

Practice Exam 1

Tables Provided on Exam 1

Table 1.5.1: Laws of propositional logic.

Idempotent laws:	$p \vee p = p$	$p \wedge p = p$
Associative laws:	$(p \vee q) \vee r = p \vee (q \vee r)$	$(p \wedge q) \wedge r = p \wedge (q \wedge r)$
Commutative laws:	$p \vee q = q \vee p$	$p \wedge q = q \wedge p$
Distributive laws:	$p \vee (q \wedge r) = (p \vee q) \wedge (p \vee r)$	$p \wedge (q \vee r) = (p \wedge q) \vee (p \wedge r)$
Identity laws:	$p \vee F = p$	$p \wedge T = p$
Domination laws:	$p \wedge F = F$	$p \vee T = T$
Double negation law:	$\neg\neg p = p$	
Complement laws:	$p \wedge \neg p = F$ $\neg T = F$	$p \vee \neg p = T$ $\neg F = T$
De Morgan's laws:	$\neg(p \vee q) = \neg p \wedge \neg q$	$\neg(p \wedge q) = \neg p \vee \neg q$
Absorption laws:	$p \vee (p \wedge q) = p$	$p \wedge (p \vee q) = p$
Conditional identities:	$p \rightarrow q = \neg p \vee q$	$p \leftrightarrow q = (p \rightarrow q) \wedge (q \rightarrow p)$

Table 1.12.1: Rules of inference known to be valid arguments.

Rule of inference	Name
$\frac{p \quad p \rightarrow q}{\therefore q}$	Modus ponens
$\frac{\neg q \quad p \rightarrow q}{\therefore \neg p}$	Modus tollens
$\frac{p}{\therefore p \vee q}$	Addition
$\frac{p \wedge q}{\therefore p}$	Simplification

Rule of inference	Name
$\frac{p \quad q}{\therefore p \wedge q}$	Conjunction
$\frac{p \rightarrow q \quad q \rightarrow r}{\therefore p \rightarrow r}$	Hypothetical syllogism
$\frac{p \vee q \quad \neg p}{\therefore q}$	Disjunctive syllogism
$\frac{p \vee q \quad \neg p \vee r}{\therefore q \vee r}$	Resolution

Table 1.13.1: Rules of inference for quantified statements

Rule of Inference	Name
c is an element (arbitrary or particular) $\forall x P(x)$ $\therefore P(c)$	Universal instantiation
c is an arbitrary element $P(c)$ ____ $\therefore \forall x P(x)$	Universal generalization
$\exists x P(x)$ $\therefore (c \text{ is a particular element}) \wedge P(c)$	Existential instantiation*
c is an element (arbitrary or particular) $P(c)$ ____ $\therefore \exists x P(x)$	Existential generalization

Table 3.6.1: Set identities.

Name	Identities	
Idempotent laws	$A \cup A = A$	$A \cap A = A$
Associative laws	$(A \cup B) \cup C = A \cup (B \cup C)$	$(A \cap B) \cap C = A \cap (B \cap C)$
Commutative laws	$A \cup B = B \cup A$	$A \cap B = B \cap A$
Distributive laws	$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$	$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$
Identity laws	$A \cup \emptyset = A$	$A \cap U = A$
Domination laws	$A \cap \emptyset = \emptyset$	$A \cup U = U$
Double Complement law	$\overline{\overline{A}} = A$	
Complement laws	$A \cap \overline{A} = \emptyset$ $\overline{\overline{U}} = U$	$A \cup \overline{A} = U$ $\overline{\emptyset} = U$
De Morgan's laws	$\overline{A \cup B} = \overline{A} \cap \overline{B}$	$\overline{A \cap B} = \overline{A} \cup \overline{B}$
Absorption laws	$A \cup (A \cap B) = A$	$A \cap (A \cup B) = A$

Set Theory

Given:

$$A = \{ 1, \{2\}, \{\{3, 4\}\} \}$$

For each of the following statements, state whether they are **true** or **false**.

- a. $1 \in A$
- b. $1 \subseteq A$
- c. $\{2\} \in A$
- d. $\{2\} \subseteq A$
- e. $\{3, 4\} \in A$
- f. $\{3, 4\} \subseteq A$
- g. $\{\{3,4\}\} \in A$
- h. $\{\{3,4\}\} \subseteq A$
- i. $\emptyset \in A$
- j. $\emptyset \subseteq A$

Solutions:

- a. T, the number 1 is an element of set A
- b. F, the number 1 is not contained within a set of braces, thus is not a subset of A
- c. T, the set containing 2 is an element of A (the number 2 is contained within a set of { } within A)
- d. F, and would only be T if $\{\{2\}\}$ was contained within A
- e. F, because $\{3,4\}$ would be need to contained within $\{\{ \}$ rather than within just $\{ \}$
- f. F, because $\{\{3,4\}\}$ would need to be an element contained within $\{\{\{ \}\}\}$
- g. T, the set containing 3 & 4 ($\{3, 4\}$) is an element of set A
- h. F
- i. F, the empty set is always a subset and not an element of the larger set
- j. T, the empty set is only a subset (always just a subset) of the larger set

1.2 Let $A = \{1, 2, 3, 4\}$. Select the statement that is **false**.

- a. $\emptyset \in P(A)$
- b. $\emptyset \subseteq P(A)$
- c. $\{2, 3\} \in A$
- d. $\{2, 3\} \subseteq A$

Solution:

C

Functions

Choose the property for which the function satisfies if well defined.

- a. Neither one-to-one, nor onto
- b. One-to-one, but not onto
- c. Onto, but not one-to-one
- d. one-to-one and onto
- e. not well defined

Given a function whose domain is the set of all integers and whose target is the set of all positive integers:

a)

$$f(x) = 2x + 1$$

Solution: e

b)

$$f(x) = |x| + 1$$

Solution: c

c)

$$f(x) = x^2 + 1$$

Solution: a

d)

$$f(x) = \{x > 0 : 2x + 1 \wedge x \leq 0 : -2x\}$$

Solution: e

e)

$$f(x) = \{x \geq 0 : 2x + 1 \wedge x < 0 : -2x + 2\}$$

Solution: b

Proofs

4.1 Direct Proof

Prove that the product of two odd integers is an odd integer.

Proof.

Assume that integers m and n are odd integers. We will show that $m \cdot n$ is also an odd integer.

Integers m and n are odd integers

$$\Rightarrow m = 2k+1 \text{ for some integer } k \text{ and } n = 2j + 1 \text{ for some integer } j$$

$$\Rightarrow m \cdot n = (2k+1)(2j+1) \text{ for some integers } k \text{ and } j$$

$$\Rightarrow \quad = 4kj + 2k + 2j + 1$$

$$\Rightarrow \quad = 2(kj + k + j) + 1$$

Since j and k are both integers, then $2(kj + k + j)$ is also an integer. Since $m \cdot n$ can be expressed as 2 times an integer plus 1, $m \cdot n$ is an odd integer.

4.2 Proof by Contrapositive

Prove that if n^2 is even, then n is even.

Consider the contraposition of the proposition which is $\sim q \rightarrow \sim p$. Show that if n is odd, then n^2 is odd.

Proof.

Assume the hypothesis of the implication to be true.

Assume that n is an odd integer.

Show the implication.

n is an odd integer

$$\Rightarrow n = 2q+1 \text{ for some integer } q$$

$$\Rightarrow n^2 = (2q + 1)^2 \text{ for some integer } q$$

$$\Rightarrow n^2 = 4q^2 + 4q + 1 \text{ for some integer } q$$

$$\Rightarrow n^2 = 2(2q^2 + 2q) + 1 \text{ for some integer } q \text{ and because } q \text{ is an integer, then } (2q^2 + 2q) \text{ is also an integer,}$$

$$\Rightarrow n^2 = 2j + 1 \text{ for some integer } j$$

Therefore n^2 is an odd integer that can be represented as $2j+1$.

Having used a direct proof of the contrapositive statement, we can conclude that if n^2 is even, then n is even.

4.2 Proof by Contradiction

Prove by contradiction that if $3n+5$ is odd, then n is even.

Proof.

Begin by assuming that $3n+5$ is odd and n is *odd*.

Integer n is odd therefore $n=2k+1$ for some integer k .

substituting $n=2k+1$ for n in $3n+5$,

$$3n+5 = 3(2k+1) + 5 \text{ for some integer } k$$

$$= 6k + 8 \text{ for some integer } k, \text{ algebra}$$

$$= 2(3k+4)$$

Let $k' = (3k+4)$ for some integer k ,

and since k is an integer, then $k' = (3k+4)$ is also an integer.

$$= 2(k')$$

Further, since $k' = (3k+4)$ is an integer, and 2 multiplied by an integer is an even integer, by definition $3(2k+1) + 5$ is even.

Thus, $3n+5$ is even and because this **contradicts** the assumption that $3n+5$ is odd then n is odd; we have proven that if $3n+5$ is odd, then n is even.

Number Systems Conversion

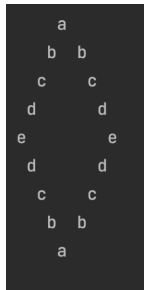
5.1 Decimal to 8-bit Two's Complement

$$(-43)_{10} = (11010101)_{8\text{-bit Two's Complement}}$$

5.2 Binary to Hexadecimal

$$(110011100)_2 = (19C)_{16}$$

Coding



Make a hollowed-out diamond made up of in-order alphabet letters. An example is below when $n=5$.

Solution:

```
#include <iostream>
using namespace std;

int main(){
    int userInput = 7;
    char letter = 'a';
    char space = ' ';

    int outsideSpaceLimit = userInput - 1;
    int insideSpaceLimit = 0;

    for (int i = 1; i < (userInput * 2); i++) {
        // outside spaces
        for (int j = 0; j < outsideSpaceLimit; j++) {
            cout << space;
        }
        // first letter
        cout << letter;

        //inside spaces
        for (int j = 0; j < insideSpaceLimit; j++) {
            cout << space;
        }

        //second letter
        if (i != 1 && i != ((userInput * 2) - 1)){
            cout << letter;
        }

        if(i < userInput){
            letter++;
            insideSpaceLimit += 2;
            outsideSpaceLimit--;
        }
    }
```



```
    else{
        letter--;
        insideSpaceLimit -= 2;
        outsideSpaceLimit++;
    }

    cout << endl;
}

return 0;
}
```

Challenge Question

For this question, you have to create the following Barn Door shape with `n=10`:

```
#####
# $      $ #
# $      $ #
# $ $ $ #
#  $ $ #
#  $ $ #
# $ $ $ #
# $      $ #
# $      $ #
#####
```

ONE answer (again, there are **many** ways to do this!):

```
#include <iostream>
using namespace std;

int main() {
    int input;

    char frame = '#';
    char planks = '$';
    char space = ' ';

    cout << "Enter positive integer: ";
    cin >> input;

    for (int i = 0; i < input; i++) { //rows
        for (int j = 0; j < input; j++) {
            // left or right or top or bottom
            if (j == 0 || j == input - 1 || i == 0 || i == input - 1) {
                cout << frame;
            }
            // left diagonal going right || right diagonal going left
            else if (i == j || i + j == input - 1) {
                cout << planks;
            }
            else {
                cout << space;
            }
        }
        cout << endl;
    }
}
```

```
    return 0;  
}
```