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Assignment 2, Problem 1. c) and d)

```
clc; close all; clear all;
epsilon = 1;
```

```
while 1
    if (epsilon+1) <= 1
        break;
    end
    epsilon = epsilon/2;
end
```

```
epsilon = epsilon*2;
fprintf('Machine Epsilon for my computer is %f\n', epsilon);
```

```
%{
In this implementation, the value of epsilon is divided by 2 through
each pass in the otherwise infinite while loop.
After each division, the if statement checks to see if adding 1 to
epsilon
would have it be less than or equal to 1. Adding a very small number
to 1
would result in the processor still equating that to 1 in the case
that
epsilon is so small, that the only value stored in the mantissa would
be
equal to 1. This kind of roundoff error can occur when adding an
extremely small number to a relatively larger error.
If epsilon+1 does equal 1, that means we divided by 2 one too many
times;
so, we break out of the while loop, and multiply by 2. The resulting
value
of epsilon is then the value of Machine Epsilon for this specific
computer.
%}
```

```
Machine Epsilon for my computer is 0.000000
```

Problem 2. Calculations only

```
clc; close all; clear all;
format long;
upper_bound_error = (.5 * 10^-6);
```

```

x = 0.3*pi;
true = sin(0.3*pi);
ml_approx = 0;
ml = [];
true_error = [];
approx_error = [-100000]; %placeholder since first approx. error = N/A
n = 0;

% I used the vectors stored in the workspace variables to fill out my
% table
% in the handwritten assignment.

while 1
    ml_approx = ml_approx + ((-1)^n)*(x^(2*n + 1)/(factorial(2*n+1)));
    ml(n+1) = ml_approx;
    true_error(n+1) = abs((true - ml_approx)/true * 100);
    if n ~= 0
        approx_error(n+1) = abs((ml(n+1) - ml(n))/ml(n+1) * 100);
        if approx_error(n+1) <= upper_bound_error
            break
        end
    end
    n = n+1;
end
fprintf('The number of iterations to approximate to 8 sig figs is %.0f\n', n+1);
%Total iterations = n+1, since n starts at 0.

```

The number of iterations to approximate to 8 sig figs is 7

Problem 3

```

clc; close all; clear all;
format long;
sumofe = 0;
sum1 = [];
for n = 0:10
    sumofe = single(sumofe + 1/factorial(n));
    sum1(n+1) = sumofe;
end
te1 = abs((exp(1) - sumofe)/exp(1))*100;
fprintf('The true percent error for e, going from 0 to 10 is %.8f\n', te1);

sumofe = 0;
sum2 = [];
for n = 10:-1:0
    sumofe = single(sumofe + 1/factorial(n));
    sum2(-n+11) = sumofe;
end
te2 = abs((exp(1) - sumofe)/exp(1))*100;
fprintf('The true percent error for e, going from 10 to 0 is %.8f\n', te2);

```

```
%{
The difference between these two percent errors can be explained if
you
consider how a computer is processing these terms. Taking a look at
the
vectors sum1 and sum2 will show that sum2 starts with numbers to the
magnitude of  $10^{-7}$ . The computer can store a certain number of
significant
digits in its mantissa, and it does not have to store all the leading
zeros
in the miniscule number, and instead saves it in scientific notation.
As
sumofe gets gradually larger, the summation is more precise, as the
computer is saving as many significant digits as possible before
rounding
off.
On the other hand, looking at sum1, we can see that from the first
term,
the computer has a digit before the leading zeros to store as well,
which
decreases the number of significant digits the computer can hold per
term,
an issue that is only exacerbated as we try and add more and more
small
yet significant terms to the sum. An example of this can be seen by
looking
at sum1(3), which should hold 1.66666666666666 as long as it can before
the
last digit is changed to a 7. However, the computer is actually
storing
this as 1.6666666626930237. As we add more and more terms to sumofe,
the
roundoff error compounds, to give a slightly larger error.
%}
```

The true percent error for e, going from 0 to 10 is 0.00000573 percent
The true percent error for e, going from 10 to 0 is 0.00000304 percent

Problem 4, part b)

```
clc; close all; clear all;
x = linspace(-pi, pi, 100);
e = exp(1);
tanhx = (e.^x - e.^(-x))./(e.^x + e.^(-x));

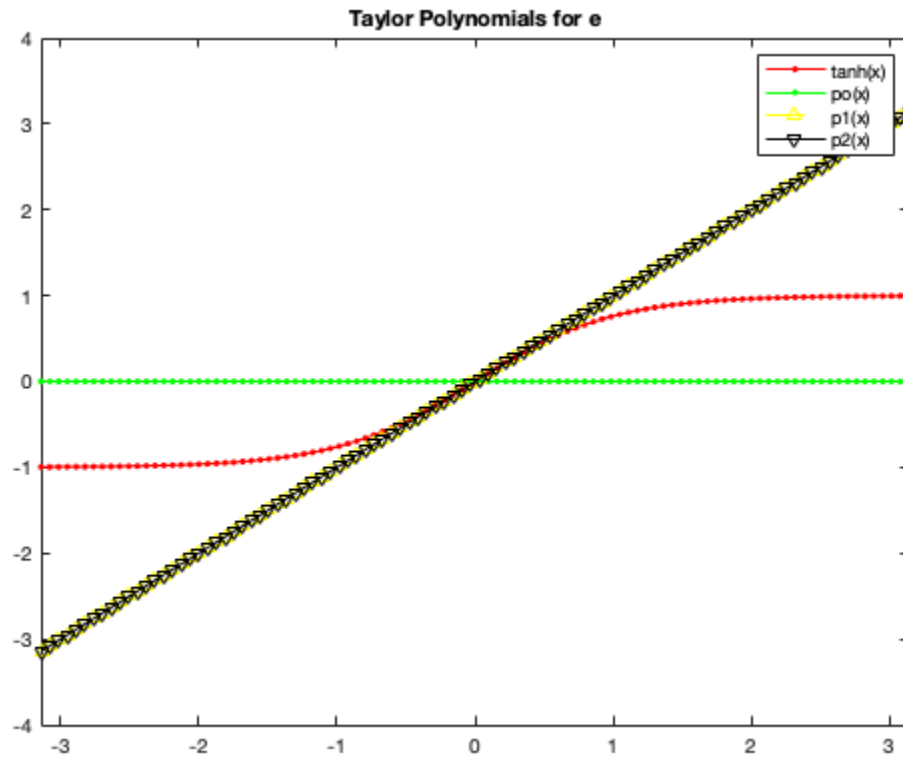
p0 = zeros([1,100]);
p1 = p0 + x;
p2 = p1 + 0.*x.^2;

plot(x,tanhx, 'r.-');
hold on
plot(x, p0, 'g.-');
```

```
plot(x,p1, 'y^--');
plot(x,p2, 'kv-');

title('Taylor Polynomials for e');
xlim([-pi pi]);

legend('tanh(x)', 'po(x)', 'p1(x)', 'p2(x)');
```



Published with MATLAB® R2020a

1. a) Machine Epsilon is a convergence criteria used as an upper bound for a given error. Machine Epsilon (ϵ) is specific to a given computer's processor, and is dependent on the number of significant digit bits the processor chip can store. In other words, the computed error should never be more than the computer's ϵ .
- b) ϵ for a 64-bit processor: $\epsilon = 2^{1-53} = 2^{-52} = 2.204 \times 10^{-16}$
- c) See Matlab code on page 3 (page 1 of MATLAB published code)
- d) See comments on Matlab Code on page 3
- e) Yes, my computer has a 64-bit processor, as my machine epsilon was equal to 2.204×10^{-16}

2. Approximate $\sin(0.3\pi)$ to 8 significant digits. **ANSWER: 7 ITERATIONS**

$$\epsilon_s = (0.5 \times 10^{2-8}) \% = 5 \times 10^{-7} \% ; \text{Continue iterating until } |\epsilon_a| \leq \epsilon_s$$

$$\sin(x) = \sum_{n=0}^{\infty} (-1)^n \frac{x^{2n+1}}{(2n+1)!} = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots$$

True Value: $\sin(0.3\pi) = 0.80901699$; See attached Matlab Code on pages 3 & 4

Iteration	Approx by ML	$\epsilon_t = \frac{\text{true} - \text{approx}}{\text{true}} \times 100\%$	$\epsilon_a = \frac{\text{current} - \text{prev}}{\text{current}}$
1	$(0.3\pi) / 1!$ $= 0.942477796$	$\frac{0.8090 - 0.9425}{0.8090} \times 100\%$ $= 16.4967\%$	N/A
2	0.80294955	0.7500%	17.377%
3	0.80914645	0.0160%	0.7659%
4	0.80901539	$1.9825 \times 10^{-4} \%$	0.0162%
5	0.80901701	$1.6047 \times 10^{-6} \%$	$1.9986 \times 10^{-4} \%$
6	0.80901699	$9.1507 \times 10^{-9} \%$	$1.6139 \times 10^{-6} \%$
7	0.80901699	$3.8727 \times 10^{-11} \%$	$9.1894 \times 10^{-9} \%$

$$4. \tanh(x) = \frac{e^x - e^{-x}}{e^x + e^{-x}}; \quad \frac{d}{dx} (\tanh(x)) = 1 - \tanh^2(x); \quad \frac{d^2}{dx^2} (\tanh(x)) = -2 \tanh(x) (1 - \tanh^2(x))$$

$$a) \text{0th term} = \tanh(0) = \tanh(0) = 0 \rightarrow p_0(x) = 0x^0$$

$$\text{1st term} = \tanh(0) + (1 - \tanh^2(0)) (x - 0) = \tanh(0) + x(1 - \tanh^2(0)) = x$$

$$\hookrightarrow p_1(x) = 0x^0 + x$$

$$\text{2nd term} = \frac{-2 \tanh(0) (1 - \tanh^2(0)) (x - 0)^2}{2!} = -x^2 \tanh(0) (1 - \tanh^2(0)) = 0$$

$$\hookrightarrow p_2(x) = 0x^0 + x^1 + 0x^2$$

$$p_0(x) = 0x^0$$

$$p_1(x) = 0x^0 + 1x^1$$

$$p_2(x) = 0x^0 + 1x^1 + 0x^2$$

b) See MATLAB code on pages 5 & 6