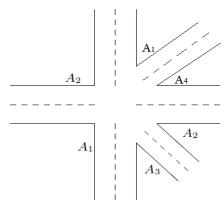


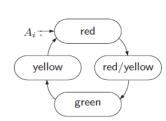
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Course Title: Formal Methods in SE Session: 2015-2019 **Course Code:** SE-364 **Semester:** Spring-2018 Engr. Afzal Ahmed **Deadline:** 18/04/2018 **Instructor:** CLO: CLO-1 Max Marks: **50**

Assignment 01

Q1. Please consider the following street junction with the specification of a traffic light as outlined on the right.





- **a)** Choose appropriate actions and label the transitions of the traffic light transition system accordingly.
- **b**) Give the transition system representation of a (reasonable) controller C that switches the green signal lamps in the following order: A_1 , A_2 , A_3 , A_4 , A_1 , A_2 , A_3 , A_4 . . . (*Hint: Choose an appropriate communication mechanism.*)
- c) Outline the transition system $A_1/|A_2/|A_3/|A_4|/C$.
- **Q2.** The following program is a mutual exclusion protocol for two processes due to Pnueli [1]. There is a single shared variable s which is either 0 or 1, and initially 1. Besides, each process has a local Boolean variable s that initially equals 0. The program text for process s s s follows:

loop forever do

begin

11: Noncritical section

12: $(y_i, s) := (1, i);$

13: **wait until** $((y_{1-i} = 0) \ V(s \neq i));$

14: Critical section

15: $y_i := 0$

end.

Here, the statement $(y_i, s) := (1, i)$; is a *multiple assignment* in which variable $y_i := 1$ and s := I is a single, atomic step.

Please answer following questions:

- a) Define the program graph of a process in Pnueli's algorithm.
- **b)** Determine the transition system for each process.
- c) Construct their parallel composition.



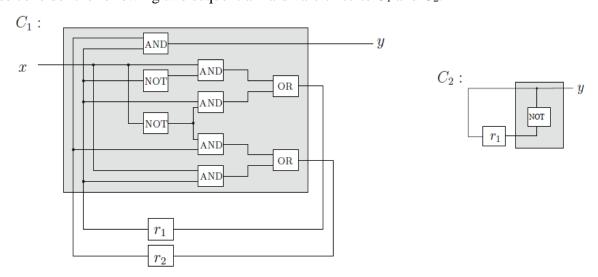
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- **d)** Check whether the algorithm ensures mutual exclusion.
- e) Check whether the algorithm ensures starvation freedom.
- f) The last two questions may be answered by inspecting the transition system.
- Q3. Please consider the following mutual exclusion algorithm that uses the shared variables y_1 and y_2 (initially both 0).

```
Process P_1:
                                                              Process P_2:
while true do
                                                              while true do
                                                                   ... noncritical section ...
     ... noncritical section . . .
     y_1 := y_2 + 1;
                                                                   y_2 := y_1 + 1;
     wait until (y_2 = 0) \ V(y_1 < y_2)
                                                                   wait until (y_1 = 0) \ V(y_2 < y_1)
     ... critical section ...
                                                                   ... critical section . . .
     y_1 := 0;
                                                                   y_2 := 0;
od
                                                              od
```

Please answer following questions:

- a) Give the program graph representations of both processes. (A pictorial representation suffices.)
- **b)** Give the reachable part of the transition system of P1 // P2 where $y1 \le 2$ and $y2 \le 2$.
- c) Describe an execution that shows that the entire transition system is infinite.
- **d)** Check whether the algorithm indeed ensures mutual exclusion.
- e) Check whether the algorithm never reaches a state in which both processes are mutually waiting for each other.
- f) Is it possible that a process that wants to enter the critical section has to wait ad infinitum?
- **Q4.** Please consider the following two sequential hardware circuits C_1 and C_2 :



- a) Give the transition system representation $TS(C_1)$ of the circuit C_1 .
- **b**) Let $TS(C_2)$ be the transition system of the circuit C_2 . Outline the transition system $TS(C_1) \otimes TS(C_2)$.



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Q5. Please consider the following leader election algorithm: For $n \in N$, n processes P_1, \ldots, P_n are located in a ring topology where each process is connected by an unidirectional channel to its neighbor in a clockwise manner.

To distinguish the processes, each process is assigned a unique identifier id $\in \{1, ..., n\}$. The aim is to elect the process with the highest identifier as the leader within the ring. Therefore each process executes the following algorithm:

```
send (id);initially set to process' idwhile (true) doreceive (m);if (m = id) then stop;process is the leaderif (m > id) then send (m);forward identifierod
```

- a) Model the leader election protocol for n processes as a channel system.
- b) Give an initial execution fragment of $TS([P_1|P_2|P_3|P_4])$ such that at least one process has executed the send statement within the body of the whileloop. Assume for $0 < i \le 3$, that process P_i has identifier $id_i = i$.

Good Luck ©



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Reference(s)

[1] W.-P. de Roever and F. S. de Boer and U. Hannemann and J. Hooman and Y. Lakhnech and M. Poel and J. Zwiers. *Concurrency Verification: Introduction to Compositional and Noncompositional Methods*. Number 54 in Cambridge Tracts in Theoretical Computer Science. Cambridge University Press, 2001.