

"Career With Rishab" Complexity Cheat Sheet

Complexity of Different Data Structures

Data Structure	Operation	Time Complexity			Explanation
		Worst Case	Average Case	Best Case	
Array	Access	$O(1)$	$O(1)$	$O(1)$	Accessing an element in an array by index takes constant time as the location can be calculated directly.
	Search	$O(n)$	$O(n)$	$O(1)$	In the worst case, searching an unsorted array requires traversing through all elements. In the best case, if the element is at the beginning, it takes constant time.
	Insertion	$O(n)$	$O(n)$	$O(1)$	In the worst case, inserting an element at the beginning of an array requires shifting all other elements.
	Deletion	$O(n)$	$O(n)$	$O(1)$	In the worst case, deleting an element from the beginning of an array requires shifting all other elements.
Singly Linked List	Access	$O(n)$	$O(n)$	$O(1)$	Accessing an element in a linked list requires traversing the list from the head to the desired element.
	Search	$O(n)$	$O(n)$	$O(1)$	Searching in a linked list involves traversing the entire list to find the desired element.
	Insertion	$O(1)$	$O(1)$	$O(1)$	Inserting an element at the head or tail of a linked list can be done in constant time by updating the necessary pointers.
	Deletion	$O(1)$	$O(1)$	$O(1)$	Deleting the head or tail of a linked list can be done in constant time by updating the necessary pointers.
Doubly Linked List	Access	$O(n)$	$O(n)$	$O(1)$	Accessing an element in a doubly linked list requires traversing the list from the head or tail to the desired element.
	Search	$O(n)$	$O(n)$	$O(1)$	Searching in a doubly linked list involves traversing the entire list to find the desired element.
	Insertion	$O(1)$	$O(1)$	$O(1)$	Inserting an element at the head or tail of a doubly linked list can be done in constant time by updating the necessary pointers.
	Deletion	$O(1)$	$O(1)$	$O(1)$	Deleting a node from a doubly linked list can be done in constant time by updating the necessary pointers.
Stack	Push	$O(1)$	$O(1)$	$O(1)$	Pushing an element onto a stack takes constant time as it involves updating the top of the stack.
	Pop	$O(1)$	$O(1)$	$O(1)$	Popping an element from a stack takes constant time as it involves updating the top of the stack.
	Peek	$O(1)$	$O(1)$	$O(1)$	Peeking at the top element of a stack takes constant time as it does not modify the stack.
Queue	Enqueue	$O(1)$	$O(1)$	$O(1)$	Enqueuing an element into a queue takes constant time as it involves updating the tail of the queue.
	Dequeue	$O(1)$	$O(1)$	$O(1)$	Dequeuing an element from a queue takes constant time as it involves updating the head of the queue.
	Peek	$O(1)$	$O(1)$	$O(1)$	Peeking at the front element of a queue takes constant time as it does not modify the queue.
Heap	Insert	$O(\log n)$	$O(\log n)$	$O(1)$	Inserting an element into a heap takes logarithmic time as it may require restructuring the heap.
	Delete	$O(\log n)$	$O(\log n)$	$O(1)$	Deleting the minimum or maximum element from a heap takes logarithmic time as it may require restructuring the heap.
	Get Min/Max	$O(1)$	$O(1)$	$O(1)$	Retrieving the minimum or maximum element from a heap can be done in constant time.
Binary Search Tree	Search	$O(n)$	$O(\log n)$	$O(1)$	In the worst case, a binary search tree can be a skewed binary search tree and can lead to linear time complexity for searching.
	Insert	$O(n)$	$O(\log n)$	$O(\log n)$	In the worst case, inserting elements in a binary search tree without balancing can lead to linear time complexity.
	Delete	$O(n)$	$O(\log n)$	$O(\log n)$	In the worst case, deleting elements in a binary search tree without balancing can lead to linear time complexity.
AVL Tree	Search	$O(\log n)$	$O(\log n)$	$O(1)$	Searching in an AVL tree takes logarithmic time as it maintains balanced height through rotations.
	Insert	$O(\log n)$	$O(\log n)$	$O(\log n)$	Inserting an element into an AVL tree requires rotations to maintain balance, resulting in logarithmic time complexity.
	Delete	$O(\log n)$	$O(\log n)$	$O(\log n)$	Deleting an element from an AVL tree requires rotations to maintain balance, resulting in logarithmic time complexity.
Red-Black Tree	Search	$O(\log n)$	$O(\log n)$	$O(\log n)$	Searching in a red-black tree takes logarithmic time as it maintains balanced height and adheres to the properties of a binary search tree.
	Insert	$O(\log n)$	$O(\log n)$	$O(\log n)$	Inserting an element into a red-black tree requires restructuring to maintain balance, resulting in logarithmic time complexity.
	Delete	$O(\log n)$	$O(\log n)$	$O(\log n)$	Deleting an element from a red-black tree requires restructuring to maintain balance, resulting in logarithmic time complexity.
Hash Table	Search	$O(1)$	$O(1)$	$O(1)$	Hash tables provide constant-time search on average, assuming a good hash function and proper handling of collisions.
	Insert	$O(1)$	$O(1)$	$O(1)$	Inserting into a hash table takes constant time on average, assuming a good hash function and proper handling of collisions.
	Delete	$O(1)$	$O(1)$	$O(1)$	Deleting from a hash table takes constant time on average, assuming a good hash function and proper handling of collisions.

n = number of elements

Complexity of Different Array Algorithms

Array Algorithm	Time Complexity			Explanation	Space Complexity	
	Worst Case	Average Case	Best Case		Worst Case	Explanation
Linear Search	$O(n)$	$O(n)$	$O(1)$	Iterate through the array	$O(1)$	Only need constant space for variables
Binary Search	$O(\log n)$	$O(\log n)$	$O(1)$	Divide and conquer approach	$O(1)$	Only need constant space for variables
Ternary Search	$O(\log^3 n)$	$O(\log^3 n)$	$O(1)$	Divide and conquer approach	$O(1)$	Only need constant space for variables
Merge Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	Divide and conquer approach	$O(n)$	Requires additional space for merging
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$	Find the minimum element	$O(1)$	Only need constant space for variables
Bubble Sort	$O(n^2)$	$O(n^2)$	$O(n)$	Compare adjacent elements	$O(1)$	Only need constant space for variables
Insertion Sort	$O(n^2)$	$O(n^2)$	$O(n)$	Insert element in sorted order	$O(1)$	Only need constant space for variables
Heap Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	Build max/min heap and extract elements	$O(1)$	Only need constant space for variables
Quick Sort	$O(n^2)$	$O(n \log n)$	$O(n \log n)$	Divide and conquer approach	$O(\log n)$	Requires space for the recursion stack
Randomized Quick Sort	$O(n^2)$	$O(n \log n)$	$O(n \log n)$	Randomly select pivot	$O(\log n)$	Requires space for the recursion stack
Bucket Sort	$O(n^2)$	$O(n + k)$	$O(n + k)$	Distribute elements into buckets	$O(n + k)$	Requires additional space for buckets
Counting Sort	$O(n+k)$	$O(n+k)$	$O(n+k)$	Count occurrences of elements	$O(n+k)$	Requires additional space for counting
Kadane's Algorithm	$O(n)$	$O(n)$	$O(n)$	Find maximum subarray sum	$O(1)$	Only need constant space for variables

Complexity of Different Graph Algorithms

Graph Algorithm	Time Complexity			Explanation	Space Complexity	
	Worst Case	Average Case	Best Case		Worst Case	Explanation
Depth First Search	$O(V + E)$	$O(V + E)$	$O(V + E)$	Traverse all vertices and edges in a graph	$O(V)$	Requires space for the recursion stack
Breadth First Search	$O(V + E)$	$O(V + E)$	$O(V + E)$	Traverse all vertices and edges in a graph	$O(V)$	Requires space for the queue
Dijkstra's Algorithm	$O(V^2)$	$O((V+E) \log V)$	$O(V \log V)$	Find shortest path from a source to all other vertices in a graph.	$O(V)$	Requires space for the priority queue
Floyd-Warshall Algorithm	$O(V^3)$	$O(V^3)$	$O(V^3)$	Find shortest path between all pairs of vertices in a graph	$O(V^2)$	Requires space for the distance matrix
Prim's Algorithm	$O((V + E) \log V)$	$O((V + E) \log V)$	$O((V + E) \log V)$	Construct a minimum spanning tree in a graph	$O(V)$	Requires space for the priority queue
Kruskal's Algorithm	$O(E \log E)$	$O(E \log E)$	$O(E \log E)$	Construct a minimum spanning tree in a graph	$O(V)$	Requires space for the disjoint set data structure
Bellman-Ford Algorithm	$O(VE)$	$O(VE)$	$O(VE)$	Find the shortest path from a source to all other vertices in a graph	$O(V)$	Requires space for the distance array
0/1 Knapsack Problem	$O(NW)$	$O(NW)$	$O(NW)$	Find the maximum value subset from a given set of items	$O(W)$	Requires space for the memoization table
Topological Sorting	$O(V + E)$	$O(V + E)$	$O(V + E)$	Sort the vertices in a directed acyclic graph	$O(V)$	Requires space for the stack

V = Number of Vertices
E = Number of Edges
W = Capacity of the Knapsack
N = Number of items available