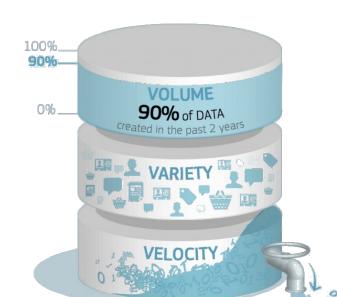




Historic view

- As of the rise of the Web 2.0:
 - Very large volumes of data being collected
 - Web logs first, then social media, web apps, etc.
 - Data from all type of sensors (phone, cars, ...)
 - Metadata from communication networks
 - Analytics on this data, of great value



- Big Data: different from earlier DBs by...
 - Volume: much larger amounts of data stored
 - Velocity: much higher data ingest rate
 - Variety: many data types, beyond relational data

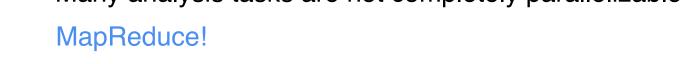


Introduction

- Large increase in the size of disks
- Access-to-disk time has not decreased proportionally
- Obvious solution: use multiple disks and work in parallel!
- **Problems:**
 - Hardware failures

HDFS!

Many analysis tasks are not completely parallelizable MapReduce!



Hadoop: General purpose storage and analysis platform for big data

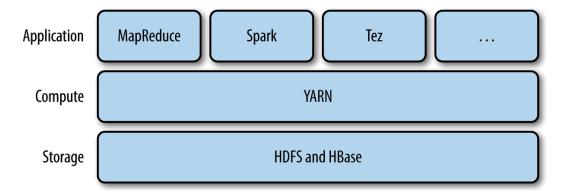




Hadoop

3 building blocks

- HDFS: distributed file system
- YARN: Resource (cluster) manager
- Map-reduce: parallel computing paradigm



- Splits data and computation across (thousands of) hosts
 - A Hadoop cluster scales computation capacity, storage capacity and I/O bandwidth by simply adding nodes.



Distributed File Systems

- Store files across many nodes while giving a single-system view
 - Need to address the distribution of files across nodes
 - Unified file system, with file names and directories
 - Users do not bother where the file actually is stored
 - Designed to store very large files
 10s MB to 100s GB
 - Highly scalable for large data-intensive applications
 - E.g., 10K nodes, 100 million files, 10 PB
 - E.g., Google File System (GFS), Hadoop File System (HDFS), CODA, Google Colossus



Distributed File Systems

- Files, divided into blocks
- Remember the concept of blocks in the context of disks:
 - Minimum amount of data that a disk can read/write.
 - Usually, disk blocks: 512 bytes; regular filesystem blocks: few KBs
- Blocks, file's chunks, are stored as independent units
 - Spread through nodes, replicated for availability
 - Much larger unit: 128 MB by default.
 - Why so large?
 - To minimize the (relative) cost of seeks
 - Unlike a regular filesystem, if the file is smaller than block-size, it doesn't occupy a full block's storage.



Distributed File Systems

- Blocks are good for:
 - Make easy to handle files larger than any disk in the cluster
 - Blocks from a file can be stored on different disks
 - Simplifying storage management (e.g., blocks have fix size) and other concerns
 - Replication, to provide fault tolerance and availability.
 - Usually 3 replicas (replication factor, user defined per file)
 - If a block becomes unavailable, a copy can be read from another location in a transparent way
 - Inaccessible blocks can be replicated from its alternative locations to other live machines to ensure replication factor
 - Multiply data transfer bandwidth



- HDFS: filesystem designed for storing very large files with streaming data access patterns, running on clusters of commodity hardware.
 - Very large files: 100+ MB, GB, TB

Data coherence

- Streaming data access: write-once, read-many pattern
 - Various analyses can be performed on the data.
 - Time to read the whole dataset, more important than the latency in reading the first record.
- Commodity hardware: expensive/highly-reliable hardware not required
 - The chance of node failure across a cluster is high
 - Designed to carry on working without a noticeable interruption to the user when this happens.



Not useful if: many small files, need low-latency data access, multiple writers and file modifications

- A HDFS cluster is composed of:
 - A single NameNode (master)
 - Manages the filesystem namespace
 - Maintains the directory system for hierarchical file organization (directories/subdirectories)
 - Stored persistently on the local disk
 - Maps filename to sequence of node-block ids
 - Not stored persistently: reconstructed from DataNodes when system starts.
 - Many DataNodes (workers)
 - Store and retrieve blocks when instructed so
 - Each block replica is represented by two files:
 - 1) One containing the data itself
 - 2) Another one with block's metadata (checksums, generation stamp, ...).
 - Report to the NameNode periodically with lists of blocks that they are storing.



Master-worker

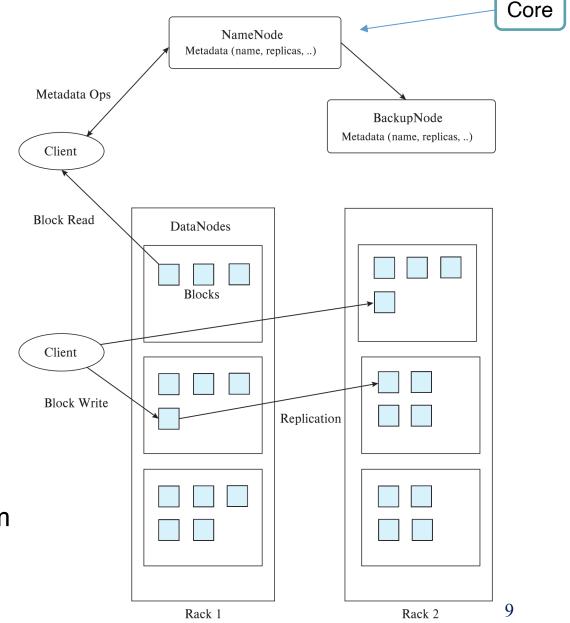
- Single Namespace
- Files are broken up into blocks

block size = 128MB

 Each block, replicated on multiple DataNodes

3 times

- NameNode stores:
 - List of blocks of a file
 - List of nodes with a copy of each block
- Access:
 - API
 - As a subdirectory of the local file system
 Connected to HDFS server



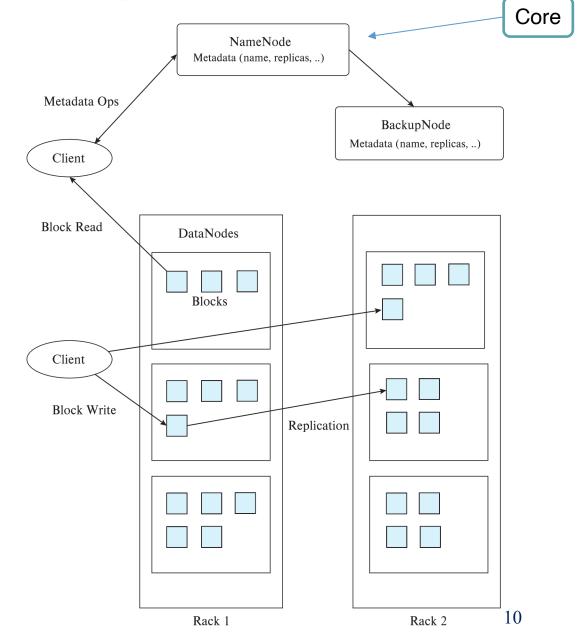


Read from a client:

- Retrieves from NameNode block and node ids
- Selects a DataNode for each block and retrieve data

Write from a client:

- Retrieves new block ids and a set of machines for each, assigned by NameNode
- Client sends block ids and data to the assigned DataNodes





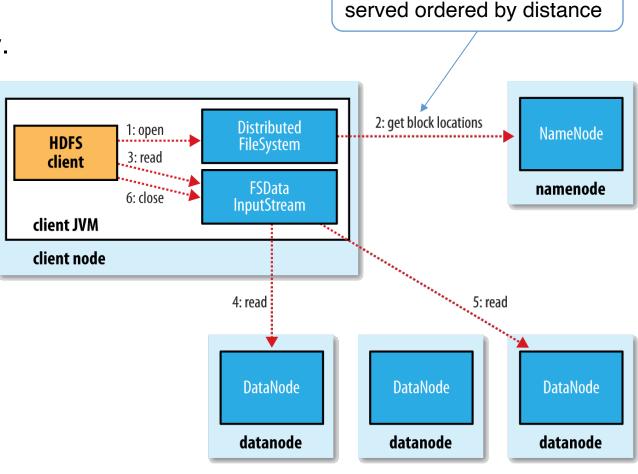
Reading a file from HDFS

Client contacts DataNodes directly.

 This design allows HDFS to scale to many concurrent clients.

 Data traffic is spread across all the cluster.

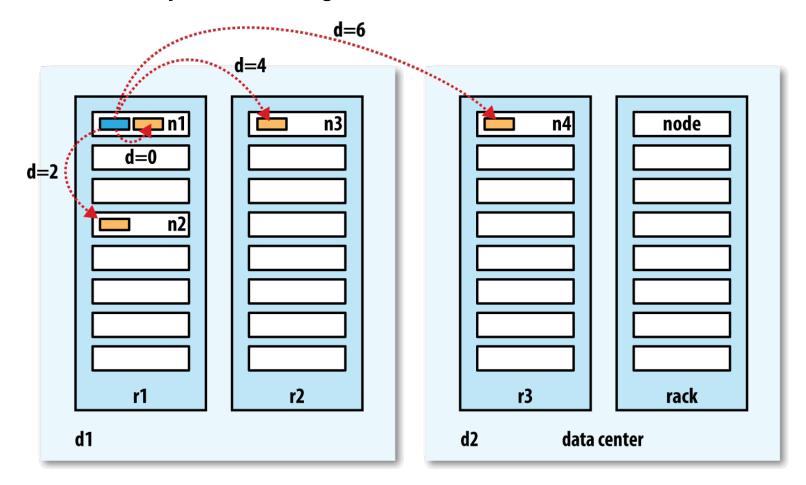
 NameNode merely serves block location requests.



DataNodes' addresses

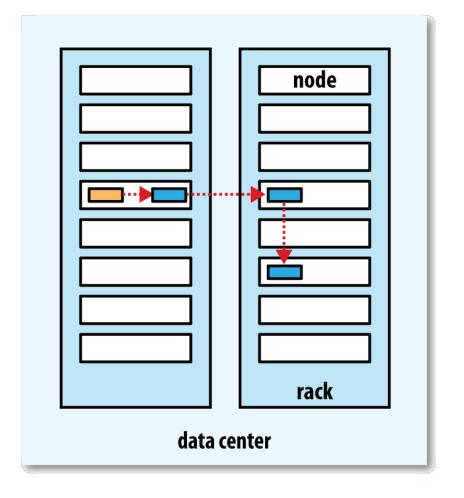


- Distance in Hadoop between nodes
 - Replication usually within a single data center





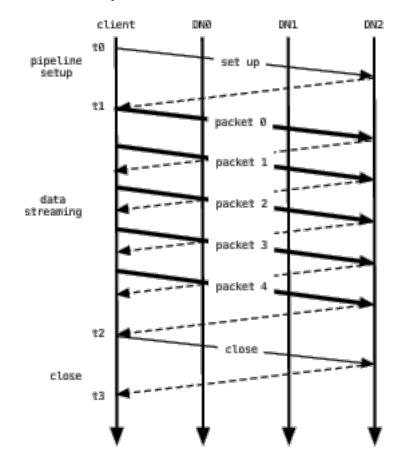
- Hadoop's replica location policy
 - 1. Client's node
 - Or randomly, if it hasn't
 - 2. Another rack
 - 3. Another node in 2nd rack



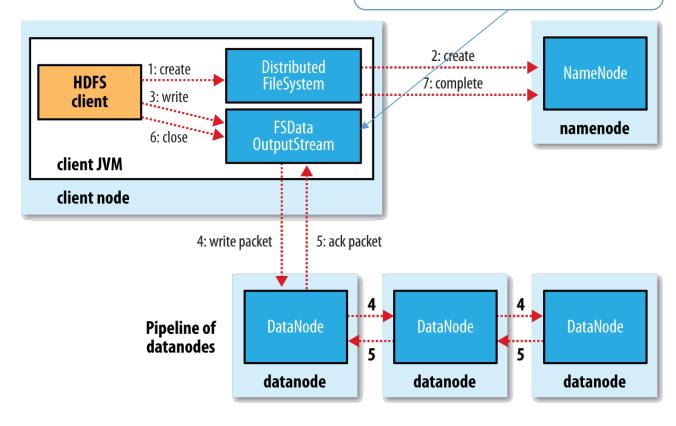


Writing a file in HDFS

Pipeline of DataNodes



DataNodes' addresses served ordered by NameNode





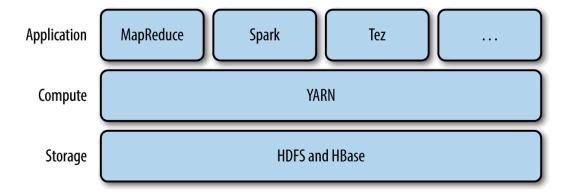
Limitations of HDFS

- Central master becomes bottleneck
 - Keep metadata in memory to avoid IO
 - Memory size limits no. files
 - Distributed master supported by Colossus file system
 - With smaller (1MB) block size
- File system directory overheads per file
 - Not appropriate for storing too many files
- Without consistency guarantees
 - File systems cache blocks locally
 - Ideal for write-once and append-only data



Hadoop

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MapReduce paradigm

- Old paradigm in parallel and functional programming
 - map(): solve a task for each element of a large list
 - reduce(): aggregates the solutions of the different elements

- E.g., word counts in a large collection of books
 - Solution:
 - Divide documents among workers
 - Each worker finds all words, and outputs (word, count) pairs
 - Partition (word, count) pairs across workers based on word
 - For each word at a worker, add up counts

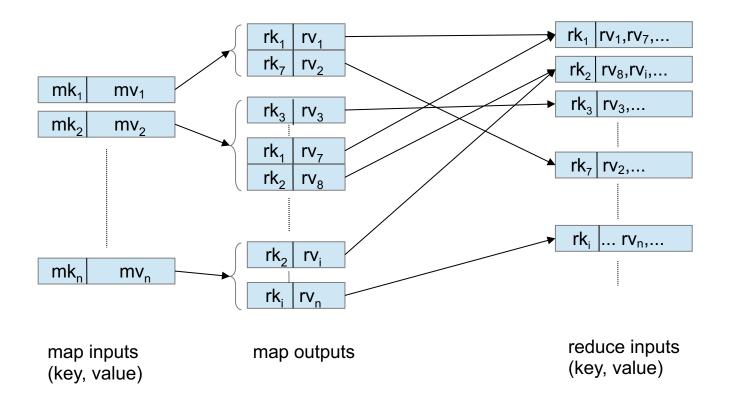
```
map(String document):
    for word in document:
        emit(word, 1);

reduce(String key, List value_list):
    int count = 0;
    for value in value_list:
        count += value;
    return(key, count);
```



MapReduce paradigm

Flow of keys and values in a map-reduce task





Workflow:

Put data in HDFS, run code, bring results back from HDFS



- Mapper and Reducer interfaces take...
 - input key, input value, output key and output value

```
public static class myMap extends Mapper < Long Writable, Text,
                                        Text, IntWritable> {
   private final static IntWritable one = new IntWritable(1);
   private Text word = new Text();
   public void map (LongWritable key, Text value, Context context)
                       throws IOException, InterruptedException {
      String line = value.toString();
      StringTokenizer tokenizer = new StringTokenizer(line);
      while (tokenizer.hasMoreTokens()) {
          word.set(tokenizer.nextToken());
          context.write(word, one);
```



- Mapper and Reducer interfaces take...
 - input key, input value, output key and output value

```
public static class myReduce extends Reducer < Text, IntWritable,
                                            Text, IntWritable> {
   public void reduce(Text key, Iterable<IntWritable> values,
                      Context context)
                      throws IOException, InterruptedException {
      int sum = 0;
      for (IntWritable val : values) {
          sum += val.get();
      context.write(key, new IntWritable(sum));
```



General class

```
public class WordCount {
    public static void main(String[] args) throws Exception {
        Configuration conf = new Configuration();
        Job job = new Job (conf, "wordcount");
        FileInputFormat.addInputPath(job, new Path(args[0]));
        FileOutputFormat.setOutputPath(job, new Path(args[1]));
        job.setMapperClass(myMap.class);
        job.setReducerClass(myReduce.class);
        job.setOutputKeyClass(Text.class);
        job.setOutputValueClass(IntWritable.class);
        System.exit(job.waitForCompletion(true)? 0:1);
```



MapReduce for a Projection operator

Explain the behavior of the mapper and the reducer



MapReduce for a Selection operator

Explain the behavior of the mapper and the reducer



MapReduce vs. Relational Queries

- MapReduce widely used for parallel processing
 - Allows procedural code in map and reduce functions
 - Allows data of any type
- Many computations...
 - of MapReduce cannot be expressed in SQL
 - are much easier to express in SQL
 - of relational operations can be expressed using MapReduce
 - select, project, join, aggregation, etc.
- SQL queries have been translated into MapReduce
 - Apache Hive SQL, Apache Pig Latin, Microsoft SCOPE



Hadoop Ecosystem



















Mapreduce (Data Processing)





HDFS (Hadoop Distributed File system)











