

## 1.4 Exercises

### 1.4.1 LEARNING ABOUT THE LANGUAGE

The first six exercises define some number-theoretic functions for you to program. Try to use recursion; it will be good practice for Chapter 2.

1.  $(\text{sigma } m \ n) = m + (m + 1) + \cdots + n$ .
2.  $(\text{exp } m \ n) = m^n$  ( $m, n \geq 0$ ).  $(\text{log } m \ n) =$  the least integer  $l$  such that  $m^{l+1} > n$  ( $m > 1, n > 0$ ).
3.  $(\text{choose } n \ k)$  is the number of ways of selecting  $k$  items from a collection of  $n$  items, without repetitions,  $n$  and  $k$  nonnegative integers. This quantity is called a *binomial coefficient*, and is notated  $\binom{n}{k}$ . It can be defined as  $\frac{n!}{k!(n-k)!}$ , but the following identities are more helpful computationally:  $\binom{n}{0} = 1$  ( $n \geq 0$ ),  $\binom{n}{n} = 1$  ( $n \geq 0$ ), and  $\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$  ( $n, k > 0$ ).
4.  $(\text{fib } m)$  is the  $m$ th Fibonacci number. The Fibonacci numbers are defined by the identities:  $(\text{fib } 0) = 0$ ,  $(\text{fib } 1) = 1$ , and for  $m > 1$ ,  $(\text{fib } m) = (\text{fib } m - 1) + (\text{fib } m - 2)$ .
5.  $(\text{prime } n) = \text{true}$  (1) if  $n$  is prime, false (0) otherwise.  $(\text{nthprime } n) =$  the  $n$ th prime number.  $(\text{sumprimes } n) =$  the sum of the first  $n$  primes.  $(\text{relprime } m \ n) = \text{true}$  if  $m$  and  $n$  are relatively prime (have no common divisors except 1), false otherwise.

## 2.7 Exercises

### 2.7.1 LEARNING ABOUT THE LANGUAGE

1. Code the following LISP functions:

- (a) `(count x l)` counts the number of occurrences of `x` at the top level of `l`, and `(countall x l)` counts the number of occurrences throughout `l`.

```
-> (count 'a '(1 b a (c a)))
```

```
1
```

```
-> (countall 'a '(1 b a (c a)))
```

```
2
```

- (b) (`reverse l`) returns a list containing the elements of `l` in reverse order. (*Hint: you will probably want to use the `append` function.*)

```
-> (reverse '(a b (c d) e))  
(e (c d) b a)
```

- (c) (`twist l`) reverses the top level of `l` and recursively twists all the items in `l`.

```
-> (twist '((a (b 5)) (c d) e))  
(e (d c) ((5 b) a))
```

- (d) (`flatten l`) constructs a list having the same atoms as `l` in the same order but in a flat list.

```
-> (flatten '((a b) ((c d) e)))  
(a b c d e)
```

- (e) (`sublist l1 l2`) and (`contig-sublist l1 l2`) determine whether the elements of `l1` are contained, and contiguously contained, respectively, in the same order in `l2`.

```
-> (sublist '(a b c) '(x a y b z c))  
T  
-> (contig-sublist '(a b c) '(x a y b z c))  
( )  
-> (contig-sublist '(a y) '(x a y b z c))  
T
```

## 2. Program the following set functions:

- (a) (`remove x s`) returns a set having the same elements as set `s` with element `x` removed.
- (b) (`subset s1 s2`) determines if `s1` is a subset of `s2`.
- (c) (`=set s1 s2`) determines if `s1` and `s2` are the same set.

## 3. Explain why `wrong-sum` (page 2.1.3) doesn't work. Then explain why it works when the last line is changed to:

```
(+ tmp (wrong-sum (cdr l))).
```

## 4. This set of questions concerns tree traversal.

- (a) Program post-order and in-order traversal for binary trees.
- (b) Modify the pre-order and level-order traversals so that, instead of the node labels being printed, they are placed in a list, which is returned as the value of the call.

- (c) Extend the pre-order and level-order traversals to trees of arbitrary degree, represented in such a way that binary trees have the same representation given them in the text. For example, '(a b c d)' represents a ternary tree whose root is labeled with a and which has three children, labeled b, c, and d, respectively, all leaf nodes. Note that there are two ways to represent leaf nodes: 'a and '(a) both represent a leaf node labeled a.

5. These problems relate to the relational data base example.

- (a) Program AND-SELECT, whose first two arguments are lists (of the same length) and which selects only those rows that have all the given values for the given attributes. For example, (AND-SELECT '(Crime Location) '(murder London) CRIMES) would select only rows representing murders in London.
- (b) Program OR-SELECT, whose first argument is an attribute name, whose second argument is a list of values, and which selects those rows which have any of the values for the given attribute.
- (c) Lift the restriction on UNION, INTER, and DIFF that the attributes of their two arguments must occur *in the same order*. These operations should check that their arguments have the same set of attributes and then choose an order of those attributes for the result.
- (d) REMOVE has the same arguments as PROJECT, but projects onto those attribute *not* in its first argument.

6. Modify the last version of eval in Section 2.4 as follows:

- (a) Add **begin** and **print**.
- (b) Add **set** and global variables.
- (c) Add local variables, as described in Exercise 12 of Chapter 1.

## 4.10 Exercises

### 4.10.1 LEARNING ABOUT THE LANGUAGE

1. Use `mapcar`, `mapc`, or `combine` to define the following functions (those not defined here were defined in Chapter 2):
  - (a) `cdr*` takes the `cdr`'s of each element of a list of lists:  
    `-> (cdr* '((a b c) (d e) (f)))`  
    `((b c) (e) ())`
  - (b) `max*` finds the maximum of a list of nonnegative integers.
  - (c) `append`.
  - (d) `addtoend` adds its first argument (an arbitrary S-expression) as the last element of its second argument (a list):  
    `-> (addtoend 'a '(b c d))`  
    `(b c d a)`

(e) `reverse` (use `addtoend`).

(f) `insertion-sort` (you may take `insert` as given).

(g) `(mkpairfn x)` is a function which, given a list, places `x` in front of each element:

```
-> ((mkpairfn 'a) '(() (b c) (d) ((e f))))  
((a) (a b c) (a d) (a (e f)))
```

2. `lex-order*` extends `lex-order` to apply to lists of differing length, as in normal alphabetical ordering:

```
-> (set alpha-order (lex-order* <))
```

```
-> (alpha-order '(4 15 7) '(4 15 7 5))
```

T

```
-> (alpha-order '(4 15 7) '(4 15 6 6))
```

()

To relate this to normal alphabetical ordering, just translate the numbers to letters: These two results say `DOG < DOGE` and `DOG  $\not<$  DOFF`. Program `lex-order*`.

3. An alternate representation for sets in SCHEME is as Boolean-valued functions, called “characteristic functions.” In particular, the null set is represented by a function that always returns *nil*, and the membership test is just application:

```
-> (set nullset (lambda (x) '()))
```

```
-> (set member? (lambda (x s) (s x)))
```

(a) Program `addelt`, `union`, `inter`, and `diff` using this new representation.

(b) Code the third approach to polymorphism (page 106) using this representation.

6. Extend the equivalence given for `letrec` in Section 4.6 to allow for multiple recursive definitions, i.e., to allow for the general form of `letrec`, as in this version of `eval`, which evaluates expressions having operations `+` and `*`, of arbitrary arity:

```
(set eval (lambda (e)
  (letrec
    ((ev (lambda (exp)
      (if (number? exp) exp
          (if (= (car exp) '+)
              (+/ (evlis (cdr exp)))
              (*/* (evlis (cdr exp)))))))
    (evlis (mapc ev)))
  (ev e))))
```

- (a) Give the extension on paper.
- (b) Translate `eval` by hand, and run it to test the translation.
- (c) Write the SCHEME functions `trans-let`, `trans-letrec`, and `trans-let*` that do the translations of each type of expression into SCHEME; for example:

```
-> (trans-letrec '(letrec ((f e1)) e2))
((lambda (f) (begin (set f e1) e2)) (quote ()))
```

- (d) Add `let` and `let*` to SCHEME `eval` (Section 4.5). (Note that adding `letrec` is much more difficult, because it requires implementing `set` in SCHEME `eval`.)