

CSAPP : datalab 实验报告

实验目标

datalab 实验要求我们使用限定的运算符 `&` `|` `<<` `>>` `!` 等，完成每一个函数的功能并满足所有限制条件。

文件包中有以下文件，注释了各自的功能。

```
Makefile      - Makes btest, fshow, and ishow
README        - This file
bits.c        - The file you will be modifying and handing in
bits.h        - Header file
btest.c       - The main btest program
  btest.h     - Used to build btest
  decl.c      - Used to build btest
  tests.c     - Used to build btest
  tests-header.c - Used to build btest
dlc*          - Rule checking compiler binary (data lab compiler)
driver.pl*    - Driver program that uses btest and dlc to autograde bits.c
Driverhdrs.pm - Header file for optional "Beat the Prof" contest
fshow.c       - Utility for examining floating-point representations
ishow.c       - Utility for examining integer representations
```

题目及解法

bitAnd

第一题较为简单 $A \& B = \sim((\sim A) \mid (\sim B))$

```
/*
 * bitAnd - x&y using only ~ and |
 *   Example: bitAnd(6, 5) = 4
 *   Legal ops: ~ |
 *   Max ops: 8rr
 *   Rating: 1
 */
```

```
int bitAnd(int x, int y) {

    int res = ~((~x)|(~y));
    return res;

}
```

getByte

题目要求返回 32 位数 `x` 中的第 `n` 个字节，通过移位运算和掩码即可。

```
/*
 * getByte - Extract byte n from word x
 *   Bytes numbered from 0 (LSB) to 3 (MSB)
 *   Examples: getByte(0x12345678,1) = 0x56
 *   Legal ops: ! ~ & ^ | + << >>
 *   Max ops: 6
 *   Rating: 2
 */
int getByte(int x, int n) {

    int shifts = n << 3;
    int res = x >> shifts;
    res &= 0xff;

    return res;

}
```

logicalShift

题目要求实现逻辑右移，我的想法是通过加法将算术右移填充的 1 溢出成 0，还可以在 `x >> 1` 之后将符号位直接清 0。

```
/*
 * logicalShift - shift x to the right by n, using a logical shift
 *   Can assume that 0 <= n <= 31
 *   Examples: logicalShift(0x87654321,4) = 0x08765432
 *   Legal ops: ! ~ & ^ | + << >>
 *   Max ops: 20
 *   Rating: 3
 */
int logicalShift(int x, int n) {

    int bias = (x >> 31) & 2;
```

```

int res, tem;
tem = (31 ^ n);
bias <=< tem;
res = bias + (x >> n);

return res;

}

```

bitCount

题目要求我们统计 32 位数字二进制表示中数字 1 出现的次数，并且限定操作符在 40 个以内。

老师上课时讲过一种统计方法，如图。因为有 $2^n - 2^{n-1} - 2^{n-2} \dots - 1 = 1$ 通过图中计算 tmp 的方式，可以每 3 位为一组，计算这一组中 1 的个数。 $(tmp + (tmp \gg 3)) \& 0x030707070707$ 将相邻红蓝块数值合并入蓝块中。因为 $1 \equiv 64 \pmod{63}$ 所以对 63 取模就得到了蓝块之和，即答案。

统计一个机器字的二进制表示中“1”的个数

```

int bitcount(unsigned int n)
{
    unsigned int tmp;

    tmp = n - ((n >> 1) & 033333333333)
           - ((n >> 2) & 011111111111);
    return ((tmp + (tmp >> 3)) & 030707070707) % 63;
}

```

The diagram illustrates the bit counting algorithm using a 32-bit word. It shows the word being processed in groups of 3 bits. The diagram uses red and blue blocks to represent the intermediate steps of the algorithm, showing how the count of 1s is accumulated and shifted. The final result is obtained by taking the modulo 63 of the accumulated value.

MIT 6.034 Lecture 4.60

L2 Chp2.35

Wu Spring 19 ©USTC

而我采用了类似线段树值合并的思想，有点类似上述方法，通过移位和加法，逐步将 1 的个数统计出来。

```

/*
 * bitCount - returns count of number of 1's in word
 * Examples: bitCount(5) = 2, bitCount(7) = 3
 */

```

```

*   Legal ops: ! ~ & ^ | + << >>
*   Max ops: 40
*   Rating: 4
*/
int bitCount(int x) {

    int bias1 = 0x55;
    int bias2 = 0x33;
    int bias3 = 0x0f;
    int bias4 = 0xff;
    int bias5 = 0xff;

    bias1 |= bias1 << 8;
    bias1 |= bias1 << 16;

    bias2 |= bias2 << 8;
    bias2 |= bias2 << 16;

    bias3 |= bias3 << 8;
    bias3 |= bias3 << 16;

    bias4 |= bias4 << 16;

    bias5 |= bias5 << 8;

    x = (x & bias1) + ((x >> 1) & bias1);
    x = (x & bias2) + ((x >> 2) & bias2);
    x = (x & bias3) + ((x >> 4) & bias3);
    x = (x & bias4) + ((x >> 8) & bias4);
    x = (x & bias5) + ((x >> 16) & bias5);

    return x;

}

```

bang

当 $x \neq 0$ 时 $!x = 0$, $x = 0$ 时 $!x = 1$ 题目要求不使用 `!` 求出 `!x` , 只需要检查 `x` 二进制表示中是否有 `1` 即可。

```

/*
* bang - Compute !x without using !
* Examples: bang(3) = 0, bang(0) = 1
* Legal ops: ~ & ^ | + << >>
* Max ops: 12
* Rating: 4
*/
int bang(int x) {

```

```

    x |= x >> 16;
    x |= x >> 8;
    x |= x >> 4;
    x |= x >> 2;
    x |= x >> 1;

    return (x & 1) ^ 1;
}

```

tmin

要求输出最小的有符号整数(int)

```

/*
 * tmin - return minimum two's complement integer
 *   Legal ops: ! ~ & ^ | + << >>
 *   Max ops: 4
 *   Rating: 1
 */
int tmin(void) {

    int res = 1 << 31;
    return res;
}

```

fitsBits

题目问 `x` 能否表示为 `n` 位二进制数。我一开始的做法十分复杂，其实只要将 `x` 截断为 `n` 位再还原回去，比较与原来的数是否相等就行。因为 `n` 位二进制数也是有符号的，所以如果有操作引起数值符号的改变导致与原本不相等的情况，正是压缩为 `n` 位之后最高位翻译成符号位的表现。

```

/*
 * fitsBits - return 1 if x can be represented as an
 *   n-bit, two's complement integer.
 *   1 <= n <= 32
 *   Examples: fitsBits(5,3) = 0, fitsBits(-4,3) = 1
 *   Legal ops: ! ~ & ^ | + << >>
 *   Max ops: 15
 *   Rating: 2
 */
int fitsBits(int x, int n) {

```

```

    int m = 33 + ~n;
    int t = x << m >> m;

    return !(t ^ x);

}

```

divpwr2

就是对 x 进行向 0 舍入的移位，当 $x < 0$ 时加一个偏置即可。

```

/*
 * divpwr2 - Compute x/(2^n), for 0 <= n <= 30
 * Round toward zero
 * Examples: divpwr2(15,1) = 7, divpwr2(-33,4) = -2
 * Legal ops: ! ~ & ^ | + << >>
 * Max ops: 15
 * Rating: 2
 */
int divpwr2(int x, int n) {

    int s = (x >> 31) & 1;
    int res = x + ((s << n) + ~0 + !s) >> n;

    return res;

}

```

negate

$-x = \sim x + 1$

```

/*
 * negate - return -x
 * Example: negate(1) = -1.
 * Legal ops: ! ~ & ^ | + << >>
 * Max ops: 5
 * Rating: 2
 */
int negate(int x) {

    return ~x + 1;

}

```

isPositive

判断 x 是否大于 0 ，注意特判 0 的情况。

```
/*
 * isPositive - return 1 if x > 0, return 0 otherwise
 * Example: isPositive(-1) = 0.
 * Legal ops: ! ~ & ^ | + << >>
 * Max ops: 8
 * Rating: 3
 */
int isPositive(int x) {

    int res = (x >> 31) & 1;
    res |= !x;
    return !res;

}
```

isLessOrEqual

询问是否 $x \leq y$ ，则可以通过先比较符号位，不等时可以直接出答案，当符号位相等时则使用减法判断大小而不会产生溢出的问题。

```
/*
 * isLessOrEqual - if x <= y then return 1, else return 0
 * Example: isLessOrEqual(4,5) = 1.
 * Legal ops: ! ~ & ^ | + << >>
 * Max ops: 24
 * Rating: 3
 */
int isLessOrEqual(int x, int y) {

    int s1 = (x >> 31) & 1;
    int s2 = (y >> 31) & 1;
    int res = (!s2) & s1;
    res |= (!(s1 ^ s2)) & (!(y + ~x + 1 >> 31));

    return res;

}
```

ilog2

题目问 `x` 的最高位 `1` 在第几位上，但给出了比较严格的限制。这题的思路有点类似二分法，运用了整数能被唯一地表示为一个二进制数的知识。（通过 `!!` 符号可以将非零整数映射为 `1`）

```
/*
 * ilog2 - return floor(log base 2 of x), where x > 0
 *   Example: ilog2(16) = 4
 *   Legal ops: ! ~ & ^ | + << >>
 *   Max ops: 90
 *   Rating: 4
 */
int ilog2(int x) {

    int res = 0;
    int t = !! (x >> 16);
    res |= t << 4;
    t = !! (x >> res + 8);
    res |= t << 3;
    t = !! (x >> res + 4);
    res |= t << 2;
    t = !! (x >> res + 2);
    res |= t << 1;
    t = !! (x >> res + 1);
    res |= t;

    return res;
}
```

float_neg

输出 `-f`，但特别注意 `NaN` 的情况。

```
/*
 * float_neg - Return bit-level equivalent of expression -f for
 *   floating point argument f.
 *   Both the argument and result are passed as unsigned int's, but
 *   they are to be interpreted as the bit-level representations of
 *   single-precision floating point values.
 *   When argument is NaN, return argument.
 *   Legal ops: Any integer/unsigned operations incl. ||, &&. also if, while
 *   Max ops: 10
 *   Rating: 2
 */
unsigned float_neg(unsigned uf) {

    if ((uf & 0x7F800000) != 0x7F800000 || !(uf & 0x007FFFFFFF))
        uf ^= 1 << 31;
}
```



```
    return uf;

}
```

float_i2f

题目要求将 `int` 类型转化为 `float` 类型。先判断 `x` 的符号，但由于 `-tmin = tmin` 所以先将其特判掉。接着按照类型转换的逻辑写下去即可，但要注意规格化问题，由于 `float` 与 `int` 精度差别，会导致舍入问题，还应考虑舍入导致的进位问题。

```
/*
 * float_i2f - Return bit-level equivalent of expression (float) x
 *   Result is returned as unsigned int, but
 *   it is to be interpreted as the bit-level representation of a
 *   single-precision floating point values.
 *   Legal ops: Any integer/unsigned operations incl. ||, &&. also if, while
 *   Max ops: 30
 *   Rating: 4
 */
unsigned float_i2f(int x) {

    unsigned res = 0;
    unsigned cnt = 0;
    unsigned y;
    int i, j;
    int zero;

    if (!x) return 0;
    if (x == 0x80000000) return 0xc0000000;

    if (x < 0) {
        x = -x;
        res = 1 << 31;
    }
    for (i = 31; i > 0; --i) {
        if (x & (1 << i)) break;
    }
    cnt = i;
    res |= cnt + 127 << 23;
    y = x << 32 - cnt;
    i = y & 0x100;
    j = y & 0x0fff;
    y >>= 9;
    res |= y;
    if (i == 0x100){
        if (j != 0 || (y & 1)) ++res;
    }
}
```

```
    return res;

}
```

float_twice

题目要求计算 $2 * u2f(f)$ 。非规格化数只需整体左移再维护符号位不变就能达到乘二的目的，因为若在指数部分产生溢出正好将数字转为规格化(隐藏最高位 1)，对于规格化数指数加一就能达到要求，而对于 NaN 及 无穷 特判返回即可。

```
/*
 * float_twice - Return bit-level equivalent of expression 2*f for
 *   floating point argument f.
 *   Both the argument and result are passed as unsigned int's, but
 *   they are to be interpreted as the bit-level representation of
 *   single-precision floating point values.
 *   When argument is NaN, return argument
 *   Legal ops: Any integer/unsigned operations incl. ||, &&. also if, while
 *   Max ops: 30
 *   Rating: 4
 */
unsigned float_twice(unsigned uf) {

    unsigned s = uf & (1 << 31);
    unsigned e = uf & 0x7f800000;
    unsigned f = uf & 0x007fffff;
    if (e == 0x7f800000) return uf;
    if (e == 0) {
        uf <<= 1;
        uf = (uf & 0x7fffffff) | s;
        return uf;
    }

    e = (e >> 23) + 1 << 23;
    return s | e | f;

}
```

实验结果

本实验所有资源在 <https://github.com/Afool1999/CSAPP-Labs/tree/master/datalab>

文件(F) 编辑(E) 查看(V) 搜索(S) 终端(T) 帮助(H)

make: 对“all”无需做任何事。

```
afool@ubuntu:~/Documents/CS:APP_LAB/datalab/datalab-handout$ ./dlc bits.c
bits.c:382: Warning: unused variable `zero'
```

Compilation Successful (1 warning)

```
afool@ubuntu:~/Documents/CS:APP_LAB/datalab/datalab-handout$ ./btest bits.c
```

Score	Rating	Errors	Function
1	1	0	bitAnd
2	2	0	getByte
3	3	0	logicalShift
4	4	0	bitCount
4	4	0	bang
1	1	0	tmin
2	2	0	fitsBits
2	2	0	divpwr2
2	2	0	negate
3	3	0	isPositive
3	3	0	isLessOrEqual
4	4	0	ilog2
2	2	0	float_neg
4	4	0	float_i2f
4	4	0	float_twice

Total points: 41/41

```
afool@ubuntu:~/Documents/CS:APP_LAB/datalab/datalab-handout$
```