

Lecture No.11

Data Structures

# Code for Simulation

```
// print the final average wait time.
```

```
double avgWait = (totalTime*1.0)/count;  
cout << "Total time:      " << totalTime << endl;  
cout << "Customer:         " << count << endl;  
cout << "Average wait: " << avgWait << endl;
```

# Priority Queue

```
#include "Event.cpp"
#define PQMAX 30

class PriorityQueue {
public:
    PriorityQueue() {
        size = 0; rear = -1;
    };
    ~PriorityQueue() {};

    int full(void)
    {
        return ( size == PQMAX ) ? 1 : 0;
    };
};
```

# Priority Queue

```
Event* remove()
{
    if( size > 0 ) {
        Event* e = nodes[0];
        for(int j=0; j < size-2; j++ )
            nodes[j] = nodes[j+1];

        size = size-1; rear=rear-1;
        if( size == 0 ) rear = -1;

        return e;
    }
    return (Event*)NULL;
    cout << "remove - queue is empty." << endl;
};
```

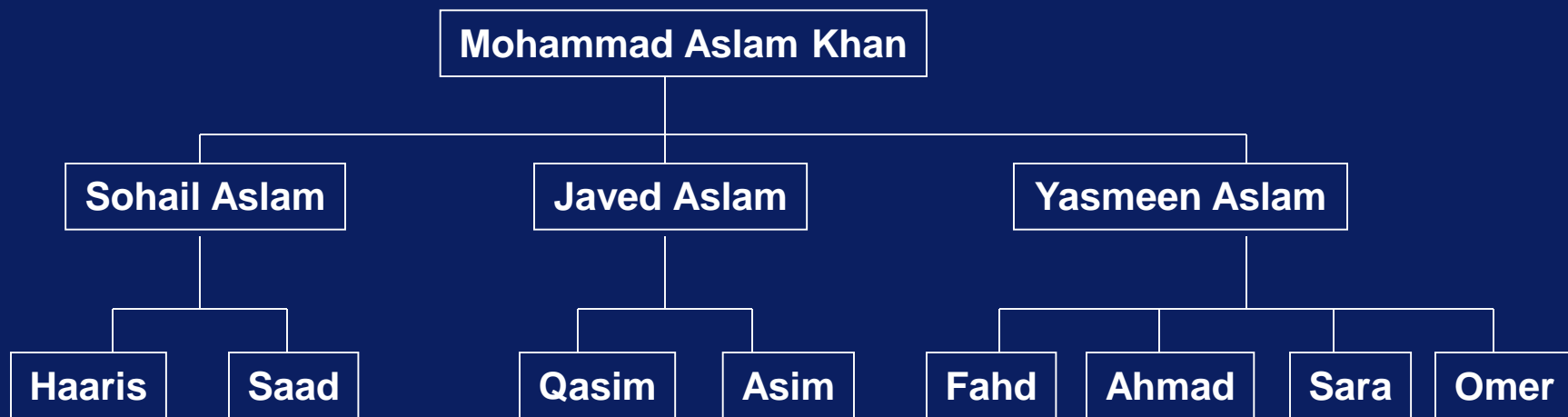
# Priority Queue

```
int insert(Event* e)
{
    if( !full() ) {
        rear = rear+1;
        nodes[rear] = e;
        size = size + 1;
        sortElements(); // in ascending order
        return 1;
    }
    cout << "insert queue is full." << endl;
    return 0;
};

int length() { return size; };
};
```

# Tree Data Structures

- There are a number of applications where linear data structures are not appropriate.
- Consider a genealogy tree of a family.



# Tree Data Structure

- A linear linked list will not be able to capture the tree-like relationship with ease.
- Shortly, we will see that for applications that require searching, linear data structures are not suitable.
- We will focus our attention on *binary trees*.

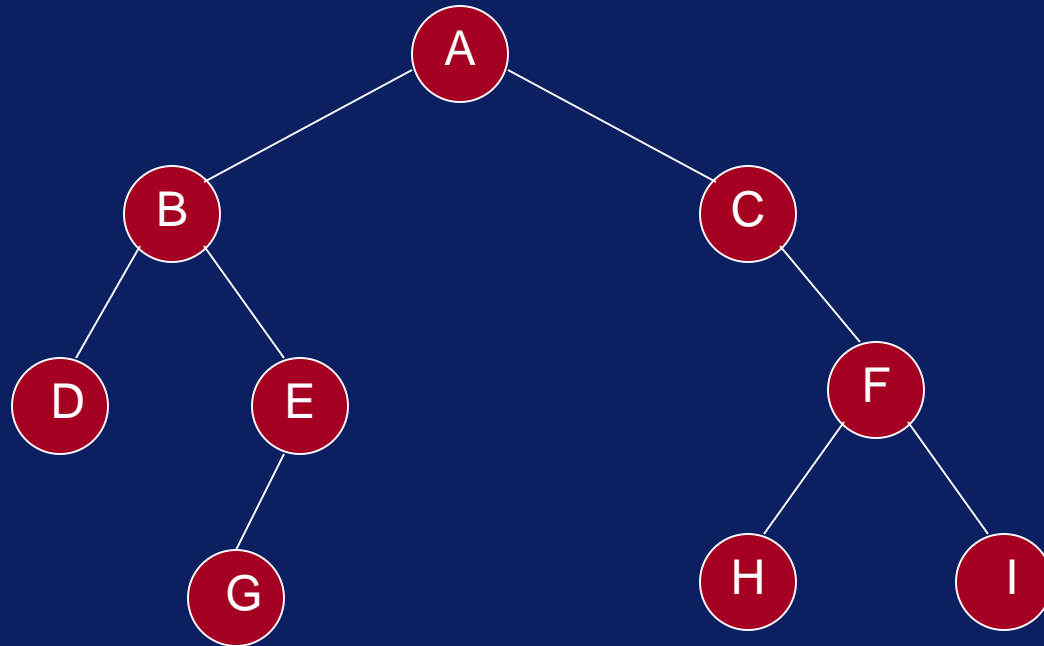
# Binary Tree

- A *binary tree* is a finite set of elements that is either empty or is partitioned into *three* disjoint subsets.
- The first subset contains a single element called the *root* of the tree.
- The other two subsets are themselves binary trees called the *left* and *right subtrees*.
- Each element of a binary tree is called a *node* of the tree.

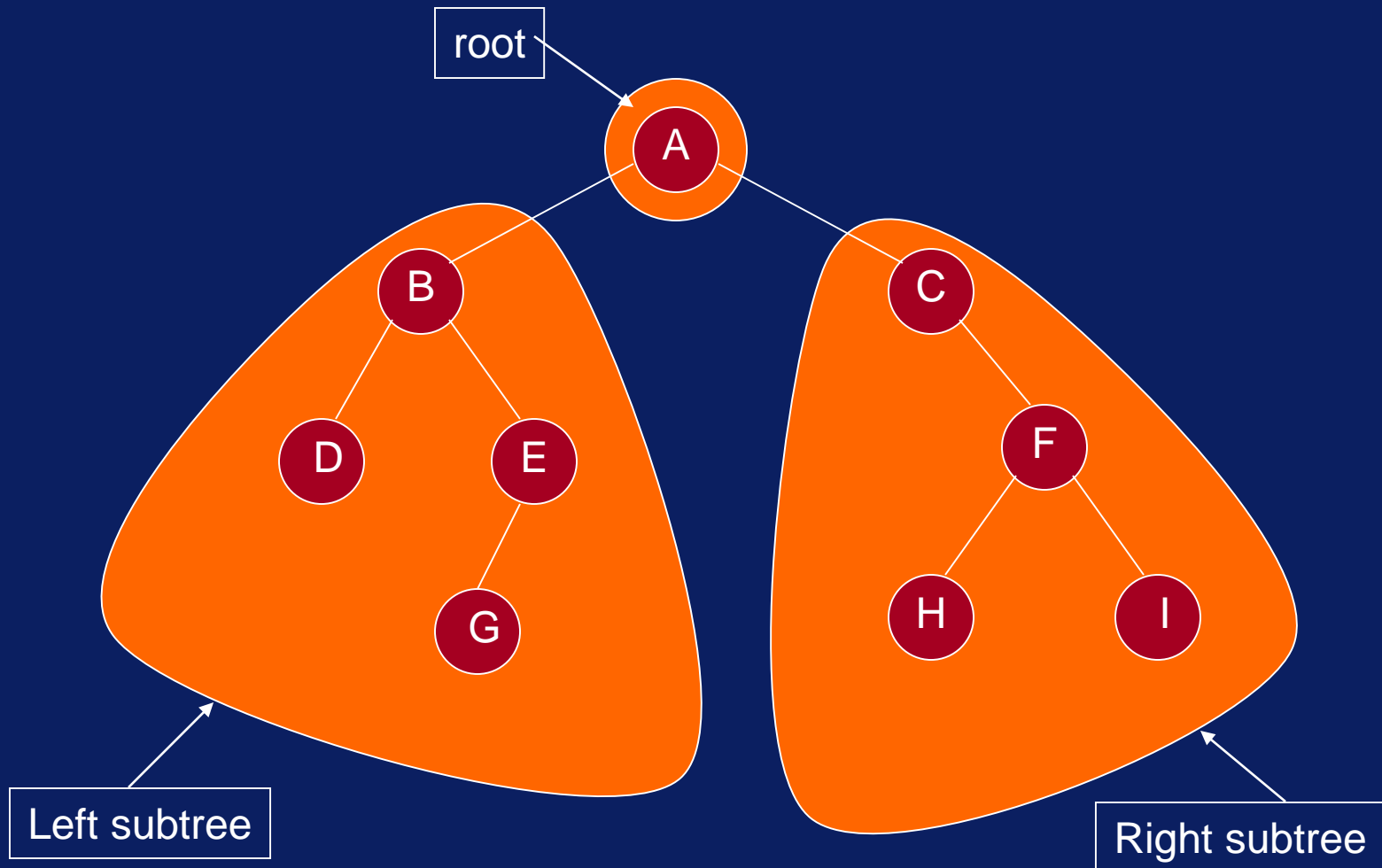


# Binary Tree

- Binary tree with 9 nodes.

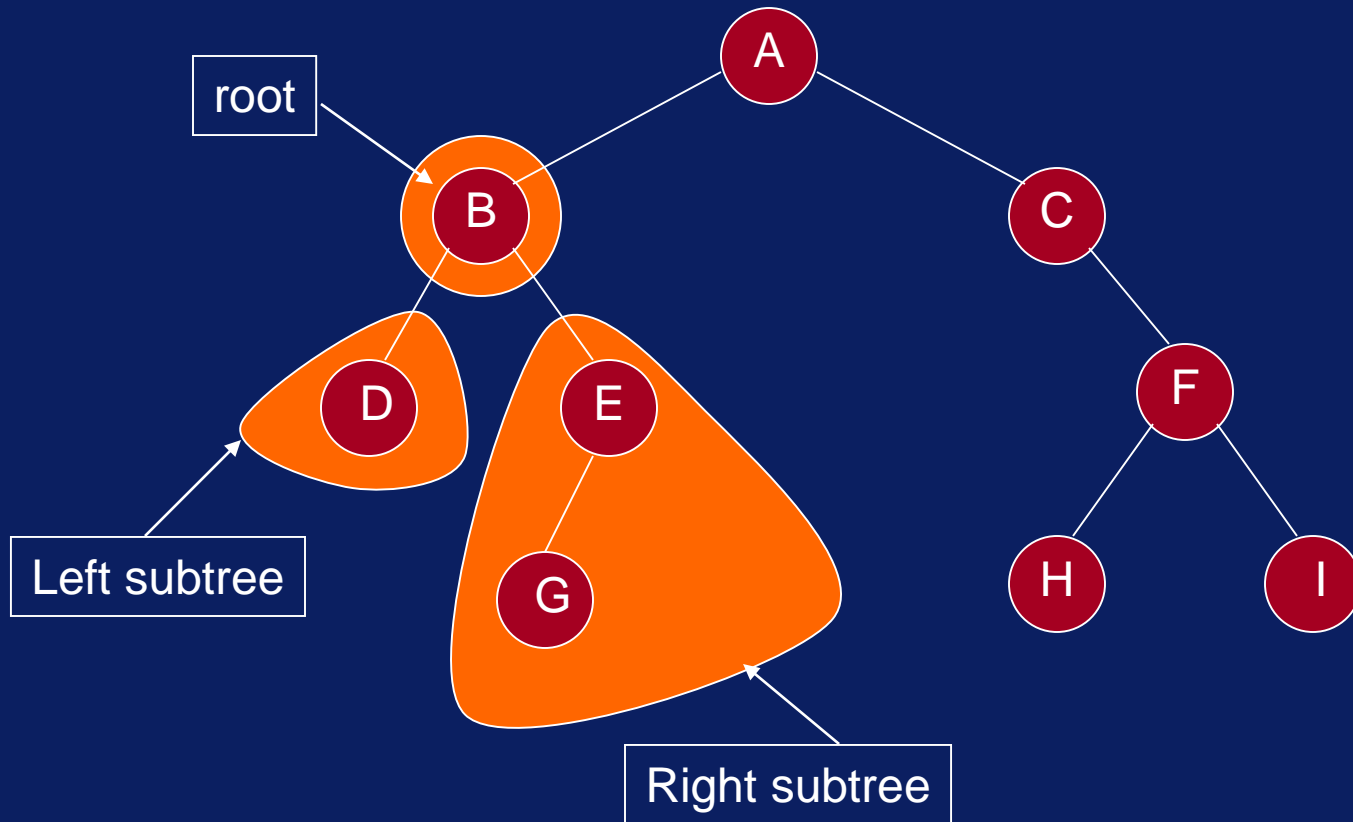


# Binary Tree



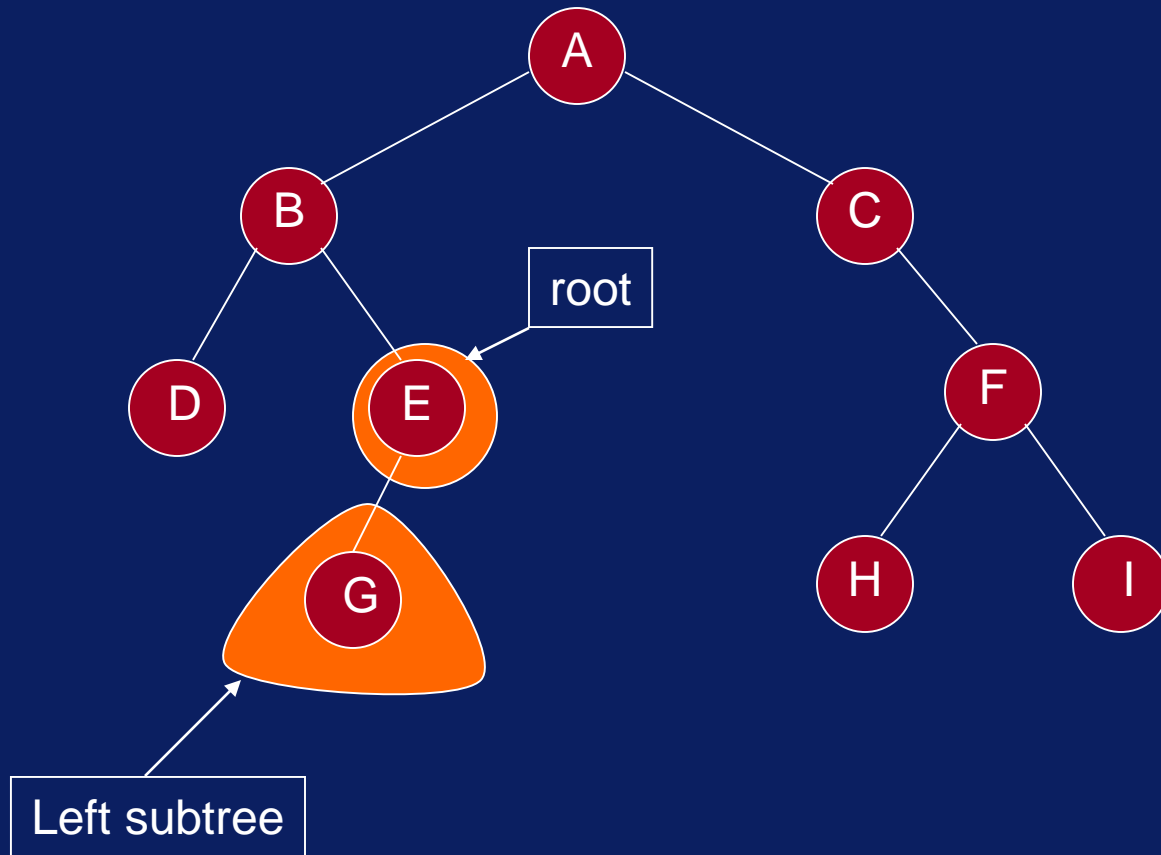
# Binary Tree

- Recursive definition



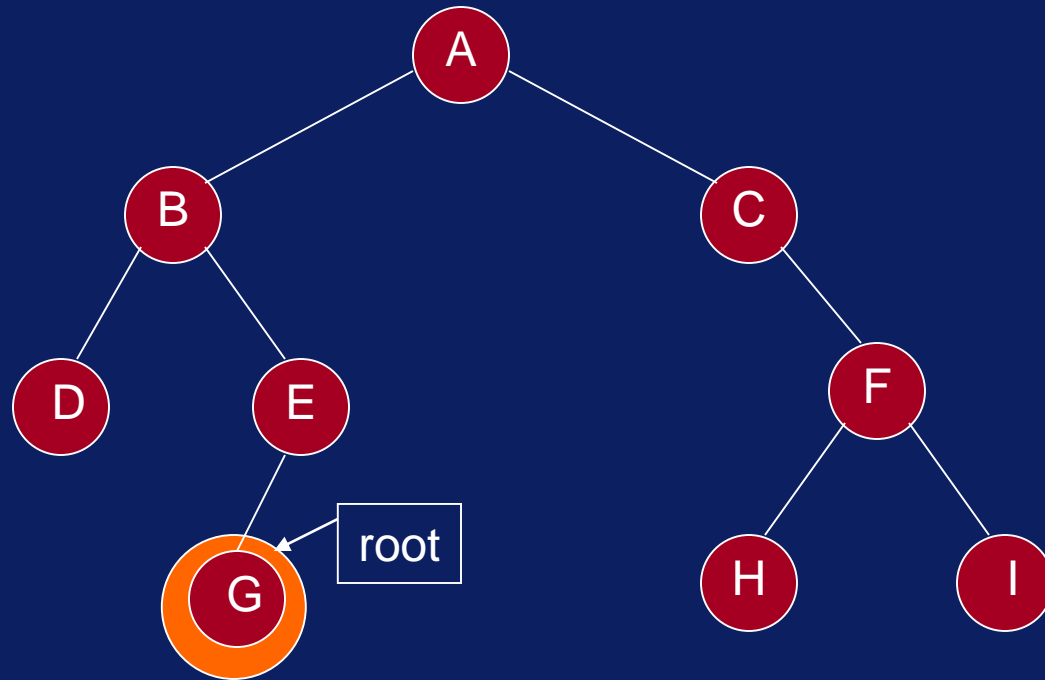
# Binary Tree

- Recursive definition



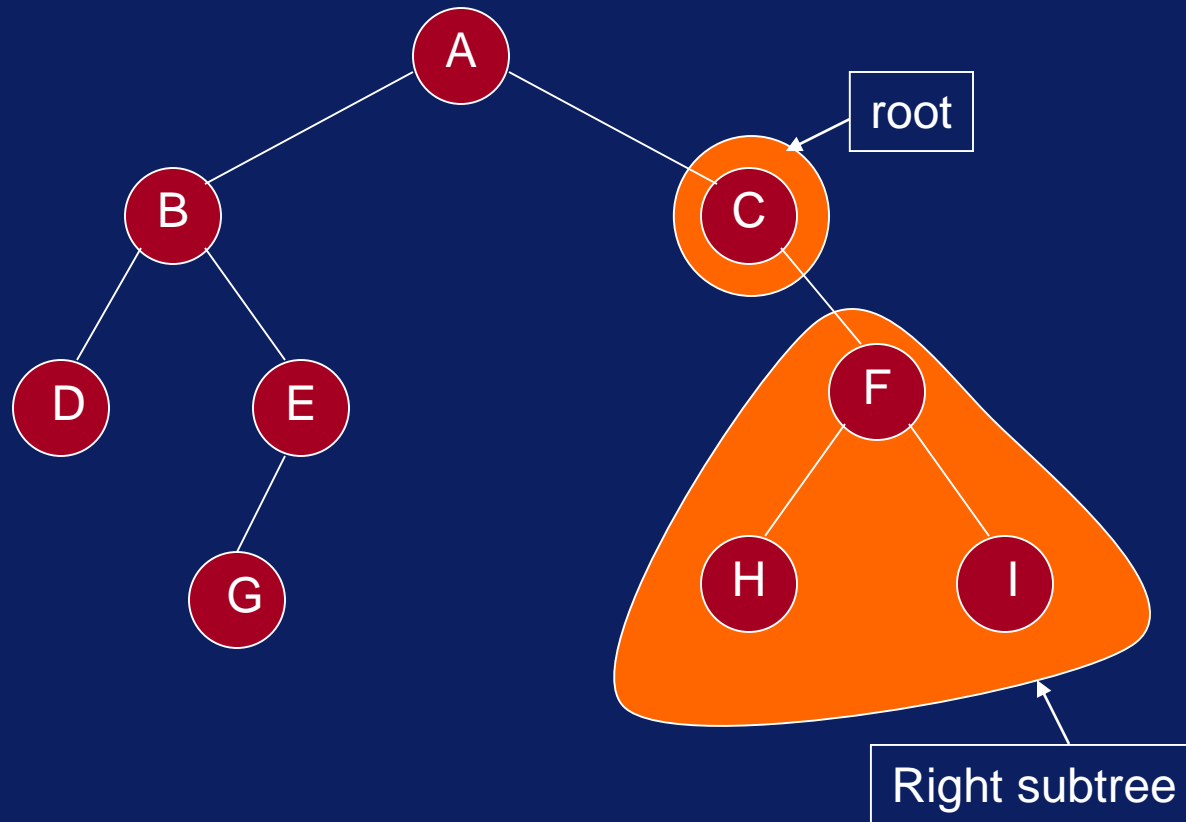
# Binary Tree

- Recursive definition



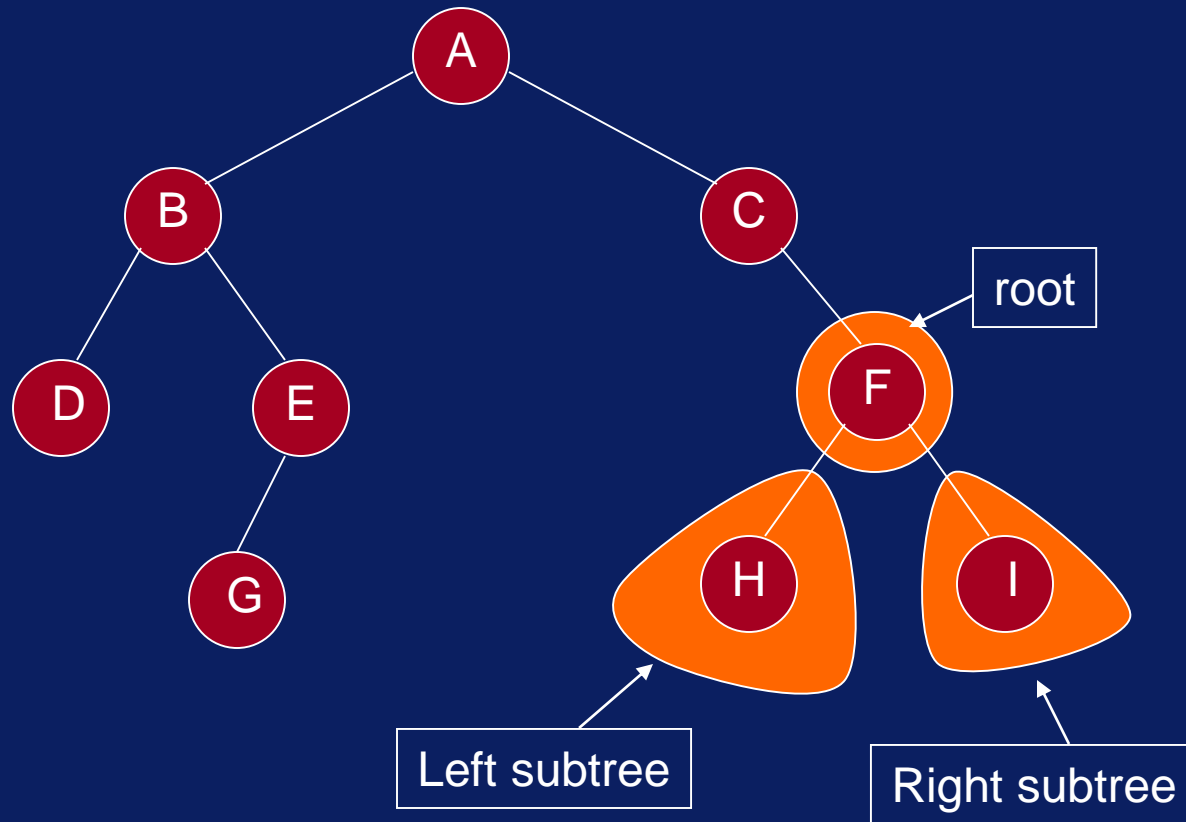
# Binary Tree

- Recursive definition



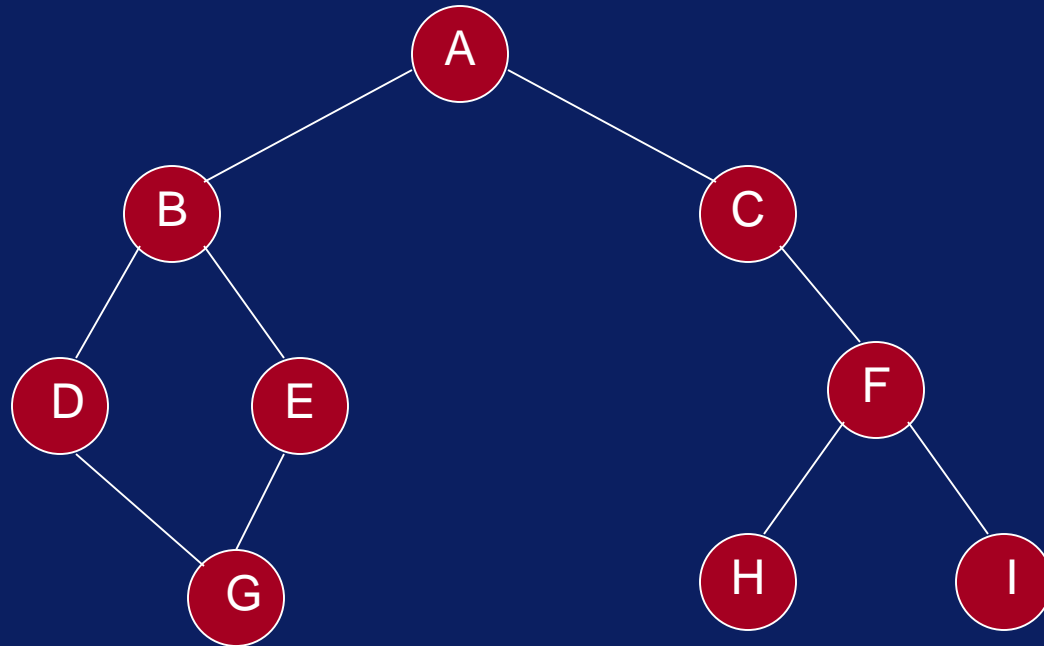
# Binary Tree

- Recursive definition



# Not a Tree

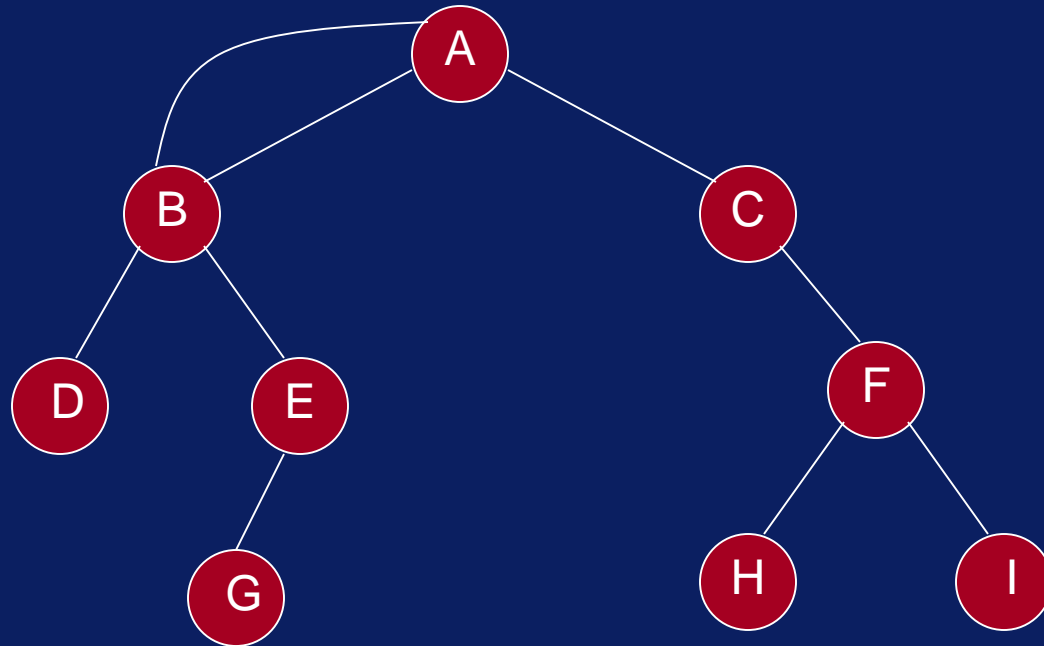
- Structures that are not trees.





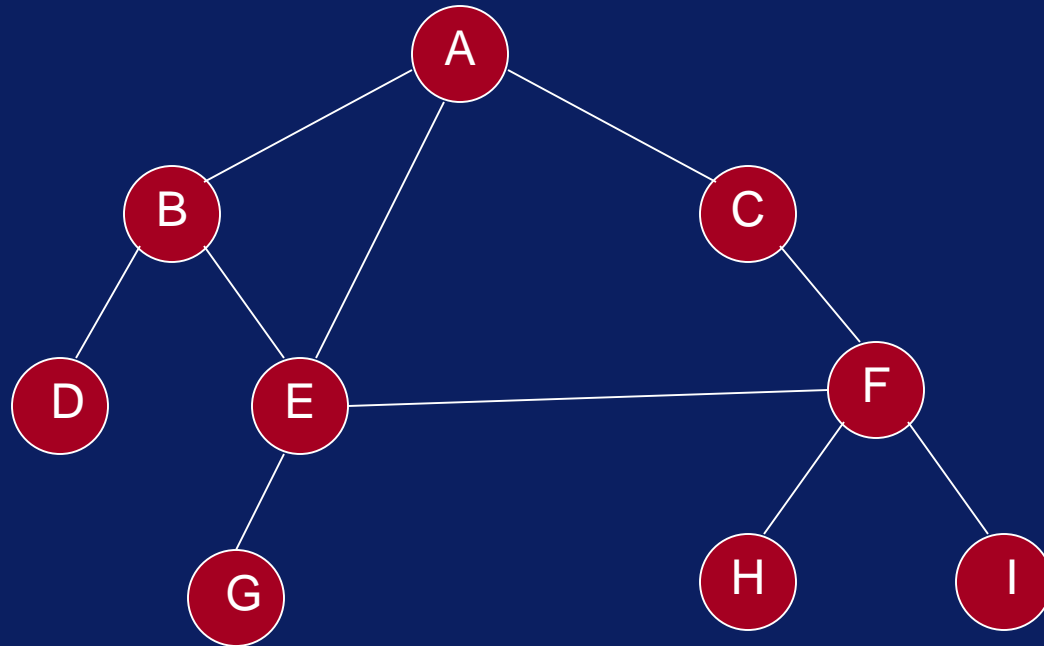
# Not a Tree

- Structures that are not trees.

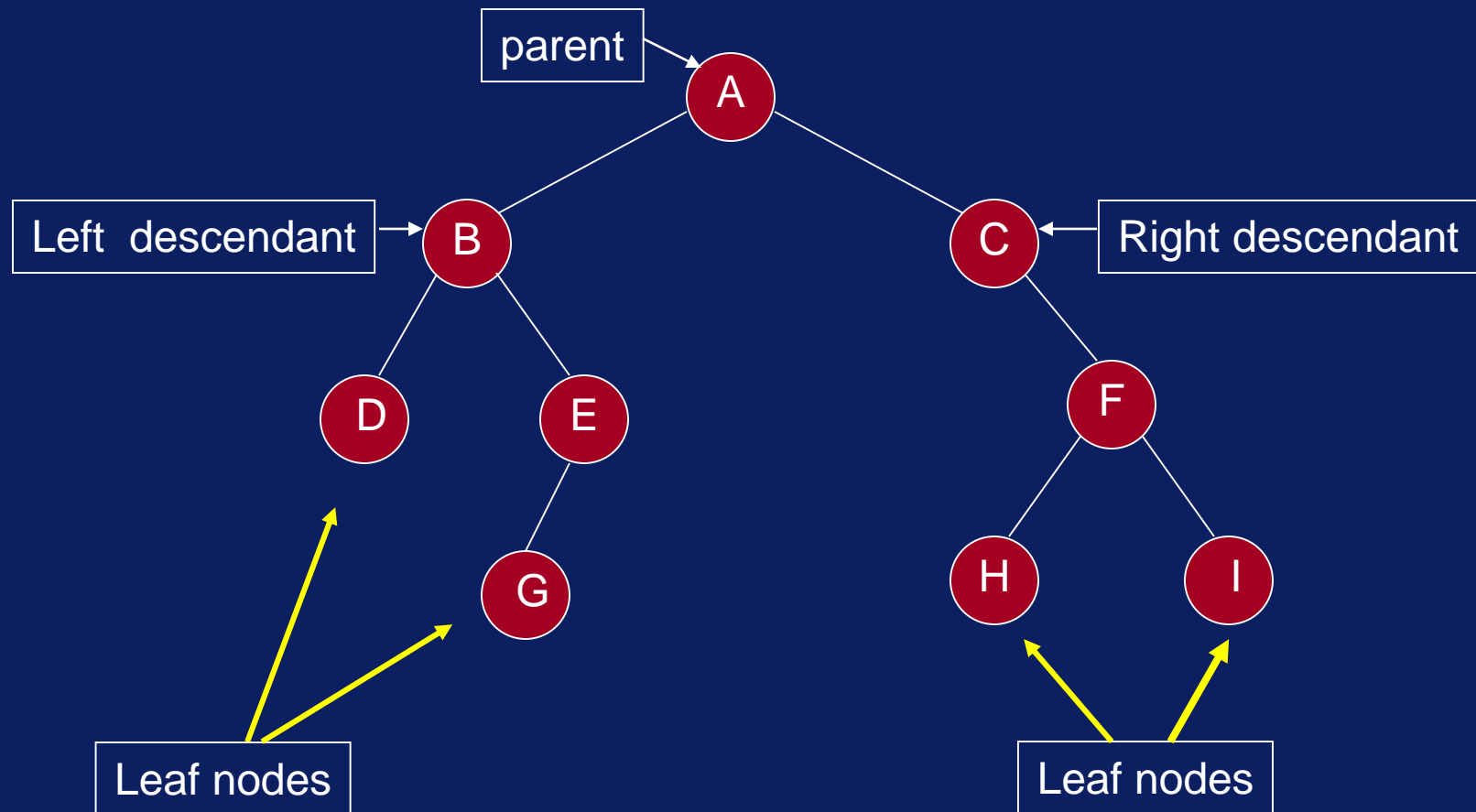


# Not a Tree

- Structures that are not trees.

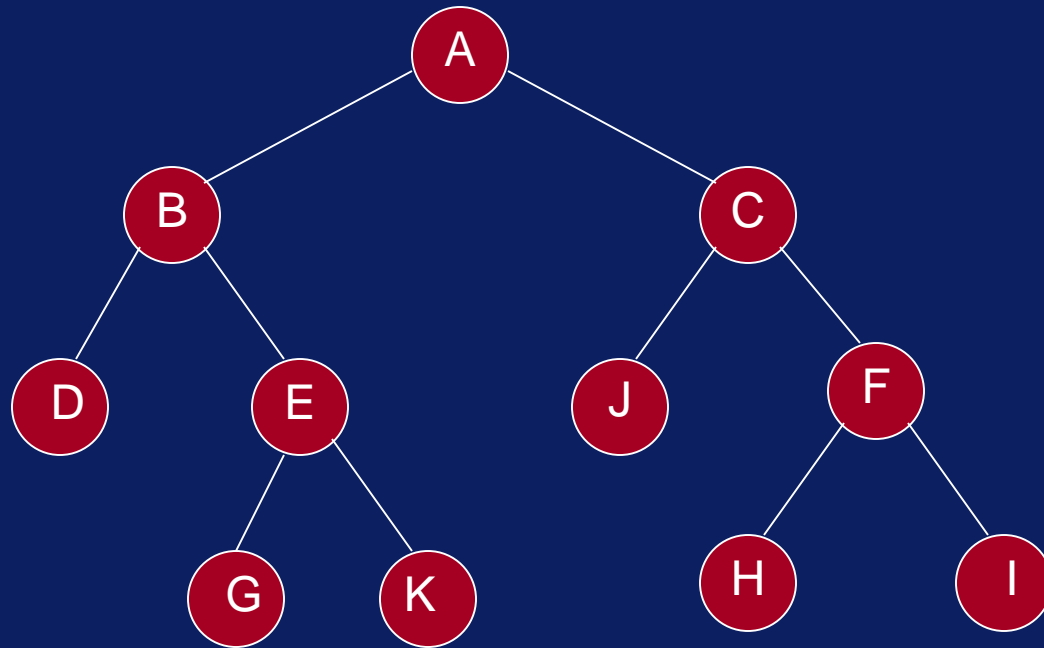


# Binary Tree: Terminology



# Binary Tree

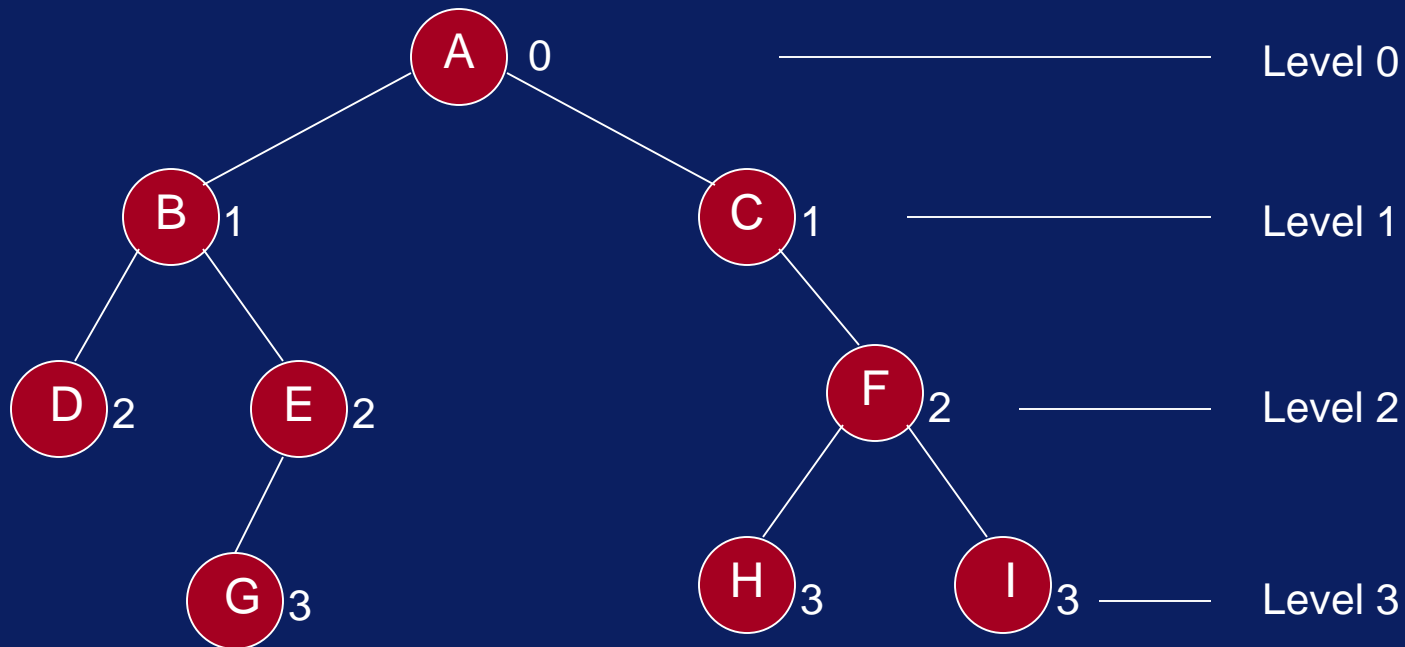
- If every non-leaf node in a binary tree has non-empty left and right subtrees, the tree is termed a *strictly binary tree*.



# Level of a Binary Tree Node

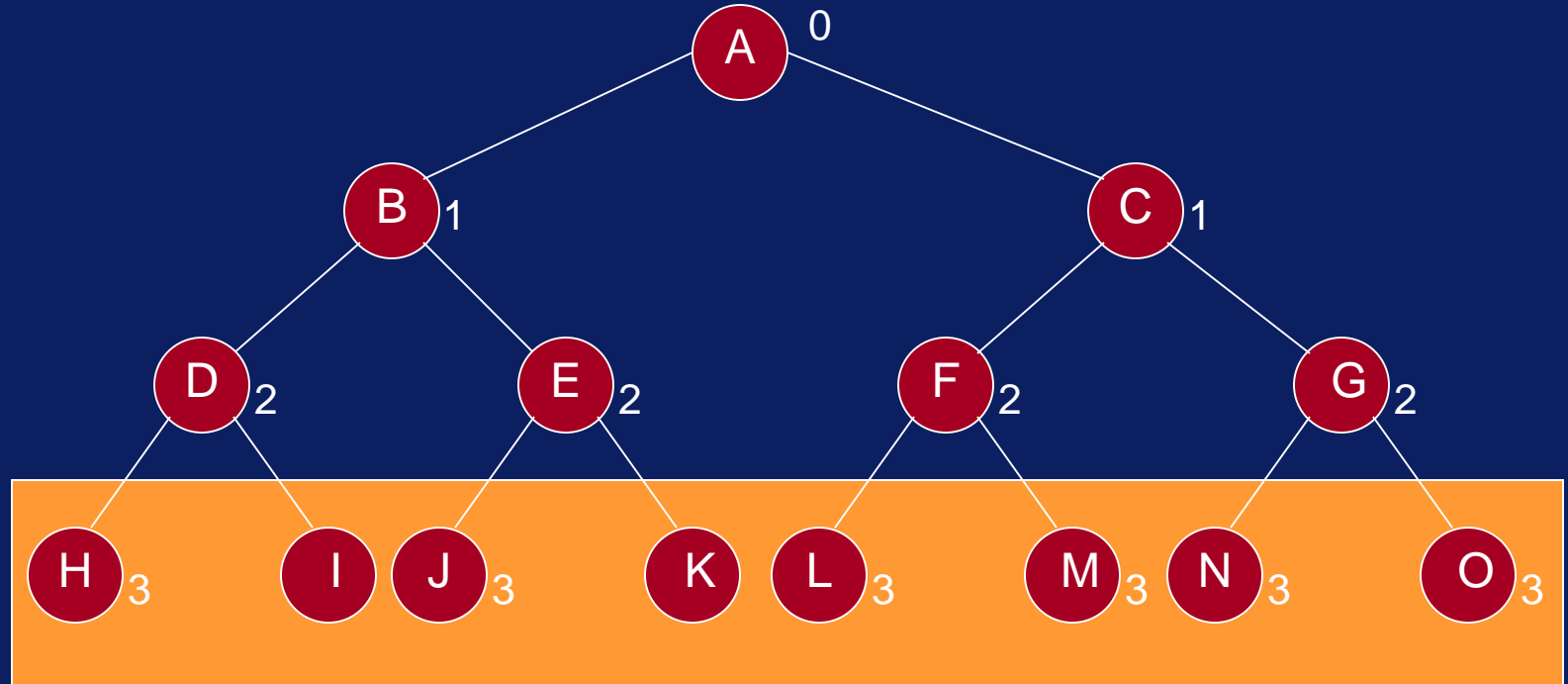
- The *level* of a node in a binary tree is defined as follows:
  - Root has level 0,
  - Level of any other node is one more than the level its parent (father).
- The *depth* of a binary tree is the maximum level of any leaf in the tree.

# Level of a Binary Tree Node

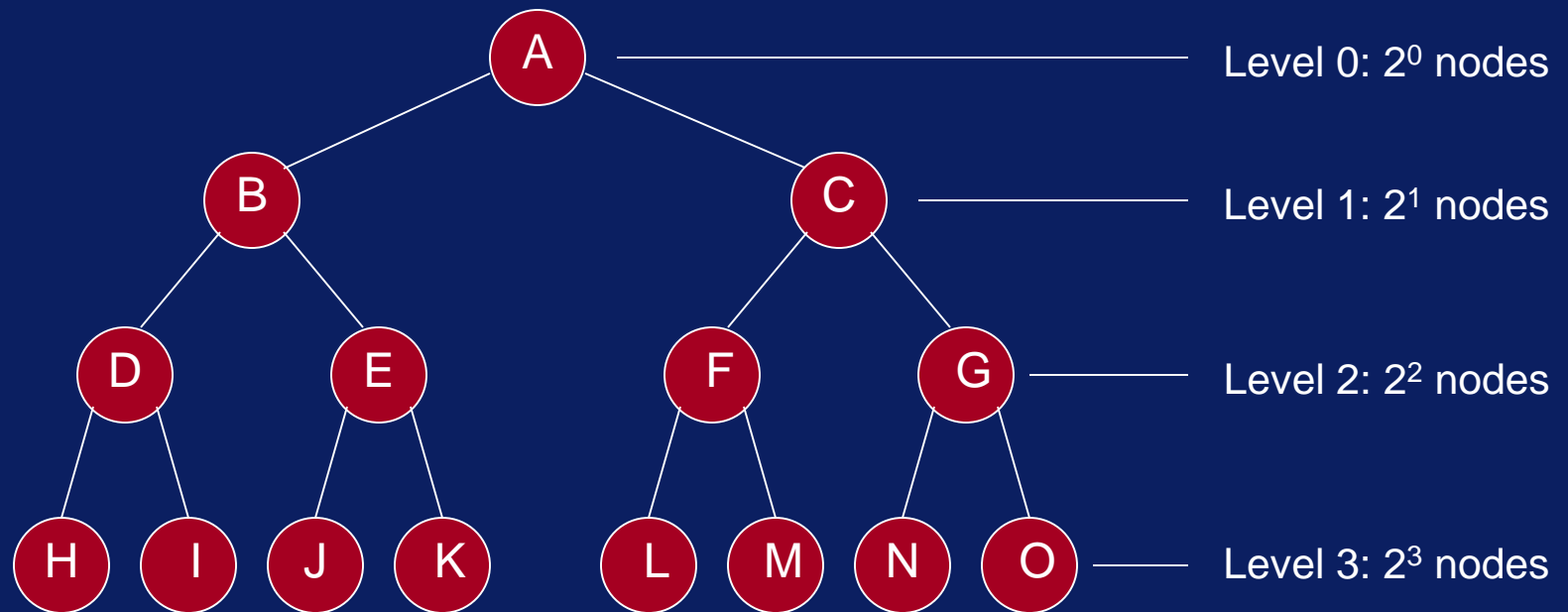


# Complete Binary Tree

- A *complete binary tree* of depth  $d$  is the strictly binary all of whose leaves are at level  $d$ .



# Complete Binary Tree





# Complete Binary Tree

- At level  $k$ , there are  $2^k$  nodes.
- Total number of nodes in the tree of depth  $d$ :

$$2^0 + 2^1 + 2^2 + \dots + 2^d = \sum_{j=0}^d 2^j = 2^{d+1} - 1$$

- In a complete binary tree, there are  $2^d$  leaf nodes and  $(2^d - 1)$  non-leaf (inner) nodes.

# Complete Binary Tree

- If the tree is built out of 'n' nodes then

$$n = 2^{d+1} - 1$$

$$\text{or } \log_2(n+1) = d+1$$

$$\text{or } d = \log_2(n+1) - 1$$

- I.e., the depth of the complete binary tree built using 'n' nodes will be  $\log_2(n+1) - 1$ .
- For example, for  $n=1,000,000$ ,  $\log_2(1000001)$  is less than 20; the tree would be 20 levels deep.
- The significance of this shallowness will become evident later.

# Operations on Binary Tree

- There are a number of operations that can be defined for a binary tree.
- If  $p$  is pointing to a node in an existing tree then
  - $\text{left}(p)$  returns pointer to the left subtree
  - $\text{right}(p)$  returns pointer to right subtree
  - $\text{parent}(p)$  returns the father of  $p$
  - $\text{brother}(p)$  returns brother of  $p$ .
  - $\text{info}(p)$  returns content of the node.