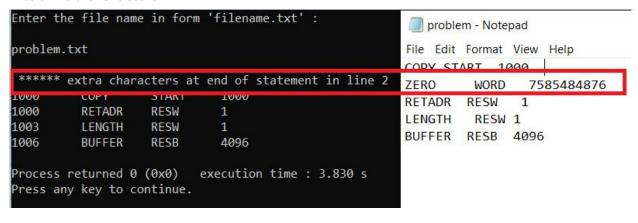
**Ex.4:** Check that the operation is existing in SIC/XE "STORE" isn't an operation in SIC/XE

```
problem - Notepad
                      "C:\Users\Kimo Store\Desktop\c_shimaa\finalAssemb\bin\Debug\finalAssemb.e
                     Enter the file name in form 'filename.txt' :
File Edit Format View
COPY START 1000
                     problem.txt
STORE LENGTH
                      1000
                                COPY
                                          START
                                                     1000
                      1000
                                          STORE
                      ****** Invalid operation Or The operand is missed in line 2
                                                  execution time : 5.579 s
                      Process returned 0 (0x0)
```

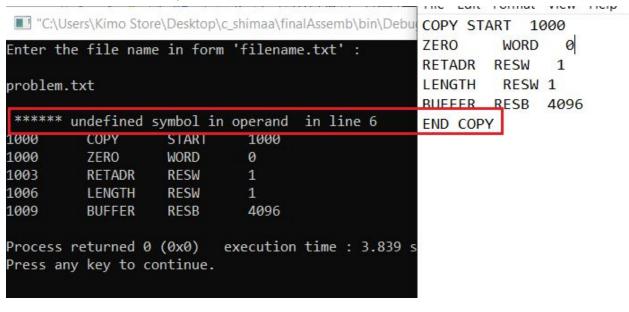
#### Ex.5: When RSUB has an operand



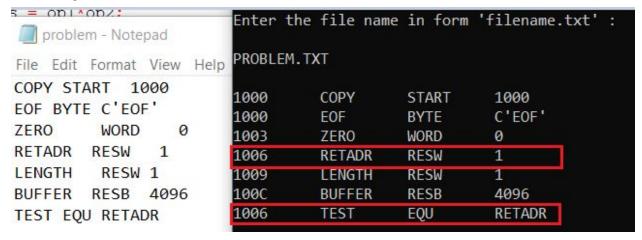
#### Ex.6: Extra characters



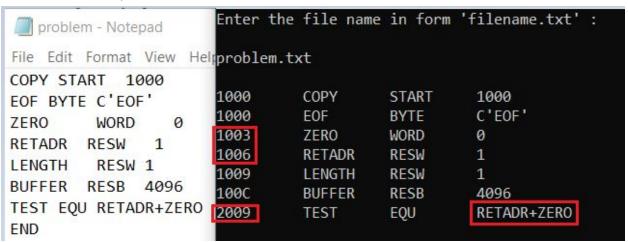
## Ex.7: When End has an Operand



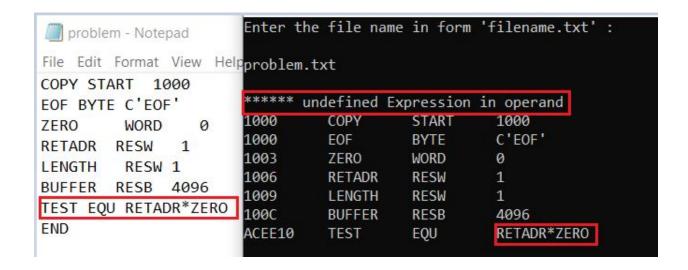
Ex.8: EQU



**Ex.9: EQU** with Expression



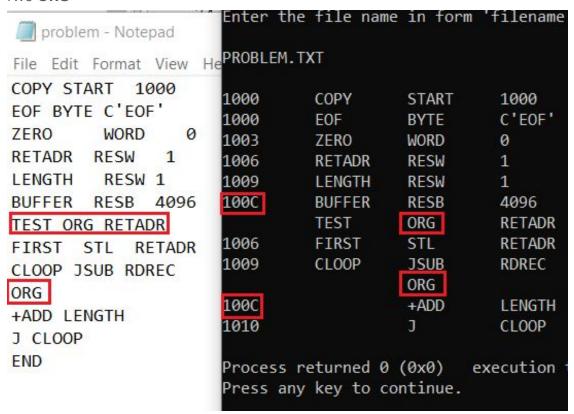
**Ex.10:** Invalid expression, multiplication is illegal for type relative



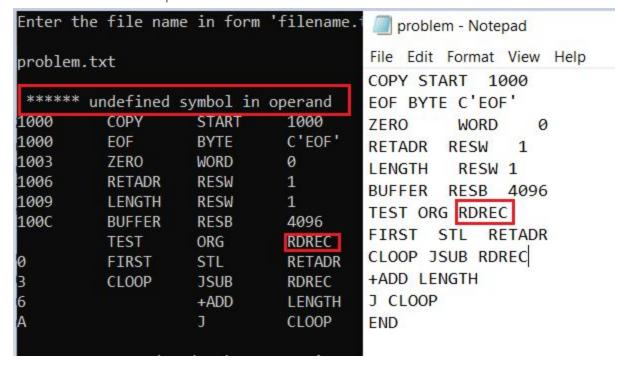
Ex.11: Test ORG

problem - Notepad	Enter t	he file nam	e in form	'filename.tx
File Edit Format View H	elp <mark>PROBLEM</mark>	.TXT		
COPY START 1000 EOF BYTE C'EOF' ZERO WORD 0 RETADR RESW 1 LENGTH RESW 1 BUFFER RESB 4096 TEST ORG RETADR FIRST STL RETADR CLOOP JSUB RDREC +ADD LENGTH J CLOOP END	1000 1000 1003 1006 1009 100C 1006 1009 100C 1010 Process	COPY EOF ZERO RETADR LENGTH BUFFER TEST FIRST CLOOP	RESW RESB ORG STL JSUB +ADD J	1000 C'EOF' 0 1 1 4096 RETADR RETADR RDREC LENGTH CLOOP

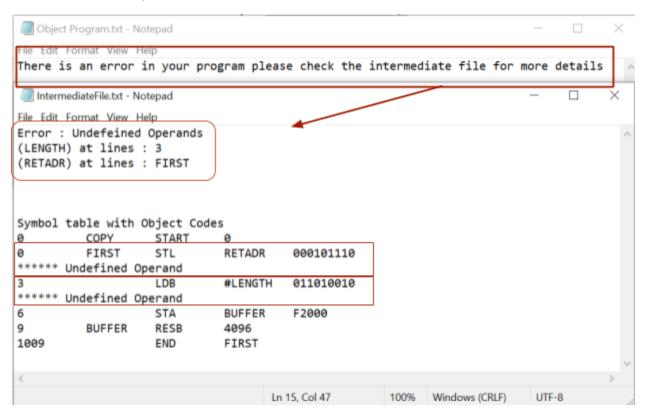
#### Two **ORG**



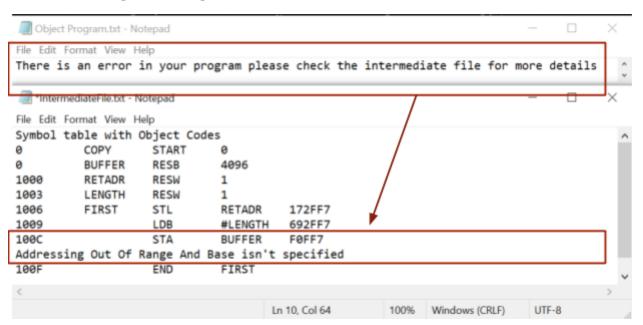
#### **ORG** with undefined Operand



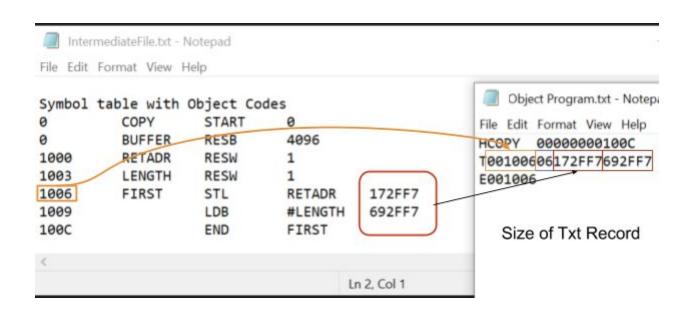
### **Ex.12:** Undefined Operands:



**Ex.13:** Addressing out of range:



**Ex.13:** Without forward referencing:



Ex.13: The Book's Example:

5	0000	COPY	START	0	
10	0000	FIRST	SIL	The same of the same	17202D
12	0003	FIRST	LDB	RETADR #LENGTH	17202D
13	0003		20000000	LENGTH	69202D
15	0006	CLOOP	BASE		AD101025
		CLUOP	+JSUB	RDREC	4B101036
20	A000		LDA	LENGTH	032026
25	000D		COMP	#0	290000
30	0010	-	JEQ	ENDFIL	332007
35	0013		+JSUB	WRREC	4B10105D
40	0017		J	CLOOP	3F2FBC
45	001A	ENDFIL	LDA	EOF	032010
50	001D	_	7STA	BUFFER	0F2016
55	0020		LDA	#3	010003
60	0023		STA	LENGTH	0F200D
65	0026		+JSUB	WRREC	4B10105D
70	002A		J	GRETADR	3E2003
80	002D	EOF	BYTE	C'EOF'	454F46
95	0030	RETADR	RESW	1	
100	0033	LENGTH	RESW	1	
105	0036	BUFFER	RESB	4096	
110		-			
115			SUBROUT	TIME TO READ RE	CORD INTO BUFFER
120					
125	1036	RDREC	CLEAR	X	B410
130	1038		CLEAR	A	B400
132	103A		CLEAR	S	B440
133	103C		+LDT	#4096	75101000
135	1049	RLOOP	TD	INPUT	E32019
140	1043		JEQ	RLOOP	332FFA
145	1046		RD	INPUT	DB2013
150	1049		COMPR	A,S	A004
155	104B		JEQ	EXIT	332008
160	104E		STCH	BUFFER, X	57C003
165	1051		TIXR	T	B850
170	1053		JLT	RLOOP	3B2FEA
175	1056	EXIT	STX	LENGTH	134000
180	1059		RSUB		4F0000
185	105C	INPUT	BYTE	X'F1'	F1
195				Tr. Tale	
200			SUBROUT	TIME TO WRITE R	ECORD FROM BUFFER
205					
210	105D	WRREC	CLEAR	X	B410
212	105F		LDT	LENGTH	774000
215	1062	WLOOP	TD	OUTPUT	E32011
220	1065	The state of the s	JEQ	WLOOP	332FFA
225	1068		LDCH	BUFFER, X	53C003
230	106B		WD	OUTPUT	DF2008
235	106E		TIXR	T	B850
240	1070		JLT *	WLOOP	3B2FEF
245	1073		RSUB		4F0000
250	1076	CUTPUT	BYTE	X'05'	05
255			END	FIRST	

