Lecture No.07

Data Structures

Infix to Postfix

Infix	Postfix
A+B	AB+
12 + 60 – 23	12 60 + 23 -
(A + B)*(C - D)	AB+CD-*
$A \uparrow B * C - D + F/F$	$AB \uparrow C*D - FF/+$

Infix to Postfix

- Note that the postfix form an expression does not require parenthesis.
- Consider '4+3*5' and '(4+3)*5'. The parenthesis are not needed in the first but they are necessary in the second.
- The postfix forms are:

- Each operator in a postfix expression refers to the previous two operands.
- Each time we read an operand, we push it on a stack.
- When we reach an operator, we pop the two operands from the top of the stack, apply the operator and push the result back on the stack.

```
Stack s;
while( not end of input ) {
   e = get next element of input
   if( e is an operand )
        s.push(e);
   else {
        op2 = s.pop();
        op1 = s.pop();
        value = result of applying operator 'e' to op1 and op2;
        s.push( value );
finalresult = s.pop();
```

```
Evaluate 6 2 3 + - 3 8 2 / + * 2 ↑ 3 +

Input op1 op2 value stack

6
```

Evaluate 6 2 3 + - 3 8 2 / + * 2 ↑ 3 +

Input op1 op2 value stack

6
2
6,2

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1
3	6	5	1	1,3

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1
3	6	5	1	1,3
8	6	5	1	1,3,8

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1
3	6	5	1	1,3
8	6	5	1	1,3,8
2	6	5	1	1,3,8,2

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1
3	6	5	1	1,3
8	6	5	1	1,3,8
2	6	5	1	1,3,8,2
1	8	2	4	1,3,4

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1
3	6	5	1	1,3
8	6	5	1	1,3,8
2	6	5	1	1,3,8,2
1	8	2	4	1,3,4
+	3	4	7	1,7

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1
3	6	5	1	1,3
8	6	5	1	1,3,8
2	6	5	1	1,3,8,2
1	8	2	4	1,3,4
+	3	4	7	1,7
*	1	7	7	7

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1
3	6	5	1	1,3
8	6	5	1	1,3,8
2	6	5	1	1,3,8,2
1	8	2	4	1,3,4
+	3	4	7	1,7
*	1	7	7	7
2	1	7	7	7,2

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1
3	6	5	1	1,3
8	6	5	1	1,3,8
2	6	5	1	1,3,8,2
1	8	2	4	1,3,4
+	3	4	7	1,7
*	1	7	7	7
2	1	7	7	7,2
↑	7	2	49	49

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1
3	6	5	1	1,3
8	6	5	1	1,3,8
2	6	5	1	1,3,8,2
1	8	2	4	1,3,4
+	3	4	7	1,7
*	1	7	7	7
2	1	7	7	7,2
\uparrow	7	2	49	49
3	7	2	49	49,3

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1
3	6	5	1	1,3
8	6	5	1	1,3,8
2	6	5	1	1,3,8,2
1	8	2	4	1,3,4
+	3	4	7	1,7
*	1	7	7	7
2	1	7	7	7,2
↑	7	2	49	49
3	7	2	49	49,3
+	49	3	52	52

Input	op1	op2	value	stack
6				6
2				6,2
3				6,2,3
+	2	3	5	6,5
-	6	5	1	1
3	6	5	1	1,3
8	6	5	1	1,3,8
2	6	5	1	1,3,8,2
1	8	2	4	1,3,4
+	3	4	7	1,7
*	1	7	7	7
2	1	7	7	7,2
↑	7	2	49	49
3	7	2	49	49,3
+	49	3	52	52

- Consider the infix expressions 'A+B*C' and ' (A+B)*C'.
- The postfix versions are 'ABC*+' and 'AB+C*'.
- The order of operands in postfix is the same as the infix.
- In scanning from left to right, the operand 'A' can be inserted into postfix expression.

- The '+' cannot be inserted until its second operand has been scanned and inserted.
- The '+' has to be stored away until its proper position is found.
- When 'B' is seen, it is immediately inserted into the postfix expression.
- Can the '+' be inserted now? In the case of 'A+B*C' cannot because * has precedence.

- In case of '(A+B)*C', the closing parenthesis indicates that '+' must be performed first.
- Assume the existence of a function 'prcd(op1,op2)' where op1 and op2 are two operators.
- Prcd(op1,op2) returns TRUE if op1 has precedence over op2, FASLE otherwise.

- prcd('*','+') is TRUE
- prcd('+','+') is TRUE
- prcd('+','*') is FALSE
- Here is the algorithm that converts infix expression to its postfix form.
- The infix expression is without parenthesis.

```
1.
     Stack s;
     While( not end of input ) {
2.
3.
       c = next input character;
4.
       if( c is an operand )
5.
          add c to postfix string;
6.
       else {
          while( !s.empty() && prcd(s.top(),c) ){
7.
8.
             op = s.pop();
9.
             add op to the postfix string;
10.
           }
           s.push( c );
11.
12.
       }
13.
       while(!s.empty()) {
14.
          op = s.pop();
          add op to postfix string;
15.
        }
16.
```

Example: A + B * C

symb postfix stack

A

```
symb postfix stack
A A +
```

symb	postfix	stack
Α	A	
+	A	+
В	AB	+

symb	postfix	stack
A	Α	
+	Α	+
В	AB	+
*	AB	+ *

symb	postfix	stack
A	Α	
+	Α	+
В	AB	+
*	AB	+ *
С	ABC	+ *

symb	postfix	stack
Α	A	
+	A	+
В	AB	+
*	AB	+ *
С	ABC	+ *
	ABC *	+

symb	postfix	stack
Α	Α	
+	Α	+
В	AB	+
*	AB	+ *
С	ABC	+ *
	ABC *	+
	ABC * +	

- Handling parenthesis
- When an open parenthesis '(' is read, it must be pushed on the stack.
- This can be done by setting prcd(op, '(') to be FALSE.
- Also, prcd('(',op) == FALSE which ensures that an operator after '(' is pushed on the stack.

- When a ')' is read, all operators up to the first '(' must be popped and placed in the postfix string.
- To do this, prcd(op,')') == TRUE.
- Both the '(' and the ')' must be discarded: prcd('(',')') == FALSE.
- Need to change line 11 of the algorithm.

```
if( s.empty() || symb != ')' )
     s.push( c );
else
    s.pop(); // discard the '('
prcd('(', op ) = FALSE for any operator
prcd( op, '(' ) = FALSE
                           for any operator
                            other than '('
prcd( op, ')' ) = TRUE
                            for any operator
                            other than '('
                            for any operator.
prcd( ')', op ) = error
```